

Intel® PXA255 Processor

Developer's Manual

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Revision History

Date	Revision	Description
March 2003	-001	Initial release

Introduction 1

This document applies to the Intel® PXA255 Processor (PXA255 processor). It is an application specific standard product (ASSP) that provides industry-leading MIPS/mW performance for handheld computing applications. The processor is a highly integrated system on a chip and includes a high-performance low-power Intel® XScaleTM microarchitecture with a variety of different system peripherals.

The PXA255 processor is a 17x17mm 256-pin PBGA package configuration for high performance. The 17x17mm package has a 32-bit memory data bus and the full assortment of peripherals.

1.1 Intel® XScale™ Microarchitecture Features

The Intel® XScaleTM microarchitecture provides these features:

- ARM* Architecture Version 5TE ISA compliant.
 - ARM* Thumb Instruction Support
 - ARM* DSP Enhanced Instructions
- Low power consumption and high performance
- Intel® Media Processing Technology
 - Enhanced 16-bit Multiply
 - 40-bit Accumulator
- 32-KByte Instruction Cache
- 32-KByte Data Cache
- 2-KByte Mini Data Cache
- 2-KByte Mini Instruction Cache
- Instruction and Data Memory Management Units
- Branch Target Buffer
- Debug Capability via JTAG Port

Refer to the *Intel*® *XScale*™ *Microarchitecture for the Intel*® *PXA255 Processor User's Manual* for more details.

1.2 System Integration Features

The processor integrates the Intel® XScaleTM microarchitecture with this peripheral set:

- Memory Controller
- · Clock and Power Controllers
- Universal Serial Bus Client



- DMA Controller
- LCD Controller
- AC97
- I²S
- MultiMediaCard
- FIR Communication
- · Synchronous Serial Protocol Port
- I²C
- General Purpose I/O pins
- UARTs
- · Real-Time Clock
- OS Timers
- Pulse Width Modulation
- Interrupt Control

1.2.1 Memory Controller

The Memory Controller provides glueless control signals with programmable timing for a wide assortment of memory-chip types and organizations. It supports up to four SDRAM partitions; six static chip selects for SRAM, SSRAM, Flash, ROM, SROM, and companion chips; support for two PCMCIA or Compact Flash slots

1.2.2 Clocks and Power Controllers

The processor functional blocks are driven by clocks that are derived from a 3.6864-MHz crystal and an optional 32.768-kHz crystal.

The 3.6864-MHz crystal drives a core Phase Locked Loop (PLL) and a Peripheral PLL. The PLLs produce selected clock frequencies to run particular functional blocks.

The 32.768-kHz crystal provides an optional clock source that must be selected after a hard reset. This clock drives the Real Time Clock (RTC), Power Management Controller, and Interrupt Controller. The 32.768-kHz crystal is on a separate power island to provide an active clock while the processor is in sleep mode.

Power management controls the transition between the turbo/run, idle, and sleep operating modes.

1.2.3 Universal Serial Bus (USB) Client

The USB Client Module is based on the Universal Serial Bus Specification, Revision 1.1. It supports up to sixteen endpoints and it provides an internally generated 48-MHz clock. The USB Device Controller provides FIFOs with DMA access to or from memory.



1.2.4 DMA Controller (DMAC)

The DMAC provides sixteen prioritized channels to service transfer requests from internal peripherals and up to two data transfer requests from external companion chips. The DMAC is descriptor-based to allow command chaining and looping constructs.

The DMAC operates in Flow-Through Mode when performing peripheral-to-memory, memory-to-peripheral, and memory-to-memory transfers. The DMAC is compatible with peripherals that use word, half-word, or byte data sizes.

1.2.5 LCD Controller

The LCD Controller supports both passive (DSTN) and active (TFT) flat-panel displays with a maximum supported resolution of 640x480x16-bit/pixel. An internal 256 entry palette expands 1, 2, 4, or 8-bit encoded pixels. Non-encoded 16-bit pixels bypass the palette.

Two dedicated DMA channels allow the LCD Controller to support single- and dual-panel displays. Passive monochrome mode supports up to 256 gray-scale levels and passive color mode supports up to 64K colors. Active color mode supports up to 64K colors.

1.2.6 AC97 Controller

The AC97 Controller supports AC97 Revision 2.0 CODECs. These CODECs can operate at sample rates up to 48 KHz. The controller provides independent 16-bit channels for Stereo PCM In, Stereo PCM Out, Modem In, Modem Out, and mono Microphone In. Each channel includes a FIFO that supports DMA access to memory.

1.2.7 Inter-IC Sound (I²S) Controller

The I²S Controller provides a serial link to standard I²S CODECs for digital stereo sound. It supports both the Normal-I²S and MSB-Justified I²S formats, and provides four signals for connection to an I²S CODEC. I²S Controller signals are multiplexed with AC97 Controller pins. The controller includes FIFOs that support DMA access to memory.

1.2.8 Multimedia Card (MMC) Controller

The MMC Controller provides a serial interface to standard memory cards. The controller supports up to two cards in either MMC or SPI modes with serial data transfers up to 20 Mbps. The MMC controller has FIFOs that support DMA access to and from memory.

1.2.9 Fast Infrared (FIR) Communication Port

The FIR Communication Port is based on the 4-Mbps Infrared Data Association (IrDA) Specification. It operates at half-duplex and has FIFOs with DMA access to memory. The FIR Communication Port uses the STUART's transmit and receive pins to directly connect to external IrDA LED transceivers.



1.2.10 Synchronous Serial Protocol Controller (SSPC)

The SSP Port provides a full-duplex synchronous serial interface that operates at bit rates from 7.2 kHz to 1.84 MHz. It supports National Semiconductor's Microwire*, Texas Instruments' Synchronous Serial Protocol*, and Motorola's Serial Peripheral Interface*. The SSPC has FIFOs with DMA access to memory.

1.2.11 Inter-Integrated Circuit (I²C) Bus Interface Unit

The I²C Bus Interface Unit provides a general purpose 2-pin serial communication port. The interface uses one pin for data and address and a second pin for clocking.

1.2.12 **GPIO**

Each GPIO pin can be individually programmed as an output or an input. Inputs can cause interrupts on rising or falling edges. Primary GPIO pins are not shared with peripherals while secondary GPIO pins have alternate functions which can be mapped to the peripherals.

1.2.13 **UARTs**

The processor provides three Universal Asynchronous Receiver/Transmitters. Each UART can be used as a slow infrared (SIR) transmitter/receiver based on the Infrared Data Association Serial Infrared (SIR) Physical Layer Link Specification.

1.2.13.1 Full Function UART (FFUART)

The FFUART baud rate is programmable up to 230 Kbps. The FFUART provides a complete set of modem control pins: nCTS, nRTS, nDSR, nDTR, nRI, and nDCD. It has FIFOs with DMA access to or from memory.

1.2.13.2 Bluetooth UART (BTUART)

The BTUART baud rate is programmable up to 921 Kbps. The BTUART provides a partial set of modem control pins: nCTS and nRTS. Other modem control pins can be implemented via GPIOs. The BTUART has FIFOs with DMA access to or from memory.

1.2.13.3 Standard UART (STUART)

The STUART baud rate is programmable up to 230 Kbps. The STUART does not provide any modem control pins. The modem control pins can be implemented via GPIOs. The STUART has FIFOs with DMA access to or from memory.

The STUART's transmit and receive pins are multiplexed with the Fast Infrared Communication Port.



1.2.13.4 Hardware UART (HWUART)

The PXA255 processor has a UART with hardware flow control. The HWUART provides a partial set of modem control pins: nCTS and nRTS. These modem control pins provide full hardware flow control. Other modem control pins can be implemented via GPIOs. The HWUART baud rate is programmable up to 921.6 Kbps.

The HWUART's pins are multiplexed with the PCMCIA control pins. Because of this, these HWUART pins operate at the same voltage as the memory bus. Also, since the PCMCIA pin nPWE is used for variable-latency input/output (VLIO), while using these pins for the HWUART, VLIO is unavailable. The HWUART pins are also available over the BTUART pins. When operating over the BTUART pins, the HWUART pins operate at the I/O voltage.

1.2.14 Real-Time Clock (RTC)

The Real-Time Clock can be clocked from either crystal. A system with a 32.768-KHz crystal consumes less power during Sleep versus a system using only the 3.6864-MHz crystal. This crystal can be removed to save system cost. The RTC provides a constant frequency output with a programmable alarm register. This alarm register can be used to wake up the processor from Sleep mode.

1.2.15 OS Timers

The OS Timers can be used to provide a 3.68-MHz reference counter with four match registers. These registers can be configured to cause interrupts when equal to the reference counter. One match register can be used to cause a watchdog reset.

1.2.16 Pulse-Width Modulator (PWM)

The PWM has two independent outputs that can be programmed to drive two GPIOs. The frequency and duty cycle are independently programmable. For example, one GPIO can control LCD contrast and the other LCD brightness.

1.2.17 Interrupt Control

The Interrupt Controller directs the processor interrupts into the core's IRQ and FIQ inputs. The Mask Register enables or disables individual interrupt sources.

1.2.18 Network Synchronous Serial Protocol Port

The PXA255 processor has an SSP port optimized for connection to other network ASICs. This NSSP adds a Hi-Z function to TXD, the ability to control when Hi-Z occurs, and swapping the TXD/RXD pins.

This port is not multiplexed with other interfaces.

Introduction



intel® System Architecture

2.1 **Overview**

The PXA255 processor is an integrated system-on-a-chip microprocessor for high performance, low power portable handheld and handset devices. It incorporates the Intel® XScale™ microarchitecture with on-the-fly frequency scaling and sophisticated power management to provide industry leading MIPs/mW performance. The PXA255 processor is ARM* Architecture Version 5TE instruction set compliant (excluding floating point instructions) and follows the ARM* programmer's model.

The processor's memory interface supports a variety of memory types to allow design flexibility. Support for the connection of two companion chips permits a glueless interface to external devices. An integrated LCD display controller provides support for displays up to 640x480 pixels, and permits 1-, 2-, 4-, and 8-bit grayscale and 8- or 16-bit color pixels. A 256 entry/512 byte palette RAM provides flexibility in color mapping.

A set of serial devices and general system resources provide computational and connectivity capability for a variety of applications. Refer to Figure 2-1 for an overview of the microprocessor system architecture.



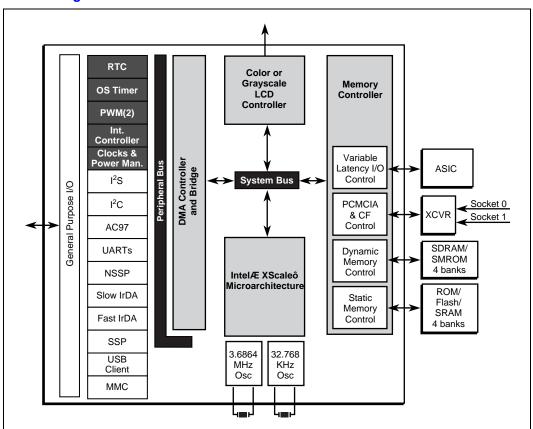


Figure 2-1. Block Diagram

2.2 Intel® XScale™ Microarchitecture Implementation Options

The processor incorporates the Intel® XScaleTM microarchitecture which is described in a separate document. This core contains implementation options which an Application Specific Standard Product (ASSP) may elect to implement or omit. This section describes those options.

Most of these options are specified within the coprocessor register space. The processor does not implement any coprocessor registers beyond those defined in the Intel® XScaleTM microarchitecture. The coprocessor registers which are ASSP specific, as stated in the *Intel® XScale*TM *Microarchitecture for the Intel® PXA255 Processor User's Manual*, order number 278793, are defined in the following sections.

2.2.1 Coprocessor 7 Register 4 - PSFS Bit

Bit 5 of this register is defined as the Power Source Fault Status bit or PSFS bit. This bit is set when either nVDD_FAULT or nBATT_FAULT pins are asserted and the Imprecise Data Abort Enable (IDAE) bit in the Power Manager Control Register (PMCR) is set.

This is a read-only register. Ignore reads from reserved bits.



Table 2-1.	CPU Core	e Fault Reg	gister Bit	Definitions
-------------------	----------	-------------	------------	--------------------

			C		roc egis		or 7 4	7					CPU Core Fault											System Architecture									
Bit	31	30	29	28	27	26	25	24	23	22	21	20 1	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
			Reserved														PSFS			Reserved													
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
		[31:	6]					_				Res Rea			fine	d.																	
		5					F	PSF\$	6			Pow 0 = 1 = Rea Clea	nV las nV 1.	DD t cle DD nly,	_FA eare _FA writ	ULT ed b ULT	Γ or y a Γ or nore	nB/ rese nB/ ed.	ATT et o	r the	€ CF \ULT	PÜ. Γpir	ı wa	as as	sser	ted							
		[4:0)]					_				Res Rea			fine	d.																	

2.2.2 Coprocessor 14 Registers 0-3 - Performance Monitoring

The processor does not define any performance monitoring features beyond those called out in the *Intel*® *XScale*TM *Microarchitecture for the Intel*® *PXA255 Processor User's Manual*, order number 278793. The interrupt generated by performance monitoring events is defined in Chapter 4, "System Integration Unit". The ASSP defined performance monitoring events (events 0x10 - 0x17), defined through the PMNC register are reserved for the processor.

2.2.3 Coprocessor 14 Register 6 and 7- Clock and Power Management

These registers allow software to use the clocking and power management modes. The valid operations are described in Table 3-23, "Coprocessor 14 Clock and Power Management Summary" on page 3-39.

2.2.4 Coprocessor 15 Register 0 - ID Register Definition

This register may be read by software to determine the device type and revision. The contents of this register for the Intel® PXA255 Processor is defined in the table below. Combined, this register must read as $0x6905\ 2X0R$ where R=0b0000 for the first stepping and then increments for subsequent steppings, and X is the revision of the Intel® XScaleTM microarchitecture present. Please see the Intel Developer Homepage at http://developer.intel.com for updates.

This is a read-only register.



Table 2-2. ID Bit Definitions

	CP15 Register 0													ID										CP15										
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14 13	12	1	1 10	9	8	7	6	5	4	3	2	1	0			
				Implementation	Trademark							Architecture	Architecture Version				Ogra	generation	(e Cor	Revision			Product	Number				Product	Revision				
Reset	0	1	1	0	1	0	0	1	0	0	0	0	0	1	0	1	0	0 1	0	0	0	0	1	0	0	0	0	0	0	0	0			
		[31:	:24]				plen Trac					Implementation trademark. 0x69 — Intel® Corporation.																						
		[23:	:16]			F	Arch Ve	itec rsic		;		ARM* Architecture version of the core. 0x05 – ARM* Architecture version 5TE																						
		[15:	:13]		Co	re G	ene	erati	ion		This field is updated when new sets of features are added to the core. This allows software that is dependent on core features to target a specific core generation: 0b001 – Intel® XScale™ core																							
		[12:	:10]			C	ore	Rev	visio	n		erra Co Ob(ata, re re 000 010	soft evisi – Fi – Ti	war ion: irst v	e w vers	orka sion sion	ach tir around of the of the	s, et core	c. e.		revi	sed	l. Di	ffere	ence	es m	ay i	nclu	ide				
		[9:	:4]			Pro	oduc	t N	umb	er					umb – P		255	proces	sor															
		[3:	:0]			Pro	duc	t Re	evis	ion		Ob010000 – PXA255 processor This field tracks the different steppings for each ASSP. Product Revision Ob0110 – A0 Stepping																						

Table 2-3. PXA255 Processor ID Values

Stepping	ARM ID	JTAG ID	
A0	0x6905_2D06	0x6926_4013	

2.2.5 Coprocessor 15 Register 1 - P-Bit

Bit 1 of this register is defined as the Page Table Memory Attribute bit or P-bit. It is not implemented in the processor and must be written as zero. Similarly, the P-bit in the page table descriptor in the MMU is not implemented and must be written to zero.



2.3 I/O Ordering

The processor uses queues that accept memory requests from the three internal masters: core, DMA Controller, and LCD Controller. Operations issued by a master are completed in the order they were received. Operations from one master may be interrupted by operations from another master. The processor does not provide a method to regulate the order of operations from different masters.

Loads and stores to internal addresses are generally completed more quickly than those issued to external addresses. The difference in completion time allows one operation to be received before another operation, but completed after the second operation.

In the following sequence, the store to the address in r4 is completed before the store to the address in r2 because the first store waits for memory in the queue while the second is not delayed.

str r1, [r2] ; store to external memory address [r2].

str r3, [r4]; store to internal (on-chip) memory address [r4].

If the two stores are control operations that must be completed in order, the recommended sequence is to insert a load to an unbuffered, uncached memory page followed by an operation that depends on data from the load:

str r1, [r2] ; first store issued

ldr r5, [r6] ; load from external unbuffered, uncached address ([r2] if possible)

mov r5, r5; nop stalls until r5 is loaded

str r3, [r4] ; second store completes in program order

2.4 Semaphores

The Swap (SWP) and Swap Byte (SWPB) instructions, as described in the ARM* Architecture reference, may be used for semaphore manipulation. No on-chip master or process can access a memory location between the load and store portion of a SWP or SWPB to the same location.

Note: Semaphore coherency may be interrupted because an external companion chip that uses the MBREQ/MBGNT handshake can take ownership of the bus during a locked sequence. To allow semaphore manipulation by external companion chips, the software must manage coherency.

2.5 Interrupts

The interrupt controller is described in detail in Section 4.2, "Interrupt Controller". All on-chip interrupts are enabled, masked, and routed to the core FIQ or IRQ. Each interrupt is enabled or disabled at the source through an interrupt mask bit. Generally, all interrupt bits in a unit are ORed together and present a single value to the interrupt controller.



Each interrupt goes through the Interrupt Controller Mask Register and then the Interrupt Controller Level Register directs the interrupt into either the IRQ or FIQ. If an interrupt is taken, the software may read the Interrupt Controller Pending Register to identify the source. After it identifies the interrupt source, software is responsible for servicing the interrupt and clearing it in the source unit before exiting the service routine.

Note: Clearing interrupts may take a delay. To allow the status bit to clear before returning from an interrupt service routine (ISR), clear the interrupt early in the routine.

2.6 Reset

The processor can be reset in any of three ways: Hardware, Watchdog, and GPIO resets. Each is described in more detail in Section 3.4, "Resets and Power Modes" on page 3-6.

- Hardware reset results from asserting the nRESET pin and forces all units into reset state.
- Watchdog reset results from a time-out in the OS Timer and may be used to recover from runaway code. Watchdog reset is disabled by default and must be enabled by software.
- GPIO reset is a "soft reset" that is less destructive than Hardware and Watchdog resets.

Each type of reset affects the state of the processor pins. Table 2-4 shows each pin's state after each type of reset.

Leaving Sleep Mode causes a Sleep Mode reset. Unlike other resets, Sleep Mode resets do not change the state of the pins.

The Reset Controller Status Register (RCSR) contains information on the type of reset, including Sleep Mode resets.

Table 2-4. Effect of Each Type of Reset on Internal Register State (Sheet 1 of 2)

Unit	Sleep Mode	GPIO Reset	Watchdog Reset	Hard Reset
Core	reset	reset	reset	reset
Memory Controller	reset	preserved	reset	reset
LCD Controller	reset	reset	reset	reset
DMA Controller	reset	reset	reset	reset
Full Function UART	reset	reset	reset	reset
Bluetooth UART	reset	reset	reset	reset
Standard UART	reset	reset	reset	reset
Hardware UART	reset	reset	reset	reset
I ² C	reset	reset	reset	reset
I ² S	reset	reset	reset	reset
AC97	reset	reset	reset	reset
USB	reset	reset	reset	reset
ICP	reset	reset	reset	reset
RTC	preserved	preserved	reset (except RTTR)	reset
OS Timer	reset	reset	reset	reset



Table 2-4. Effect of Each Type of Reset on Internal Register State (Sheet 2 of 2)

Unit	Sleep Mode	GPIO Reset	Watchdog Reset	Hard Reset
PWM	reset	reset	reset	reset
Interrupt Controller	reset	reset	reset	reset
GPIO	reset	reset	reset	reset
Power Manager	preserved	reset	reset	reset
SSP	reset	reset	reset	reset
NSSP	reset	reset	reset	reset
MMC	reset	reset	reset	reset
Clocks	preserved (except CP14)	preserved (except CP14)	reset (except OSCC)	reset

2.7 Internal Registers

All internal registers are mapped in physical memory space on 32-bit address boundaries. Use word access loads and stores to access internal registers. Internal register space must be mapped as non-cacheable.

Byte and halfword accesses to internal registers are not permitted and yield unpredictable results.

Register space where a register is not specifically mapped is defined as reserved space. Reading or writing reserved space causes unpredictable results.

The processor does not use all register bit locations. The unused bit locations are marked reserved and are allocated for future use. Write reserved bit locations as zeros. Ignore the values of these bits during reads because they are unpredictable.

2.8 Selecting Peripherals vs. General Purpose I/O

Most peripherals connect to the external pins through GPIOs. To use a peripheral connected through a GPIO, the software must first configure the GPIO so that the desired peripheral is connected to its pins. The default state of the pins is GPIO inputs.

To allocate a peripheral to a pin, disable the GPIO function for that pin, then map the peripheral function onto the pin by selecting the proper alternate function for the pin. Some GPIOs have multiple alternate functions. After a function is selected for a pin, all other functions are excluded. For this reason some peripherals are mapped to multiple GPIOs, as shown in Section 4.1.2, "GPIO Alternate Functions" on page 4-2. Multiple mapping does not mean multiple instances of a peripheral - only that the peripheral is connected to the pins in several ways.



2.9 Power on Reset and Boot Operation

Before the device that uses the processor is powered on, the system must assert nRESET and nTRST. To allow the internal clocks to stabilize, all power supplies must be stable for a specified period before nRESET or nTRST are deasserted. When nRESET is asserted, nRESET_OUT is driven active and can be used to reset other devices in the system. For additional information, see the *Intel*® *PXA255 Processor Design Guide*.

When the system deasserts nRESET and nTRST, the processor deasserts nRESET_OUT a specified time later and the device attempts to boot from physical address location 0x0000_0000.

The BOOT_SEL[2:0] pins are sampled when reset is deasserted and let the user specify the type and width of memory device from which the processor attempts to boot. The software can read the pins as described in Section 6.10.2, "Boot Time Defaults" on page 6-72.

2.10 Power Management

The processor offers a number of modes to manage power in the system. These range widely in level of power savings and level of functionality. The following modes are supported:

- Turbo Mode: low latency (nanoseconds) switch between two preprogrammed frequencies.
- Run Mode: normal full function mode.
- Idle Mode: core clocks are stopped resume through an interrupt.
- Sleep Mode: low power mode that does not save state but keeps I/Os powered. The RTC, Power Manager, and Clock modules are saved, except for Coprocessor 14.

Note: In low power modes, ensure that input pins are not floating and output pins are not driven by an external device that opposes how the processor is driving that pin. In either case, the system will draw excess current. Current draw that varies in sleep mode or varies greatly between parts is typically a sign of floating pins.

Section 3.4, "Resets and Power Modes" describes the modes in detail.

2.11 Pin List

Some of the processor pins can be connected to multiple signals. The signal connected to the pin is determined by the GPIO Alternate Function Select Registers (GAFRn m). Some signals can go to multiple pins. The signal must be routed to only one pin by using the GAFRn m registers. Because this is true, some pins are listed twice, once in each unit that can use the pin.

Table 2-5. Processor Pin Types

Туре	Function
IC	CMOS input
OC	CMOS output
OCZ	CMOS output, Hi-Z
ICOCZ	CMOS bidirectional, Hi-Z



Table 2-5. Processor Pin Types

Туре	Function
IA	Analog Input
OA	Analog output
IAOA	Analog bidirectional
SUP	Supply pin (either VCC or VSS)

Table 2-6 describes the PXA255 processor pins.

Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 1 of 9)

Pin Name	Type	Signal Descriptions	Reset State	Sleep State	
Memory Controller Pins					
MA[25:0]	ocz	Memory address bus. (output) Signals the address requested for memory accesses.	Driven Low	Driven Low	
MD[15:0]	ICOCZ	Memory data bus. (input/output) Lower 16 bits of the data bus.	Hi-Z	Driven Low	
MD[31:16]	ICOCZ	Memory data bus. (input/output) Used for 32-bit memories.	Hi-Z	Driven Low	
nOE	OCZ	Memory output enable. (output) Connect to the output enables of memory devices to control data bus drivers.	Driven High	Note [4]	
nWE	OCZ	Memory write enable. (output) Connect to the write enables of memory devices.	Driven High	Note [4]	
nSDCS[3:0]	OCZ	SDRAM CS for banks 3 through 0. (output) Connect to the chip select (CS) pins for SDRAM. For the PXA255 processor nSDCS0 can be Hi-Z, nSDCS1-3 cannot.	Driven High	Note [5]	
DQM[3:0]	OCZ	SDRAM DQM for data bytes 3 through 0. (output) Connect to the data output mask enables (DQM) for SDRAM.	Driven Low	Driven Low	
nSDRAS	OCZ	SDRAM RAS. (output) Connect to the row address strobe (RAS) pins for all banks of SDRAM.	Driven High	Driven High	
nSDCAS	OCZ	SDRAM CAS. (output) Connect to the column address strobe (CAS) pins for all banks of SDRAM.	Driven High	Driven High	
SDCKE[0]	ос	Synchronous Static Memory clock enable. (output) Connect to the CKE pins of SMROM. The memory controller provides control register bits for deassertion.	Driven Low	Driven Low	
SDCKE[1]	ос	SDRAM and/or Synchronous Static Memory clock enable. (output) Connect to the clock enable pins of SDRAM. It is deasserted during sleep. SDCKE[1] is always deasserted upon reset. The memory controller provides control register bits for deassertion.	Driven Low	Driven Low	
SDCLK[0]	ос	Synchronous Static Memory clock. (output) Connect to the clock (CLK) pins of SMROM. It is driven by either the internal memory controller clock, or the internal memory controller clock divided by 2. At reset, all clock pins are free running at the divide by 2 clock speed and may be turned off via free running control register bits in the memory controller. The memory controller also provides control register bits for clock division and deassertion of each SDCLK pin. SDCLK[0] control register assertion bit defaults to on if the boot-time static memory bank 0 is configured for SMROM.			



Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 2 of 9)

Pin Name	Type	Signal Descriptions	Reset State	Sleep State
SDCLK[1]	OCZ	SDRAM Clocks (output) Connect SDCLK[1] and	Driven Low	Driven Low
SDCLK[2]	ос	SDCLK[2] to the clock pins of SDRAM in bank pairs 0/1 and 2/3, respectively. They are driven by either the internal memory controller clock, or the internal memory controller clock divided by 2. At reset, all clock pins are free running at the divide by 2 clock speed and may be turned off via free running control register bits in the memory controller. The memory controller also provides control register bits for clock division and deassertion of each SDCLK pin. SDCLK[2:1] control register assertion bits are always deasserted upon reset.	Driven Low	Driven Low
nCS[5]/ GPIO[33]	ICOCZ			
nCS[4]/ GPIO[80]	ICOCZ	Statio chin calcate (output) Chin calcate to statio		
nCS[3]/ GPIO[79]	ICOCZ	Static chip selects. (output) Chip selects to static memory devices such as ROM and Flash. Individually programmable in the memory configuration registers.	Hi-Z - Note [1]	Note [4]
nCS[2]/ GPIO[78]	ICOCZ	nCS[5:0] can be used with variable latency I/O devices.		
nCS[1]/ GPIO[15]	ICOCZ			
nCS[0]	ICOCZ	Static chip select 0. (output) Chip select for the boot memory. nCS[0] is a dedicated pin.	Driven High	Note [4]
RD/nWR	OCZ	Read/Write for static interface. (output) Signals that the current transaction is a read or write.	Driven Low	Holds last state
RDY/ GPIO[18]	ICOCZ	Variable Latency I/O Ready pin. (input) Notifies the memory controller when an external bus device is ready to transfer data.	Hi-Z - Note [1]	Note [3]
L_DD[8]/ GPIO[66]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. Memory Controller alternate bus master request. (input) Allows an external device to request the system bus from the Memory Controller.	Hi-Z - Note [1]	Note [3]
L_DD[15]/ GPIO[73]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. Memory Controller grant. (output) Notifies an external device that it has been granted the system bus.	Hi-Z - Note [1]	Note [3]
MBGNT/ GP[13]	ICOCZ	Memory Controller grant. (output) Notifies an external device that it has been granted the system bus.	Hi-Z - Note [1]	Note [3]
MBREQ/ GP[14]	ICOCZ	Memory Controller alternate bus master request. (input) Allows an external device to request the system bus from the Memory Controller.	Hi-Z - Note [1]	Note [3]
PCMCIA/CF Co	ontrol Pins			·
nPOE/ GPIO[48]	ICOCZ	PCMCIA output enable. (output) Reads from PCMCIA memory and to PCMCIA attribute space.	Hi-Z - Note [1]	Note [5]
nPWE/ GPIO[49]	ICOCZ	PCMCIA write enable. (output) Performs writes to PCMCIA memory and to PCMCIA attribute space. Also used as the write enable signal for Variable Latency I/O.	Hi-Z - Note [1]	Note [5]



Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 3 of 9)

Pin Name	Type	Signal Descriptions	Reset State	Sleep State
nPIOW/ GPIO[51]	ICOCZ	PCMCIA I/O write. (output) Performs write transactions to PCMCIA I/O space.	Hi-Z - Note [1]	Note [5]
nPIOR/ GPIO[50]	ICOCZ	PCMCIA I/O read. (output) Performs read transactions from PCMCIA I/O space.	Hi-Z - Note [1]	Note [5]
nPCE[2]/ GPIO[53]	ICOCZ	PCMCIA card enable 2. (output) Selects a PCMCIA card. nPCE[2] enables the high byte lane and nPCE[1] enables the low byte lane. MMC clock. (output) Clock signal for the MMC Controller.	Hi-Z - Note [1]	Note [5]
nPCE[1]/ GPIO[52]	ICOCZ	PCMCIA card enable 1. (outputs) Selects a PCMCIA card. nPCE[2] enables the high byte lane and nPCE[1] enables the low byte lane.	Hi-Z - Note [1]	Note [5]
nIOIS16/ GPIO[57]	ICOCZ	IO Select 16. (input) Acknowledge from the PCMCIA card that the current address is a valid 16 bit wide I/O address.	Hi-Z - Note [1]	Note [5]
nPWAIT/ GPIO[56]	ICOCZ	PCMCIA wait. (input) Driven low by the PCMCIA card to extend the length of the transfers to/from the PXA255 processor.	Hi-Z - Note [1]	Note [5]
PSKTSEL/ GPIO[54]	ICOCZ	PCMCIA socket select. (output) Used by external steering logic to route control, address, and data signals to one of the two PCMCIA sockets. When PSKTSEL is low, socket zero is selected. When PSKTSEL is high, socket one is selected. Has the same timing as the address bus.	Hi-Z - Note [1]	Note [5]
nPREG/ GPIO[55]	ICOCZ	PCMCIA Register select. (output) Indicates that the target address on a memory transaction is attribute space. Has the same timing as the address bus.	Hi-Z - Note [1]	Note [5]
LCD Controller F	Pins			
L_DD(7:0)/ GPIO[65:58]	ICOCZ	LCD display data. (outputs) Transfers pixel information from the LCD Controller to the external LCD panel.	Hi-Z - Note [1]	Note [3]
L_DD[8]/ GPIO[66]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. Memory Controller alternate bus master request. (input) Allows an external device to request the system bus from the Memory Controller.	Hi-Z - Note [1]	Note [3]
L_DD[9]/ GPIO[67]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. MMC chip select 0. (output) Chip select 0 for the MMC Controller.	Hi-Z - Note [1]	Note [3]
L_DD[10]/ GPIO[68]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. MMC chip select 1. (output) Chip select 1 for the MMC Controller.	Hi-Z - Note [1]	Note [3]
L_DD[11]/ GPIO[69]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. MMC clock. (output) Clock for the MMC Controller.	Hi-Z - Note [1]	Note [3]
L_DD[12]/ GPIO[70]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. RTC clock. (output) Real time clock 1 Hz tick.	Hi-Z - Note [1]	Note [3]
				



Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 4 of 9)

Pin Name	Туре	Signal Descriptions	Reset State	Sleep State		
L_DD[13]/ GPIO[71]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. 3.6864 MHz clock. (output) Output from 3.6864 MHz oscillator.	Hi-Z - Note [1]	Note [3]		
L_DD[14]/ GPIO[72]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. 32 kHz clock. (output) Output from the 32 kHz oscillator.	Hi-Z - Note [1]	Note [3]		
L_DD[15]/ GPIO[73]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. Memory Controller grant. (output) Notifies an external device it has been granted the system bus.	Hi-Z - Note [1]	Note [3]		
L FCLK/ GPIO[74]	ICOCZ	LCD frame clock. (output) Indicates the start of a new frame. Also referred to as Vsync.	Hi-Z - Note [1]	Note [3]		
L LCLK/ GPIO[75]	ICOCZ	LCD line clock. (output) Indicates the start of a new line. Also referred to as Hsync.	Hi-Z - Note [1]	Note [3]		
L PCLK/ GPIO[76]	ICOCZ	LCD pixel clock. (output) Clocks valid pixel data into the LCD's line shift buffer.	Hi-Z - Note [1]	Note [3]		
L BIAS/ GPIO[77]	ICOCZ	AC bias drive. (output) Notifies the panel to change the polarity for some passive LCD panel. For TFT panels, this signal indicates valid pixel data.	Hi-Z - Note [1]	Note [3]		
Full Function UA	ART Pins					
FFRXD/ GPIO[34]	ICOCZ	Full Function UART Receive. (input) MMC chip select 0. (output) Chip select 0 for the MMC Controller.	Hi-Z - Note [1]	Note [3]		
FFTXD/ GPIO[39]	ICOCZ	Full Function UART Transmit. (output) MMC chip select 1. (output) Chip select 1 for the MMC Controller.	Hi-Z - Note [1]	Note [3]		
FFCTS/ GPIO[35]	ICOCZ	Full Function UART Clear-to-Send. (input)	Hi-Z - Note [1]	Note [3]		
FFDCD/ GPIO[36]	ICOCZ	Full Function UART Data-Carrier-Detect. (input)	Hi-Z - Note [1]	Note [3]		
FFDSR/ GPIO[37]	ICOCZ	Full Function UART Data-Set-Ready. (input)	Hi-Z - Note [1]	Note [3]		
FFRI/ GPIO[38]	ICOCZ	Full Function UART Ring Indicator. (input)	Hi-Z - Note [1]	Note [3]		
FFDTR/ GPIO[40]	ICOCZ	Full Function UART Data-Terminal-Ready. (output)	Hi-Z - Note [1]	Note [3]		
FFRTS/ GPIO[41]	ICOCZ	Full Function UART Request-to-Send. (output)	Hi-Z - Note [1]	Note [3]		
Bluetooth UART	Bluetooth UART Pins					
BTRXD/ GPIO[42]	ICOCZ	Bluetooth UART Receive. (input)	Hi-Z - Note [1]	Note [3]		
BTTXD/ GPIO[43]	ICOCZ	Bluetooth UART Transmit. (output)	Hi-Z - Note [1]	Note [3]		



Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 5 of 9)

Pin Name	Type	Signal Descriptions	Reset State	Sleep State
BTCTS/ GPIO[44]	ICOCZ	Bluetooth UART Clear-to-Send. (input)	Hi-Z - Note [1]	Note [3]
BTRTS/ GPIO[45]	ICOCZ	Bluetooth UART Data-Terminal-Ready. (output)	Hi-Z - Note [1]	Note [3]
Standard UART	and ICP F	rins		
IRRXD/ GPIO[46]	ICOCZ	IrDA receive signal. (input) Receive pin for the FIR function. Standard UART receive. (input)	Hi-Z - Note [1]	Note [3]
IRTXD/ GPIO[47]	ICOCZ	IrDA transmit signal. (output) Transmit pin for the Standard UART, SIR and FIR functions. Standard UART transmit. (output)	Hi-Z - Note [1]	Note [3]
HWUART Pins				
HWTXD/ GPIO[48]	ICOCZ	Hardware UART Transmit Data.	Pulled High Note [1]	Note [3]
HWRXD/ GPIO[49]	ICOCZ	Hardware UART Receive Data.	Pulled High Note [1]	Note [3]
HWCTS/ GPIO[50]	ICOCZ	Hardware UART Clear-To-Send.	Pulled High Note [1]	Note [3]
HWRTS/ GPIO[51]	ICOCZ	Hardware UART Request-to-Send.	Pulled High Note [1]	Note [3]
MMC Controller	Pins			
MMCMD	ICOCZ	Multimedia Card Command. (bidirectional)	Hi-Z	Hi-Z
MMDAT	ICOCZ	Multimedia Card Data. (bidirectional)	Hi-Z	Hi-Z
nPCE[2]/ GPIO[53]	ICOCZ	PCMCIA card enable 2. (outputs) Selects a PCMCIA card. Bit one enables the high byte lane and bit zero enables the low byte lane. MMC clock. (output) Clock signal for the MMC Controller.	Hi-Z - Note [1]	Note [5]
L_DD[9]/ GPIO[67]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. MMC chip select 0. (output) Chip select 0 for the MMC Controller.	Hi-Z - Note [1]	Note [3]
L_DD[10]/ GPIO[68]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. MMC chip select 1. (output) Chip select 1 for the MMC Controller.	Hi-Z - Note [1]	Note [3]
L_DD[11]/ GPIO[69]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. MMC clock. (output) Clock for the MMC Controller.	Hi-Z - Note [1]	Note [3]
FFRXD/ GPIO[34]	ICOCZ	Full Function UART Receive. (input) MMC chip select 0. (output) Chip select 0 for the MMC Controller.	Hi-Z - Note [1]	Note [3]
FFTXD/ GPIO[39]	ICOCZ	Full Function UART Transmit. (output) MMC chip select 1. (output) Chip select 1 for the MMC Controller.	Hi-Z - Note [1]	Note [3]



Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 6 of 9)

Pin Name	Туре	Signal Descriptions	Reset State	Sleep State
MMCCLK/ GP[6]	ICOCZ	MMC clock. (output) Clock signal for the MMC Controller.	Hi-Z - Note [1]	Note [3]
MMCCS0/ GP[8]	ICOCZ	MMC chip select 0. (output) Chip select 0 for the MMC Controller.	Hi-Z - Note [1]	Note [3]
MMCCS1/ GP[9]	ICOCZ	MMC chip select 1. (output) Chip select 1 for the MMC Controller.	Hi-Z - Note [1]	Note [3]
SSP Pins				
SSPSCLK/ GPIO[23]	ICOCZ	Synchronous Serial Port Clock. (output)	Hi-Z - Note [1]	Note [3]
SSPSFRM/ GPIO[24]	ICOCZ	Synchronous Serial Port Frame. (output)	Hi-Z - Note [1]	Note [3]
SSPTXD/ GPIO[25]	ICOCZ	Synchronous Serial Port Transmit. (output)	Hi-Z - Note [1]	Note [3]
SSPRXD/ GPIO[26]	ICOCZ	Synchronous Serial Port Receive. (input)	Hi-Z - Note [1]	Note [3]
SSPEXTCLK/ GPIO[27]	ICOCZ	Synchronous Serial Port External Clock. (input)	Hi-Z - Note [1]	Note [3]
Network SSP pi	ns			
NSSPSCLK/ GPIO[81]	ICOCZ	Network Synchronous Serial Port Clock.	Pulled High Note [1]	Note [3]
NSSPSFRM/ GPIO[82]	ICOCZ	Network Synchronous Serial Port Frame Signal.	Pulled High Note [1]	Note [3]
NSSPTXD/ GPIO[83]	ICOCZ	Network Synchronous Serial Port Transmit.	Pulled High Note [1]	Note [3]
NSSPRXD/ GPIO[84]	ICOCZ	Network Synchronous Serial Port Receive.	Pulled High Note [1]	Note [3]
USB Client Pins	;			
USB P	IAOAZ	USB Client Positive. (bidirectional)	Hi-Z	Hi-Z
USB N	IAOAZ	USB Client Negative pin. (bidirectional)	Hi-Z	Hi-Z
AC97 Controller	and I ² S C	Controller Pins		
BITCLK/ GPIO[28]	ICOCZ	AC97 Audio Port bit clock. (input) AC97 clock is generated by Codec 0 and fed into the PXA255 processor and Codec 1. AC97 Audio Port bit clock. (output) AC97 clock is generated by the PXA255 processor. I²S bit clock. (input) I2S clock is generated externally and fed into PXA255 processor. I²S bit clock. (output) I²S clock is generated by the PXA255 processor.	Hi-Z - Note [1]	Note [3]
SDATA_IN0/ GPIO[29]	ICOCZ	AC97 Audio Port data in. (input) Input line for Codec 0. I ² S data in. (input) Input line for the I ² S Controller.	Hi-Z - Note [1]	Note [3]
SDATA_IN1/ GPIO[32]	ICOCZ	AC97 Audio Port data in. (input) Input line for Codec 1. I ² S system clock. (output) System clock from I ² S Controller.	Hi-Z - Note [1]	Note [3]
	1	I	1	L



Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 7 of 9)

Pin Name	Type	Signal Descriptions	Reset State	Sleep State
SDATA_OUT/ GPIO[30]	ICOCZ	AC97 Audio Port data out. (output) Output from the PXA255 processor to Codecs 0 and 1. 128 data out. (output) Output line for the I2S Controller.	Hi-Z - Note [1]	Note [3]
SYNC/ GPIO[31]	ICOCZ	AC97 Audio Port sync signal. (output) Frame sync signal for the AC97 Controller. I ² S sync. (output) Frame sync signal for the I ² S Controller.	Hi-Z - Note [1]	Note [3]
nACRESET	ОС	AC97 Audio Port reset signal. (output)	Driven Low	Driven Low
I ² C Controller P	ins			
SCL	ICOCZ	I ² C clock. (bidirectional)	Hi-Z	Hi-Z
SDA	ICOCZ	I ² C data. (bidirectional).	Hi-Z	Hi-Z
PWM Pins	<u> </u>			
PWM[1:0]/ GPIO[17:16]	ICOCZ	Pulse Width Modulation channels 0 and 1. (outputs)	Hi-Z - Note [1]	Note [3]
DMA Pins	•			
DREQ[1:0]/ GPIO[19:20]	ICOCZ	DMA Request. (input) Notifies the DMA Controller that an external device requires a DMA transaction. DREQ[1] is GPIO[19]. DREQ[0] is GPIO[20].	Hi-Z - Note [1]	Note [3]
GPIO Pins	<u> </u>			
GPIO[1:0]	ICOCZ	General Purpose I/O. Wakeup sources on both rising and falling edges on nRESET.	Hi-Z - Note [1]	Note [3]
GPIO[14:2]	ICOCZ	General Purpose I/O. More wakeup sources for sleep mode.	Hi-Z - Note [1]	Note [3]
GPIO[22:21]	ICOCZ	General Purpose I/O. Additional General Purpose I/O pins.	Hi-Z - Note [1]	Note [3]
Crystal and Clo	ck Pins			
PXTAL	IA	3.6864 Mhz crystal input. No external caps are required.	Note [2]	Note [2]
PEXTAL	OA	3.6864 Mhz crystal output. No external caps are required.	Note [2]	Note [2]
TXTAL	IA	32 Khz crystal input. No external caps are required.	Note [2]	Note [2]
TEXTAL	OA	32 Khz crystal output. No external caps are required.	Note [2]	Note [2]
L_DD[12]/ GPIO[70]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. RTC clock. (output) Real time clock 1 Hz tick.	Hi-Z - Note [1]	Note [3]
L_DD[13]/ GPIO[71]	ICOCZ	LCD display data. (output) Transfers the pixel information from the LCD Controller to the external LCD panel. 3.6864 MHz clock. (output) Output from 3.6864 MHz oscillator.	Hi-Z - Note [1]	Note [3]
L_DD[14]/ GPIO[72]	ICOCZ	LCD display data. (output) Transfers pixel information from the LCD Controller to the external LCD panel. 32 kHz clock. (output) Output from the 32 kHz oscillator.	Hi-Z - Note [1]	Note [3]
48MHz/GP[7]	ICOCZ	48 MHz clock. (output) Peripheral clock output derived from the PLL. NOTE: This clock is only generated when the USB unit clock enable is set.	Hi-Z - Note [1]	Note [3]



Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 8 of 9)

Pin Name	Туре	Signal Descriptions	Reset State	Sleep State
RTCCLK/ GP[10]	ICOCZ	Real time clock. (output) 1 Hz output derived from the 32kHz or 3.6864MHz output.	Hi-Z - Note [1]	Note [3]
3.6MHz/GP[11]	ICOCZ	3.6864 MHz clock. (output) Output from 3.6864 MHz oscillator.	Hi-Z - Note [1]	Note [3]
32kHz/GP[12]	ICOCZ	32 kHz clock. (output) Output from the 32 kHz oscillator.	Hi-Z - Note [1]	Note [3]
Miscellaneous P	ins			
BOOT_SEL [2:0]	IC	Boot select pins. (input) Indicates type of boot device.	Input	Input
PWR_EN	ОС	Power Enable for the power supply. (output) When negated, it signals the power supply to remove power to the core because the system is entering sleep mode.	Driven High	Driven low while entering sleep mode. Driven high when sleep exit sequence begins.
nBATT_FAULT	IC	Main Battery Fault. (input) Signals that main battery is low or removed. Assertion causes PXA255 processor to enter sleep mode or force an Imprecise Data Exception, which cannot be masked. PXA255 processor will not recognize a walk-up event while this signal is asserted. Minimum assertion time for nBATT_FAULT is 1 ms.	Input	Input
nVDD_FAULT	IC	VDD Fault. (input) Signals that the main power source is going out of regulation. nVDD_FAULT causes the PXA255 processor to enter sleep mode or force an Imprecise Data Exception, which cannot be masked. nVDD_FAULT is ignored after a walk-up event until the power supply timer completes (approximately 10 ms). Minimum assertion time for nVDD_FAULT is 1 ms.	Input	Input
nRESET	IC	Hard reset. (input) Level sensitive input used to start the processor from a known address. Assertion causes the current instruction to terminate abnormally and causes a reset. When nRESET is driven high, the processor starts execution from address 0. nRESET must remain low until the power supply is stable and the internal 3.6864 MHz oscillator has stabilized.	Input	Input. Driving low during sleep will cause normal reset sequence and exit from sleep mode.
nRESET_OUT	ос	Reset Out. (output) Asserted when nRESET is asserted and deasserts after nRESET is deasserted but before the first instruction fetch. nRESET_OUT is also asserted for "soft" reset events: sleep, watchdog reset, or GPIO reset.	Driven low during any reset sequence - driven high prior to first fetch.	Driven Low
JTAG and Test F	Pins			
nTRST	IC	JTAG Test Interface Reset. Resets the JTAG/Debug port. If JTAG/Debug is used, drive nTRST from low to high either before or at the same time as nRESET. If JTAG is not used, nTRST must be either tied to nRESET or tied low.	Input	Input
TDI	IC	JTAG test data input. (input) Data from the JTAG controller is sent to the PXA255 processor using this pin. This pin has an internal pull-up resistor.	Input	Input
TDO	OCZ	JTAG test data output. (output) Data from the PXA255 processor is returned to the JTAG controller using this pin.	Hi-Z	Hi-Z
TMS	IC	JTAG test mode select. (input) Selects the test mode required from the JTAG controller. This pin has an internal pull-up resistor.	Input	Input



Table 2-6. Pin & Signal Descriptions for the PXA255 Processor (Sheet 9 of 9)

Pin Name	Туре	Signal Descriptions	Reset State	Sleep State
TCK	IC	JTAG test clock. (input) Clock for all transfers on the JTAG test interface.	Input	Input
TEST	IC	Test Mode. (input) Reserved. Must be grounded.	Input	Input
TESTCLK	IC	Test Clock. (input) Reserved. Must be grounded.	Input	Input
Power and Grou	nd Pins			
VCC	SUP	Positive supply for internal logic. Must be connected to the low voltage supply on the PCB.	Powered	Note [6]
VSS	SUP	Ground supply for internal logic. Must be connected to the common ground plane on the PCB.	Grounded	Grounded
PLL_VCC	SUP	Positive supply for PLLs and oscillators. Must be connected to the common low voltage supply.	Powered	Note [6]
PLL_VSS	SUP	Ground supply for the PLL. Must be connected to common ground plane on the PCB.	Grounded	Grounded
VCCQ	SUP	Positive supply for all CMOS I/O except memory bus and PCMCIA pins. Must be connected to the common 3.3v supply on the PCB.	Powered	Note [7]
VSSQ	SUP	Ground supply for all CMOS I/O except memory bus and PCMCIA pins. Must be connected to the common ground plane on the PCB.	Grounded	Grounded
VCCN	SUP	Positive supply for memory bus and PCMCIA pins. Must be connected to the common 3.3v or 2.5v supply on the PCB.	Powered	Note [7]
VSSN	SUP	Ground supply for memory bus and PCMCIA pins. Must be connected to the common ground plane on the PCB.	Grounded	Grounded

Table 2-7. Pin Description Notes (Sheet 1 of 2)

Note	Description				
[1]	GPIO Reset Operation: Configured as GPIO inputs by default after any reset. The input buffers for these pins are disabled to prevent current drain and the pins are pulled high with 10K to 60K internal resistors. The input paths must be enabled and the pullups turned off by clearing the Read Disable Hold (RDH) bit described in Section 3.5.7, "Power Manager Sleep Status Register (PSSR)" on page 3-29. Even though sleep mode sets the RDH bit, the pull-up resistors are not re-enabled by sleep mode.				
[2]	Crystal oscillator pins: These pins are used to connect the external crystals to the on-chip oscillators. Refer to Section 3.3.1, "32.768 kHz Oscillator" on page 3-4 and Section 3.3.2, "3.6864 MHz Oscillator" on page 3-4 for details on Sleep Mode operation.				
[3]	GPIO Sleep operation: During the transition into sleep mode, the state of these pins is determined by the corresponding PGSRn. See Section 3.5.10, "Power Manager GPIO Sleep State Registers (PGSR0, PGSR1, PGSR2)" and Section 4.1.3.2, "GPIO Pin Direction Registers (GPDR0, GPDR1, GPDR2)" on page 4-8. If selected as an input, this pin does not drive during sleep. If selected as an output, the value contained in the Sleep State Register is driven out onto the pin and held there while the PXA255 processor is in Sleep Mode. GPIOs configured as inputs after exiting sleep mode cannot be used until PSSR[RDH] is cleared.				
[4]	Static Memory Control Pins: During Sleep Mode, these pins can be programmed to either drive the value in the Sleep State Register or to be placed in Hi-Z. To select the Hi-Z state, software must set the FS bit in the Power Manager General Configuration Register. If PCFR[FS] is not set, then during the transition to sleep these pins function as described in [3], above. For nWE, nOE, and nCS[0], if PCFR[FS] is not set, they are driven high by the Memory Controller before entering sleep. If PCFR[FS] is set, these pins are placed in Hi-Z.				



Table 2-7. Pin Description Notes (Sheet 2 of 2)

Note	Description
[5]	PCMCIA Control Pins: During Sleep Mode: Can be programmed either to drive the value in the Sleep State Register or to be placed in Hi-Z. To select the Hi-Z state, software must set PCFR[FP]. If it is not set, then during the transition to sleep these pins function as described in [3], above.
[6]	During sleep, this supply may be driven low. This supply must never be high impedance.
[7]	Remains powered in sleep mode.

2.12 Memory Map

Figure 2-2 and Figure 2-3 show the full processor memory map.

Any unused register space from 0x4000_0000 to 0x4BFF_FFFF is reserved.

Note: Accessing reserved portions of the memory map will give unpredictable results.

The PCMCIA interface is divided into Socket 0 and Socket 1 space. These two sockets are each subdivided into I/O, memory and attribute space. Each socket is allocated 256 MB of memory space.



Figure 2-2. Memory Map (Part One) — From 0x8000_0000 to 0xFFFF FFFF

2. Memory Map (Fart One) — From		
0xFFFF_FFFF		
0xFC00_0000	Reserved (64 MB)	
0xF800_0000	Reserved (64 MB)	
0xF400_0000	Reserved (64 MB)	
0xF000_0000	Reserved (64 MB)	
0xEC00_0000	Reserved (64 MB)	
0xE800_0000	Reserved (64 MB)	
0xE400_0000	Reserved (64 MB)	
0xE000_0000	Reserved (64 MB)	
0xDC00_0000	Reserved (64 MB)	
0xD800_0000	Reserved (64 MB)	
0xD400_0000	Reserved (64 MB)	
0xD000_0000	Reserved (64 MB)	
0xCC00_0000	Reserved (64 MB)	
0xC800_0000	Reserved (64 MB)	
0xC400_0000	Reserved (64 MB)	
0xC000_0000	Reserved (64 MB)	
0xBC00_0000	Reserved (64 MB)	
0xB800_0000	Reserved (64 MB)	
0xB400_0000	Reserved (64 MB)	
0xB000_0000	Reserved (64 MB)	
0xAC00_0000	SDRAM Bank 3 (64 MB)	
0xA800_0000	SDRAM Bank 2 (64 MB)	
0xA400_0000	SDRAM Bank 1 (64 MB)	
0xA000_0000	SDRAM Bank 0 (64 MB)	
0x9C00 0000	Reserved (64 MB)	
0x9800_0000	Reserved (64 MB)	
0x9400_0000	Reserved (64 MB)	
0x9000_0000	Reserved (64 MB)	
0x8C00_0000	Reserved (64 MB)	
0x8800_0000	Reserved (64 MB)	
0x8400_0000	Reserved (64 MB)	
0x8000_0000	Reserved (64 MB)	



Figure 2-3. Memory Map (Part Two) — From 0x0000_0000 to 0x7FFF FFFF

0x7FFF FFFF		_
0x7C00_0000	Reserved (64 MB)	
0x7800_0000	Reserved (64 MB)	
0x7400_0000	Reserved (64 MB)	
0x7000_0000	Reserved (64 MB)	
0x6C00_0000	Reserved (64 MB)	
0x6800_0000	Reserved (64 MB)	
0x6400_0000	Reserved (64 MB)	
0x6000_0000	Reserved (64 MB)	
0x5C00_0000	Reserved (64 MB)	
0x5800_0000	Reserved (64 MB)	
0x5400_0000	Reserved (64 MB)	
0x5000_0000	Reserved (64 MB)	
0x4C00_0000	Reserved (64 MB)	
0x4800_0000	Memory Mapped registers (Memory Ctl)	
0x4400_0000	Memory Mapped registers (LCD)	
0x4000_0000	Memory Mapped registers (Peripherals)	
0x3C00_0000		
0x3800_0000	PCMCIA/CF- Slot 1 (256 MB)	
0x3400_0000		
0x3000_0000		
0x2C00_0000		
0x2800_0000	PGMGLA/GE, GLAQ/Q5Q4B)	
0x2400_0000	PCMCIA/CF - Slot 0 (256MB)	
0x2000_0000		
0x1C00_0000	Reserved (64 MB)	
0x1800_0000	Reserved (64 MB)	
0x1400_0000	Static Chip Select 5 (64 MB)	
0x1000_0000	Static Chip Select 4 (64 MB)	
0x0C00_0000	Static Chip Select 3 (64 MB)	
0x0800_0000	Static Chip Select 2 (64 MB)	
0x0400_0000	Static Chip Select 1 (64 MB)	
0x0000_0000	Static Chip Select 0 (64 MB)	



2.13 System Architecture Register Summary

Table 2-8. System Architecture Register Address Summary (Sheet 1 of 12)

Unit	Address	Register Symbol	Register Description
DMA Controller	0x4000_0000		
	0x4000_0000	DCSR0	DMA Control / Status Register for Channel 0
	0x4000_0004	DCSR1	DMA Control / Status Register for Channel 1
	0x4000_0008	DCSR2	DMA Control / Status Register for Channel 2
	0x4000_000C	DCSR3	DMA Control / Status Register for Channel 3
	0x4000_0010	DCSR4	DMA Control / Status Register for Channel 4
	0x4000_0014	DCSR5	DMA Control / Status Register for Channel 5
	0x4000_0018	DCSR6	DMA Control / Status Register for Channel 6
	0x4000_001C	DCSR7	DMA Control / Status Register for Channel 7
	0x4000_0020	DCSR8	DMA Control / Status Register for Channel 8
	0x4000_0024	DCSR9	DMA Control / Status Register for Channel 9
	0x4000_0028	DCSR10	DMA Control / Status Register for Channel 10
	0x4000_002C	DCSR11	DMA Control / Status Register for Channel 11
	0x4000_0030	DCSR12	DMA Control / Status Register for Channel 12
	0x4000_0034	DCSR13	DMA Control / Status Register for Channel 13
	0x4000_0038	DCSR14	DMA Control / Status Register for Channel 14
	0x4000_003C	DCSR15	DMA Control / Status Register for Channel 15
	0x4000_00f0	DINT	DMA Interrupt Register
	0x4000_0100	DRCMR0	Request to Channel Map Register for DREQ 0
	0x4000_0104	DRCMR1	Request to Channel Map Register for DREQ 1
	0x4000_0108	DRCMR2	Request to Channel Map Register for I2S receive Request
	0x4000_010C	DRCMR3	Request to Channel Map Register for I2S transmit Request
	0x4000_0110	DRCMR4	Request to Channel Map Register for BTUART receive Request
	0x4000_0114	DRCMR5	Request to Channel Map Register for BTUART transmit Request.
	0x4000_0118	DRCMR6	Request to Channel Map Register for FFUART receive Request
	0x4000_011C	DRCMR7	Request to Channel Map Register for FFUART transmit Request
	0x4000_0120	DRCMR8	Request to Channel Map Register for AC97 microphone Request
	0x4000_0124	DRCMR9	Request to Channel Map Register for AC97 modem receive Request
	0x4000_0128	DRCMR10	Request to Channel Map Register for AC97 modem transmit Request
	0x4000_012C	DRCMR11	Request to Channel Map Register for AC97 audio receive Request
	0x4000_0130	DRCMR12	Request to Channel Map Register for AC97 audio transmit Request
	0x4000_0134	DRCMR13	Request to Channel Map Register for SSP receive Request
	0x4000_0138	DRCMR14	Request to Channel Map Register for SSP transmit Request
	0x4000_013C	DRCMR15	Request to Channel Map Register for NSSP receive Request
	0x4000_0140	DRCMR16	Request to Channel Map Register for NSSP transmit Request
	0x4000_0144	DRCMR17	Request to Channel Map Register for ICP receive Request
	0x4000_0148	DRCMR18	Request to Channel Map Register for ICP transmit Request



Table 2-8. System Architecture Register Address Summary (Sheet 2 of 12)

Unit	Address	Register Symbol	Register Description
	0x4000_014C	DRCMR19	Request to Channel Map Register for STUART receive Request
	0x4000_0150	DRCMR20	Request to Channel Map Register for STUART transmit Request
	0x4000_0154	DRCMR21	Request to Channel Map Register for MMC receive Request
	0x4000_0158	DRCMR22	Request to Channel Map Register for MMC transmit Request
	0x4000_015C	DRCMR23	Reserved
	0x4000_0160	DRCMR24	Reserved
	0x4000_0164	DRCMR25	Request to Channel Map Register for USB endpoint 1 Request
	0x4000_0168	DRCMR26	Request to Channel Map Register for USB endpoint 2 Request
	0x4000_016C	DRCMR27	Request to Channel Map Register for USB endpoint 3 Request
	0x4000_0170	DRCMR28	Request to Channel Map Register for USB endpoint 4 Request
	0x4000_0174	DRCMR29	Request to Channel Map Register for HWUART receive Request
	0x4000_0178	DRCMR30	Request to Channel Map Register for USB endpoint 6 Request
	0x4000_017C	DRCMR31	Request to Channel Map Register for USB endpoint 7 Request
	0x4000_0180	DRCMR32	Request to Channel Map Register for USB endpoint 8 Request
	0x4000_0184	DRCMR33	Request to Channel Map Register for USB endpoint 9 Request
	0x4000_0188	DRCMR34	Request to Channel Map Register for HWUART transmit Request
	0x4000_018C	DRCMR35	Request to Channel Map Register for USB endpoint 11 Request
	0x4000_0190	DRCMR36	Request to Channel Map Register for USB endpoint 12 Request
	0x4000_0194	DRCMR37	Request to Channel Map Register for USB endpoint 13 Request
	0x4000_0198	DRCMR38	Request to Channel Map Register for USB endpoint 14 Request
	0x4000_019C	DRCMR39	Reserved
	0x4000_0200	DDADR0	DMA Descriptor Address Register Channel 0
	0x4000_0204	DSADR0	DMA Source Address Register Channel 0
	0x4000_0208	DTADR0	DMA Target Address Register Channel 0
	0x4000_020C	DCMD0	DMA Command Address Register Channel 0
	0x4000_0210	DDADR1	DMA Descriptor Address Register Channel 1
	0x4000_0214	DSADR1	DMA Source Address Register Channel 1
	0x4000_0218	DTADR1	DMA Target Address Register Channel 1
	0x4000_021C	DCMD1	DMA Command Address Register Channel 1
	0x4000_0220	DDADR2	DMA Descriptor Address Register Channel 2
	0x4000_0224	DSADR2	DMA Source Address Register Channel 2
	0x4000_0228	DTADR2	DMA Target Address Register Channel 2
	0x4000_022C	DCMD2	DMA Command Address Register Channel 2
	0x4000_0230	DDADR3	DMA Descriptor Address Register Channel 3
	0x4000_0234	DSADR3	DMA Source Address Register Channel 3
	0x4000_0238	DTADR3	DMA Target Address Register Channel 3
	0x4000_023C	DCMD3	DMA Command Address Register Channel 3
	0x4000_0240	DDADR4	DMA Descriptor Address Register Channel 4
	0x4000_0244	DSADR4	DMA Source Address Register Channel 4
	0x4000_0248	DTADR4	DMA Target Address Register Channel 4



Table 2-8. System Architecture Register Address Summary (Sheet 3 of 12)

Unit	Address	Register Symbol	Register Description
	0x4000_024C	DCMD4	DMA Command Address Register Channel 4
	0x4000_0250	DDADR5	DMA Descriptor Address Register Channel 5
	0x4000_0254	DSADR5	DMA Source Address Register Channel 5
	0x4000_0258	DTADR5	DMA Target Address Register Channel 5
	0x4000_025C	DCMD5	DMA Command Address Register Channel 5
	0x4000_0260	DDADR6	DMA Descriptor Address Register Channel 6
	0x4000_0264	DSADR6	DMA Source Address Register Channel 6
	0x4000_0268	DTADR6	DMA Target Address Register Channel 6
	0x4000_026C	DCMD6	DMA Command Address Register Channel 6
	0x4000_0270	DDADR7	DMA Descriptor Address Register Channel 7
	0x4000_0274	DSADR7	DMA Source Address Register Channel 7
	0x4000_0278	DTADR7	DMA Target Address Register Channel 7
	0x4000_027C	DCMD7	DMA Command Address Register Channel 7
	0x4000_0280	DDADR8	DMA Descriptor Address Register Channel 8
	0x4000_0284	DSADR8	DMA Source Address Register Channel 8
	0x4000_0288	DTADR8	DMA Target Address Register Channel 8
	0x4000_028C	DCMD8	DMA Command Address Register Channel 8
	0x4000_0290	DDADR9	DMA Descriptor Address Register Channel 9
	0x4000_0294	DSADR9	DMA Source Address Register Channel 9
	0x4000_0298	DTADR9	DMA Target Address Register Channel 9
	0x4000_029C	DCMD9	DMA Command Address Register Channel 9
	0x4000_02A0	DDADR10	DMA Descriptor Address Register Channel 10
	0x4000_02A4	DSADR10	DMA Source Address Register Channel 10
	0x4000_02A8	DTADR10	DMA Target Address Register Channel 10
	0x4000_02AC	DCMD10	DMA Command Address Register Channel 10
	0x4000_02B0	DDADR11	DMA Descriptor Address Register Channel 11
	0x4000_02B4	DSADR11	DMA Source Address Register Channel 11
	0x4000_02B8	DTADR11	DMA Target Address Register Channel 11
	0x4000_02BC	DCMD11	DMA Command Address Register Channel 11
	0x4000_02C0	DDADR12	DMA Descriptor Address Register Channel 12
	0x4000_02C4	DSADR12	DMA Source Address Register Channel 12
	0x4000_02C8	DTADR12	DMA Target Address Register Channel 12
	0x4000_02CC	DCMD12	DMA Command Address Register Channel 12
	0x4000_02D0	DDADR13	DMA Descriptor Address Register Channel 13
	0x4000_02D4	DSADR13	DMA Source Address Register Channel 13
	0x4000_02D8	DTADR13	DMA Target Address Register Channel 13
	0x4000_02DC	DCMD13	DMA Command Address Register Channel 13
	0x4000_02E0	DDADR14	DMA Descriptor Address Register Channel 14
	0x4000_02E4	DSADR14	DMA Source Address Register Channel 14
	0x4000_02E8	DTADR14	DMA Target Address Register Channel 14



Table 2-8. System Architecture Register Address Summary (Sheet 4 of 12)

Unit	Address	Register Symbol	Register Description
	0x4000_02EC	DCMD14	DMA Command Address Register Channel 14
	0x4000_02F0	DDADR15	DMA Descriptor Address Register Channel 15
	0x4000_02F4	DSADR15	DMA Source Address Register Channel 15
	0x4000_02F8	DTADR15	DMA Target Address Register Channel 15
	0x4000_02FC	DCMD15	DMA Command Address Register Channel 15
Full Function UART	0x4010_0000		
	0x4010_0000	FFRBR	Receive Buffer Register (read only)
	0x4010_0000	FFTHR	Transmit Holding Register (write only)
	0x4010_0004	FFIER	Interrupt Enable Register (read/write)
	0x4010_0008	FFIIR	Interrupt ID Register (read only)
	0x4010_0008	FFFCR	FIFO Control Register (write only)
	0x4010_000C	FFLCR	Line Control Register (read/write)
	0x4010_0010	FFMCR	Modem Control Register (read/write)
	0x4010_0014	FFLSR	Line Status Register (read only)
	0x4010_0018	FFMSR	Modem Status Register (read only)
	0x4010_001C	FFSPR	Scratch Pad Register (read/write)
	0x4010_0020	FFISR	Infrared Selection Register (read/write)
	0x4010_0000	FFDLL	Divisor Latch Low Register (DLAB = 1) (read/write)
	0x4010_0004	FFDLH	Divisor Latch High Register (DLAB = 1) (read/write)
Bluetooth UART	0x4020_0000		
	0x4020_0000	BTRBR	Receive Buffer Register (read only)
	0x4020_0000	BTTHR	Transmit Holding Register (write only)
	0x4020_0004	BTIER	Interrupt Enable Register (read/write)
	0x4020_0008	BTIIR	Interrupt ID Register (read only)
	0x4020_0008	BTFCR	FIFO Control Register (write only)
	0x4020_000C	BTLCR	Line Control Register (read/write)
	0x4020_0010	BTMCR	Modem Control Register (read/write)
	0x4020_0014	BTLSR	Line Status Register (read only)
	0x4020_0018	BTMSR	Modem Status Register (read only)
	0x4020_001C	BTSPR	Scratch Pad Register (read/write)
	0x4020_0020	BTISR	Infrared Selection Register (read/write)
	0x4020_0000	BTDLL	Divisor Latch Low Register (DLAB = 1) (read/write)
	0x4020_0004	BTDLH	Divisor Latch High Register (DLAB = 1) (read/write)
I2C	0x4030_0000		
	0x4030 1680	IBMR	I2C Bus Monitor Register - IBMR
	0x4030 1688	IDBR	I2C Data Buffer Register - IDBR
	0x4030 1690	ICR	I2C Control Register - ICR
	0x4030 1698	ISR	I2C Status Register - ISR
	0x4030 16A0	ISAR	I2C Slave Address Register - ISAR



Table 2-8. System Architecture Register Address Summary (Sheet 5 of 12)

4040_0000 4040_0000 4040_0004 4040_0008 4040_0010 4040_0014 4040_0018 4040_001C through 4040_0060 4040_0060 4040_0060 4040_007C 4040_0080	SACRO SACR1 SASRO SAIMR SAICR SADIV	Global Control Register Serial Audio I ² S/MSB-Justified Control Register Reserved Serial Audio I ² S/MSB-Justified Interface and FIFO Status Register Reserved Serial Audio Interrupt Mask Register Serial Audio Interrupt Clear Register Reserved Audio Clock Divider Register.
4040_0004 4040_0008 4040_000C 4040_0010 4040_0014 4040_0018 4040_001C through 4040_005C 4040_0060 4040_0064 through 4040_007C	SACR1 — SASR0 — SAIMR SAICR —	Serial Audio I ² S/MSB-Justified Control Register Reserved Serial Audio I ² S/MSB-Justified Interface and FIFO Status Register Reserved Serial Audio Interrupt Mask Register Serial Audio Interrupt Clear Register Reserved
4040_0008 4040_000C 4040_0010 4040_0014 4040_0018 4040_001C through 4040_005C 4040_0060 4040_0064 through 4040_007C	SASRO SAIMR SAICR	Reserved Serial Audio I ² S/MSB-Justified Interface and FIFO Status Register Reserved Serial Audio Interrupt Mask Register Serial Audio Interrupt Clear Register Reserved
4040_000C 4040_0010 4040_0014 4040_0018 4040_001C through 4040_005C 4040_0060 4040_0064 through 4040_007C	— SAIMR SAICR —	Serial Audio I ² S/MSB-Justified Interface and FIFO Status Register Reserved Serial Audio Interrupt Mask Register Serial Audio Interrupt Clear Register Reserved
4040_0010 4040_0014 4040_0018 4040_001C through 4040_005C 4040_0060 4040_0064 through 4040_007C	— SAIMR SAICR —	Reserved Serial Audio Interrupt Mask Register Serial Audio Interrupt Clear Register Reserved
4040_0014 4040_0018 4040_001C through 4040_005C 4040_0060 4040_0064 through 4040_007C	SAICR —	Serial Audio Interrupt Mask Register Serial Audio Interrupt Clear Register Reserved
4040_0018 4040_001C through 4040_005C 4040_0060 4040_0064 through 4040_007C	SAICR —	Serial Audio Interrupt Clear Register Reserved
4040_001C through 4040_005C 4040_0060 4040_0064 through 4040_007C	-	Reserved
through 4040_005C 4040_0060 4040_0064 through 4040_007C	— SADIV —	
4040_005C 4040_0060 4040_0064 through 4040_007C	SADIV	
4040_0060 4040_0064 through 4040_007C	SADIV	Audio Clock Divider Register.
4040_0064 through 4040_007C	SADIV —	Audio Clock Divider Register.
through 4040_007C	_	
4040_007C	_	
_		Reserved
4040 0080		
	SADR	Serial Audio Data Register (TX and RX FIFO access Register).
4050_0000		
(4050_0000	POCR	PCM Out Control Register
(4050_0004	PICR	PCM In Control Register
4050_0008	MCCR	Mic In Control Register
4050_000C	GCR	Global Control Register
4050_0010	POSR	PCM Out Status Register
4050_0014	PISR	PCM In Status Register
4050_0018	MCSR	Mic In Status Register
4050_001C	GSR	Global Status Register
(4050_0020	CAR	CODEC Access Register
4050_0024		
through	_	Reserved
4050_003C		
4050_0040	PCDR	PCM FIFO Data Register
4050_0044		
ŭ	_	Reserved
	MODE	Min in EIEO Data Danistan
_	MCDR	Mic-in FIFO Data Register
_		Poconyad
Ŭ	_	Reserved
	MOCR	Modem Out Control Register
		Reserved
_	MICR	Modern In Control Register
_		Reserved
_	MOSR	Modem Out Status Register
	4040_0080 4050_0000 4050_0000 4050_0004 4050_0008 4050_0010 4050_0014 4050_0018 4050_001C 4050_0020 4050_0024 through 4050_003C 4050_0040	4040_0080 SADR 4050_0000 POCR 4050_0004 PICR 4050_0008 MCCR 4050_000C GCR 4050_0010 POSR 4050_0014 PISR 4050_0016 GSR 4050_001C GSR 4050_0020 CAR 4050_0020 CAR 4050_003C PCDR 4050_003C PCDR 4050_0044 Through — 4050_005C MCDR 4050_0060 MCDR 4050_006C MCDR 4050_00100 MOCR 4050_0100 MCCR 4050_0100 MICR 4050_0100 MICR 4050_0100 MICR 4050_0100 MICR



Table 2-8. System Architecture Register Address Summary (Sheet 6 of 12)

Unit	Address	Register Symbol	Register Description
	0x4050_0114	_	Reserved
	0x4050_0118	MISR	Modem In Status Register
	0x4050_011C		
	through	_	Reserved
	0x4050_013C		
	0x4050_0140	MODR	Modem FIFO Data Register
	0x4050_0144		
	through 0x4050_01FC	_	Reserved
	0x4050_0200 through	_	Primary Audio CODEC registers
	0x4050_02FC		I filling Additional Control of the
	0x4050_0300		
	through	_	Secondary Audio CODEC registers
	0x4050_03FC		
	0x4050_0400		
	through	_	Primary Modem CODEC registers
	0x4050_04FC		
	0x4050_0500		
	through	_	Secondary Modem CODEC registers
LIDC	0x4050_05FC		
UDC	0x4060_0000	LIDCCD	LIDC Control Posicios
	0x4060_0000	UDCCR	UDC Control Register
	0x4060_0010	UDCCS0	UDC Endpoint 0 Control/Status Register
	0x4060_0014	UDCCS1	UDC Endpoint 1 (IN) Control/Status Register
	0x4060_0018	UDCCS2	UDC Endpoint 2 (OUT) Control/Status Register
	0x4060_001C	UDCCS3	UDC Endpoint 3 (IN) Control/Status Register
	0x4060_0020	UDCCS4	UDC Endpoint 4 (OUT) Control/Status Register
	0x4060_0024	UDCCS5	UDC Endpoint 5 (Interrupt) Control/Status Register
	0x4060_0028	UDCCS6	UDC Endpoint 6 (IN) Control/Status Register
	0x4060_002C	UDCCS7	UDC Endpoint 7 (OUT) Control/Status Register
	0x4060_0030	UDCCS8	UDC Endpoint 8 (IN) Control/Status Register
	0x4060_0034	UDCCS9	UDC Endpoint 9 (OUT) Control/Status Register
	0x4060_0038	UDCCS10	UDC Endpoint 10 (Interrupt) Control/Status Register
	0x4060_003C	UDCCS11	UDC Endpoint 11 (IN) Control/Status Register
	0x4060_0040	UDCCS12	UDC Endpoint 12 (OUT) Control/Status Register
	0x4060_0044	UDCCS13	UDC Endpoint 13 (IN) Control/Status Register
	0x4060_0048	UDCCS14	UDC Endpoint 14 (OUT) Control/Status Register
	0x4060_004C	UDCCS15	UDC Endpoint 15 (Interrupt) Control/Status Register
	0x4060_0060	UFNRH	UDC Frame Number Register High
	0x4060_0064	UFNRL	UDC Frame Number Register Low
_	0x4060_0068	UBCR2	UDC Byte Count Register 2



Table 2-8. System Architecture Register Address Summary (Sheet 7 of 12)

Unit	Address	Register Symbol	Register Description
	0x4060_006C	UBCR4	UDC Byte Count Register 4
	0x4060_0070	UBCR7	UDC Byte Count Register 7
	0x4060_0074	UBCR9	UDC Byte Count Register 9
	0x4060_0078	UBCR12	UDC Byte Count Register 12
	0x4060_007C	UBCR14	UDC Byte Count Register 14
	0x4060_0080	UDDR0	UDC Endpoint 0 Data Register
	0x4060_0100	UDDR1	UDC Endpoint 1 Data Register
	0x4060_0180	UDDR2	UDC Endpoint 2 Data Register
	0x4060_0200	UDDR3	UDC Endpoint 3 Data Register
	0x4060_0400	UDDR4	UDC Endpoint 4 Data Register
	0x4060_00A0	UDDR5	UDC Endpoint 5 Data Register
	0x4060_0600	UDDR6	UDC Endpoint 6 Data Register
	0x4060_0680	UDDR7	UDC Endpoint 7 Data Register
	0x4060_0700	UDDR8	UDC Endpoint 8 Data Register
	0x4060_0900	UDDR9	UDC Endpoint 9 Data Register
	0x4060_00C0	UDDR10	UDC Endpoint 10 Data Register
	0x4060_0B00	UDDR11	UDC Endpoint 11 Data Register
	0x4060_0B80	UDDR12	UDC Endpoint 12 Data Register
	0x4060_0C00	UDDR13	UDC Endpoint 13 Data Register
	0x4060_0E00	UDDR14	UDC Endpoint 14 Data Register
	0x4060_00E0	UDDR15	UDC Endpoint 15 Data Register
	0x4060_0050	UICR0	UDC Interrupt Control Register 0
	0x4060_0054	UICR1	UDC Interrupt Control Register 1
	0x4060_0058	USIR0	UDC Status Interrupt Register 0
	0x4060_005C	USIR1	UDC Status Interrupt Register 1
Standard UART	0x4070_0000		
	0x4070_0000	STRBR	Receive Buffer Register (read only)
	0x4070_0000	STTHR	Transmit Holding Register (write only)
	0x4070_0004	STIER	Interrupt Enable Register (read/write)
	0x4070_0008	STIIR	Interrupt ID Register (read only)
	0x4070_0008	STFCR	FIFO Control Register (write only)
	0x4070_000C	STLCR	Line Control Register (read/write)
	0x4070_0010	STMCR	Modem Control Register (read/write)
	0x4070_0014	STLSR	Line Status Register (read only)
	0x4070_0018	STMSR	Reserved
	0x4070_001C	STSPR	Scratch Pad Register (read/write)
	0x4070_0020	STISR	Infrared Selection Register (read/write)
	0x4070_0000	STDLL	Divisor Latch Low Register (DLAB = 1) (read/write)
	0x4070_0004	STDLH	Divisor Latch High Register (DLAB = 1) (read/write)



Table 2-8. System Architecture Register Address Summary (Sheet 8 of 12)

Unit	Address	Register Symbol	Register Description
ICP	0x4080_0000	Trogressor Gymbol	
ICF	0x4080_0000	ICCR0	ICP Control Register 0
	0x4080_0000 0x4080_0004	ICCR0	ICP Control Register 1
		ICCR1	ICP Control Register 2
	0x4080_0008	ICCR2	
	0x4080_000C	ICDR	ICP Data Register
	0x4080_0010		Reserved
	0x4080_0014	ICSR0	ICP Status Register 0
DTO	0x4080_0018	ICSR1	ICP Status Register 1
RTC	0x4090_0000		
	0x4090_0000	RCNR	RTC Count Register
	0x4090_0004	RTAR	RTC Alarm Register
	0x4090_0008	RTSR	RTC Status Register
	0x4090_000C	RTTR	RTC Timer Trim Register
OS Timer	0x40A0_0000		
	0x40A0_0000	OSMR<0>	
	0x40A0_0004	OSMR<1>	OS Timer Match registers<3:0>
	0x40A0_0008	OSMR<2>	OS Timer Mator registers<5.02
	0x40A0_000C	OSMR<3>	
	0x40A0_0010	OSCR	OS Timer Counter Register
	0x40A0_0014	OSSR	OS Timer Status Register
	0x40A0_0018	OWER	OS Timer Watchdog Enable Register
	0x40A0_001C	OIER	OS Timer Interrupt Enable Register
PWM 0	0x40B0_0000		
	0x40B0_0000	PWM_CTRL0	PWM 0 Control Register
	0x40B0_0004	PWM_PWDUTY0	PWM 0 Duty Cycle Register
	0x40B0_0008	PWM_PERVAL0	PWM 0 Period Control Register
PWM 1	0x40C0_0000		
	0x40C0_0000	PWM_CTRL1	PWM 1Control Register
	0x40C0_0004	PWM_PWDUTY1	PWM 1 Duty Cycle Register
	0x40C0_0008	PWM_PERVAL1	PWM 1 Period Control Register
Interrupt Control	0x40D0_0000		
	0x40D0_0000	ICIP	Interrupt Controller IRQ Pending Register
	0x40D0_0004	ICMR	Interrupt Controller Mask Register
	0x40D0_0008	ICLR	Interrupt Controller Level Register
	0x40D0_000C	ICFP	Interrupt Controller FIQ Pending Register
	0x40D0_0010	ICPR	Interrupt Controller Pending Register
	0x40D0_0014	ICCR	Interrupt Controller Control Register
GPIO	0x40E0_0000		
	0x40E0_0000	GPLR0	GPIO Pin-Level Register GPIO<31:0>
	0x40E0_0004	GPLR1	GPIO Pin-Level Register GPIO<63:32>
			<u> </u>



Table 2-8. System Architecture Register Address Summary (Sheet 9 of 12)

Unit	Address	Register Symbol	Register Description
	0x40E0_0008	GPLR2	GPIO Pin-Level Register GPIO<80:64>
	0x40E0_000C	GPDR0	GPIO Pin Direction Register GPIO<31:0>
	0x40E0_0010	GPDR1	GPIO Pin Direction Register GPIO<63:32>
	0x40E0_0014	GPDR2	GPIO Pin Direction Register GPIO<80:64>
	0x40E0_0018	GPSR0	GPIO Pin Direction Register GPIO<31:0>
	0x40E0_001C	GPSR1	GPIO Pin Output Set Register GPIO<63:32>
	0x40E0_0020	GPSR2	GPIO Pin Output Set Register GPIO<80:64>
	0x40E0_0024	GPCR0	GPIO Pin Output Clear Register GPIO<31:0>
	0x40E0_0028	GPCR1	GPIO Pin Output Clear Register GPIO <63:32>
	0x40E0_002C	GPCR2	GPIO Pin Output Clear Register GPIO <80:64>
	0x40E0_0030	GRER0	GPIO Rising-Edge Detect Register GPIO<31:0>
	0x40E0_0034	GRER1	GPIO Rising-Edge Detect Register GPIO<63:32>
	0x40E0_0038	GRER2	GPIO Rising-Edge Detect Register GPIO<80:64>
	0x40E0_003C	GFER0	GPIO Falling-Edge Detect Register GPIO<31:0>
	0x40E0_0040	GFER1	GPIO Falling-Edge Detect Register GPIO<63:32>
	0x40E0_0044	GFER2	GPIO Falling-Edge Detect Register GPIO<80:64>
	0x40E0_0048	GEDR0	GPIO Edge Detect Status Register GPIO<31:0>
	0x40E0_004C	GEDR1	GPIO Edge Detect Status Register GPIO<63:32>
	0x40E0_0050	GEDR2	GPIO Edge Detect Status Register GPIO<80:64>
	0x40E0_0054	GAFR0_L	GPIO Alternate Function Select Register GPIO<15:0>
	0x40E0_0058	GAFR0_U	GPIO Alternate Function Select Register GPIO<31:16>
	0x40E0_005C	GAFR1_L	GPIO Alternate Function Select Register GPIO<47:32>
	0x40E0_0060	GAFR1_U	GPIO Alternate Function Select Register GPIO<63:48>
	0x40E0_0064	GAFR2_L	GPIO Alternate Function Select Register GPIO<79:64>
	0x40E0_0068	GAFR2_U	GPIO Alternate Function Select Register GPIO 80
Power Manager and Reset Control	0x40F0_0000		
	0x40F0_0000	PMCR	Power Manager Control Register
	0x40F0_0004	PSSR	Power Manager Sleep Status Register
	0x40F0_0008	PSPR	Power Manager Scratch Pad Register
	0x40F0_000C	PWER	Power Manager Wake-up Enable Register
	0x40F0_0010	PRER	Power Manager GPIO Rising-Edge Detect Enable Register
	0x40F0_0014	PFER	Power Manager GPIO Falling-Edge Detect Enable Register
	0x40F0_0018	PEDR	Power Manager GPIO Edge Detect Status Register
	0x40F0_001C	PCFR	Power Manager General Configuration Register
	0x40F0_0020	PGSR0	Power Manager GPIO Sleep State Register for GP[31-0]
	0x40F0_0024	PGSR1	Power Manager GPIO Sleep State Register for GP[63-32]
	0x40F0_0028	PGSR2	Power Manager GPIO Sleep State Register for GP[84-64]
	0x40F0_002C		Reserved



Table 2-8. System Architecture Register Address Summary (Sheet 10 of 12)

Unit	Address	Register Symbol	Register Description
	0x40F0_002C	_	Reserved
	0x40F0_0030	RCSR	Reset Controller Status Register
SSP	0x4100_0000		
	0x4100_0000	SSCR0	SSP Control Register 0
	0x4100_0004	SSCR1	SSP Control Register 1
	0x4100_0008	SSSR	SSP Status Register
	0x4100_000C	SSITR	SSP Interrupt Test Register
	0x4100_0010	SSDR (Write / Read)	SSP Data Write Register/SSP Data Read Register
MMC Controller	0x4110_0000		
	0x4110_0000	MMC_STRPCL	Control to start and stop MMC clock
	0x4110_0004	MMC_STAT	MMC Status Register (read only)
	0x4110_0008	MMC_CLKRT	MMC clock rate
	0x4110_000C	MMC_SPI	SPI mode control bits
	0x4110_0010	MMC_CMDAT	Command/response/data sequence control
	0x4110_0014	MMC_RESTO	Expected response time out
	0x4110_0018	MMC_RDTO	Expected data read time out
	0x4110_001C	MMC_BLKLEN	Block length of data transaction
	0x4110_0020	MMC_NOB	Number of blocks, for block mode
	0x4110_0024	MMC_PRTBUF	Partial MMC TXFIFO FIFO written
	0x4110_0028	MMC_I_MASK	Interrupt Mask
	0x4110_002C	MMC_I_REG	Interrupt Register (read only)
	0x4110_0030	MMC_CMD	Index of current command
	0x4110_0034	MMC_ARGH	MSW part of the current command argument
	0x4110_0038	MMC_ARGL	LSW part of the current command argument
	0x4110_003C	MMC_RES	Response FIFO (read only)
	0x4110_0040	MMC_RXFIFO	Receive FIFO (read only)
	0x4110_0044	MMC_TXFIFO	Transmit FIFO (write only)
Clocks Manager	0x4130_0000		
	0x4130_0000	CCCR	Core Clock Configuration Register
	0x4130_0004	CKEN	Clock Enable Register
	0x4130_0008	OSCC	Oscillator Configuration Register
Network SSP	0x4140_0000		
	0x4140_0000	NSSCR0	NSSP Control Register 0
	0x4140_0004	NSSCR1	NSSP Control Register 1
	0x4140_0008	NSSSR	NSSP Status Register
	0x4140_000C	NSSITR	NSSP Interrupt Test Register
	0x4140_0010	NSSDR	NSSP Data Read/Write Register
	0x4140_0028	NSSTO	NSSP Time Out Register



Table 2-8. System Architecture Register Address Summary (Sheet 11 of 12)

Unit	Address	Register Symbol	Register Description
	0x4140_002C	NSSPSP	NSSP Programmable Serial Protocol
Hardware UART	0x4160_0000		
	0x4160_0000	HWRBR	Receive Buffer Register (read only)
	0x4160_0000	HWTHR	Transmit Holding Register (write only)
	0x4160_0004	HWIER	Interrupt Enable Register (read/write)
	0x4160_0008	HWIIR	Interrupt ID Register (read only)
	0x4160_0008	HWFCR	FIFO Control Register (write only)
	0x4160_000C	HWLCR	Line Control Register (read/write)
	0x4160_0010	HWMCR	Modem Control Register (read/write)
	0x4160_0014	HWLSR	Line Status Register (read only)
	0x4160_0018	HWMSR	Modem Status Register (read only)
	0x4160_001C	HWSPR	Scratch Pad Register (read/write)
	0x4160_0020	HWISR	Infrared Selection Register (read/write)
	0x4160_0024	HWFOR	FIFO Occupancy Register (read only)
	0x4160_0028	HWABR	Auto-Baud Control Register (read/write)
	0x4160_002C	HWACR	Auto-Baud Count Register
	0x4160_0000	HWDLL	Divisor Latch Low Register (DLAB = 1) (read/write)
	0x4160_0000	HWDLH	Divisor Latch High Register (DLAB = 1) (read/write)
LCD Controller	0x4400_0000		
	0x4400_0000	LCCR0	LCD Controller Control Register 0
	0x4400_0004	LCCR1	LCD Controller Control Register 1
	0x4400_0008	LCCR2	LCD Controller Control Register 2
	0x4400_000C	LCCR3	LCD Controller Control Register 3
	0x4400_0200	FDADR0	DMA Channel 0 Frame Descriptor Address Register
	0x4400_0204	FSADR0	DMA Channel 0 Frame Source Address Register
	0x4400_0208	FIDR0	DMA Channel 0 Frame ID Register
	0x4400_020C	LDCMD0	DMA Channel 0 Command Register
	0x4400_0210	FDADR1	DMA Channel 1 Frame Descriptor Address Register
	0x4400_0214	FSADR1	DMA Channel 1 Frame Source Address Register
	0x4400_0218	FIDR1	DMA Channel 1 Frame ID Register
	0x4400_021C	LDCMD1	DMA Channel 1 Command Register
	0x4400_0020	FBR0	DMA Channel 0 Frame Branch Register
	0x4400_0024	FBR1	DMA Channel 1 Frame Branch Register
	0x4400_0038	LCSR	LCD Controller Status Register
	0x4400_003C	LIIDR	LCD Controller Interrupt ID Register
	0x4400_0040	TRGBR	TMED RGB Seed Register
	0x4400_0044	TCR	TMED Control Register
Memory Controller	0x4800_0000		



Table 2-8. System Architecture Register Address Summary (Sheet 12 of 12)

Unit	Address	Register Symbol	Register Description
	0x4800_0000	MDCNFG	SDRAM Configuration Register 0
	0x4800_0004	MDREFR	SDRAM Refresh Control Register
	0x4800_0008	MSC0	Static Memory Control Register 0
	0x4800_000C	MSC1	Static Memory Control Register 1
	0x4800_0010	MSC2	Static Memory Control Register 2
	0x4800_0014	MECR	Expansion Memory (PCMCIA/Compact Flash) Bus Configuration Register
	0x4800_001C	SXCNFG	Synchronous Static Memory Control Register
	0x4800_0024	SXMRS	MRS value to be written to SMROM
	0x4800_0028	MCMEM0	Card interface Common Memory Space Socket 0 Timing Configuration
	0x4800_002C	MCMEM1	Card interface Common Memory Space Socket 1 Timing Configuration
	0x4800_0030	MCATT0	Card interface Attribute Space Socket 0 Timing Configuration
	0x4800_0034	MCATT1	Card interface Attribute Space Socket 1 Timing Configuration
	0x4800_0038	MCIO0	Card interface I/O Space Socket 0 Timing Configuration
	0x4800_003C	MCIO1	Card interface I/O Space Socket 1 Timing Configuration
	0x4800_0040	MDMRS	MRS value to be written to SDRAM
	0x4800_0044	BOOT_DEF	Read-only Boot-Time Register. Contains BOOT_SEL and PKG SEL values.
	0x4800_0058	MDMRSLP	Low Power SDRAM Mode Register Set Configuration Register
	0x4800_0064	SA1111CR	SA1111 Compatibility Register



Clocks and Power Manager

The Clocks and Power Manager for the PXA255 processor controls the clock frequency to each module and manages transitions between the different power manager (PM) operating modes to optimize both computing performance and power consumption.

3.1 Clock Manager Introduction

The Clocks and Power Manager provides fixed clocks for each peripheral unit. Many of the devices' peripheral clocks can be disabled using the Clock Enable Register (CKEN), or through bits in the peripheral's control registers. To minimize power consumption, turn off the clock to any unit that is not being used. The Clocks and Power Manager also provides the programmable-frequency clocks for the LCD Controller, Memory Controller, and CPU. These clocks are related to each other because they come from the same internal Phase Locked Loop (PLL) clock source. To program the PLL's frequency, follow these steps (for information on the factors L, M, and N, see Section 3.6.1, "Core Clock Configuration Register (CCCR)" on page 3-34):

- 1. Determine the fastest synchronous memory requirement (SDRAM frequency).
- 2. If the SDRAM frequency is less than 99.5 MHz, the Memory Frequency must be twice the SDRAM Frequency and the SDRAM clock ratio in the Memory Controller must be set to two. If the SDRAM frequency is 99.5 MHz, the Memory Frequency is equal to the SDRAM frequency.
- 3. Round the Memory Frequency down to the nearest value of 99.5 MHz (L=0x1B), 118.0 MHz (L=0x20), 132.7 MHz (L=0x24), 147.5 MHz (L=0x28), or 165.9 MHz (L=0x2D), and program the value of L in the Core Clock Configuration register. This frequency (or half, if the SDRAM clock ratio is 2) is the External Synchronous Memory Frequency.
- 4. Determine the required Core Frequency for normal (Run Mode) operation. This mode is used during normal processing, when the application must make occasional fetches to external memory. The possible values are one, two, or four times the Memory Frequency. Program this value (M) in the Core Clock Configuration register.
- 5. Determine the required Core Frequency for Turbo Mode operation. This mode is generally used when the application runs entirely from the caches, because any fetches to external memory slow the Core's performance. This value is a multiple (1.0, 1.5, 2.0, or 3.0) of the Run Mode Frequency. Program the value (N) in the Core Clock Configuration register.
- 6. Configure the LCD Controller and Memory Controller for the new Memory Frequency and enter the Frequency Change Sequence (described in Section 3.4.7, "Frequency Change Sequence" on page 3-11).

Note: Not all frequency combinations are valid. See Section 3.3.3, "Core Phase Locked Loop" for valid combinations.



3.2 Power Manager Introduction

The Clocks and Power Manager can place the processor in one of three resets.

- Hardware Reset (nRESET asserted) is a nonmaskable total reset. It is used at power up or when no system information requires preservation.
- Watchdog Reset is asserted through the Watchdog Timer and resets the system except the Clocks and Power Manager. This reset is used as a code monitor. If code fails to complete a specified sequence, the processor assumes a fatal system error has occurred and causes a Watchdog Reset.
- GPIO Reset is enabled through the GPIO alternate function registers. It is used as an alternative to Hardware Reset that preserves the Memory Controller registers and a few critical states in the Clocks and Power Manager and the Real Time Clock (RTC).

The Clocks and Power Manager also controls the entry into and exit from any of the low power or special clocking modes on processor. These modes are:

- Turbo Mode: the Core runs at its peak frequency. In this mode, make very few external memory accesses because the Core must wait on the external memory.
- Run Mode: the Core runs at its normal frequency. In this mode, the Core is assumed to be
 doing frequent external memory accesses, so running slower is optimum for the best power/
 performance trade-off.
- Idle Mode: the Core is not being clocked, but the rest of the system is fully operational. This mode is used during brief lulls in activity, when the external system must continue operation but the Core is idle.
- Sleep Mode: places the processor in its lowest power state but maintains I/O state, RTC, and
 the Clocks and Power Manager. Wake-up from Sleep Mode requires re-booting the system,
 since most internal state was lost. The core power must be grounded in sleep to prevent current
 leakage.

The Clocks and Power Manager also controls the processor's actions during the Frequency Change Sequence. The Frequency Change Sequence is a sequence that changes the Core Frequency (Run and Turbo) and Memory Frequency from the previously stored values to the new values in the Core Clock Configuration register. This sequence takes time to complete due to PLL relock time, but it allows dynamic frequency changes without compromising external memory integrity. Any peripherals that rely on the Core or Memory Controller must be configured to withstand a data flow interruption.

3.3 Clock Manager

The processor's clocking system incorporates five major clock sources:

- 32.768 kHz Oscillator
- 3.6864 MHz Oscillator
- Programmable Frequency Core PLL
- 95.85 MHz Fixed Frequency Peripheral PLL
- 147.46 MHz Fixed Frequency PLL

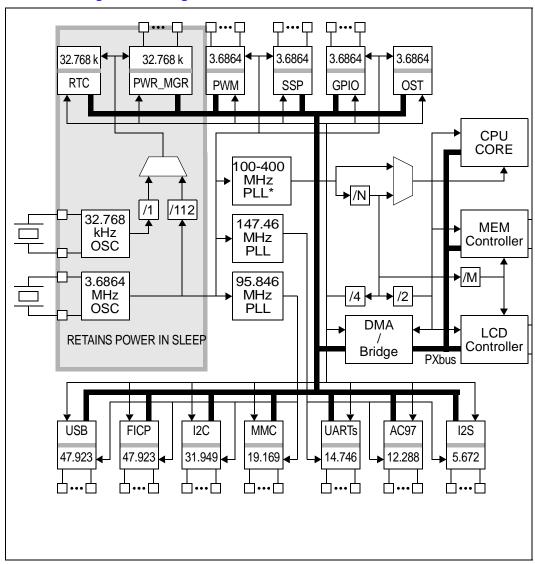


The clocks manager also contains clock gating for power reduction.

Figure 3-1 shows a functional representation of the clocking network. "L" is in the core PLL.

The PXbus is the internal bus between the Core, the DMA/Bridge, the LCD Controller, and the Memory Controller as shown in Figure 3-1. This bus is clocked at 1/2 the run mode frequency. For optimal performance, the PXbus should be clocked as fast as possible. For example, if a target core frequency of 200 MHz is desired use 200 MHz run mode instead of 200 MHz turbo mode with run at 100 MHz. Increasing the PXbus frequency may help reduce the latency involved in accessing non-cacheable memory.

Figure 3-1. Clocks Manager Block Diagram





3.3.1 32.768 kHz Oscillator

The 32.768 kHz oscillator is a low power, low frequency oscillator that clocks the RTC and Power Manager. This oscillator is disabled out of Hardware Reset and the RTC and Power Manager blocks use the 3.6864 MHz oscillator instead. Software writes the Oscillator On bit in the Oscillator Configuration Register to enable the 32.768 kHz. This configures the RTC and Power Manager to use the 32.768 kHz oscillator after it stabilizes.

32.768 kHz oscillator use is optional and provides the lowest power consumption during Sleep Mode. In less power-sensitive applications, disable the 32.768 kHz oscillator in the Oscillator Configuration Register (OSCC) and leave the external pins floating (no external crystal required) for cost savings. If the 32.768 kHz oscillator is not in the system, the frequency of the RTC and Power Manager will be 3.6864 MHz divided by 112 (32.914 kHz). In Sleep, the 3.6864 MHz oscillator consumes hundreds of microamps of extra power when it stays enabled. See Section 3.5.2, "Power Manager General Configuration Register (PCFR)" on page 3-24 for information on The Oscillator Power Down Enable (OPDE) bit, which determines if the 3.6864 MHz oscillator is enabled in Sleep Mode. No external capacitors are required.

3.3.2 3.6864 MHz Oscillator

The 3.6864 MHz oscillator provides the primary clock source for the processor. The on-chip PLL frequency multipliers, Synchronous Serial Port (SSP), Pulse Width Modulator (PWM), and the Operating System Timer (OST) use the 3.6864 MHz oscillator as a reference. Out of Hardware Reset, the 3.6864 MHz oscillator also drives the RTC and Power Manager (PM). The user may then enable the 32.768 kHz oscillator, which will drive the RTC and PM after it is stabilized. The 3.6864 MHz oscillator can be disabled during Sleep Mode by setting the OPDE (see Section 3.5.2) bit but only if the 32.768 kHz oscillator is enabled and stabilized (both the OON and OOK bits in the OSCC set). See Section 3.6.3 for more information. No external capacitors are required.

3.3.3 Core Phase Locked Loop

The Core PLL is the clock source of the CPU Core, the Memory Controller, the LCD Controller, and DMA Controller. The Core PLL uses the 3.6864 MHz oscillator as a reference and multiplies its frequency by the following variables:

- L: Crystal Frequency to Memory Frequency Multiplier, set to 27, 32, 36, 40, or 45.
- M: Memory Frequency to Run Mode Frequency Multiplier, set to 1 or 2.
- N: Run Mode Frequency to Turbo Mode Frequency Multiplier, set to 1.0, 1.5, 2.0, or 3.0.

The output frequency selections are shown in Table 3-1, "Core PLL Output Frequencies for 3.6864 MHz Crystal". See Section 3.6.1 for programming information on the L, M, and N factors. See Section 3.6.1, "Core Clock Configuration Register (CCCR)" for the hexadecimal settings.

Do not choose a combination that generates a frequency that is not supported in the voltage range and package in which the processor is operating.

SDCLK must not be greater than 100 MHz. If MEMCLK is greater than 100 MHz, the SDCLK to MEMCLK ratio must be set to 1:2 in the Memory Controller.



L M		Turbo Mode Frequency (MHz) for Values "N" and Core Clock Configuration Register (CCCR[15:0]) programming for Values of "N"				PXbus Frequency (MHz)	MEM, LCD Frequency (MHz)	SDRAM max Frequency
		1.00 (Run)	1.50	2.00	3.00	(MHZ)	(WITIZ)	(MHz)
27	1	99.5 @1.0 V	_	199.1 @1.0 V	298.6 @1.1 V	50	99.5	99.5
36	1	132.7 @1.0 V	_	_	_	66	132.7	66
27	2	199.1 @1.0 V	298.6 @1.1 V	398.1 @1.3 V	_	99.5	99.5	99.5
36	2	265.4 @1.1 V	_	_	_	132.7	132.7	66
45	2	331.8 @1.3 V	_	_	_	165.9	165.9	83
27	4	398.1 @1.3 V	_	_	_	196	99.5	99.5

Table 3-1. Core PLL Output Frequencies for 3.6864 MHz Crystal

3.3.4 95.85 MHz Peripheral Phase Locked Loop

The 95.85 MHz PLL is the clock source for many of the peripheral blocks' external interfaces. These interfaces require ~48 MHz (UDC/USB, FICP), ~33 MHz ($\rm I^2C$), and ~20 MHz (MMC). The generated frequency is not exactly the required frequency due to the chosen crystal and the lack of a perfect Least Common Multiple between the units. The chosen frequencies keep each unit's clock frequency within the unit's clock tolerance. If a crystal other than 3.6864 MHz is used, the clock frequencies to the peripheral blocks' interfaces may not yield the desired baud rates (or protocol's rate).

Table 3-2. 95.85 MHz Per	ipheral PLL Out	put Frequencies f	or 3.6864 MHz Cry	/stal
--------------------------	-----------------	-------------------	-------------------	-------

Unit Name	Nominal Frequency	Actual Frequency
USB (UDC)	48 MHz	47.923 MHz
FICP	48 MHz	47.923 MHz
I ² C	33 MHz	31.949 MHz
MMC	20 MHz	19.169 MHz

3.3.5 147.46 MHz Peripheral Phase Locked Loop

The 147.46 MHz PLL is the clock source for many of the peripheral blocks' external interfaces. These interfaces require ~14.75 MHz (UARTs), 12.288 MHz (AC97), and variable frequencies (I²S). The generated frequency may not exactly match the required frequency due to the choice of crystal and the lack of a perfect Least Common Multiple between the units. The chosen frequencies keep each unit's clock frequency within the unit's clock tolerance. If a crystal other than 3.6864 MHz is used, the clock frequencies to the peripheral blocks' interfaces may not yield the desired baud rates (or other protocol's rate)



Table 3-3. 147.46 MHz Peripheral PLL Output Frequencies for 3.6864 MHz Crystal

Unit Name	Nominal Frequency	Actual Frequency
UARTs	14.746 MHz	14.746 MHz
AC97	12.288 MHz	12.288 MHz
I ² S	146.76 MHz	147.46 MHz

3.3.6 Clock Gating

The Clocks Manager contains the CKEN register. This register contains configuration bits that can disable the clocks to individual units. The configuration bits are used when a module is not being used. After Hardware Reset, any module that is not being used must have its clock disabled. If a module is temporarily quiescent but does not have clock gating functionality, the CKEN register can be used to disable the unit's clock.

When a module's clock is disabled, the registers in that module are still readable and writable. The AC97 is an exception and is completely inaccessible if the clock is disabled.

3.4 Resets and Power Modes

The Clocks and Power Manager Unit determines the processor's Resets, Power Sequences and Power Modes. Each behaves differently during operation and has specific entry and exit sequences. The resets and power modes are:

- Hardware Reset
- · Watchdog Reset
- GPIO Reset
- Run Mode
- Turbo Mode
- · Idle Mode
- Frequency Change Sequence
- Sleep Mode

3.4.1 Hardware Reset

To invoke the Hardware Reset and reset all units in the processor to a known state, assert the nRESET pin. Hardware Reset is only intended to be used for power up and complete resets.

3.4.1.1 Invoking Hardware Reset

Hardware Reset is invoked when the nRESET pin is pulled low by an external source. The processor does not provide a method of masking or disabling the propagation of the external pin value. When the nRESET pin is asserted, Hardware Reset is invoked, regardless of the mode of operation. The nRESET_OUT pin is asserted when the nRESET pin is asserted. To enter Hardware



Reset, nRESET must be held low for t_{DHW_NRESET} to allow the system to stabilize and the reset state to propagate. Refer to the *Intel® PXA255 Processor Electrical, Mechanical, and Thermal Specification* for details.

3.4.1.2 Behavior During Hardware Reset

During Hardware Reset, all internal registers and units are held at their defined reset conditions. While the nRESET pin is asserted, nothing inside the processor is active except the 3.6864 MHz oscillator. The internal clocks are stopped and the chip is static. All pins return to their reset conditions and the nBATT_FAULT and nVDD_FAULT pins are ignored. Because the memory controller receives a full reset, all dynamic RAM contents are lost during Hardware Reset.

3.4.1.3 Completing Hardware Reset

To complete Hardware Reset, deassert the nRESET pin. All power supplies must be stable for t_{D_NRESET} before nRESET is deasserted. Refer to the Intel® PXA255 Processor *Electrical*, *Mechanical*, *and Thermal Specification* for details. After the nRESET pin is deasserted, the following sequence occurs:

- 1. The 3.6864 MHz oscillator and internal PLL clock generators wait for stabilization.
- 2. The nRESET_OUT pin is deasserted.
- 3. The normal boot-up sequence begins. All processor units return to their predefined reset conditions. Software must examine the Reset Controller Status register (RCSR) to determine the cause for the boot.

3.4.2 Watchdog Reset

Watchdog Reset is invoked when software fails to properly prevent the Watchdog Time-out Event from occurring. It is assumed that Watchdog Resets are only generated when software is not executing properly and has potentially destroyed data. In Watchdog Reset all units in the are reset except the Clocks and Power Manager.

3.4.2.1 Invoking Watchdog Reset

Watchdog Reset is invoked when the Watchdog Enable bit (WE) in the OWER is set and the OSMR[3] matches the OS timer counter. When these conditions are met, they invoke Watchdog Reset, regardless of the previous mode of operation. Watchdog Reset asserts nRESET_OUT.

3.4.2.2 Behavior During Watchdog Reset

During Watchdog Reset, all units except the Real Time Clock and parts of the Clocks and Power Manager maintain their defined reset conditions. All pins except the oscillator pins assume their reset conditions and the nBATT_FAULT and nVDD_FAULT pins are ignored. All dynamic RAM contents are lost during Watchdog Reset because the memory controller receives a full reset.

Refer to Table 2-6, "Pin & Signal Descriptions for the PXA255 Processor" for the pin states during Watchdog and other Resets.



3.4.2.3 Completing a Watchdog Reset

Watchdog resets immediately revert to hardware resets when the nRESET pin is asserted. Otherwise, the completion sequence for watchdog reset is:

- 1. The 3.6864 MHz oscillator and internal PLL clock generators wait for stabilization. The 32.768 kHz oscillator's configuration and status are not affected by watchdog reset.
- 2. The nRESET_OUT pin is deasserted after t_{DHW_OUT}. Refer to the *Intel® PXA255 Processor Electrical, Mechanical, and Thermal Specification*.
- 3. The normal boot-up sequence begins. All processor units except the RTTR in the RTC and parts of the Clocks and Power Manager return to their predefined reset conditions. Software must examine the RCSR to determine the cause for the reboot.

3.4.3 GPIO Reset

A GPIO Reset is invoked when GP[1] is properly configured as a reset source and is asserted low for greater than four 3.6864-MHz clock cycles. In GPIO Reset all processor units except the RTC, parts of the Clocks and Power Manager, and the Memory Controller return to their predefined, known states.

3.4.3.1 Invoking GPIO Reset

To use the GPIO Reset function, set it up through the GPIO Controller. The GP[1] pin must be configured as an input and set to its alternate GPIO Reset function in the GPIO Controller. The GPIO Reset alternate function is level-sensitive and not edge-triggered. To ensure no spurious resets are generated when the alternate GPIO Reset function is set, follow these steps:

- 1. GP[1] must be set up as an output with its data register set to a 1.
- 2. Externally drive the GP[1] pin to a high state.
- 3. Configure GP[1] as an input.
- 4. Configure GP[1] for its Alternate (Reset) Function.

The previous mode of operation does not affect a GPIO Reset. When performing a GPIO Reset, nRESET_OUT is asserted. If GP[1] is asserted for less than four 3.6864-MHz clock cycles, the processor may remain in its previous mode or enter into a GPIO reset.

GPIO Reset does not function in Sleep Mode because all GPIO pins' Alternate Function Inputs are disabled. External wake-up sources must be routed through one of the enabled GPIO wake-up sources (see Section 3.5.3 for details) during Sleep Mode. GP[1] may be enabled as a wake-up source.

3.4.3.2 Behavior During GPIO Reset

During GPIO Reset, most, but not all, internal registers and processes are held at their defined reset conditions. The exceptions are the RTC, the Clocks and Power Manager (unless otherwise noted), and the Memory Controller. During GPIO Reset, the clocks unit continues to operate with its previously programmed values, so the processor enters and exits GPIO Reset with the same clock configurations. All pins except the oscillator and Memory Controller pins return to their reset conditions and the nBATT_FAULT and nVDD_FAULT pins are ignored.



GPIO Reset does not reset the Memory Controller Configuration registers. This creates the possibility that the contents of external memories may be preserved if the external memories are properly configured before GPIO Reset is entered. To preserve SDRAM contents during a GPIO Reset, software must correctly configure the Memory Control and the time spent in GPIO Reset must be shorter than the SDRAM refresh interval. The amount of time spent in GPIO Reset depends on the CPU's mode before GPIO Reset. See Section 6, "Memory Controller" for details.

Refer to Table 2-6, "Pin & Signal Descriptions for the PXA255 Processor" for the states of all the PXA255 processor pins during GPIO reset and other resets.

3.4.3.3 Completing GPIO Reset

GPIO Reset immediately reverts to Hardware Reset when the nRESET pin is asserted. Otherwise, the completion sequence for GPIO Reset is:

- 1. The GPIO Reset Source is deasserted because the internal reset has propagated to the GPIO Controller and its registers, which are set back to their reset states.
- 2. The nRESET_OUT pin is deasserted.
- 3. The normal boot-up sequence begins. All processor units except the Real Time Clock, parts of the Clocks and Power Manager, and the Memory Controller return to their predefined reset conditions. Software must examine the RCSR to determine the cause for the reset.

3.4.4 Run Mode

Run Mode is the processor's normal operating mode. All power supplies are enabled and all functionally enabled clocks are running. Run Mode is entered after any power mode, power sequence, or reset completes its sequence. Run Mode is exited when any other power mode, power sequence, or reset begins.

3.4.5 Turbo Mode

Turbo Mode allows the user to clock the processor core at a higher frequency during peak processing requirements. It allows a synchronous switch in frequencies without disrupting the Memory Controller, LCD Controller, or any peripheral.

3.4.5.1 Entering Turbo Mode

Turbo Mode is invoked when software sets the TURBO bit in the Clock Config (CCLKCFG) Register (See Section 3.7.1). After software sets the TURBO bit, the CPU waits for all instructions currently in the pipeline to complete. When the instructions are completed, the CPU resumes operation at the higher Turbo Mode Frequency.

Software can set or clear other bits in the CCLKCFG in the same write that sets the TURBO bit. The other bits in the register take precedence over Turbo Mode, so, if another bit is set, that mode's sequence is followed before the CPU enters Turbo Mode. When the CPU exits the other mode, it enters either Run or Turbo Mode, based on the state of the CCLKCFG [TURBO] bit.

Do not confuse the CCLKCFG Register, which is in Coprocessor 14, with the CCCR (See Section 3.6.1), which is in the processor's Clocks and Power Manager.



3.4.5.2 Behavior in Turbo Mode

The processor's behavior in Turbo Mode is identical to its behavior in Run Mode, except that the processor's clock frequency relative to the memory and peripherals is increased by N, the value in the CCCR (see Section 3.6.1). Turbo mode is intended for use during peak processing, when there are very few accesses to external memory. The higher Core to external memory clock ratio increases the relative delay for each external memory access. This increased delay lowers the processor's power efficiency. For optimum performance, software must load applications in the caches in Run Mode and execute them in Turbo Mode.

3.4.5.3 Exiting Turbo Mode

To exit Turbo Mode, software clears the TURBO bit in the CCLKCFG Register. After software clears the TURBO bit, the CPU waits for all instructions in the pipeline to complete. When the instructions are completed, the CPU enters Run Mode.

Other bits in the CCLKCFG may be set or cleared in the write that clears CCLKCFG [TURBO]. All other bits in the register take precedence over Turbo Mode, so the new mode's proper sequence is followed.

Idle, Sleep, Frequency Change Sequence, and Reset have precedence over Turbo Mode and cause the processor to exit Turbo Mode. When the CPU exits of one of these modes, it enters either Run or Turbo Mode, based on the state of CCLKCFG [TURBO].

3.4.6 Idle Mode

Idle Mode allows the user to stop the CPU core clock during periods of processor inactivity and continue to monitor on- and off-chip interrupt service requests. Idle mode does not change clock generation, so when an interrupt occurs the CPU is quickly reactivated in the state that preceded Idle Mode.

During Idle mode these resources are active:

- System unit modules (real-time clock, operating system timer, interrupt controller, general-purpose I/O, and clocks and power manager)
- Peripheral unit modules (DMA controller, LCD controller, and all other peripheral units)
- Memory Controller resources

3.4.6.1 Entering Idle Mode

During Idle Mode, the clocks to the CPU core stop. All critical applications must be finished and peripherals must be set up to generate interrupts when they require CPU attention. To enter the Idle Mode, software selects Idle Mode in PWRMODE[M] (See Section 3.7.2). An interrupt immediately aborts Idle Mode and normal processing resumes. After software selects Idle Mode, the CPU waits until all instructions in the pipeline are completed. When the instructions are completed, the CPU clock stops and Idle Mode begins. In Idle Mode, interrupts are recognized as wake-up sources.



3.4.6.2 Behavior in Idle Mode

In Idle Mode the CPU clocks are stopped, but the remainder of the processor operates normally. For example, the LCD controller can continue refreshing the screen with the same frame buffer data in memory.

When ICCR[DIM] is cleared, any enabled interrupt wakes up the processor. When ICCR[DIM] is set, only unmasked interrupts cause wake-up.

Enabled interrupts are interrupts that are allowed at the unit level. Masked interrupts are interrupts that are prevented from interrupting the core based on the Interrupt Controller Mask Register.

3.4.6.3 Exiting Idle Mode

Idle Mode exits when any Reset is asserted. Reset entry and exit sequences take precedence over Idle Mode. When the Reset exit sequence is completed, the CPU is not in Idle Mode. If the Watchdog Timer is enabled, software must set the Watchdog Match Registers before it sets Idle Mode to ensure that another interrupt will bring the processor out of Idle Mode before the Watchdog Reset is asserted. Use an RTC alarm or another OS timer channel for this purpose.

Any enabled interrupt causes Idle Mode to exit. When ICCR[DIM] is cleared, the Interrupt Controller Mask register (ICMR) is ignored during Idle Mode. This means that an interrupt does not have to be unmasked to cause Idle Mode to exit. Idle Mode exits in the following sequence:

- 1. A valid, enabled Interrupt asserts.
- 2. The CPU clocks restart and the CPU resumes operation at the state indicated by CCLKCFG [TURBO].

Idle Mode also exits when the nBATT_FAULT or nVDD_FAULT pin is asserted. When either pin is asserted, Idle Mode exits in the following sequence:

- 1. The nBATT_FAULT or nVDD_FAULT pin is asserted.
- 2. If the Imprecise Data Abort Enable (IDAE) bit in the Power Manager Control Register (PMCR) is clear (not recommended), the processor enters Sleep Mode immediately.
- 3. If the IDAE bit is set, the nBATT_FAULT or nVDD_FAULT assertion is treated as a valid interrupt to the clocks module and Idle Mode exits using its normal, interrupt-driven sequence. Software must then shut down the system and enter Sleep Mode. See Section 3.4.9.3, "Entering Sleep Mode" for more details.

3.4.7 Frequency Change Sequence

The Frequency Change Sequence is used to change the processor clock frequency. During the Frequency Change Sequence, the CPU, Memory Controller, LCD Controller, and DMA clocks stop. The other peripheral units continue to function during the Frequency Change Sequence. This mode is intended to be used to change the frequency from the default condition at initial boot-up. It may also be used as a power-saving feature used to allow the processor to run at the minimum required frequency when the software requires major changes in frequency.

3.4.7.1 Preparing for a Frequency Change Sequence

Software must complete the following steps before it initiates the Frequency Change Sequence:



- 1. Configure the Memory Controller to ensure SDRAM contents are maintained during the Frequency Change Sequence. The Memory Controller's refresh timer must be programmed to match the maximum refresh time associated with the slower of two frequencies (current and desired). The SDRAM divide by two must be set to a value that prevents the SDRAM frequency from exceeding the specified frequency. For example, to change from 100/100 to 133/66, the SDRAM bus must be set to divide by two before the frequency change. To change from 133/66 to 100/100, the SDRAM must be set to one-to-one after the frequency change sequence is completed. See Section 6, "Memory Controller" for more details.
- 2. Disable the LCD Controller or configure it to avoid the effects of an interruption in the LCD clocks and data from the processor.
- 3. Configure peripheral units to handle a lack of DMA service for up to 500 µs. If a peripheral unit can not function for 500 µs without DMA service, it must be disabled.
- 4. Disable peripheral units that can not accommodate a 500 μs interrupt latency. The interrupts generated during the Frequency Change Sequence are serviced when the sequence exits.
- 5. Program the CCCR (Section 3.6.1, "Core Clock Configuration Register (CCCR)") to reflect the desired frequency.

3.4.7.2 Invoking the Frequency Change Sequence

To invoke the Frequency Change Sequence, software must set FCS in the CCLKCFG (See Section 3.7.1). When software sets FCS, it may also set or clear other bits in CCLKCFG. If software sets the TURBO bit in the same write, the CPU enters Turbo Mode when the Frequency Change Sequence exits.

After software sets the FCS:

- 1. The CPU clock stops and interrupts to the CPU are gated.
- 2. The Memory Controller completes all outstanding transactions in its buffers and from the CPU. New transactions from the LCD or DMA controllers are ignored.
- 3. The Memory Controller places the SDRAM in self-refresh mode.

Note: Program the Memory Controller to ensure the correct self-refresh time for SDRAM, given the slower of the current and desired clock frequencies.

3.4.7.3 Behavior During the Frequency Change Sequence

In the frequency change sequence, the processor's PLL clock generator is in the process of locking to the correct frequency and cannot be used. This means that interrupts cannot be processed. Interrupts that occur during the frequency change sequence are serviced after the processor's PLL has locked. The 95.85 MHz and 147.46 MHz PLL clock generators are active and peripherals, except the memory, LCD, and DMA controllers, may continue to operate normally, provided they can accommodate the inability to process DMA or interrupt requests. DMA or interrupt requests are not recognized until the frequency change sequence is complete.

The Imprecise Data Abort is also not recognized and if nVDD_FAULT or nBATT_FAULT is asserted, the assertion is ignored until the Frequency Change Sequence exits. This means that the processor does not enter Sleep Mode until the Frequency Change Sequence is complete.



3.4.7.4 Completing the Frequency Change Sequence

The Frequency Change Sequence exits when any Reset is asserted. In Hardware and Watchdog Resets, the Reset entry and exit sequences take precedence over the Frequency Change Sequence and the PLL resumes in its Reset condition. In GPIO Reset, the Reset exit sequence is delayed while the PLL relocks and the frequency is set to the desired frequency of the Frequency Change Sequence.

If the Watchdog Timer is enabled during the Frequency Change Sequence, set the Watchdog Match Register to ensure that the Frequency Change Sequence completes before the Watchdog Reset is asserted.

If Hardware or Watchdog Reset is asserted during the Frequency Change Sequence, the DRAM contents are lost because all states, including Memory Controller configuration and information about the previous Frequency Change Sequence, are reset. If GPIO Reset is asserted during the Frequency Change Sequence, the SDRAM contents will be lost during the GPIO Reset exit sequence if the SDRAM is not in self-refresh mode and the exit sequence exceeds the refresh interval.

Normally, the Frequency Change Sequence exits in the following sequence:

1. The processor's PLL clock generator is reprogrammed with the desired values, which are in the CCCR, and begins to relock to those values.

Note: This sequence occurs even if the before and after frequencies are the same.

- 2. The internal PLL clock generator for the processor clock waits for stabilization. Refer to the Intel® PXA250 and PXA210 Application Processors Electrical, Mechanical, and Thermal Specification for details.
- 3. The CPU clocks restart and the CPU resumes operation at the state indicated by the TURBO bit (either Run or Turbo Mode). Interrupts to the CPU are no longer gated.
- 4. The FCS bit is not automatically cleared. To prevent an accidental return to the Frequency Change Sequence, software must not immediately clear the FCS bit. The bit must be cleared on the next required write to the register.
- 5. Values may be written to the CCCR, but they are ignored until the Frequency Change Sequence is re-entered.
- 6. The SDRAM must transition out of self-refresh mode and into its idle state. See Section 6, "Memory Controller" for details on configuring the SDRAM interface.

3.4.8 33-MHz Idle Mode

33-MHz idle mode has the lowest power consumption of any idle mode. The run mode frequency selected in the Core Clock Configuration Register (CCCR) directly affects the processor idle mode power consumption. Faster run mode frequencies consume more power. 33-MHz idle mode places the processor a special low speed run mode before entering idle. This is similar to normal idle since the CPU core clock can be stopped during periods of processor inactivity and continue to monitor on- and off-chip interrupt service requests. 33-MHz idle limitations are:

- Peripherals will not function correctly and should be disabled before entering this mode.
- A Frequency Change Sequence must be performed upon entry to and exit from 33-MHz idle mode.



- SDRAM is placed in self refresh before entering 33-MHz idle mode, because SDRAM cannot be refreshed correctly in 33-MHz idle mode. Carefully consider the processor interrupt behavior when the SDRAM in self refresh. To allow the interrupts to occur while SDRAM is in self refresh, set the I and F bits in the CPSR. This allows interrupts to wake the processor from idle mode without jumping to the interrupt handler. When the system's SDRAM is no longer in self refresh, the I and F bits can be cleared and the interrupt is handled.
- Because nBATT_FAULT and nVDD_FAULT can cause a data abort interrupt, the function of these pins in 33-MHz idle mode also needs special consideration. Either the Imprecise Data Abort Enable (IDAE) bit in the Power Manager Control Register (PMCR) must be clear, (causing the processor to immediately enter sleep mode if either nBATT_FAULT or nVDD_FAULT are asserted) or take software precautions to avoid starting execution in or trying to use SDRAM while it is in self refresh.

During 33-Mhz idle mode these system unit modules are functional:

- · Real-time clock
- · Operating system timer
- Interrupt controller
- · General purpose I/O
- Clocks and power manager
- Flash ROM/SRAM

Unlike normal idle mode, in 33-MHz idle mode all other peripheral units cannot be used, including SDRAM, LCD and DMA controllers.

3.4.8.1 Entering 33-MHz Idle Mode

During idle mode, the processor core clocks stop. Before the clocks stop, all critical applications must be finished and peripherals turned off. If software is executing from SDRAM, the last three of the following steps must be loaded into the cache before being performed.

- 1. Set the I and F bits in the CPSR register to mask all interrupts
- 2. Place the SDRAM into self refresh mode
- 3. Perform a frequency change sequence to 33MHz mode. The CCCR value for this mode is 0x13F
- 4. Enter idle mode by selecting the PWRMODE[M] bit (refer to Section 3.7.2)

3.4.8.2 Behavior in 33-MHz Idle Mode

In 33-MHz idle mode the CPU clocks are stopped. While in 33-MHz idle mode these features of the processor all operate normally: the RTC timer, the OS timers including the watchdog timer, and the GPIO interrupt capabilities.

When ICCR[DIM] is cleared, any enabled interrupt wakes up the processor. When ICCR[DIM] is set, only unmasked interrupts cause wake-up.

Enabled interrupts are interrupts that are allowed at the unit level. Masked interrupts are interrupts that are prevented from interrupting the core based on the Interrupt Controller Mask Register (ICMR).



3.4.8.3 Exiting 33-MHz Idle Mode

The 33-MHz idle mode exit procedure is the same as the exit procedure for normal idle mode. However, because the I and F bits are set in the CPSR, the processor does not immediately jump to the interrupt vector. Instead processing continues with the instruction following the last executed instruction before 33-MHz idle mode was entered. If execution occurs from SDRAM, steps 1 and 2 must have been previously loaded into the instruction cache. The steps below are then taken:

- 1. Perform a frequency change to a supported run mode frequency, greater or equal to 100 MHz.
- 2. Take the SDRAM out of self refresh.
- 3. Clear the I and F bits in the CPSR. Execution immediately jumps to the pending interrupt handler.

3.4.9 Sleep Mode

Sleep Mode offers lower power consumption at the expense of the loss of most of the internal processor state. In Sleep Mode, the processor goes through an orderly shut-down sequence and power is removed from the core. The Power Manager watches for a wake-up event and, after it receives one, re-establishes power and goes through a reset sequence. During Sleep Mode, the RTC and Power Manager continue to function. Pin states can be controlled throughout Sleep Mode and external SDRAM is preserved because it is in self-refresh mode.

Because all activity on the processor except the RTC stops when Sleep Mode starts, peripherals must be disabled to allow an orderly shutdown. When Sleep Mode exits, the processor's state resets and processing resumes in a boot-up mode.

3.4.9.1 Sleep Mode External Voltage Regulator Requirements

To implement Sleep Mode in the simplest manner, the External Voltage Regulator, which supplies power to the processor's internal elements, must have the following characteristics:

- A power enable input pin that enables the primary supply output connected to VCC and PLL_VCC. This pin must be connected to the processor's PWR_EN pin. To support fast sleep walk-up by maintaining power during sleep, the regulator should be software configurable to ignore PWR_EN. When PWR_EN is not used, VCC and PLL_VCC may be powered on before or simultaneously with VCCN and VCCQ. In this configuration, when PWR_EN is deasserted, the core regulator must be able to maintain regulation when the load power is as little as 0.5 mW. Core supply current during sleep will vary with voltage and temperature.
- When core power is enabled during sleep, the power management IC or logic that generates nVDD_FAULT must assert this signal when any supply including VCC and PLL_VCC falls below the lower regulation limit during sleep. nVDD_FAULT must not be deasserted until all supplies are in regulation again since there is no power supply stabilization delay during the fast sleep walk-up sequence. If nVDD_FAULT is asserted during fast sleep walk-up, then the processor returns to Sleep Mode.
- When configured to disable the core supply to save power during sleep, the core regulator's output must be driven to ground when PWR_EN goes low.
- Higher-voltage outputs connected to VCCQ and VCCN are continuously driven and do not change when the PWR_EN pin is asserted.



3.4.9.2 Preparing for Sleep Mode

Before Sleep Mode starts, software must take the following steps:

- 1. The Memory Controller must be configured to ensure SDRAM contents are maintained during Sleep Mode. See Section 6, "Memory Controller" for details.
- 2. If a graceful shutdown is required for a peripheral, the peripheral must be disabled before Sleep Mode asserts. This includes monitoring DMA transfers to and from peripherals or memories to ensure they are completed. All other peripherals need not be disabled, since they are held in their reset states internally during Sleep Mode.
- 3. The following Power Manager registers must be set up for proper sleep entry and exit:
 - PM GPIO Sleep State registers (PGSR0, PGSR1, PGSR2). To avoid contention on the bus when the processor attempts to wake up, ensure that the chip selects are not set to 0 during sleep mode. If a GPIO is used as an input, it must not be allowed to float during sleep mode. The GPIO can be pulled up or down externally or changed to an output and driven with the unasserted value.
 - PM General Configuration Register Float bits [FS/FP] must be configured appropriately for the system. The General Configuration Register Float bits must be cleared on wake up. To avoid contention on the bus when the processor attempts to wake up, ensure that the chip selects are not set to 0 during sleep mode. The PCFR[OPDE] bit must be cleared to leave the 3.6864 MHz enabled during sleep if the fast walk-up sleep configuration is selected by setting the PMFW[FWAKE] bit.
 - PMFW configuration register must be set to select between the standard and fast sleep wakeup configurations. Set PMFW[FWAKE] to 1 to disable the 10 ms power supply stabilization delay during sleep wakeup if power is maintained during sleep. This configuration reduces the sleep wakeup time to approximately 650 μs.
- 4. Before the IDAE bit is set, software must configure an imprecise data abort exception handler to put the processor into sleep mode when a data abort occurs in response to nVDD_FAULT or nBATT_FAULT assertion. This abort exception event indicates that the processor is in peril of losing its main power supply.
- 5. The following Power Manager registers must be set up to detect wake-up sources and oscillator activity:
 - PM GPIO Sleep State registers (PGSR0, PGSR,1 and PGSR2).
 - PM Wake-up Enable register (PWER)
 - PM GPIO Falling-edge Detect Enable and PM GPIO Rising-edge Detect Enable registers (PFER and PRER)
 - OPDE bit in the Power Manager Configuration Register (PCFR)
 - IDAE bit in PMCR

Note: The PCFR[OPDE] bit must be cleared to enable the 3.6864 MHz oscillator during sleep when fast sleep wakeup is selected by setting the PMFW[FWAKE] bit.

3.4.9.3 Entering Sleep Mode

Software uses the PWRMODE register to enter sleep mode (See Section 3.7.2).



If the external voltage regulator is failing or the main battery is low or missing, some systems must enter sleep mode quickly. When nBATT_FAULT or nVDD_FAULT is asserted, the system is required to shut down immediately.

To allow the assertion of nVDD_FAULT or nBATT_FAULT to cause an imprecise data abort, set the Imprecise Data Abort Enable (IDAE) bit in the PMCR. Setting the IDAE bit in the PMCR will result in software executing the data abort handler routine as part of entering sleep mode. If the IDAE bit is clear, the processor enters sleep mode immediately without executing the abort handler routine.

Note:

Using an exception handler to invoke sleep in response to a power fault event is advantageous because software can clear the PMFW[FWAKE] bit and configure the power management IC to use PWR_EN to disable the core power supply during sleep to minimize power consumption from a critically low battery.

PSSR[VFS] and PSSR[BFS] can not be used prior to entering Sleep Mode to determine which type of fault occurred, VDD fault or battery fault, respectively. If either nVDD_FAULT or nBATT_FAULT signals are asserted or if both are asserted at the same time (and the IDAE bit of the PMCR is set), the software data abort handler will be called. Since there is only one common data abort handler, software must first determine if one of the two nVDD_FAULT or nBATT_FAULT assertion events resulted in an imprecise data abort by reading Coprocessor 7, Register 4, Bit 5 (PSFS). If the PSFS bit is cleared, neither a nVDD_FAULT or nBATT_FAULT assertion occurred and the data abort handler was called for some other reason. If the PSFS bit is set, this indicates either a nVDD_FAULT or nBATT_FAULT assertion occurred, but it is not possible to determine which of the two faults was asserted. For either case, nVDD_FAULT or nBATT_FAULT assertion, software should shut the system down as quickly as possible by performing the steps outlined below to enter Sleep Mode.

Note

All addresses (data and instruction) used in the abort handler routines should be resident and accessible in the memory page tables, i.e. system software developers should ensure no further aborts occur while executing an abort handler. The processor does not support recursive (nested) aborts. The system must not assert nBATT_FAULT or nVDD_FAULT signals more than once before nRESET_OUT is asserted. System software can not return to normal execution following a nBATT_FAULT or nVDD_FAULT. If a battery or VDD fault occurs while executing in the abort mode, the abort handler is reentered. This condition of a recursive abort occurrence can be detected in software by reading the Saved Program Status Register (SPSR) to see if the previous context was executing in abort mode.

To enter Sleep Mode, software must complete the following sequence:

- 1. Software uses external memory and the Power Manager Scratch Pad Register (PSPR) to preserve critical states.
- 2. Software sets Sleep Mode in PWRMODE[M]. An interrupt immediately aborts Sleep Mode and normal processing resumes.
- 3. The CPU waits until all instructions in the pipeline are complete.
- 4. The Memory Controller completes outstanding transactions in its buffers and from the CPU. New transactions from the LCD or DMA controllers are ignored.
- 5. The Memory Controller places the SDRAM in self-refresh mode.
- 6. The Power Manager switches the GPIO output pins to their sleep state. This sleep state is programmed in advance by loading the Power Manager GPIO Sleep State registers (PGSR0, PGSR1, and PGSR2). To avoid contention on the bus when the processor attempts to wake up, ensure that the chip selects are not set to 0 during sleep mode.



- 7. The CPU clock stops and power is removed from the Core.
- 8. PWR EN is deasserted.

When the Power Manger get the indication from the Memory Controller that it has finished its outstanding transactions and has put the SDRAM into self-refresh, there are eight core clock cycles before the GPIOs latch the PGSR values and four core clock cycles after that, nRESET_OUT asserts low.

In some systems the Imprecise Data Abort latency lasts longer than the residual charge in the failed power supply can sustain operation. This normally only occurs when the processor is in a Power Mode or Sequence that requires that the processor exit before Sleep Mode starts. Frequency Change Sequence is an example of such a Power Sequence. In these Power Modes and Sequences, the IDAE bit must not be set. This allows the processor to enters Sleep Mode immediately but any critical states in the processor are lost.

If the IDAE bit is not set and the nVDD_FAULT or nBATT_FAULT pin is asserted, the Sleep Sequence begins at Step 4.

3.4.9.4 Behavior in Sleep Mode

In Sleep Mode, all processor and peripheral clocks are disabled, except the RTC. The processor does not recognize interrupts or external pin transitions except valid wake-up signals, Reset signals, and the nBATT_FAULT signal.

If the nBATT_FAULT signal is asserted while in Sleep Mode, GPIO[1:0] are set as the only valid wake-up signals.

The Power Manager watches for wake up events programmed by the CPU before Sleep Mode starts or set by the Power Manager it detects a fault condition. In order to detect a rising-edge or falling-edge on a GPIO pin, the rising- or falling-edge must be held for more than one full 32.768 kHz clock cycle. The Power Manager takes three 32.768 kHz clock cycles to acknowledge the GPIO edge and begin the wake up sequence.

Refer to Table 2-6, "Pin & Signal Descriptions for the PXA255 Processor" on page 2-9 for the PXA255 processor pin states during sleep mode reset and other resets.

3.4.9.5 Exiting Sleep Mode

Sleep Mode exits when Hardware Reset is asserted. Hardware Reset's entry and exit sequences take precedence over Sleep Mode.

Note: If Hardware Reset is asserted during Sleep Mode, the DRAM contents are lost because all states, including Memory Controller configuration and information about the previous Sleep Mode, are reset.

Normally, Sleep Mode exits in the following sequence. Any time the nBATT_FAULT pin is asserted, the processor returns to Sleep Mode. The nVDD_FAULT pin is ignored until the external power supply stabilization timer expires.

1. A pre-programmed wake up event from an enabled GPIO or RTC source occurs. If the nBATT FAULT pin is asserted, the wake up source is ignored.



- 2. The PWR_EN signal is asserted and the Power Manager waits for the external power supply to stabilize. If nVDD_FAULT is asserted after the external power supply timer expires, the processor returns to Sleep Mode.
- 3. If PCFR[OPDE] and OSCC[OON] were set when Sleep Mode started, the 3.6864 MHz oscillator is enabled and stabilizes. Otherwise, the 3.6864 MHz oscillator is already stable and this step is bypassed.
- 4. The processor's PLL clock generator is reprogrammed with the values in the CCCR and stabilizes.
- 5. The Sleep Mode configuration in PWRMODE[M] is cleared.
- 6. The processor's internal reset is deasserted and the CPU begins a normal boot sequence. When the normal boot sequence begins, all of the processor's units, except the RTC and portions of the Clocks and Power Manager and the Memory Controller, return to their predefined reset settings.
- 7. The nRESET_OUT pin is deasserted. This indicates that the processor is about to perform a fetch from the Reset vector.
- 8. Clear PSSR[PH] before accessing GPIOs, including chip selects that are muxed with GPIOs.
- 9. Clear PCFR[FS] and PCFR[FP] if either was set before Sleep Mode was triggered.
- 10. The SDRAM must transition out of self-refresh mode and into its idle state. See Section 6, "Memory Controller" for details on configuring the SDRAM interface.
- 11. Software must examine the RCSR, to determine what caused the reboot, and the Power Manager Sleep Status register (PSSR), to determine what triggered Sleep Mode.
- 12. If the PSPR was used to preserve any critical states during Sleep Mode, software can now recover the information.

If the nVDD_FAULT or nBATT_FAULT pin is asserted during the Sleep Mode exit sequence, the system re-enters Sleep Mode in the following sequence:

- 1. Regardless of the state of the IDAE bit:
 - All GPIO edge detects and the RTC alarm interrupt are cleared.
 - The Power Manager wake-up source registers (PWER, PRER, and PFER) are loaded with 0x0000 0003, their wake-up fault state. This limits the potential wake-up sources to a rising or falling edge on GPIO[0] or GPIO[1]. The wake-up fault state prevents spurious events from causing an unwanted wake-up while the battery is low or the power supply is at risk. The fault state is also the default state after a Hardware Reset.
- 2. The PLL clock generators are disabled.
- 3. If the OPDE bit in the PCFR is set and the OON bit in the OSCC is set, the 3.6864 MHz oscillator is disabled. If the oscillator is disabled, Sleep Mode consumes less power. If it is enabled, Sleep Mode exits more quickly.
- 4. An internal reset is generated to the core and most peripheral modules. This reset asserts the nRESET_OUT pin.
- 5. The PWR_EN pin is deasserted. If PMFW[FWAKE] is cleared, the system must respond by grounding the VCC and PLL_VCC power supplies to minimize power consumption.



3.4.10 Power Mode Summary

Table 3-4 shows the actions that occur when a Power Mode is entered. Table 3-5 shows the actions that occur when a Power Mode is exited. In the tables, an empty cell means that the power mode skips that step. Table 3-6 shows the expected behavior for power supplies in each power mode.

Table 3-4. Power Mode Entry Sequence Table

Step	Description of Action	Turbo	Run (from Turbo)	əlpi	Freq Change	Sleep	Fault ¹ Sleep
1	Software writes a bit in CP14	Х	х	Х	х	Х	
2	The CPU waits until all instructions to be completed	Х	Х	Х	Х	Х	
3	Wake up sources are cleared and limited to GP[1:0]						Х
4	The PM places GPIOs in their sleep states					х	Х
5	The Memory Controller finishes all outstanding transactions				х	х	Х
6	The Memory Controller places SDRAMs in self-refresh				х	х	Х
7	The PLL is disabled				х	х	Х
8	If OPDE and OOK bits are set, disable 3.6864 MHz oscillator					х	Х
9	Internal Reset to most modules. nRESET_OUT asserted					х	Х
10	PWR_EN is deasserted. Power is cut off					х	Х
11	Power to most I/O pins is cut off						

^{1:} Fault Sleep Mode starts if IDAE is clear and nBATT_FAULT or nVDD_FAULT is asserted.

Table 3-5. Power Mode Exit Sequence Table (Sheet 1 of 2)

Step	Description of Action	Turbo	Run (from Turbo)	ldle	Freq Change	Sleep	Fault ¹ Sleep
1	Wake up source or Interrupt is received			Х		Х	Х
2	Power to I/O pins restored						
3	PWR_EN is asserted					Х	х
4	External power ramp					Х	х
5	Enable 3.6864 MHz oscillator if OPDE and OOK are set					Х	х
6	Wait for 3.6864 MHz oscillator to stabilize if OPDE and OOK are set					х	х
7	Enable PLL with new frequency				х	Х	х
8	Wait for PLL stabilization				х	х	х
9	Wait for internal stabilization					х	х
10	Clear CP14 bit			х		х	



Table 3-5. Power Mode Exit Sequence Table (Sheet 2 of 2)

Step	Description of Action	Turbo	Run (from Turbo)	Idle	Freq Change	Sleep	Fault ¹ Sleep
11	Deassert nRESET_OUT					Х	Х
12	Restart CPU clocks, enable interrupts	Х	х	Х	Х	Х	Х

^{1:} Fault Sleep Mode starts if IDAE is clear and nBATT_FAULT or nVDD_FAULT is asserted.



Table 3-6. Power and Clock Supply Sources and States During Power Modes

						P	ower	Mod	le			
Module	Supp	ly Source	Tui	rbo	R	un	ld	le		eq inge	Sle	ер
	Pw	Ck	Pw	Ck	Pw	Ck	Pw	Ck	Pw	Ck	Pw	Ck
CPU, Caches, Buffers		Run/ Turbo (R/T)		Т		R		Off				
Memory Controller										changing		
LCD Controller	VCC	Mem	On		On		On		On	cha	Off	Off
DMA Controller				On		On		On				
General Periphs.		PLL										
OS timer		3.686								On		
Interrupts		MHz Osc										
Real Time Clock	VCC/ Reg	32.768	V	On	V	On	V	On	V	On	ı	On
Power Manager	(V/R)	kHz Osc	•	On	•	On	v	5	v	On		011
GP[3:0], PM pads, Osc pads	HV/ Batt (H/B)	Dynamic/ Static (D/S)	Н	D	Н	D	Н	D	Н	D	Н	S
General IO	Н											

KEY:

- T: Turbo clock
- R: Run clock
- V: Module powered off VCC.
- I: Module powered off internal regulator
- H: Module powered off VCCQ or VCCN
- D: Module is dynamic or actively clocked
- S: Module is static or clocks are gated.

3.5 Power Manager Registers

This section describes the 32-bit registers that control the Power Manager.



3.5.1 Power Manager Control Register (PMCR)

The PMCR is used to select the manner in which Sleep Mode is entered when the nVDD_FAULT or the nBATT_FAULT pin is asserted low. When the IDAE bit is set, an Imprecise Data Abort indication is sent to the CPU. The CPU then performs an abort routine. Software must ensure that the abort routine sets the Sleep Mode configuration in the PWRMODE register (see Section 3.7.2, "Power Mode Register (PWRMODE)"). The IDAE bit is cleared in any Reset and when Sleep Mode exits. Software may also clear the IDAE bit when necessary. The PMCR must be protected through Memory Management Unit (MMU) permissions.

Table 3-7. PMCR Bit Definitions

				0x4	0F0	0_0	000								PM	CR							Clo	cks	an	d P	owe	er M	ana	ger		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
																Reserved																IDAE
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	scr	ipti	on										
		[31	:1]			_	_			serv ad u	ed. Inde	fine	ed a	nd r	nus	t alv	vays	s be	wri	tten	with	n ze	roe	S.								
		C)			ID/	ΑE		0 -	Al as Fo	ise I low ssert orce nen d on	imn ed. imp nV[nedi orec DD_	ate ise (FAL	entr data JLT	y to ab or n	ort s ort s	sign	al to FAL	CF JLT	PU to	o all sser	ow ted	soft . Re	war	e to	ent ende	er s	leep ode	mo		



3.5.2 Power Manager General Configuration Register (PCFR)

The PCFR contains bits used to configure functions in the processor. When the OPDE bit is set, it allows the 3.6864 MHz oscillator to be disabled during Sleep Mode. The OPDE bit is cleared in Hardware, Watchdog, and GPIO Resets. The Float PCMCIA (FP) and Float Static Memory (FS) bits control the state of the PCMCIA control pins and the static memory control pins during Sleep Mode.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 3-8. PCFR Bit Definitions

				0x4	OF	0_001	С							PC	FR							Clo	cks	and	d Po	owe	r M	ana	ager		
Bit	31	30	29	28	27	26 2	5 2	4 :	23 22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														Reserved															FS	FP	OPDE
Reset	0	0	0	0	0	0 () ()	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Nam	е											De	esci	ripti	on										
		[31	1:3]			_			Reser Read			ed a	nd r	mus	t alv	vays	s be	wri	itten	wit	h ze	roe	S.								
	[31:3] — Read undefined and must always be written with zeroes. Float Static Chip Selects during Sleep Mode. 0 = Static Chip Select pins are not floated in Sleep Mode. nCS[5:1] are driver state of the appropriate PGSR register bits. nCS[1], nWE, and nOE are d 1 = Static Chip Select pins are floated in Sleep Mode. The pins nCS[5:0], nW nOE are affected. Cleared on Hardware, Watchdog, and GPIO Resets. Float PCMCIA controls during Sleep Mode. 0 = PCMCIA pins are not floated in Sleep Mode. They are driven to the state															drive	n hi	gh.													
		,	1			FP			0 = F a 1 = T S	PCM ppro he F leep ne st	CIA pria PCM Mo ate	pins Ite F ICIA Ide. of th	are PGS sig nPS ne a	e no R re nals SKT	t flo egist s: nF SEL ess	ated ter to OE and bus	d in bits. i, nF d nF dur	Slee PWE PRE ring	ep N E, nl :G a Sle	PIO re d ep N	W, r eriv Mod	PIC ed fi)R, a	and	nP(CE[2	2:1]	are		ted	
		(0			OPD	E		3.686 If the Config disabl 0 = E 11 = S	32.7 gura ed. o no	686 tion ot st the	kHz Reg op tl 3.68	cry giste he c	/sta er is oscil MH:	l is o 0, C lato z os	disa DPD r du cilla	bled E is ring tor	d be s ign g Sle duri	cau iore eep ng \$	d ar Mod Slee	nd th de. p M	ne 3. ode	.686						s not		

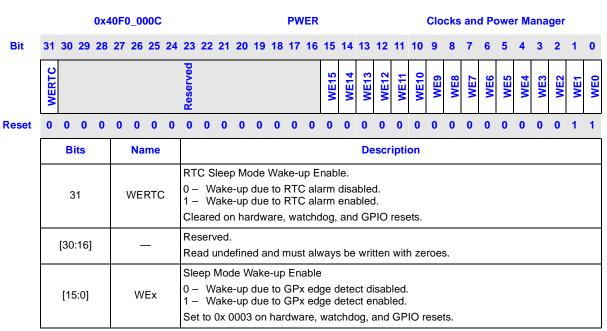


3.5.3 Power Manager Wake-Up Enable Register (PWER)

Table 3-9 shows the location of all wake up source enable bits in the Power Manager Wake-Up Enable Register (PWER). If a GPIO is to be used as a wake up source from Sleep, it must be programmed as an input in the GPDR and either one or both of the corresponding bits in the PRER and PFER must be set. When the IDAE bit is zero and a fault condition is detected on the nVDD_FAULT or nBATT_FAULT pin, PWER is set to 0x0000 0003 and only allows GP[1:0] as wake-up sources. When the IDAE bit is set, fault conditions on the nVDD_FAULT or nBATT_FAULT pins do not affect wake-up sources. PWER is also set to 0x0000 0003 in Hardware, Watchdog, or GPIO Resets.

Software should enable wakeups only for those GPIO pins that are configured as inputs during sleep. Any GPIO pins that are configured as outputs during sleep, should have their associated wake enable bits set to logic zero in all three PMU wake enable registers (PWER, PRER, and PFER).

Table 3-9. PWER Bit Definitions





3.5.4 Power Manager Rising-Edge Detect Enable Register (PRER)

The PRER, shown in Table 3-10, determines whether the GPIO pin enabled with the PWER register causes a wake up from sleep mode on that GPIO pin's rising edge. When PWER[IDAE] is zero and a fault condition is detected on the nVDD_FAULT or nBATT_FAULT pin, PRER is set to 0x0000_0003. This enables rising edges on GP[1:0] to act as wake up sources. When PWER[IDAE] is set, fault conditions on the nVDD_FAULT or nBATT_FAULT pins do not affect wake-up sources. PRER is also set to 0x0000_0003 in hardware, watchdog, and GPIO resets.

Software should enable wakeups only for those GPIO pins that are configured as inputs during sleep. Any GPIO pins that are configured as outputs during sleep, should have their associated wake enable bits set to logic zero in all three PMU wake enable registers (PWER, PRER, and PFER).

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 3-10. PRER Bit Definitions

				0x4	IOF	0_00	010								PR	ER							Clo	cks	and	d Po	owe	r M	ana	ger		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								Doynood									RE15	RE14	RE13	RE12	RE11	RE10	RE9	RE8	RE7	RE6	RE5	RE4	RE3	RE2	RE1	RE0
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1
		В	its			Na	me												De	scr	ipti	on										
		[31	:16]			_	_		Res Rea			efine	ed a	nd ı	mus	t alv	vays	be	writ	tten	with	n ze	roes	S.								
		[15	5:0]			RI	Ξx		Slee 0 – 1 – Set	W W	ake ake	up up	due due	to to	GPx GPx	risi risi	ng-e ng-e	edge edge	e de e de	tect	ena	able	d.	s.								



3.5.5 Power Manager Falling-Edge Detect Enable Register (PFER)

The PFER, Table 3-11, determines if the GPIO pin enabled with the PWER causes a wake up from sleep mode on that GPIO pin's falling edge. When PWER[IDAE] is zero and a fault condition is detected on the nVDD_FAULT or nBATT_FAULT pin, PFER is set to 0x0000_0003. This enables falling edges on GP[1:0] to act as wake up sources. When PWER[IDAE] is set, fault conditions on the nVDD_FAULT or nBATT_FAULT pins do not affect wake-up sources. PFER is also set to 0x0000_0003 during hardware, watchdog, and GPIO resets.

Software should enable wakeups only for those GPIO pins that are configured as inputs during sleep. Any GPIO pins that are configured as outputs during sleep, should have their associated wake enable bits set to logic zero in all three PMU wake enable registers (PWER, PRER, and PFER).

Table 3-11. PFER Bit Definitions

				0x4	IOF(0_0	014								PF	ER							Clo	cks	and	d Po	owe	r M	ana	ger		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								Reserved									FE15	FE14	FE13	FE12	FE11	FE10	FE9	FE8	FE7	FE6	FE5	FE4	FE3	FE2	FE1	FE0
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1
		В	its			Na	me												De	scr	ipti	on										
		[31	:16]			_	_		Res Rea			fine	d a	nd r	nus	t alv	vays	be	writ	ten	with	n ze	roes	S.								
		[15	5:0]			FI	Εx			W W	ake ake	up (due due	to to	GPx GPx	fall fall	ing- ing-	edg edg	e de e de	etec	t en	able	ed.	S.								



3.5.6 Power Manager GPIO Edge Detect Status Register (PEDR)

The PEDR, Table 3-12, indicates which of the GPIO pins enabled through the PWER, PRER, and PFER registers caused a wake up from sleep mode. The bits in PEDR can only be set on a rising or falling edge on a given GPIO pin. If PRER is set, the bits in PEDR can only be set on a rising edge. If PFER is set, the bits in PEDR can only be set on a falling edge. To reset a bit in PEDR to zero, write a 1 to it. The PEDR bits are reset to zero in hardware, watchdog, and GPIO resets.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 3-12. PEDR Bit Definitions

				0x4	10F(0_00	18								PE	DR							Clo	cks	and	d Po	owe	r M	ana	ger		
Bit	31	30	29	28	27	26	25	24	23 2	2 2	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								Received									ED15	ED14	ED13	ED12	ED11	ED10	ED9	ED8	ED7	ED6	ED5	ED4	ED3	ED2	ED1	ED0
Reset	0	0	0	0	0	0	0	0	0 0	(0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Naı	me												De	scr	iptio	on										
		[31	:16]			_	_		Rese Read			fine	d a	nd r	nus	t alv	vays	s be	writ	tten	with	ı ze	roes	S.								
		[15	5:0]			E	Эx		Sleep 0 - \ 1 - \ Clear	Val Val	ke ke	up (on (due	GPx to (not edge	det e on	tecte GP	x de				ets. (Clea	ared	by	writ	ing	a 1.				



3.5.7 Power Manager Sleep Status Register (PSSR)

The PSSR, shown in Table 3-13, contains the following status flags:

- Software Sleep Status (SSS) flag is set when the sleep mode configuration in the PWRMODE register is set and sleep mode starts (see Section 3.7.2).
- Battery Fault Status (BFS) bit is set after wake up any time the nBATT_FAULT pin is asserted (even when the processor is already in sleep mode).
- VDD Fault Status (VFS) bit is set after wake up when the nVDD_FAULT pin is asserted and
 causes the processor to enter sleep mode. The VFS bit is not set if software starts the sleep
 mode and then the nVDD FAULT pin is asserted.
- Peripheral Control Hold (PH) bit is set when sleep mode starts and indicates that the GPIO pins are retaining their sleep mode state values.
- Read Disable Hold (RDH) bit is set in hardware, GPIO, and watchdog resets and sleep mode. The RDH bit indicates that all the processor's GPIO input paths are disabled. To allow a GPIO input pin to be enabled, software must reset the RDH bit by writing a one to it. Clearing RDH also disables the 10 K to 60 K GPIO pullup resistors that are present during and after hardware, GPIO and watchdog reset. Sleep mode disables the GPIO input path, but the pullup resisters are not re-enabled in this case.

To clear a status flag write a 1 to it. Writing a 0 to a status bit has no effect. Hardware, watchdog, and GPIO resets clear or set the PSSR bits.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 3-13. PSSR Bit Definitions (Sheet 1 of 2)

				0×4	OE	00	10.4								DC	SR							Cla	oko	an	4 D	owe	- NA	ono	aar		
				UX-	юг		<i>7</i> 04								-	JON							Cio	CKS	alli	u F	owe	ı ıvı	ana	yeı		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												ı	ese	rve	d												RDH	ЬН	reserved	VFS	BFS	SSS
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0	0
		Bi	its			Na	me												De	escr	ipti	on										
		[31	:6]			_	_		rese	rve	ed																					
		Ę	5			RE	ЭН		Rea 0 - 1 -	G Re ha cle bu	PIO ecei ardw eare at be	pin iver vare ed b efor	s ar s of e, GI y th e th	e co all (PIO e pi ey a	GPI , or roce are	O pi wate sso usec	ns t chdo r aft l.	hat og r er tl	can eset he p	act t, int erip	as i terna hera	inpu al G al aı	ts a PIO nd C	re d pul SPIC	lisak II-up D int	oled s ar terfa	and re ad aces	ctive are	e. M e coi	ust I nfigu	oe ired	
		2	1			P	Н		Peri 0 – 1 –	G G M	PIO PIO ust onfig	pin pin be d gure	s ar s ar clea d bu	e core be red	onfigeing by to efor	hel the p	d in proc ey a	the ess ire a	ir sl or a actu	eep fter ally	mo the use	de s peri d by	tate phe the	e. Se eral i e pro	et wi nter	hen face ssor	slee es h	ave	bee		arts.	
		3	3			_	_		rese	rve	ed																					



Table 3-13. PSSR Bit Definitions (Sheet 2 of 2)

											•					•																	
				0x4	IOF(0_00	004								PS	SR							С	lock	(S a	and	l Po	owe	r M	ana	ger		
Bit	31	30	29	28	27	26	25	24	23 2	22	21	20	19	18	17	16	15	14	13	12	2 11	1) 9	8 (3	7	6	5	4	3	2	1	0
												re	sei	rve	d													RDH	ЬН	reserved	VFS	BFS	SSS
Reset	0	0	Description VDD Fault Status. O – nVDD_FAULT pin has not been asserted since it was last cleared by a rescPU.															0	0	0													
		Bi	Bits Name Description VDD Fault Status. 0 - nVDD_FAULT pin has not been asserted since it was last cleared by a reset																														
	VDD Fault Status. 0 - nVDD_FAULT pin has not been asserted since it was last cleared by a re CPU. 2 VFS 1 - nVDD_FAULT pin was asserted in Run or idle mode and caused the chip sleep mode; bit is set only after wake up. This bit is not set when nVDD_FAULT is asserted while in sleep mode. Cleared by hardware, watchdog, and GPIO resets.																																
		1	1			BF	-s		0 – 1 – This	nB CF nB bit	AT PU. AT cai	ult S T_FA T_FA n be r hare	AUL AUL set	T p T p wh	in h	as l nB/	oee	n as _FA	ser ULT	ted is	l; bit ass	is s erte	et o	only	aft	er v	wak	e u	p.	a re	eset	or tl	he
		()			SS	SS		0 – 1 –	So wa Ch	oftwa ns la nip v	Sleep are h ast cl was p hare	nas lear plac	not ed ced	en by a	a re slee	set (or th	ne C by	PL set	J. tting	the	sle		Ċ			e bit	sind	ce th	ne S	SS	

3.5.8 Power Manager Scratch Pad Register (PSPR)

The PM contains a 32-bit register that can be used to save processor configuration information in any desired format. The PSPR, shown in Table 3-14, is a holding register that is powered during sleep mode and is reset by hardware, watchdog, and GPIO resets. During run and turbo modes, any value can be written to PSPR. The value can be read after sleep mode exits. The value in PSPR can be used to represent the processor's configuration before sleep mode is invoked.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 3-14. PSPR Bit Definitions

Ia	DIE	3-	14.	PS	אר	B	IT L	еп	niti	ion	S																					
				0x4	40F(0_0	800								PS	PR							Clo	cks	an	d Po	owe	r M	ana	ger		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 SP																												
Reset	0	0	SP 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0															0	0													
		Bi	its			Na	me												De	escr	ipti	on										
		[31	:0]			S	Р		Scr 32-			-	res	erve	ed ir	ı sle	ep n	nod	e.													
									Cle	ared	by	har	dwa	re,	wate	chdo	og, a	and	GPI	O re	eset	s.										



3.5.9 Power Manager Fast Sleep Walk-up Configuration Register (PMFW)

The PSPR, shown in Table 3-15, provides a single bit called FWAKE which is used to select between the standard and fast sleep walk-up sequences. The PMFW register is reset by a hardware reset, GPIO reset, watchdog reset, but is not cleared by the sleep walk-up sequence. Using an exception handler to invoke sleep in response to a power fault event is advantageous because software can clear the PMFW[FWAKE] bit and configure the power management IC to use PWR_EN to disable the core power supply during sleep to minimize power consumption from a critically low battery. Also, the PCFR[OPDE] bit must be cleared to enable the 3.6864 MHz oscillator during sleep when fast sleep walk-up is selected by setting the PMFW[FWAKE] bit.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Power Manager Fast Sleep Wakeup 0x40F0 0034 Power Manager Configuration Register (PMFW) 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 Reset Reserved [31:3] Read undefined and must always be written with zeroes. FAST WAKEUP ENABLE 0 - Selects the standard sleep wakeup sequence with a 10 ms power supply stabilization delay when power is disabled during sleep. **FWAKE** [1] Selects the fast sleep wakeup sequence without a power supply stabilization delay when power is maintained during sleep. Cleared by hardware reset.

Read undefined and must always be written with zeroes.

Table 3-15. PMFW Register Bitmap and Bit Definitions

3.5.10 Power Manager GPIO Sleep State Registers (PGSR0, PGSR1, PGSR2)

PGSR0, PGSR1, and PGSR2, shown in Table 3-16, Table 3-17, and Table 3-18 allow software to select the output state of each GPIO pin when the processor goes into sleep mode. When a transition to sleep mode is required (through software or the nBATT_FAULT or nVDD_FAULT pin), the contents of the PGSR registers are loaded into the GPIO output data registers that software normally controls through the GPSR and GPCR registers. Only pins that are already configured as outputs reflect the new state. All bits in the output registers are loaded. When the processor reenters the run mode, these GPIO pins retain the programmed sleep state until software resets PSSR[PH]. If a pin is reconfigured from an input to an output, the register's last contents are driven onto the pin.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

[0]



Table 3-16. PGSR0 Bit Definitions

				0x4	0F0	0_00	20								PGS	R0						(Clo	cks	and	d Po	owe	r Ma	ana	ger		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	SS31	SS30	SS29	SS28	SS27	SS26	SS25	SS24	SS23	SS22	SS21	SS20	SS19	SS18	SS17	SS16	SS15	SS14	SS13	SS12	SS11	SS10	888	888	SS7	SS6	SS5	SS4	SS3	SS2	SS1	SS0
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	scri	iptio	on										
		[31	:0]			SS	Sx		0 – 1 –	Pi Pi	amr n is n is d by	driv driv	en t en t	o a	zero one	du du	ıring ring	sle	ep r ep n	mod node	e e											

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 3-17. PGSR1 Bit Definitions

				0x4	OF	0_00)24								PGS	R1						•	Clo	cks	and	d Po	owe	r Ma	ana	ger		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	SS63	SS62	SS61	SS60	SS59	SS58	SS57	SS56	SS55	SS54	SS53	SS52	SS51	SS50	SS49	SS48	SS47	SS46	SS45	SS44	SS43	SS42	SS41	SS40	SS39	SS38	SS37	SS36	SS35	SS34	SS33	SS32
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Naı	me												De	scr	iptio	on										
		[31	:0]			SS	Sx		0 – 1 –	Pii Pii	n is n is	driv driv	en t en t	o a	zero one	du du	ring ring	sle	ep i	mod nod												



Table 3-18. PGSR2 Bit Definitions

				0x4	l0F(0_0)28								PGS	SR2						(Clo	cks	and	l Po	we	r Ma	ana	ger		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					res	serv	ed					SS84	SS83	SS82	SS81	SS80	SS79	SS78	SS77	SS76	SS75	SS74	SS73	SS72	SS71	SS70	8869	SS68	SS67	SS66	SS65	SS64
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	scr	iptio	on										
		[31:	:17]			_	_		res	erve	ed																					
		[16	5:0]			SS	Зx		0 – 1 –	Pi Pi	n is n is	driv driv	en t	to a to a	zer	out, so du e dui tchd	ring ring	sle	ep i	mod	le e											

3.5.11 Reset Controller Status Register (RCSR)

The CPU uses the RCSR, shown in Table 3-19, to determine a reset's last cause or causes. The processor can be reset in four ways:

- Hardware reset
- Watchdog reset
- Sleep mode
- GPIO reset

Refer to Table 2-4, "Effect of Each Type of Reset on Internal Register State" on page 2-6 for details of the behavior of different modules during each type of reset.

Each RCSR status bit is set by a different reset source and can be cleared by writing a 1 back to the bit. The RCSR status bits for watchdog reset, sleep mode, and GPIO resets have a hardware reset state of zero.



Table 3-19. RCSR Bit Definitions

Idb		J- 1	٠. ١	101	5 1 \	Dit		C 1 11	IIIIO	113																					
				0x4	OF	0_00	30							RC	SF	₹						Clo	ocks	an	d P	owe	r M	ana	ger		
Bit	31	30	29	28	27	26	25	24	23 2	2 21	20	1	9 18	17	16	6 15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	erve	d													GPR	SMR	WDR	HWR
Reset	0	0	0	0	0	0	0	0	0	0	0	(0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
		Bi	its			Nan	ne											De	esci	ipti	on										
		[31	:4]			_			rese	ved																					
		3	3			GP	R		1 –	GPIC this b GPIC bit.	res it.	et	has	occı	ırre	ed sir	nce t	the	last	time	e th										
		2	2			SM	R		1 –	Sleep this b Sleep this b	o mo oit. o mo oit.	ode	e has	occ	urr	ed s	ince	the	e las	t tim	ne t										
		1	1			WD	R		1 –	Watc clear Watc this b	hdo ed tl hdo it.	g r his g r	reset s bit. reset	has	oc	curre	d si	nce	the	last	tim										ed
		()			HW	'R		Hard 0 – 1 – Set b	Hard Hard	ware	e r	eset eset	has	OC	curre	d si	nce	the	last	t tin				_						

3.6 Clocks Manager Registers

The Clocks Manager contains three registers:

- Core Clock Configuration Register (CCCR)
- Clock Enable Register (CKEN)
- Oscillator Configuration Register (OSCC)

3.6.1 Core Clock Configuration Register (CCCR)

The CCCR, shown in Table 3-20, controls the core clock frequency, from which the core, memory controller, LCD controller, and DMA controller frequencies are derived. The crystal frequency to memory frequency multiplier (L), memory frequency to run mode frequency multiplier (M), and run mode frequency to turbo mode frequency multiplier (N) are set in this register. The clock frequencies are shown below.



Memory frequency = 3.6864 MHz crystal freq. * crystal frequency to memory frequency multiplier (L)

Run mode frequency = Memory frequency * memory frequency to run mode frequency multiplier (M)

Turbo mode frequency = run mode frequency * run mode frequency to turbo mode frequency multiplier (N)

The value for L is chosen based on external memory or LCD requirements and can be constant while M and N change to allow run and turbo mode frequency changes without disrupting memory settings. The value for M is chosen based on bus bandwidth requirements and minimum core performance requirements. The value for N is chosen based on peak core performance requirements.

Table 3-20. CCCR Bit Definitions

				0x4	413 (0_0	000			Co	re Clo	ck C		nfigu CCCF		n R	egis	ster				C	loc	ks I	Man	age	r			
Bit	31	30	29	28	27	26	25	24	23 22	21	20 1	18	1	7 16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										rese	rved											N		N	1			L		
Reset	0	0	0	0	0	0	0	0	0 0	0	0 0	0	(0 0	0	0	0	0	0	0	0	1	0	0	1	0	0	0	0	1
		Bi	le			Na	me										D	000	riptio	on.										
		31:							reserve	<u> </u>								C 3C	ipti	<u> </u>										
		31.	10						Run M		F=0.0011		+-	Turk	- N	- d -	F=-	~	2011	1114	منامن									
												,						•	,	viuit	ipiie	: I								
									Turbo			. = r	tur	1 mod	ie ir	eque	ency	/ " IN												
									1 – 000																					
									001 – 1																					
		ro.	71						010 – 1		•																			
		[9:	/]				N		011 – N																					
									100 – 1		•	2																		
									101 – 1			^																		
									110 – N		•	3																		
									111 – r			.l			4			4_												
									Set to							•				:1:	_									
									Memor	-	•	-				Free	que	ncy	wuit	ipiie	r									
									Memor	•	•	rys	aı	Freq	. " L															
			-1						00 – re			(D																		
		[6:	5]			ı	M		01 – M			`				•	•		•			•			• /	,				
									10 – M			•					-						-			• /				
									11 – M			•					•		mes	the	me	mor	y fre	eque	ency	')				
									Set to							·			41 11											
									Crystal 00000		•	•	IVIE	HIOT	rre	que	псу	iviul	прпе	#1										
													,	11.10	- m :			:	- 00	E 2 1	41 I —		~ ^	600		11	ow. 15	١٨١		
									00001 00010		•		1 (iviem	огу	-req	luer	icy I	5 99	۱۱۵C.	viHZ	ror	11 3.	080	4 IVI	ΠZ (crys	ai)		
		۲4۰	∩1										c /	Mar	on, l	-ro~		· · · · ·	. 12	2 74	N/LI	- fra	·m ·	0 60	61 N	/ LI-	on.	ato!\		
		[4:	υJ				L		00011		•		о (iviem	ory I	-req	uen	icy is	5 13	2./1	IVIH	ZIIC	om 3	ა.ხგ	04 N	vi⊟Z	сгу	stai)		
									00100	– re	served																			

00110 to 11111 - reserved

Set to 00001 on hardware and watchdog resets.

00101 – Multiplier = 45 (Memory Frequency is 165.89MHz from 3.6864 MHz crystal)



3.6.2 Clock Enable Register (CKEN)

CKEN, shown in Table 3-21, enables or disables the clocks to most of the peripheral units. For lowest power consumption, the clock to any unit that is not being used must be disabled by writing a zero to the appropriate bit.

Table 3-21. CKEN Bit Definitions (Sheet 1 of 2)

				0x	413	0_00	004							СК	EN								С	loci	ks N	lan	age	r			
Bit	31	30	29	28	27	26	25	24	23 2	22 21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							res	erv	ed						CKEN16	reserved	CKEN14	CKEN13	CKEN12	CKEN11	reserved	reserved	CKEN8	CKEN7	CKEN6	CKEN5	reserved	CKEN3	CKEN2	CKEN1	CKENO
Reset	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	1	0	1	1	1	1	1	0	1	1	1	1	0	1	1	1	1
		В	its			Naı	me											De	scri	ptic	on										
		[31	:17]			_	_		rese	rved																					
									LCD	Unit	Clo	ck E	nab	le																	
		1	16			CKE	N16	6		Clock																					
										by ha							eset	s													
		1	15			_	_		rese	rved																					
									I2C	Unit C	Cloc	k En	able)																	
		1	14			CKE	N14	ļ	-	Clock																					
										by ha							eset	s													
									FICE	² Unit	Clo	ock E	nab	ole																	
		1	13			CKE	N13	3	-	Clock																					
									Thes	se bits	ar	e set	by	har	dwa	re re	set	or v	vatc	hdo	g re	set									
									MMC	C Uni	Cle	ock E	Enal	ole																	
		1	12			CKE	N12	2	-	Clock Clock																					
									Thes	se bits	ar	e set	by	har	dwa	re re	set	or v	vatc	hdo	g re	set									
									USB	Unit	Clo	ck E	nab	le																	
		1	11			CKE	N11		-	Clock Clock																					
										oy ha						•															
									This	bit m	ust	be s	et to	all	ow t	he 4	·8M	hz c	lock	ou	tput	on	GP7	7 Alt	tern	ate	Fun	ctio	n 1.		
		1	10			_	_		rese																						
										P Un						اء ما															
		9	9			CKE	EN9		-	Clock Clock																					
									Set I	oy ha	rdw	are a	and	wat	chdo	g re	eset	s													



Tab	le :	3- 2 ′	1. (CK	EN	Bit	t De	efii	nition	s (Shee	t	2 o	f 2	2)																
				0x4	4130	0_00	004							CK	EN								С	loci	ks I	/lan	age	r			
Bit	31	30	29	28	27	26	25	24	23 22	21	20 1	9	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															9	ed	14	13	12	=	ed	ed	8	4	16	22	ed	2	2	5	9
							res	erv	ed						CKEN16	reserved	CKEN14	CKEN13	CKEN12	CKEN11	reserved	eserved	CKEN8	CKEN7	CKEN6	CKEN5	reserved	CKEN3	CKEN2	CKEN1	CKEN0
Reset	_	0	0	0	0	0	0	0	0 0	_	0	<u> </u>	0	0	1	0	1	1	1	1	1	0	1	1	1	1	0	1	1	1	1
Reset					_			•	0 0		•						•	_			_		•	•			•	_	•	•	•
		Bit	ts			Na	me											De	escr	ipti	on										
		_							12S Ur 0 - C	-					lisabl	ed															
		8	3			CKE	EN8		1 – C	ock	to the	u	nit	is e	nabl	ed.															
									Set by							g re	eset	S													
									BTUAI 0 - C							ed															
		7	•			CKE	EN7		1 – C	lock	to the	u	nit	is e	nabl	ed.															
									These				<u> </u>			e re	eset	or	wato	chdo	og re	set									
									FFUAF 0 - C							ed															
		6	6			CKE	EN6		1 – C	ock	to the	u	nit	is e	nabl	ed.															
									Set by							g re	eset	s													
									STUAI 0 - C							ed															
		5	•			CKE	=N5		1 – C	ock	to the	u	nit	is e	nabl	ed.															
									Set by							g re	eset	s													
									HWUA							ed															
		4				CKE	EN4		1 – C	lock	to the	u	nit	is e	nabl	ed.															
									Set by						chdc	g re	eset	S													
		_							SSP U						lisabl	ed															
		3	3			CKE	=N3		1 – C	ock	to the	u	nit	is e	nabl	ed.															
									Set by						chdc	g re	eset	S													
		_							AC97						lisabl	ed															
		2	2			CKE	:N2		1 – C	lock	to the	u	nit	is e	nabl	ed.															
									Set by					wat	chdc	g re	eset	S													
						OK	-		PWM1 0 - C					is d	lisabl	ed															
		1				CKE	=IN 1		1 – C	lock	to the	u	nit	is e	nabl	ed.															
									Set by	nar	dware	a	nd \	wat	chdc	g re	eset	S													

CKEN0

0

PWM0 Clock Enable

0 - Clock to the unit is disabled1 - Clock to the unit is enabled.

Set by hardware and watchdog resets



3.6.3 Oscillator Configuration Register (OSCC)

The OSCC, shown in Table 3-22, controls the 32.768 kHz oscillator configuration. It contains two bits, the set-only 32.768 kHz OSCC[OON] and the read-only 32.768 kHz OSCC[OOK]. OSCC[OON] enables the external 32.768 kHz oscillator and can only be set by software. When the oscillator is enabled, it takes up to 10 seconds for to stabilize. When the oscillator is stabilized, the processor sets OSCC[OOK].

When OSCC[OOK] is set, the RTC and PM are clocked from the 32.768 KHz oscillator. Otherwise, the 3.6864 MHz oscillator is used. The OPDE bit, which allows the 3.6864 MHz oscillator to be disabled in sleep mode, is ignored (treated as if it were clear) if OSCC[OOK] is clear. OSCC[OOK] can only be reset by a hardware reset.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 3-22. OSCC Bit Definitions 0x4130_0008 OSCC **Clocks and Power Manager** Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 reserved Reset **Bits** Name **Description** [31:2] reserved 32.768 KHz OON (write-once only bit) 0 - 32.768 KHz oscillator is disabled. The 3.6864 MHz oscillator (divided by 112) clocks the RTC and PM. OON 32.768 KHz oscillator is enabled. OON can not be cleared once written except by hardware reset. Cleared by hardware reset. 32.768 kHz OOK (read-only bit)

0 - 32.768 KHz oscillator is disabled or not stable. The 3.6864 MHz oscillator (divided

32.768 KHz oscillator has been enabled (OON=1) and stabilized. It will clock the

3.7 Coprocessor 14: Clock and Power Management

by 112) clocks the RTC and PM.

RTC and PM.
Cleared by hardware reset.

Coprocessor 14 contains two registers that control the power modes and sequences:

• CP14 register 6 – CCLKCFG register

OOK

0

• CP14 register 7 – PWRMODE register



Table 3-23. Coprocessor 14 Clock and Power Management Summary

Function	Data in Rd	Instruction
Read CCLKCFG	_	MRC p14, 0, Rd, c6, c0, 0
Enter turbo mode	TURBO = 1	MCR p14, 0, Rd, c6, c0, 0
Enter frequency change sequence	FCS = 1 (Turbo mode bit may be set or cleared in the same write)	MCR p14, 0, Rd, c6, c0, 0
Enter idle mode	M = 1	MCR p14, 0, Rd, c7, c0, 0
Enter sleep mode	M = 3	MCR p14, 0, Rd, c7, c0, 0

3.7.1 Core Clock Configuration Register (CCLKCFG)

The CCLKCFG register (CP14 register 6), shown in Table 3-24, is used to enter the turbo mode and frequency change sequence. To enter the mode or sequence, software executes the appropriate function from Table 3-23. All core-initiated memory requests are completed before the Clocks and Power Manager initiates the desired mode or sequence.

To ensure that CCLKCFG[TURBO] does not change when entering the frequency change sequence, software must do a read-modify-write.

Table 3-24. CCLKCFG Bit Definitions

				R	CP egis		6							С	CL	(CF	G						Clo	cks	an	d P	owe	r M	ana	ger		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														r	ese	rve	d														FCS	TURBO
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bit	s			Naı	me												De	esci	ipti	on										
		[31:	2]			_	-		res	erve	ed																					
		1				FC	s		0 – 1 –	Do Er	o no nter	Char ot en freq n ha	ter Juer	freq	uen cha	icy o	char sec	luen	ice .			et ar	nd v	vher	n sle	еер	mod	de e	xits	-		
		0			-	TUR	ВО		1 –	Do Er	o no	le ot en turb n ha	o m	ode	Э							et aı	nd v	vher	n sle	еер	mod	de e	xits			



3.7.2 Power Mode Register (PWRMODE)

The PWRMODE register (CP14, register 7), shown in Table 3-25, is used to enter idle and sleep modes. To select a mode, software writes to PWRMODE[M]. All core-initiated memory requests are completed before the Clocks and Power Manager initiates the desired mode.

Table 3-25. PWRMODE Bit Definitions

				R	_	14 ster	7							P۱	ΝF	RMO	DE									CF	214					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	7 16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														r	es	erve	ed														ı	/ 1
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												D	esci	ipti	on										
		[31	:2]			_			res	erve	ed																					
		[1:	0]			N	Л		01 10 11	– Ri – Id – re – SI	un/ti le m serv eep	urbo node /ed mo	o mo																			

3.8 External Hardware Considerations

The Clocks and Power Manager controls the timing in and out of resets and the voltage ramp and stabilization. As a result, the hardware that is used with the processor must meet certain requirements to operate properly. This section describes those requirements.

3.8.1 Power-On-Reset Considerations

The nRESET and nTRST pins must be held low while the power supplies initialize and for a fixed time after power is stable. This can be controlled with an external Power-On-Reset device or another circuit.

To ensure that the internal ESD protection devices do not activate during power up, a minimum rise time must be observed. Refer to the *Intel® PXA255 Processor Electrical, Mechanical, and Thermal Specification* for details.

3.8.2 Power Supply Connectivity

The processor requires two or three externally-supplied voltage levels. VCCQ requires 3.3 V (+/-10%), VCCN requires 3.3 V (+/-10%) or 2.5 V (+15/-5%), and VCC and PLL_VCC shall be connected together and require .85-1.3 V. PLL_VCC must be separated from other low voltage supplies. Depending on the availability of independent regulator outputs and the desired memory voltage, VCCQ may have to be separated from VCCN.



3.8.3 Driving the Crystal Pins from an External Clock Source

The information in this section is provided as a guideline. The electrical specifications for the crystal oscillator pins are in *Intel® PXA255 Processor Electrical, Mechanical, and Thermal Specification*.

A 3.6864 MHz crystal must be connected between the PXTAL and PEXTAL pins. A 32.768 kHz crystal is normally connected between the TXTAL and TEXTAL pins. This configuration gives the lowest overall power consumption because the crystal's resonant nature provides better power efficiency than an external source that drives the crystal pins. Some applications have other clock sources of the same frequency and can reduce overall cost by driving the crystal pins externally.

Refer to the Oscillator Electrical Specifications in the *Intel® PXA255 Processor Design Guide* for more information.

Note: No external capacitors are required.

3.8.4 Noise Coupling Between Driven Crystal Pins and a Crystal Oscillator

The two pairs of crystal pins are located near each other on the processor. When crystal oscillators are connected to the pins, this proximity leads to low signal swings and slow edges that result in limited noise coupling between the pins. If one of the crystal oscillators is replaced by an independent signal source and the other is not, the noise coupling may increase. To limit this effect, reduce the slew rate on the pins driven by the independent source.

3.9 Clocks and Power Manager Register Summary

3.9.1 Clocks Manager Register Locations

Table 3-26 shows the registers associated with the clocks manager and the physical addresses used to access them.

Table 3-26. Clocks Manager Register Summary

Address	Name	Description
0x4130_0000	CCCR	Core Clock Configuration Register
0x4130_0004	CKEN	Clock Enable Register
0x4130_0008	OSCC	Oscillator Configuration Register

3.9.2 Power Manager Register Summary

Table 3-27 shows the registers associated with the PM and the physical addresses used to access them.



Table 3-27. Power Manager Register Summary

		-
Address	Name	Description
0x40F0_0000	PMCR	Power Manager Control register
0x40F0_0004	PSSR	Power Manager Sleep Status register
0x40F0_0008	PSPR	Power Manager Scratch Pad register
0x40F0_000C	PWER	Power Manager Wake-up Enable register
0x40F0_0010	PRER	Power Manager GPIO Rising-edge Detect Enable register
0x40F0_0014	PFER	Power Manager GPIO Falling-edge Detect Enable register
0x40F0_0018	PEDR	Power Manager GPIO Edge Detect Status register
0x40F0_001C	PCFR	Power Manager General Configuration register
0x40F0_0020	PGSR0	Power Manager GPIO Sleep State register for GP[31-0]
0x40F0_0024	PGSR1	Power Manager GPIO Sleep State register for GP[63-32]
0x40F0_0028	PGSR2	Power Manager GPIO Sleep State register for GP[84-64]
0x40F0_0030	RCSR	Reset controller status register



This chapter describes the System Integration Unit (SIU) for the PXA255 processor. The SIU controls several processor-wide system functions. The units contained in the SIU are:

- General-purpose I/O (GPIO) ports
- Interrupt controller
- Real-time clock (RTC)
- Operating system timer (OS timer)
- Pulse Width modulator

General-Purpose I/O 4.1

The PXA255 processor enables and controls its 81 GPIO pins through the use of 27 registers which configure the pin direction (input or output), pin function, pin state (outputs only), pin level detection (inputs only), and selection of alternate functions. A portion of the GPIOs can be used to bring the processor out of Sleep mode. Take care when choosing which GPIO pin is assigned as a GPIO function because many of the GPIO pins have alternate functions and can be configured to support processor peripherals.

Configure all unused GPIOs as outputs to minimize power consumption.

4.1.1 **GPIO Operation**

The PXA255 processor provides 81 GPIO pins for use in generating and capturing applicationspecific input and output signals. Each pin can be programmed as either an input or output. When programmed to be an input, a GPIO can also serve as an interrupt source. All 81 pins are configured as inputs during the assertion of all resets, and remain as inputs until they are configured otherwise.

Use the GPIO Pin Direction Register (GPDR) to set whether the GPIO pins are outputs or inputs. When programmed as an output, the pin can be set high by writing to the GPIO Pin Output Set Register (GPSR) and cleared low by writing to the GPIO Pin Output Clear Register (GPCR). The set and clear registers can be written to regardless of whether the pin is configured as an input or an output. If a pin is configured as an input, the programmed output state occurs when the pin is reconfigured to be an output.

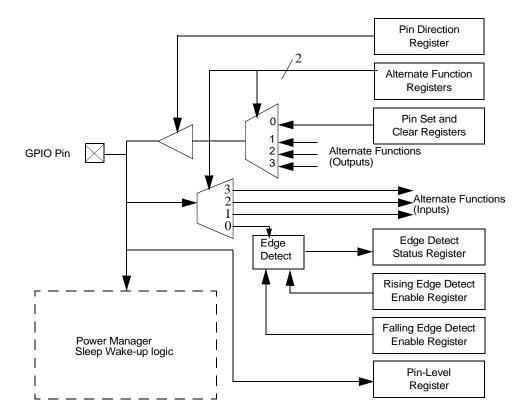
Validate each GPIO pin's state by reading the GPIO Pin Level Register (GPLR). You can read this register any time to confirm the state of a pin. In addition, use the GPIO Rising Edge Detect Enable Register (GRER) and GPIO Falling Edge Detect Enable Register (GFER) to detect either a rising edge or falling edge on each GPIO pin. Use the GPIO Edge Detect Status register (GEDR) to read edge detect state. These edge detects can be programmed to generate interrupts (see Section 4.2). Also use GPIO[15:0] to generate wake-up events that bring the PXA255 processor out of sleep mode (refer to Section 3.4.9.5, "Exiting Sleep Mode" on page 3-18).



When the processor enters sleep mode, the contents of the Power Manager Sleep State registers (PGSR0, PGSR1 and PGSR2) are loaded into the output data registers. If the particular pin is programmed as an output, then the value in the PGSR is driven onto the pin before entering sleep mode. When the processor exits sleep mode, these values remain driven until the GPIO pins are reprogrammed by writing to the GPDR, GPSR or GPCR, and setting the GPIO bit in the Power Manager Sleep Status register (PSSR) to indicate that the GPIO registers have been re-initialized after sleep mode. This is necessary since the GPIO logic loses power during sleep mode

Most GPIO pins can also serve an alternate function within the processor. Certain modes within the serial controllers and LCD controller require extra pins. These functions are hardwired into specific GPIO pins and their use is described in the following paragraphs. Even though a GPIO pin is used for an alternate function, you must still program the proper direction of that pin through the GPDR. Details on alternate functions are provided in Section 4.1.2. Figure 4-1 shows a block diagram of a single GPIO pin.

Figure 4-1. General-Purpose I/O Block Diagram



4.1.2 **GPIO Alternate Functions**

GPIO pins are capable of having as many as six alternate functions that can be set to enable additional functionality within the processor. If a GPIO is used for an alternate function, then it cannot be used as a GPIO at the same time. GPIO[0] is reserved because of its special use during sleep mode and is not available for alternate functions. GPIO[15:0] are used for walk-up from sleep mode. The walk-up functionality is described in Section 3.4.9.5, "Exiting Sleep Mode" on page 3-18. Table 4-1 shows each GPIO pin and its corresponding alternate functions.



For more information on alternate functions, refer to the Source Unit column in Table 4-1 for the appropriate section of this document.

Table 4-1. GPIO Alternate Functions (Sheet 1 of 4)

Pin	Alternate Function Name	Alternate Function Assignment	AF{n} encoding	Source Unit	Signal Description and comments
GP1	GP_RST	ALT_FN_1_IN	01	Clocks & Power Manager Unit	Active low GP_reset
GP6	MMCCLK	ALT_FN_1_OUT	01	Multimedia Card (MMC) Controller	MMC Clock
GP7	48 MHz clock [†]	ALT_FN_1_OUT	01	Clocks & Power Manager Unit	48 MHz clock output
GP8	MMCCS0	ALT_FN_1_OUT	01	Multimedia Card	MMC Chip Select 0
GP9	MMCCS1	ALT_FN_1_OUT	01	(MMC) Controller	MMC Chip Select 1
GP10	RTCCLK	ALT_FN_1_OUT	01	System Integation Unit	real time clock (1 Hz)
GP11	3.6 MHz	ALT_FN_1_OUT	01	Clocks & Power	3.6 MHz oscillator out
GP12	32 kHz	ALT_FN_1_OUT	01	Manager Unit	32 kHz out
GP13	MBGNT	ALT_FN_2_OUT	10		memory controller grant
GP14	MBREQ	ALT_FN_1_IN	01		memory controller alternate bus master request
GP15	nCS_1	ALT_FN_2_OUT	10	Memory Controller	Active low chip select 1
GP16	PWM0	ALT_FN_2_OUT	10		PWM0 output
GP17	PWM1	ALT_FN_2_OUT	10		PWM1 output
GP18	RDY	ALT_FN_1_IN	01		Ext. Bus Ready
GP19	DREQ[1]	ALT_FN_1_IN	01		External DMA Request
GP20	DREQ[0]	ALT_FN_1_IN	01		External DMA Request
GP23	SCLK	ALT_FN_2_OUT	10		SSP clock
GP24	SFRM	ALT_FN_2_OUT	10	SSP Serial Port	SSP Frame
GP25	TXD	ALT_FN_2_OUT	10		SSP transmit
GP26	RXD	ALT_FN_1_IN	01		SSP receive
GP27	EXTCLK	ALT_FN_1_IN	01		SSP ext_clk
GP28	BITCLK	ALT_FN_1_IN	01	AC97 Controller Unit	AC97 bit_clk
	BITCLK	ALT_FN_2_IN	10	- I2S Controller	I2S bit_clk
	BITCLK	ALT_FN_1_OUT	01		I2S bit_clk
GP29	SDATA_IN0	ALT_FN_1_IN	01	AC97 Controller Unit	AC97 Sdata_in0
	SDATA_IN	ALT_FN_2_IN	10	I2S Controller	I2S Sdata_in
GP30	SDATA_OUT	ALT_FN_1_OUT	01	I2S Controller	I2S Sdata_out
	SDATA_OUT	ALT_FN_2_OUT	10	AC97 Controller Unit	AC97 Sdata_out
GP31	SYNC	ALT_FN_1_OUT	01	I2S Controller	I2S sync
	SYNC	ALT_FN_2_OUT	10	AC97 Controller Unit	AC97 sync
GP32	SDATA_IN1	ALT_FN_1_IN	01	AC97 Controller Unit	AC97 Sdata_in1
	SYSCLK	ALT_FN_1_OUT	01	I2S Controller	I2S System Clock



Table 4-1. GPIO Alternate Functions (Sheet 2 of 4)

GP41 RTS ALT_FN_2_OUT 10 UARTS FFUART request to send GP42 BTRXD ALT_FN_1_IN 01 UARTS BTUART receive data HWRXD ALT_FN_3_IN 11 HWUART HWUART receive data GP43 BTTXD ALT_FN_2_OUT 10 UARTS BTUART transmit data HWTXD ALT_FN_3_OUT 11 HWUART HWUART transmit data HWTXD ALT_FN_3_OUT 11 HWUART HWUART clear to send HWCTS ALT_FN_3_IN 11 HWUART HWUART clear to send HWCTS ALT_FN_3_IN 11 HWUART HWUART request to send HWRTS ALT_FN_3_OUT 10 UARTS BTUART request to send HWRTS ALT_FN_3_OUT 11 HWUART HWUART request to send GP46 ICP_RXD ALT_FN_1_IN 01 UARTS STD_UART request to send GP47 TXD ALT_FN_2_IN 10 UARTS STD_UART receive data GP47 TXD ALT_FN_2_OUT		Iternate unction Name	Alternate Function Assignment	AF{n} encoding	Source Unit	Signal Description and comments
MMCCS0	GP33 n0	CS[5]	ALT_FN_2_OUT	10	Memory Controller	Active low chip select 5
MMCCS0	FI	FRXD	ALT_FN_1_IN	01	UARTs	FFUART receive
GP36 DCD ALT_FN_1_IN 01 UARTS FFUART Data carrier detect FFUART data set ready FFUART data set ready FFUART data set ready FFUART Ring Indicator FFUART data set ready FFUART Ring Indicator FFUART Ring Indicator MMC Chip Select 1 MMC Chi		IMCCS0	ALT_FN_2_OUT	10		MMC Chip Select 0
DSR	GP35 C	TS	ALT_FN_1_IN	01		FFUART Clear to send
GP36 RI	GP36 D	OCD	ALT_FN_1_IN	01	UARTs	FFUART Data carrier detect
MMCCS1	GP37 D	SR	ALT_FN_1_IN	01		FFUART data set ready
MMCCS1	GP38 R	N.	ALT_FN_1_IN	01		FFUART Ring Indicator
ALT_FN_2_OUT 10	GP39	IMCCS1	ALT_FN_1_OUT	01		MMC Chip Select 1
BTRXD	FI	FTXD	ALT_FN_2_OUT	10	UARTs	FFUART transmit data
GP41 RTS ALT_FN_2_OUT 10 FFUART request to send GP42 BTRXD ALT_FN_1_IN 01 UARTS BTUART receive data HWRXD ALT_FN_3_IN 11 HWUART HWUART receive data BTXD ALT_FN_2_OUT 10 UARTS BTUART transmit data HWTXD ALT_FN_3_OUT 11 HWUART HWUART transmit data BTCTS ALT_FN_1_IN 01 UARTS BTUART clear to send HWCTS ALT_FN_3_IN 11 HWUART HWUART clear to send HWRTS ALT_FN_2_OUT 10 UARTS BTUART request to send HWRTS ALT_FN_3_OUT 11 HWUART HWUART request to send GP46 ICP_RXD ALT_FN_1_IN 01 UARTS STD_UART receive data GP47 TXD ALT_FN_2_IN 10 UARTS STD_UART transmit data GP47 ICP_TXD ALT_FN_2_OUT 10 Memory Controller Output Enable for Card Space GP48 IWTXD ALT_FN_2_OUT 1	GP40 D	TR	ALT_FN_2_OUT	10	UARTs	FFUART data terminal Ready
HWRXD	GP41 R	RTS	ALT_FN_2_OUT	10		FFUART request to send
HWRXD		TRXD	ALT_FN_1_IN	01	UARTs	BTUART receive data
HWTXD		IWRXD	ALT_FN_3_IN	11	HWUART	HWUART receive data
HWTXD		TTXD	ALT_FN_2_OUT	10	UARTs	BTUART transmit data
HWCTS		IWTXD	ALT_FN_3_OUT	11	HWUART	HWUART transmit data
HWCTS		TCTS	ALT_FN_1_IN	01	UARTs	BTUART clear to send
HWRTS		IWCTS	ALT_FN_3_IN	11	HWUART	HWUART clear to send
HWRTS		TRTS	ALT_FN_2_OUT	10	UARTs	BTUART request to send
CP_RXD		IWRTS	ALT_FN_3_OUT	11	HWUART	HWUART request to send
TXD	GP46	CP_RXD	ALT_FN_1_IN	01		ICP receive data
ICP_TXD	R	XD	ALT_FN_2_IN	10	UARTs	STD_UART receive data
ICP_TXD	T	XD	ALT_FN_1_OUT	01	UARTs	STD_UART transmit data
GP48 HWTXD ALT_FN_1_OUT 01 HWUART HWUART transmit data GP49 nPWE ALT_FN_2_OUT 10 Memory Controller Write Enable for Card Space HWRXD ALT_FN_1_IN 01 HWUART HWUART receive data GP50 nPIOR ALT_FN_2_OUT 10 Memory Controller I/O Read for Card Space HWCTS ALT_FN_1_IN 01 HWUART HWUART clear to send GP51 nPIOW ALT_FN_2_OUT 10 Memory Controller I/O Write for Card Space HWRTS ALT_FN_1 OUT 01 HWUART HWUART request to send GP52 nPCE[1] ALT_FN_2_OUT 10 Memory Controller Card Enable for Card Space GP53 nPCE[2] ALT_FN_2_OUT 10 Memory Controller Card Enable for Card Space		CP_TXD	ALT_FN_2_OUT	10		ICP transmit data
HWTXD		POE	ALT_FN_2_OUT	10	Memory Controller	Output Enable for Card Space
HWRXD	H'	IWTXD	ALT_FN_1_OUT	01	HWUART	HWUART transmit data
HWRXD	CP40	PWE	ALT_FN_2_OUT	10	Memory Controller	Write Enable for Card Space
GP50 HWCTS ALT_FN_1_IN 01 HWUART HWUART clear to send GP51 nPIOW ALT_FN_2_OUT 10 Memory Controller I/O Write for Card Space HWRTS ALT_FN_1 OUT 01 HWUART HWUART request to send GP52 nPCE[1] ALT_FN_2_OUT 10 Memory Controller Card Enable for Card Space GP53 nPCE[2] ALT_FN_2_OUT 10 Card Enable for Card Space		IWRXD	ALT_FN_1_IN	01	HWUART	HWUART receive data
HWCTS	ODEO nf	PIOR	ALT_FN_2_OUT	10	Memory Controller	I/O Read for Card Space
HWRTS ALT_FN_1 OUT 01 HWUART HWUART request to send GP52 nPCE[1] ALT_FN_2_OUT 10 GP53 nPCE[2] ALT_FN_2_OUT 10 Memory Controller Card Enable for Card Space Card Enable for Card Space	Grou H	IWCTS	ALT_FN_1_IN	01	HWUART	HWUART clear to send
HWRTS ALT_FN_1 OUT 01 HWUART HWUART request to send GP52 nPCE[1] ALT_FN_2_OUT 10 GP53 nPCE[2] ALT_FN_2_OUT 10 Memory Controller Card Enable for Card Space Card Enable for Card Space	ODE1	PIOW	ALT_FN_2_OUT	10	Memory Controller	I/O Write for Card Space
GP53 nPCE[2] ALT_FN_2_OUT 10 Memory Controller Card Enable for Card Space	GP51 H	IWRTS	ALT_FN_1 OUT	01	HWUART	HWUART request to send
GP53 nPCE[2] ALT_FN_2_OUT 10 Card Enable for Card Space	GP52 nf	PCE[1]	ALT_FN_2_OUT	10	Memory Controller	Card Enable for Card Space
	GP53 nf	PCE[2]	ALT_FN_2_OUT	10		Card Enable for Card Space
GP53 MMCCLK ALT_FN_1_OUT 01 Multimedia Card (MMC) Controller MMC Clock	GP53 M	IMCCLK	ALT_FN_1_OUT	01	Multimedia Card (MMC) Controller	MMC Clock



Table 4-1. GPIO Alternate Functions (Sheet 3 of 4)

Pin	Alternate Function Name	Alternate Function Assignment	AF{n} encoding	Source Unit	Signal Description and comments
GP54	MMCCLK	ALT_FN_1_OUT	01	Multimedia Card (MMC) Controller	MMC Clock
GP54	nPSKTSEL	ALT_FN_2_OUT	10	Memory Controller	Socket Select for Card Space
GP55	nPREG	ALT_FN_2_OUT	10	Marray Cantrallar	Card Address bit 26
GP56	nPWAIT	ALT_FN_1_IN	01	Memory Controller	Wait signal for Card Space
GP57	nIOIS16	ALT_FN_1_IN	01	Memory Controller	Bus Width select for I/O Card Space
GP58	LDD[0]	ALT_FN_2_OUT	10		LCD data pin 0
GP59	LDD[1]	ALT_FN_2_OUT	10		LCD data pin 1
GP60	LDD[2]	ALT_FN_2_OUT	10		LCD data pin 2
GP61	LDD[3]	ALT_FN_2_OUT	10		LCD data pin 3
GP62	LDD[4]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 4
GP63	LDD[5]	ALT_FN_2_OUT	10		LCD data pin 5
GP64	LDD[6]	ALT_FN_2_OUT	10		LCD data pin 6
GP65	LDD[7]	ALT_FN_2_OUT	10	1	LCD data pin 7
	LDD[8]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 8
GP66	MBREQ	ALT_FN_1_IN	01	Memory Controller	memory controller alternate bus master req
	LDD[9]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 9
GP67	MMCCS0	ALT_FN_1_OUT	01	Multimedia Card (MMC) Controller	MMC Chip Select 0
GP68	MMCCS1	ALT_FN_1_OUT	01	Multimedia Card (MMC) Controller	MMC Chip Select 1
	LDD[10]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 10
GP69	MMCCLK	ALT_FN_1_OUT	01	Multimedia Card (MMC) Controller	MMC_CLK
	LDD[11]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 11
GP70	RTCCLK	ALT_FN_1_OUT	01	System Integation Unit	Real Time clock (1 Hz)
	LDD[12]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 12
GP71	3.6 MHz	ALT_FN_1_OUT	01	Clocks & Power Manager Unit	3.6 MHz Oscillator clock
	LDD[13]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 13
GP72	32 kHz	ALT_FN_1_OUT	01	Clocks & Power Manager Unit	32 kHz clock
	LDD[14]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 14
CD72	LDD[15]	ALT_FN_2_OUT	10	LCD Controller	LCD data pin 15
GP73	MBGNT	ALT_FN_1_OUT	01	Memory Controller	Memory controller grant



Table 4-1. GPIO Alternate Functions (Sheet 4 of 4)

Pin	Alternate Function Name	Alternate Function Assignment	AF{n} encoding	Source Unit	Signal Description and comments
GP74	LCD_FCLK	ALT_FN_2_OUT	10		LCD Frame clock
GP75	LCD_LCLK	ALT_FN_2_OUT	10	LCD Controller	LCD line clock
GP76	LCD_PCLK	ALT_FN_2_OUT	10	LCD Controller	LCD Pixel clock
GP77	LCD_ACBIAS	ALT_FN_2_OUT	10		LCD AC Bias
GP78	nCS[2]	ALT_FN_2_OUT	10	Memory Controller	Active low chip select 2
GP79	nCS[3]	ALT_FN_2_OUT	10	Memory Controller	Active low chip select 3
GP80	nCS[4]	ALT_FN_2_OUT	10	Memory Controller	Active low chip select 4
GP81	NSSPSCLK	ALT_FN_1_IN	01		NSSP Serial clock is input
GP61	NSSPSCLK	ALT_FN_1_OUT	01		NSSP Serial clock is output
CDOO	NSSPSFRM	ALT_FN_1_IN	01		NSSP frame is input
GP82	NSSPSFRM	ALT_FN_1_OUT	01	Natural CCD	NSSP frame is output
CDOO	NSSPTXD	ALT_FN_1_OUT	01	Network SSP	NSSP transmit
GP83	NSSPRXD	ALT_FN_2_IN	10	1	NSSP receive
CD04	NSSPTXD	ALT_FN_1_OUT	01	7	NSSP transmit
GP84	NSSPRXD	ALT_FN_2_IN	10	1	NSSP receive

[†] CKEN[11] - USB Unit Clock Enable bit must be enabled to allow the 48 MHz clock output on GP7

4.1.3 **GPIO** Register Definitions

There are twenty-seven 32-bit registers within the GPIO control block. There are nine distinct register functions and there are three sets of each of the nine registers to serve the 81 GPIOs. The various functions of the nine registers corresponding to each GPIO pin are described here:

- Three monitor pin state (GPLR)
- Six control output pin state (GPSR, GPCR)
- Three control pin direction (GPDR)
- Six control whether rising edges and/or falling edges are detected (GRER & GFER)
- Three indicate when specified edge types have been detected on pins (GEDR).
- Six determine whether a pin is used as a normal GPIO or whether it is to be taken over by one of three possible alternate functions (GAFR_L, GAFR_U).

Table 4-2. GPIO Register Definitions (Sheet 1 of 2)

Register Type	Register Function	GPIO[15:0]	GPIO[31:16]	GPIO[47:32]	GPIO[63:48]	GPIO[79:64]	GPIO[80]
GPLR	Monitor Pin State	GPI	LR0	GPI	LR1	GPI	LR2
GPSR	Control Output	GPS	SR0	GPS	SR1	GPS	SR2
GPCR	Pin State	GPO	CR0	GPO	CR1	GPO	CR2
GPDR	Set Pin Direction	GPI	DR0	GPI	DR1	GPI	DR2



Table 4-2. GPIO Register Definitions (Sheet 2 of 2)

Register Type	Register Function	GPIO[15:0]	GPIO[31:16]	GPIO[47:32]	GPIO[63:48]	GPIO[79:64]	GPIO[80]
GRER	Detect Rising/Falling	GRI	ER0	GRI	ER1	GRI	ER2
GFER	Edge	GF	ER0	GF	ER1	GF	R2
GEDR	Detect Edge Type	GEI	OR0	GEI	DR1	GEI	DR2
GAFR	Set Alternate Functions	GAFR0_L	GAFR0_U	GAFR1_L	GAFR1_U	GAFR2_L	GAFR2_U

NOTE: For the alternate function registers, the designator _L signifies that the lower 16 GPIOs' alternate functions are configured by that register and _U designates that the upper 16 GPIOs' alternate functions are configured by that register.

Note: GPLR2[31:17], GPSR2[31:17], GPCR2[31:17], GPDR2[31:17], GRER2[31:17], GFER2[31:17], GEDR2[31:17] and GAFR2_U[31:2] are reserved bits. Write zeros to these bits and ignore all reads from these bits.

Note: All GPIO registers are initialized to 0x0 at reset, which results in all GPIO pins being initialized as inputs.

4.1.3.1 GPIO Pin-Level Registers (GPLR0, GPLR1, GPLR2)

Check the state of each of the GPIO pins by reading the GPIO Pin Level register (GPLR). Each bit in the GPLR corresponds to one pin in the GPIO. GPLR0[31:0] correspond to GPIO[31:0], GPLR1[31:0] correspond to GPIO[63:32] and GPLR2[16:0] correspond to GPIO[80:64]. Use the GPLR0—2 read-only registers to determine the current value of a particular pin (regardless of the programmed pin direction). For reserved bits (GPLR2[31:17]), reads return zero.

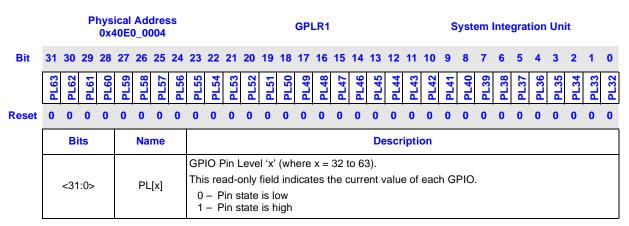
Table 4-3. GPLR0 Bit Definitions

				nysi 0x4				s							GPL	.R0							S	yste	em I	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
																PL1	PL0															
Reset																0	0															
		Bi	ts			Naı	me												De	scr	iptio	on										
		<31	:0>			PL	.[x]		-	s rea		nly tate	field is lo	ind w				,		alue	e of	eacl	n GI	PIO.	-							

This is read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 4-4. GPLR1 Bit Definitions



This is read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 4-5. GPLR2 Bit Definitions

						Add 0_00		SS							GPL	.R2							S	yste	em I	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
																PL66	PL65	PL64														
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Nar	ne												De	scr	iptio	on										
		<31:	21>			_	-		rese	erve	ed																					
		<20	:0>			PL	[x]		-	re:	ad-c in s		field is lo	ind w					,	alue	e of	each	n Gl	PIO.								

4.1.3.2 GPIO Pin Direction Registers (GPDR0, GPDR1, GPDR2)

GPDR0, GPDR1, GPDR2, shown in Table 4-6, Table 4-7, and Table 4-8, control whether a pin is an input or an output. The GPDR contain one direction control bit for each of the 81 GPIO pins. If a direction bit is programmed to a one, the GPIO is an output. If it is programmed to a zero, it is an input. Reserved bits (GPDR2[31:17]), must be written to zeros and reads to the reserved bits must be ignored.

Note: A reset clears all bits in the GPDR0-2 registers and configures all GPIO pins as inputs.



Table 4-6. GPDR0 Bit Definitions

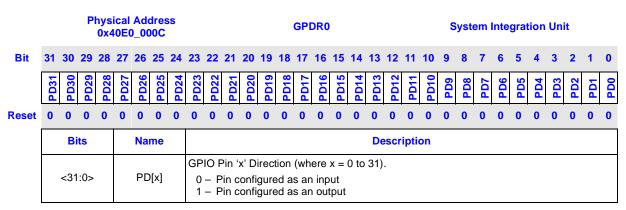


Table 4-7. GPDR1 Bit Definitions

						Add 0_00		S							GPI	DR1						S	yste	em I	nte	grat	ion	Uni	it		
Bit																1	0														
	PD63 PD64 PD60 PD60 PD60 PD60 PD60 PD60 PD63 PD44 PD44 PD44 PD44 PD44 PD44 PD44 PD33 PD34 PD35 PD36 PD37 PD36 PD37															PD33	PD32														
Reset																0	0														
		Bi	ts			Nar	ne											De	scr	iptio	on										
		<31	:0>			PD	[x]		GPI 0 1	– P	in c	c' Di onfi onfi	gure	d a	` s an	inp	ut.	o 63	3).												

Table 4-8. GPDR2 Bit Definitions

			P	hysi 0x4		Ad: 0_0(SS							GPE	R2							S	yste	em I	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 reserved Res															PD65	PD64															
Reset																0	0	0														
		В	its			Na	me												De	scr	iptic	on										
		<31	:21>	•		_	_		res	erve	ed																					
		<2	0:0>			PC)[x]		0	– F	in c	x' Di onfiç	gure	d as	` s an	inp	ut.	64 to	o 80).												

4.1.3.3 GPIO Pin Output Set Registers (GPSR0, GPSR1, and GPSR2) and Pin Output Clear Registers (GPCR0, GPCR1, GPCR2)



When a GPIO is configured as an output, the state of the pin can be controlled by writing to either the GPSR or GPCR. An output pin is set high by writing a one to its corresponding bit within the GPSR. To clear an output pin, a one is written to the corresponding bit within the GPCR. GPSR and GPCR are write-only registers. Reads return unpredictable values.

Writing a zero to any of the GPSR or GPCR bits has no effect on the state of the pin. Writing a one to a GPSR or GPCR bit corresponding to a pin that is configured as an input is effective only after the pin is configured as an output. Reserved bits (GPSR2[31:17] and GPCR2[31:17]), must be written with zeros and reads must be ignored.

Table 4-9, Table 4-10, and Table 4-11 show the bit definitions of GPSR0, GPSR1, and GPSR2. Table 4-12, Table 4-13, and Table 4-14 show the bit definitions of GPCR0, GPCR1, and GPCR2.

Table 4-9. GPSR0 Bit Definitions

				nysi 0x4			dres 018	S							GPS	R0							S	yste	em I	nte	grat	ion	Uni	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	PS30 PS29 PS29 PS28 PS28 PS26 PS27 PS27 PS27 PS27 PS27 PS27 PS27 PS27															PS2	PS1	PS0														
Reset																0	0	0														
		Bi	ts			Na	me												De	escr	ipti	on										
		<31	:0>			PS	[x]		0	– P	in le	x' O evel con	una	ffec	ted.	•					Ū	,		e).								

Table 4-10. GPSR1 Bit Definitions

		-	•	Τ.	•																											
						Add 0_00	dres 1C	S							GPS	R1							S	yste	em I	nte	grat	ion	Uni	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	PS63	PS62	PS61	PS60	PS59	PS58	PS57	PS56	PS55	PS54	PS53	PS52	PS51	PS50	PS49	PS48	PS47	PS46	PS45	PS44	PS43	PS42	PS41	PS40	PS39	PS38	PS37	PS36	PS35	PS34	PS33	PS32
Reset																0	0	0	0	0	0	0	0	0	0							
		Bi	ts			Naı	me												De	scr	iptio	on										
		<31	:0>			PS	[x]		-	– P	in le	vel	una	ffec		`					Ü	,		e).								



Table 4-11. GPSR2 Bit Definitions

				hysi 0x4				SS						•	GPS	R2							S	yste	em I	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					res	erv	ed					PS84	PS83	PS82	PS81	PS80	PS79	PS78	PS77	PS76	PS75	PS74	PS73	PS72	PS71	PS70	PS69	PS68	PS67	PS66	PS65	PS64
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	scri	iptio	on										
		<31	21>			_	-		res	erve	ed																					
		<20):0>			PS	6[x]		0	– F	Pin le	evel	una	ffec	ted.	`					J	80) nigh		e).								

Table 4-12. GPCR0 Bit Definitions

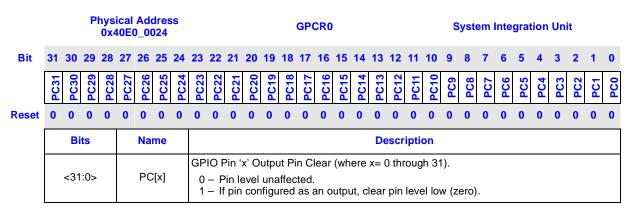


Table 4-13. GPCR1 Bit Definitions

				•		Add 0_00		S						•	GPO	CR1							S	yste	em I	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	PC63	PC62	PC61	PC60	PC59	PC58	PC57	PC56	PC55	PC54	PC53	PC52	PC51	PC50	PC49	PC48	PC47	PC46	PC45	PC44	PC43	PC42	PC41	PC40	PC39	PC38	PC37	PC36	PC35	PC34	PC33	PC32
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Naı	me												De	scr	iptio	on										
		<31	:0>			РС	[x]		-	– P	in le	evel	una	ffec	ted.		`				rou leve	•	,	ero)								



Table 4-14. GPCR2 Bit Definitions

				•		Ad 0_00								•	GPO	CR2							S	yste	em l	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					res	serv	ed					PC84	PC83	PC82	PC81	PC80	PC79	PC78	PC77	PC76	PC75	PC74	PC73	PC72	PC71	PC70	PC69	PC68	PC67	PC66	PC65	PC64
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Na	me												De	scr	iptic	n										
		<31	:21>	•		-	_		res	erve	ed																					
		<20):0>			PC	(x)		0	– F	Pin le	evel	una	ffec	ted.		`		x= 6 ear				,	ero).								

4.1.3.4 GPIO Rising Edge Detect Enable Registers (GRER0, GRER1, GRER2) and Falling Edge Detect Enable Registers (GFER0, GFER1, GFER2)

Each GPIO can also be programmed to detect a rising-edge, falling-edge, or either transition on a pin. When an edge is detected that matches the type of edge programmed for the pin, a status bit is set. The interrupt controller can be programmed so that an interrupt is signalled to the core when any of these status bits are set. Additionally, the interrupt controller can be programmed so that a subset of the status bits causes the processor to wake from Sleep mode when they are set. Refer to Section 3.4.9, "Sleep Mode" on page 3-15 and Section 3.5.6, "Power Manager GPIO Edge Detect Status Register (PEDR)" on page 3-28 for more information on which status bits can cause a wake up from Sleep mode.

Use the GRER and the GFER to select the type of transition on a GPIO pin that causes a bit within the GPIO Edge Detect Enable Status register (GEDR) to be set. For a given GPIO pin, its corresponding GRER bit is set causing a GEDR status bit to be set when the pin transitions from logic level zero to logic level one. Likewise, the GFER is used to set the corresponding GEDR status bit when a transition from logic level one to logic level zero occurs. When the corresponding bits are set in both registers, either a falling- or a rising-edge transition causes the corresponding GEDR status bit to be set.

Note: The minimum pulse width duration to guarantee edge detection is $1\mu S$.

Table 4-15 through Table 4-17 show the bitmaps of the GRER0, GRER1, and GRER2. Table 4-18 through Table 4-20 show the bitmaps of the GFER, GFER1, and GFER2.

Note: For reserved bits in GRER2 and GFER2, writes must be zeros and reads must be ignored.



Table 4-15. GRER0 Bit Definitions

				•		Add 0_00		s						•	GRE	ER0							S	yste	em I	nte	grat	ion	Uni	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	RE31	RE30	RE29	RE28	RE27	RE26	RE25	RE24	RE23	RE22	RE21	RE20	RE19	RE18	RE17	RE16	RE15	RE14	RE13	RE12	RE11	RE10	RE9	RE8	RE7	RE6	RE5	RE4	RE3	RE2	RE1	RE0
Reset	THE STATE OF THE S															0	0	0														
		<31	:0>			RE	[x]		0	– D	isal	x' Ri ole ri orre	isin	g-ed	ge o	dete	ct e	nab	le.					Ū		,	cted	on	the	GPI	O pi	ın

Table 4-16. GRER1 Bit Definitions

				•		Add 0_00		S							GRE	R1							S	yste	em I	nte	grat	ion	Uni	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	RE63	RE62	RE61	RE60	RE59	RE58	RE57	RE56	RE55	RE54	RE53	RE52	RE51	RE50	RE49	RE48	RE47	RE46	RE45	RE44	RE43	RE42	RE41	RE40	RE39	RE38	RE37	RE36	RE35	RE34	RE33	RE32
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	et 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																		De	scr	ipti	on										
		<31	:0>			RE	[x]			- D	isab	ole ri	ising	g-ed	lge d	dete	ct e	nab	le.			= 3		_		,	cted	on	the	GPI	O pi	n

Table 4-17. GRER2 Bit Definitions

				nysi 0x4				S						•	GRE	R2							S	yste	em l	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					res	erv	ed					RE84	RE83	RE82	RE81	RE80	RE79	RE78	RE77	RE76	RE75	RE74	RE73	RE72	RE71	RE70	RE69	RE68	RE67	RE66	RE65	RE64
Reset	RE84 RE73 RE73 RE73 RE73 RE73 RE74 RE77 RE77 RE78 RE77 RE77 RE77 RE77 RE77															0	0	0														
		Bi	ts			Na	me												De	scr	iptic	n										
		<31:	21>			_			res	erve	ed																					
		<20	:0>			RE	[x]		0	– C	Disal	ole r	ising	, g-ed	ge o	dete	ct e	nabl	e. `		re x en a					,	cted	on ·	the	GPI	O pi	n



Table 4-18. GFER0 Bit Definitions

				•		Add 0_00		ss							GFE	ER0							S	yste	em I	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	FE31	FE30	FE29	FE28	FE27	FE26	FE25	FE24	FE23	FE22	FE21	FE20	FE19	FE18	FE17	FE16	FE15	FE14	FE13	FE12	FE11	FE10	FE9	FE8	FE7	FE6	FE5	FE4	FE3	FE2	FE1	FE0
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Nai	me												De	escr	iptio	on										
		<31	:0>			FE	[x]			– D	isat	k' Fa ole fa orre	allin	g-e	dge	dete	ect e	enat	ole.	`				Ū		,	ected	d on	the	GP	IO p	oin

Table 4-19. GFER1 Bit Definitions

				nysi 0x4				S						•	GFE	R1							S	yste	em I	nte	grat	ion	Uni	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	FE63	FE62	FE61	FE60	FE59	FE58	FE57	FE56	FE55	FE54	FE53	FE52	FE51	FE50	FE49	FE48	FE47	FE46	FE45	FE44	FE43	FE42	FE41	FE40	FE39	FE38	FE37	FE36	FE35	FE34	FE33	FE32
Reset																0	0	0														
		Bi	ts			Naı	me												De	scr	iptic	n										
		<31	:0>			FE	[x]			– D	isat	ole fa	allin	g-ed	dge	dete	ect e	nab	le.					•	,	,	ected	d on	the	GPI	IO p	oin

Table 4-20. GFER2 Bit Definitions

			P	-		Ad 0_0		SS							GFE	R2							S	yste	em I	nte	grat	ion	Uni	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					res	serv	/ed					FE84	FE83	FE82	FE81	FE80	FE79	FE78	FE77	FE76	FE75	FE74	FE73	FE72	FE71	FE70	FE69	FE68	FE67	FE66	FE65	FE64
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Na	me												De	scr	iptio	on										
		<31	:21>			-	_		res	erve	ed																					
		<20):0>			FE	[x]		0	– C	Disal	x' Fa ble fa corre	allin	g-e	dge	dete	ect e	nab	ole.							,	cted	d on	the	GPI	Ор	in



4.1.3.5 GPIO Edge Detect Status Register (GEDR0, GEDR1, GEDR2)

GEDR1, GEDR2, shown in Table 4-21, Table 4-22, and Table 4-23, contain a total of 81 status bits that correspond to the 81 GPIO pins. When an edge detect occurs on a pin that matches the type of edge programmed in the GRER and/or GFER registers, the corresponding status bit is set in GEDR. Once a GEDR bit is set by an edge event, the bit remains set until the user clears it by writing a one to the status bit. Writing a zero to a GEDR status bit has no effect.

Each edge detect that sets the corresponding GEDR status bit for GPIO[80:0] can trigger an interrupt request. GPIO[80:2] together form a group that can cause one interrupt request to be triggered when any one of GEDR[80:2] are set. GPIO[0] and GPIO[1] cause independent first-level interrupts. Refer to Section 4.2, for a description of the programming of GPIO interrupts.

Table 4-21 through Table 4-23 show the bitmaps of the GEDR0, GEDR1, and GEDR2.

Table 4-21. GEDR0 Bit Definitions

				nysi 0x4			dres 048	S						•	GED	R0							S	yste	em l	nte	grat	ion	Uni	it		
Bit	ED23 ED28 ED29 ED26 ED26 ED26 ED27 ED26 ED27 ED27 ED27 ED27 ED28 ED19 ED19 ED19 ED19 ED19 ED19 ED19 ED19															1	0															
																ED2	ED1	ED0														
Reset	ot 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0															0	0	0														
	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																															
		<31:	:0>			ED)[x]		REA 0 1 WR 0	IO F AD - N - E ITE - N - C	o ed dge o ef	dge det	dete ect	ect h has	nas (occu	urred ed or	d on n pir	pin	as	spe	cifie	d in	GR					R.			

Table 4-22. GEDR1 Bit Definitions

				nysi 0x4			dres 04C	S						•	GEI	R1							S	yste	em I	nte	grat	ion	Uni	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	ED63	ED62	ED61	ED60	ED59	ED58	ED57	ED56	ED55	ED54	ED53	ED52	ED51	ED50	ED49	ED48	ED47	ED46	ED45	ED44	ED43	ED42	ED41	ED40	ED39	ED38	ED37	ED36	ED35	ED34	ED33	ED32
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	Bits Name Description GPIO Pin 'x' Edge Detect Status (where x= 32 through																															
		<31	:0>			ED)[x]		REA 0 1 WR 0	AD - N - E	o ed dge o ef	dge det	dete	ect h has	nas (occ	occu	urred ed oi	d on n pir	pin	as	spe	cifie	d in	GR	ER a				R.			



Table 4-23. GEDR2 Bit Definitions

	Bits Name Description <31:21> — reserved GPIO Pin 'x' Edge Detect Status (where x=64 through 80).																															
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2															ED65	ED64															
Reset	0x40E0_0050 t 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 reserved The standard of the standa															0	0	0														
	Bits Name Description																															
	reserved 20																															
		<20	:0>			EC)[x]		RE 0 1 WF	AD - N - E RITE - N	lo e dge	dge	dete	ect h has	nas (occ	occı urre	urred ed or	d on n pir	ı pin	as	spe	cifie	d in	, GRI					R.			

4.1.3.6 GPIO Alternate Function Register (GAFR0_L, GAFR0_U, GAFR1_L, GAFR1_U, GAFR2_L, GAFR2_U)

GAFR0_L, GAFR0_U, GAFR1_L, GAFR1_U, GAFR2_L, GAFR2_U, shown in Table 4-24, Table 4-25, Table 4-26, Table 4-27, Table 4-28, and Table 4-29, contain select bits that correspond to the 81 GPIO pins. Each GPIO can be configured to be either a generic GPIO pin, one of 3 alternate input functions, or one of 3 alternate output functions. To select any of the alternate functions, the GPDR register must configure the GPIO to be an input. Similarly, only GPIOs configured as outputs by the GPDR can be configured for alternate output functions. Each GPIO pin has a pair of bits assigned to it whose values determine which function (normal GPIO, alternate function 1, alternate function 2 or alternate function 3) the GPIO performs. The function selected is determined by writing the GAFR bit pair as below:

- "00" indicates normal GPIO function
- "01" selects alternate input function 1 (ALT_FN_1_IN) or alternate output function 1 (ALT_FN_1_OUT)
- "10" selects alternate input function 2 (ALT_FN_2_IN) or alternate output function 2 (ALT_FN_2_OUT)
- "11" selects alternate input function 3 (ALT_FN_3_IN) or alternate output function 3 (ALT_FN_3_OUT)



Caution: Configuring a GPIO to map to an alternate function that is not available causes indeterminate results.

Table 4-24. GAFR0 L Bit Definitions

Physical Address GAFR0_L **System Integration Unit** 0x40E0_0054 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 AF14 AF13 AF12 AF11 AF10 AF9 AF8 AF7 AF6 AF5 AF3 AF2 AF1 AF0 Reset 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0

Bits	Name	Description
<31:0>	AF[x]	GPIO Pin 'x' Alternate Function Select Bits (where x=0 through 15). A bit-pair in this register determines the corresponding GPIO pin's functionality as one of the alternate functions that is mapped to it or as a generic GPIO pin. 00 – The corresponding GPIO pin (GPIO[x]) is used as a general purpose I/O. 01 – The corresponding GPIO pin (GPIO[x]) is used for its alternate function 1. 10 – The corresponding GPIO pin (GPIO[x]) is used for its alternate function 2.
		11 – The corresponding GPIO pin (GPIO[x]) is used for its alternate function 3.

Table 4-25. GAFR0 U Bit Definitions

Physical Address 0x40E0_0058 GAFR0_U System Integration Unit

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 AF31 AF30 AF29 AF28 **AF27** AF26 AF25 AF24 AF23 AF22 AF21 **AF20** AF19 AF18 **AF17** AF16

Reset 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 **Bits Description** Name GPIO Pin 'x' Alternate Function Select Bits (where x=16 through 31). A bit-pair in this register determines the corresponding GPIO pin's functionality as one of the alternate functions that is mapped to it or as a generic GPIO pin. 00 - The corresponding GPIO pin (GPIO[x]) is used as a general purpose I/O. <31:0> AF[x] 01 – The corresponding GPIO pin (GPIO[x]) is used for its alternate function 1. 10 – The corresponding GPIO pin (GPIO[x]) is used for its alternate function 2.

11 – The corresponding GPIO pin (GPIO[x]) is used for its alternate function 3.



Table 4-26. GAFR1_L Bit Definitions

			-				ss					0X40E0_005C it 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 AF37 AF36 AF35 AF34															tion	Un	it		
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 AF37 AF36 AF35 AF34 AF35 Set 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0															-33	AF	32														
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
Bits Name Description																															
	<31	:0>			AF	[x]		A bit the a 00 – 01 –	t-pa alte - Th - Th	ir in rnate e co e co	this e fu rres	s reg incti spoi spoi	giste ons ndin	er de tha ng G ng G	eterr t is r PIO	mine map pin pin	es the ped (GF (GF	ne co I to i PIO[PIO[orre t or [x]) i [x]) i	spo as a s us	ndin a ge sed a sed 1	g G neri as a for i	PIO c GI ger ts al	pin PIO nera tern	's fu pin. I pu ate	incti rpos func	se I/	O. 11.	as o	ine (of
	AF	AF47 0 0 Bi	31 30 29 AF47 AF 0 0 0	0x4 81 30 29 28 AF47 AF46 0 0 0 0 Bits	0x40E0 81 30 29 28 27 AF47 AF46 AF 0 0 0 0 0 Bits	0x40E0_00 81 30 29 28 27 26 AF47 AF46 AF45 0 0 0 0 0 0 0 Bits Na	0x40E0_005C 81 30 29 28 27 26 25 AF47 AF46 AF45 AF 0 0 0 0 0 0 0 0 Bits Name	0x40E0_005C 31 30 29 28 27 26 25 24 AF47 AF46 AF45 AF44 0 0 0 0 0 0 0 0 0 0 Bits Name <31:0> AF[x]	0x40E0_005C 31 30 29 28 27 26 25 24 23 AF47 AF46 AF45 AF44 AF4 0 0 0 0 0 0 0 0 0 0 0 Bits Name GPI A bithe <31:0> AF[x] 000- 01-	0x40E0_005C 31 30 29 28 27 26 25 24 23 22 AF47 AF46 AF45 AF44 AF43 0 0 0 0 0 0 0 0 0 0 0 0 Bits Name GPIO F A bit-pa the alte 00 - Th 01 - Th	0x40E0_005C 31 30 29 28 27 26 25 24 23 22 21 3 AF47 AF46 AF45 AF44 AF43 AF4 0 0 0 0 0 0 0 0 0 0 0 0 0 Bits Name GPIO Pin 'x A bit-pair in the alternate of the control	0x40E0_005C 31 30 29 28 27 26 25 24 23 22 21 20 AF47 AF46 AF45 AF44 AF43 AF42 0 0 0 0 0 0 0 0 0 0 0 0 0 0 Bits Name GPIO Pin 'x' AI A bit-pair in this the alternate for 00 - The corre 01 - The corre 01 - The corre	0x40E0_005C 31 30 29 28 27 26 25 24 23 22 21 20 19 AF47 AF46 AF45 AF44 AF43 AF42 AF 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 Bits Name GPIO Pin 'x' Altern A bit-pair in this reg the alternate functi 00 - The correspon 01 - The correspon	0x40E0_005C 31 30 29 28 27 26 25 24 23 22 21 20 19 18 AF47 AF46 AF45 AF44 AF43 AF42 AF41 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 Bits Name GPIO Pin 'x' Alternate A bit-pair in this registe the alternate functions 00 - The correspondin 01 - The correspondin 01 - The correspondin	0x40E0_005C 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 Bits Name GPIO Pin 'x' Alternate Fur A bit-pair in this register de the alternate functions that 00 - The corresponding G 01 - The corresponding G	Ox40E0_005C GAFR1_ 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 Bits Name GPIO Pin 'x' Alternate Function A bit-pair in this register determing the alternate functions that is a colspan="8">ON — The corresponding GPIO 01 — The correspondi	0x40E0_005C GAFR1_L 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0x40E0_005C GAFR1_L GAFR1_L 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0x40E0_005C GAFR1_L 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	Ox40E0_005C GAFR1_L 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0x40E0_005C GAFR1_L GAFR1_L	0x40E0_005C GAFR1_L 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 AF37 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0x40E0_005C GAFR1_L State	Ox40E0_005C GAFRI_L System 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 AF37 AF36 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0x40E0_005C GAFR1_L System 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 AF37 AF36 AF 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0x40E0_005C GAFR1_L System Interval	Ox40E0_005C GAFR1_L System Integral 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 AF37 AF36 AF35 AF 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	Ox40E0_005C GAFR1_L System Integration Syst	0x40E0_005C GAFR1_L System integration of the content of the corresponding GPIO pin 'x' Alternate functions that is mapped to it or as a general purpose I/O. 01 – The corresponding GPIO pin (GPIO[x]) is used for its alternate function 1.	0x40E0_005C GAFR1_L System Integration Unit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 AF37 AF36 AF35 AF34 AF33 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	Ox40E0_005C GAFR1_L System Integration Unit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 AF47 AF46 AF45 AF44 AF43 AF42 AF41 AF40 AF39 AF38 AF37 AF36 AF35 AF34 AF33 AF 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0

Table 4-27. GAFR1_U Bit Definitions

	AF63 AF62 AF61 AF60 AF59 AF58 AF57 AF56 AF55 FA54 AF53 AF52 AF51 AF50 AF49 Bits Name Description GPIO Pin 'x' Alternate Function Select Bits (where x=48 through 63). A bit-pair in this register determines the corresponding GPIO pin's functionality as or																															
Bit	Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 AF63 AF62 AF61 AF60 AF59 AF58 AF57 AF56 AF55 FA54 AF53 AF52 AF51 AF50 AF50 Bits Name Description GPIO Pin 'x' Alternate Function Select Bits (where x=48 through 63).														3	2	1	0														
•	AF	63	AF	62	AF	61	AF	60	AF	59	AF	58	AF	57	AF	56	AF	55	FA	54	AF	53	Al	-52	A	F51	AF	5 0	AF	49	A	-48
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
•		Bi	ts			Na	me												De	esci	ipti	on										
		<31	:0>			AF	·[x]		A b the 00 01 10	it-pa alte – Th – Th – Th	air in ernat ne co ne co	thise fu orre orre	s reg unction spor spor	giste ons ndin ndin	er de tha ig G ig G ig G	eteri t is i PIO PIO	minemap pin pin pin pin	es the ped (GF (GF (GF	ne c I to i PIO[PIO[PIO[orre t or [x]) i [x]) i [x]) i	spo as a s us s us	ndin a ge sed : sed :	nerias a for i	PIC c G gei ts al	pir PIO nera terr	n's fu pin al pu nate nate	uncti irpo: func func	se I/ ction	O. 11. 12.	as o	ne (of



Table 4-28. GAFR2_L Bit Definitions

			P	•	cal A 0E0_			S						G	AFI	R2_	L						S	yst	em	Inte	gra	tion	Un	it		
Bit	31	30	29	28	27 2	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	AF	79	AF	78	AF7	7	AF	76	AF	75	AF	74	AF	73	AF	72	AF	71	AF	70	AF	69	AF	68	AF	67	AF	66	AF	65	AF	64
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts		N	Nar	ne												De	escr	ipti	on										
		<31	:0>		ı	AF	[x]		A b the 00 01 10	oit-pa alte – Th – Th – Th	air ir erna ne c ne c	n thi te fu orre orre	s reguncti spo spo	giste ions ndir ndir ndir	er de thang G ng G ng G	eterr t is r PIO PIO PIO	mine map pin pin pin	es the ped (GI (GI (GI	ne c to i PIO PIO	orre t or [x]) i [x]) i	spo as a s us s us	e x= ndin a ge sed a sed f	g G neri as a or it or it	PIO c GI ger s al	pin PIO nera tern tern	's fu pin. I pu ate ate	rpos func	se I/ction	O. 1. 2.	as o	ne d	of

Table 4-29. GAFR2 U Bit Definitions

			PI	•		Add 0_0(S						G	AF	R2_	U						S	yst	em	Inte	gra	tion	Un	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										r	ese	rve	d										AF	84	AF	83	AF	82	AF	81	AF	80
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	escr	iptio	n										
		<31:	:10>			-	_		res	erve	ed																					
		<9:	:0>			AF	·[x]		A b the 00 01 10	it-pa alte – Th – Th – Th	air ir erna ne c ne c ne c	thiste fu orre orre	s re incti spo spo spo	giste ions ndir ndir ndir	er do tha ig G ig G ig G	eteri t is i PIC PIC	mine map) pin) pin) pin	es the ped (GF (GF (GF	ne co I to i PIO[PIO[PIO[orre t or [x]) i [x]) i	where sport as a sustant as us a sustant as us a sustant as us as us a sustant as us a sus	ndin ger ed a ed f	g G neri as a or it or it	PIO c GI ger s al	pin PIO nera tern tern	i's fu pin. Il pu ate ate	rpos func func	se I/ction	0. 1. 2.	as o	ne d	of

4.1.3.7 Example Procedure for Configuring the Alternate Function Registers

In this example, GP0 is used as a generic GPIO and GP[15:1] are configured as their alternate functions. Refer to Table 4-1 for the list of alternate functions. No other GPIOs are configured. After the de-assertion of any RESET, GPDR0[15:0] configures GPIO pins in this example to be inputs. GAFR00[31:0] will be 0x0000_0000 to indicate normal GPIO function. For simplicity, assume that GP[16-31] are inputs configured as normal GPIOs.

In this example,

• GPIO[0] is configured as a normal GPIO input



- GPIO[1] is an input configured to alternate function 1 (ALT_FN_1_IN)
- GPIO[5:2] are reserved and configured as normal GPIOs inputs
- GPIO[12:6] are outputs configured to alternate function 1 (ALT_FN_1_OUT)
- GPIO[13] is an output configured to alternate function 2 (ALT_FN_2_OUT)
- GPIO[14] is an input configured to alternate function 1 (ALT FN 1 IN)
- GPIO[15] is an output configured to alternate function 2 (ALT_FN_2_OUT)

This programming sequence is required for programming the GPIO alternate functions out of reset:

- **1. WRITE GPSR0** 0x0000_8000 this sets GPIO15 (active low chip select) when it is configured as an output.
- **2. WRITE GPDR0** 0x0000_BFC0 GPIO[12:6], GPIO[13] and GPIO[15] as outputs. This drives GPIO[15] high until the alternate function information is programmed. This is required for active low outputs.
- **3. WRITE GAFRO_L** 0x9955_5004 this maps the alternate functions of GPIO[15:0]

For GPIOs that need to be configured as outputs, you must first program the GPSR and GPCR signals so the pin direction is changed. Change pin direction by setting the bit in the GPDR register—a '0' is driven for active high signals and '1' for active low signals.

Note: For more information on alternate functions, refer to the Source Unit column in Table 4-1 for the appropriate section of this document.

Table 4-24 through Table 4-29 show the bitmaps of the GPIO Alternate Function registers.

4.2 Interrupt Controller

The Interrupt Controller controls the interrupt sources available to the processor and also contains the location to determine the first level source of all interrupts. It also determines whether interrupts cause an IRQ or an FIQ to occur and masks the interrupts. The interrupt controller only supports a single priority level, however, interrupts can be routed to either IRQs or FIQ, with FIQs having priority over IRQs.

4.2.1 Interrupt Controller Operation

The interrupt controller provides masking capability for all interrupt sources and generates either an FIQ or IRQ processor interrupt. The interrupt hierarchy of the processor is a two-level structure.

- The first level identifies the interrupts from all the enabled and unmasked interrupt sources in the Interrupt Controller Mask Register (ICMR). First level interrupts are controlled by these registers:
 - Interrupt Controller Pending Register (ICPR) identifies all the active interrupts within the system
 - Interrupt Controller IRQ Pending Register (ICIP) contains the interrupts from all sources that can generate an IRQ interrupt. The Interrupt Controller Level Register (ICLR) is programmed to send interrupts to the ICIP to generate an IRQ.



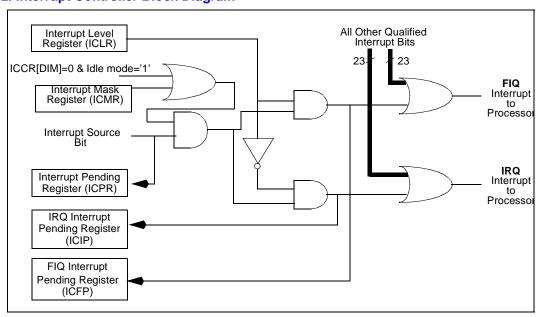
- Interrupt Controller FIQ Pending Register (ICFP) contains the interrupts from all sources that can generate an FIQ interrupt. The Interrupt Controller Level register (ICLR) is programmed to send interrupts to the ICFP to generate an FIQ.
- The second level uses registers contained in the source device (the device generating the first-level interrupt bit). The second-level interrupt status gives additional information about the interrupt and is used inside the interrupt service routine. In general, multiple second-level interrupts are OR'ed to produce a first-level interrupt bit.

In most cases, the root cause of an interrupt can be determined by reading two register locations: the ICIP for an IRQ interrupt or the ICFP for an FIQ interrupt to determine the interrupting device. You then read the status register within that device to find the exact function requesting service.

When the ICCR[DIM] bit is zero, the Interrupt Mask Register is ignored during Idle mode, and all enabled interrupts cause the processor to exit from idle mode. Otherwise, only unmasked interrupts cause the processor to exit from idle mode. The reset state of ICCR[DIM] is zero.

Figure 4-2 shows a block diagram of the Interrupt Controller.

Figure 4-2. Interrupt Controller Block Diagram



4.2.2 Interrupt Controller Register Definitions

The interrupt controller contains the following registers:

- Interrupt Controller IRQ Pending register (ICIP)
- Interrupt Controller FIQ Pending register (ICFP)
- Interrupt Controller Pending register (ICPR)
- Interrupt Controller Mask register (ICMR)
- Interrupt Controller Level register (ICLR)
- Interrupt Controller Control register (ICCR)



After a reset, the FIQ and IRQ interrupts are disabled within the CPU, and the states of all of the interrupt controller registers are set to 0x0. The interrupt controller registers must be initialized by software before interrupts are again enabled within the CPU.

4.2.2.1 Interrupt Controller Mask Register (ICMR)

The ICMR, shown in Table 4-30, contains one mask bit per pending interrupt bit (22 total). The mask bits control whether a pending interrupt bit generates a processor interrupt (IRQ or FIQ). When a pending interrupt becomes active, it is only processed by the CPU if the corresponding ICMR mask bit is set to 1. While in Idle mode, ICCR[DIM] must be set for the mask to be effective, otherwise any interrupt source that makes a request sets the corresponding pending bit and the interrupt is automatically processed, regardless of the state of its mask bit.

Mask bits allow periodic software polling of interruptible sources while preventing them from actually causing an interrupt. The ICMR is initialized to zero at reset, indicating that all interrupts are masked and the ICMR has to be configured by the user to select the desired interrupts.

Table 4-36 describes the available first-level interrupts and their location in the ICPR.

Table 4-30. ICMR Bit Definitions **Physical Address ICMR System Integration Unit** 0x40D0_0004 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 **IM27** reserved 0 0 Reset 0 0 0 0 0 0 Bits **Description** Name Interrupt Mask 'x' (where x= 8 through 14 and 17 through 31). 0 - Pending interrupt is masked from becoming active (interrupts are NOT sent to CPU or Power Manager). <31:8> IM[x] Pending interrupt is allowed to become active (interrupts are sent to CPU and Power

4.2.2.2 Interrupt Controller Level Register (ICLR)

The ICLR register, shown in Table 4-31, controls whether a pending interrupt generates an FIQ or an IRQ interrupt. If a pending interrupt is unmasked, the corresponding ICLR bit field is decoded to select which processor interrupt is asserted. If the interrupt is masked, then the corresponding bit in the ICLR has no effect. At reset the ICLR is initialized to all zeros, and software must configure the ICLR to reflect the normal operation value.

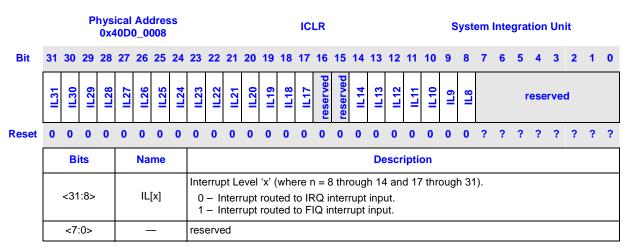
NOTE: In idle mode, the IM bits are ignored if ICCR[DIM] is cleared.

Table 4-36 describes the available first-level interrupts and their location in the ICPR.

<7:0>



Table 4-31. ICLR Bit Definitions



4.2.2.3 Interrupt Controller Control Register (ICCR)

The ICCR, shown in Table 4-32, contains a single control bit, Disable Idle Mask (DIM). In normal Idle mode any enabled interrupt can bring the processor out of Idle mode regardless of the value in ICMR. If this bit is set, then the interrupts that can bring the processor out of Idle mode are defined by the ICMR.

Note: This register is cleared during all resets.

Table 4-36 describes the available first-level interrupts and their location in the ICPR.

Table 4-32. ICCR Bit Definitions

					ical IOD(S							IC	CR							S	yst	em	Inte	gra	tion	Un	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															res	serv	ed															DIM
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0
		В	its			Na	me												De	escr	ipti	on										
		<31	:1>			_	_		res	erve	ed																					
		<()>			DI	М		0	– C	all ei Only node	nabl ena	ed i	nter d an	d ur	nma		d (as				out o					e pr	oces	ssor	out	of i	dle



4.2.2.4 Interrupt Controller IRQ Pending Register (ICIP) and FIQ Pending Register (ICFP)

The ICIP and the ICFP, shown in Table 4-33 and Table 4-34, contain one bit per interrupt (22 total.) These bits indicate an interrupt request has been made by a unit. Inside the interrupt service routine, read the ICIP and ICFP to determine the interrupt source. In general, software then reads status registers within the interrupting device to determine how to service the interrupt. Bits within the ICPR (see Section 4.2.2.5) are read only, and represent the logical OR of the status bits in the ICIP and ICFP for a given interrupt. Once an interrupt has been serviced, the handler writes a one to the required status bit, clearing the pending interrupt at the source.

Clearing the interrupt status bit at the source, automatically clears the corresponding ICIP or ICFP flag, provided there are no other interrupt status bits set within the source unit.

Table 4-36 describes the available first-level interrupts and their location in the ICPR.

Table 4-33. ICIP Bit Definitions

				•	ical IOD(S							IC	IP							S	yste	em	Inte	gra	tion	Un	it		
Bit	31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11															10	9	8	7	6	5	4	3	2	1	0						
	P31 P30 P30 P29 P28 P24 P24 P21 P21															IP8			r	ese	rve	d										
Reset	et 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0															0	0	0	?	?	?	?	?	?	?	?						
•		Bi	ts			Naı	me												De	scr	ipti	on										
•		<31	:8>			ΙΡΙ	[x]		0	– IF	1 QS	TON	rec	nere ques ed b	ted	by a	any	ena	bled	l so		•	gh 3	1).								
•		<7:	0>			_	_		rese	erve	ed																					

Table 4-34. ICFP Bit Definitions

		Physical Address 0x40D0_000C ICFP 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 \[\begin{array}{c c c c c c c c c c c c c c c c c c c																S	yste	em	Inte	gra	tion	Un	it							
Bit	31	0x40D0_000C 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 1 \[\begin{array}{c c c c c c c c c c c c c c c c c c c															11	10	9	8	7	6	5	4	3	2	1	0				
																FP11	FP10	FP9	FP8			r	ese	rve	d							
Reset																0	0	0	0	?	?	?	?	?	?	?	?					
•		Bi	ts			Naı	me												De	scr	ipti	on										
		<31	:8>			FP	[x]		0	– F	ndir IQ N IQ r	TO	rec	ues	ted	by a	any (enat	oled	sou		roug	gh 3	1).								
•		<7:	0>			_	_		rese	erve	ed																					



4.2.2.5 Interrupt Controller Pending Register (ICPR)

The ICPR, shown in Table 4-35, is a 32-bit read-only register that shows all active interrupts in the system. These bits are not affected by the state of the mask register (ICMR). Clearing the interrupt status bit at the source, automatically clears the corresponding ICPR flag, provided there are no other interrupt status bits set within the source unit.

Table 4-36 shows the pending interrupt source assigned to each bit position in the ICPR. Also included in the table are the source units for the interrupts and the number of second-level interrupts associated with each. For more information on the second-level interrupts, see the section that corresponds to its name in the Source Unit column.

Table 4-35. ICPR Bit Definitions (Sheet 1 of 3)

			PI			Add 0_00		is							ICI	PR							s	yst	em	Inte	egra	tion	Ur	nit		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	IS31	1830	IS29	1528	1827	1826	1825	IS24	IS23	IS22	IS21	1820	IS19	IS18	1817			1S14	IS13	1S12	IS11	1S10	IS9	IS8				recerved				
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Nai	me												De	escr	iptio	on										
	RTC Alarm Match Register Interrupt Pending 0 – Interrupt NOT pending due to RTC Alarm Match Register. 1 – Interrupt pending due to RTC Alarm Match Register. RTC HZ Clock Tick Interrupt Pending 30> IS30 0 – Interrupt NOT pending due to RTC HZ Clock Tick. 1 – Interrupt pending due to RTC HZ Clock Tick. 1 – Interrupt pending due to RTC HZ Clock Tick.																															
	1 - Interrupt pending due to RTC Alarm Match Register. RTC HZ Clock Tick Interrupt Pending																															
		<28	8>			IS	28		0	– Ir	nterr	upt	NO.	Трє	endii	ng d		o O	S Ti	imer					r 2							
		<2	7>			IS	27		0	– Ir	nterr	upt	NO.	Трє	endii	ng d	terru due t OS	o O	S Ti	imer					r 1							
		<20	6>			IS	26		0	– Ir	nterr	upt	NO.	Трє	endii	ng d	terru due t	to O	S Ti	imer					r O							
		<2	5>			IS	25		0	– Ir	nterr	upt	NO.	Тре	endii	ng d	st In due t	o D	ΜA	Cha	nne	l Se				uest.						
		<2	4>			IS	24		0	– Ir	nterr	upt	NO.	Трє	endii	ng d	ot Pe	o S	SP :				lues	st.								
		<23	3>			IS	23		0	– Ir	nterr	upt	NO.	Тре	endii	ng d	nteri due 1	:о М	МС	Sta	tus/l				tioı	n.						



Physical Address	Та	ble	4-3	5. IC	P	RE	3it [Def	fin	ition	s (\$	Sh	eet	2 (of 3	3)																	
Reset Bits				•					•						ı	CPF	2						;	Syst	en	n Int	eg	rati	ion	Uni	t		
Bits	Bit	31	30 2	29 28	3 2	27 2	26 2	5 2	24	23 22	21	2	0 1	9 18	3 1	7 1	6 15	14	13	12	11	10	9	8	7	7 6		5	4	3	2	1	0
Bits Name Description		IS31	1830	IS29	0761	IS27	1526	1825	IS24	IS23	1821		1820	200	2	IS17	reserved	IS14	IS13	IS12	1811	1810	681	88					reserved				
S22	Reset	0	0	0 0		0	0 ()	0	0 0	0	(0 0	0	(0	0	0	0	0	0	0	0	0	0	0		0	0	0	0	0	0
S22 IS22 0 - Interrupt NOT pending due to FFUART Transmit/Receive/Error. 1 - Interrupt pending due to FFUART Transmit/Receive/Error.			Bits	5		1	Nam	е											De	escr	ipti	on											
S21 S21 0 - Interrupt NOT pending due to BTUART Transmit/Receive/Error. 1 - Interrupt pending due to BTUART Transmit/Receive/Error. STUART Transmit/Receive/Error.			<22	>			IS22	2		0 –	Inter	ru	pt N	OT p	en	ding	due	to FI	FUA	RT	Trai	nsr					or.						
S20			<21	>			IS21			0 –	Inter	ru	pt N	OT p	en	ding	due	to B	TU <i>P</i>	ART	Tra	nsr					or.						
S19			<20	>			IS20)		0 –	Inter	ICPR System Integration Unit 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 Description Transmit/Receive/Error Interrupt Pending errupt NOT pending due to FFUART Transmit/Receive/Error. errupt pending due to FFUART Transmit/Receive/Error. errupt pending due to BTUART Transmit/Receive/Error. errupt pending due to BTUART Transmit/Receive/Error. errupt pending due to BTUART Transmit/Receive/Error. errupt pending due to STUART Transmit/Receive/Error. errupt pending due to STUART Transmit/Receive/Error. errupt pending due to STUART Transmit/Receive/Error. errupt pending due to ICP Transmit/Receive/Error. errupt pending due to ICC Service Request. errupt pending due to ICC Service Request. errupt pending due to ICC Controller Service Request. errupt pending due to ICD Controller Service Request. errupt pending due to ICD Controller Service Request. errupt pending due to ICD Controller Service Request. errupt pending due to Network SSP Service Request. errupt NOT pending due to Network SSP Service Request. errupt Pending errupt NOT pending due to AC97 unit errupt Pending errupt NOT pending due to ICS unit errupt pending due to PMU unit. errupt NOT pending due to PMU unit. errupt NOT pending due to PMU unit. errupt Pending errupt NOT pending due to PMU unit. errupt Pending errupt NOT pending due to PMU unit. errupt pending due to PMU unit. errupt Pending errupt NOT pending due to PMU unit.																					
S18			<19	>	Physical Address																												
S17			<18	System Integration Unit System Integrati																													
S16			<17	>			IS17	,		0 –	Inter	ru	pt N	ОТ р	en	CPR System Integration Unit 7 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 Description ive/Error Interrupt Pending ding due to FFUART Transmit/Receive/Error. due to FFUART Transmit/Receive/Error. due to BTUART Transmit/Receive/Error. due to BTUART Transmit/Receive/Error. due to STUART Transmit/Receive/Error. due to STUART Transmit/Receive/Error. due to STUART Transmit/Receive/Error. due to STUART Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/Receive/Error. due to ICP Transmit/R																	
AC97 Interrupt Pending O = Interrupt NOT pending due to AC97 unit 1 = Interrupt pending due to AC97 unit. IS13			<16	>			IS16	6		0 – In	erru	pt	NO	ICPR System Integration Unit 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 Description Smit/Receive/Error Interrupt Pending NOT pending due to FFUART Transmit/Receive/Error. pending due to BTUART Transmit/Receive/Error. Smit/Receive/Error Interrupt Pending NOT pending due to BTUART Transmit/Receive/Error. Smit/Receive/Error Interrupt Pending NOT pending due to STUART Transmit/Receive/Error. Smit/Receive/Error Interrupt Pending NOT pending due to STUART Transmit/Receive/Error. Smit/Receive/Error Interrupt Pending NOT pending due to ICP Transmit/Receive/Error. Sequest Interrupt Pending NOT pending due to ICP Transmit/Receive/Error. Sequest Interrupt Pending NOT pending due to ICP Service Request. Pending due to ICC Service Request. SP SERVICE REQUEST INTERRUPT PENDING: OT pending due to Network SSP Service Request. SP SERVICE REQUEST INTERRUPT PENDING: OT pending due to Network SSP Service Request. SP SERVICE REQUEST INTERRUPT PENDING: OT pending due to AC97 unit pending due to AC97 unit pending due to AC97 unit. Service Request Interrupt Pending NOT pending due to AC97 unit pending due to PMU unit. Service Request to PMU unit. Service Request.																			
 IS14			<15	>						reserv	ed																						
<13> IS13			<14	>			IS14	ļ		0 -	Inter	ru	Description Descr																				
<12> IS12			<13	>			IS13	3		0 –	Inter	ru	pt N	OT p	en					nit													
<11> IS11 0 – Interrupt NOT pending due to USB service request.			<12	>			IS12	2		0 –	Inter	ru	pt N	ΟΤ μ	en	ding	due	to Pl	MU	•		ndir	ng										
			<11:	>			IS11			0 –	Inter	ru	pt N	ОТ р	en	ding	due						ues	t.									

GPIO[80:2] Edge Detect Interrupt Pending

0 – Interrupt NOT pending due to edge detect on one (or more) of GPIO[80:2].
1 – Interrupt pending due to edge detect on one (or more) of GPIO[80:2].

<10>

IS10



Table 4-35. ICPR Bit Definitions (Sheet 3 of 3)

	Physical Address 0x40D0_0010 IC 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17															PR							S	yste	em I	Inte	grat	ion	Uni	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	IS31	IS30	IS29	IS28	IS27	1826	1825	IS24	IS23	IS22	IS21	IS20	IS19	IS18	1817	reserved		IS14	IS13	1812	IS11	IS10	<mark>6S</mark> 1	IS8				rosorvod				
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bit	s			Na	me												De	scr	iptio	on										
		<9	>			IS	9		0	– Ir	- nterr	upt	NO.	Тре	endi	rupt ng d e to	ue t	o ed	dge					[1].								
		<8	>			IS	88		0	– Ir	- nterr	upt	NO.	Тре	endi	rupt ng d e to	ue t	o ed	dge					[0].								
		<7	>			IS	57		0 –	Inte	rrup	t N	TC	pend	ding	Requ due to Ha	to	Har	dwa	re L	JAR [®]	T Se			•	ıest.						
		<6:0	0>						res	erve	h																					

Table 4-36. List of First-Level Interrupts (Sheet 1 of 2)

Bit Position	Source Unit	# of Level 2 Sources	Bit Field Description
IS<31>	Real-time clock	1	RTC equals alarm register.
IS<30>	Real-time clock	1	One Hz clock TIC occurred.
IS<29>		1	OS timer equals match register 3.
IS<28>	Operating system	1	OS timer equals match register 2.
IS<27>	timer	1	OS timer equals match register 1.
IS<26>		1	OS timer equals match register 0.
IS<25>	DMA controller	16	DMA Channel service request.
IS<24>	Synchronous Serial Port	3	SSP service request.
IS<23>	MUlti Media Card	9	MMC status / error detection
IS<22>	FFUART	5	x-mit, receive, error in FFUART.
IS<21>	BTUART	5	x-mit, receive, error in BTUART
IS<20>	STUART	4	x-mit, receive, error in STUART
IS<19>	ICP	6	x-mit, receive, error in ICP.
IS<18>	I2C	6	I2C service request.
IS<17>	LCD controller	15	LCD controller service request.
IS<16>	Network SSP	4	Network SSP service request
IS<15>			reserved
IS<14>	AC97	10	AC97 interrupt



Table 4-36. List of First-Level Interrupts (Sheet 2 of 2)

Bit Position	Source Unit	# of Level 2 Sources	Bit Field Description
IS<13>	I2S	5	I2S interrupt
IS<12>	Core	1	PMU (Performance Monitor) interrupt
IS<11>	USB	7	USB interrupt
IS<10>		79	"OR" of GPIO edge detects 80-2
IS<9>		1	GPIO<1> edge detect
IS<8>		1	GPIO<0> edge detect
IS<7>	Hardware UART	7	Hardware UART service request
IS<6>			reserved
IS<5>			reserved
IS<4>			reserved
IS<3>			reserved
IS<2>			reserved
IS<1>			reserved
IS<0>			reserved
	Total level 2 interrupt sources	179	

Several units have more than one source per interrupt signal. When an interrupt is signalled from one of these units, the interrupt handler routine identifies which interrupt was signalled using the interrupt controller's pending register. This identifies the unit that made the request, but not the exact source. The handler then reads the interrupting unit's status register to identify which source within the unit signalled the interrupt. For all interrupts that have one corresponding source, the interrupt handler routine needs to use only the interrupt controller's registers to identify the exact cause of the interrupt. ICPR[16:15] and ICPR[7:0] are reserved bits and must be written as zeros. Reads to these bits must be ignored.

4.3 Real-Time Clock (RTC)

Use the RTC to configure a clock source with a wide range of frequencies. Typically, the RTC is set to be a 1 Hz output and is utilized as a system time keeper. There is also an alarm feature that enables an interrupt or a wake up event when the RTC output clock increments to a pre-set value.

4.3.1 Real-Time Clock Operation

The RTC provides a general-purpose real-time reference for your design. The RTC Counter register (RCNR) is initialized to zero after a hardware reset or a watchdog reset. It is a free running counter that starts incrementing the count value after the deassertion of reset. The counter is incremented one 32kHz cycle after the rising edge of the Hz clock. Since the high phase of the 1 Hz clock is one 32kHz cycle wide, it appears to increment on the falling edge of the 1 Hz clock. Set this counter to the desired value. If the counter is set to a value other than zero, write the desired value to the RCNR. The value of the counter is unaffected by transitions into and out of Sleep or Idle mode.



In addition to the RCNR, the RTC incorporates a 32-bit, RTC Alarm register (RTAR). The RTAR may be programmed with a value that is compared against the RCNR. One 32-kHz cycle after each rising edge of the HZ clock, the counter is incremented and then compared to the RTAR. If the values match, and the enable bit is set, then the RTC Status register (RTSR) alarm match bit (RTSR[AL]) is set. This status bit is also routed to the interrupt controller and may be unmasked in the interrupt controller to generate a processor interrupt. Another available interruptible status bit that can be set whenever the HZ clock transitions is the RTSR. By writing a one to the AL or HZ bit in the RTSR, the status bit is cleared.

The HZ clock is generated by dividing one of two selectable clock sources, both approximately 32.768 kHz in frequency. The first source is the output of the 3.6864 MHz crystal oscillator further divided by 112 to approximately 32.914 kHz. The other source is the optional 32.768 kHz crystal oscillator output itself. Your system may be built with both the 32.768 kHz crystal oscillator and the 3.6864 MHz crystal oscillator. ALternately, your system may only use the 3.6864 MHz crystal oscillator, if the additional power consumption during sleep mode is acceptable.

The divider logic for generating the HZ clock is programmable. This lets you trim the counter to adjust for inherent inaccuracies in the crystal's frequency and the inaccuracy caused by the division of the 3.6864 MHz oscillator which yields only an approximate 32 kHz frequency. The trimming mechanism lets you adjust the RTC to an accuracy of +/- 5 seconds per month. The trimming procedure is described in a later paragraph.

All registers in the RTC, with the exception RTTR, are reset by hardware reset and the watchdog reset. The trim register, RTTR is reset only by hardware reset.

4.3.2 RTC Register Definitions

The following sections provide register descriptions for the RTC.

4.3.2.1 RTC Trim Register (RTTR)

Program the RTTR to set the frequency of the HZ clock. The reset value of this register (0x0000_7FFF) (assuming a perfect 32.768 kHz crystal) would produce an HZ clock output of exactly 1 Hz. However, by using values other than 0x0000_7FFF, a different HZ clock frequency is possible. Additionally, developers can use a crystal that is not exactly 32.768 kHz and compensate by writing a value other than 0x0000_7FFF to the RTTR. Section 4.3.3 describes how to calculate the value in this register. A write to the RTTC will increment the RTC Count Register (RCNR) by one. RTTC[LCK] does not prevent the RCNR from incrementing.

All reserved bits must be written to zeros and reads to these bits must be ignored. You can only reset the RTTR with a hardware reset. To safeguard the validity of the data written into the trim register, bit 31 is used as a Lock Bit. The data in RTTR may be changed only if RTTR[LCK] is cleared. Once, RTTR[LCK] is set to be a one, only a hardware reset can clear the RTTR.



Table 4-37. RTTR Bit Definitions

			PI			Add	dres	S							RT	TR							S	yst	em	Inte	grat	tion	Un	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•	LCK		res	erv	ed						DI	EL												CK_	DIV	,						
Reset	0	?	?	?	?	?	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
		Bi	ts			Na	me												De	escr	ipti	on										
·		<3	1>			LC	CK		0	king – R – R	RTTF	R va	lue	is a	llow	ed t					l.											
		<30:	26>			-	_		res	erve	ed																					
·	<25:16> DEL									n de s va				ents	the	num	ber	of 3	32 kl	Hz d	clock	ks to	de	lete	whe	en cl	ock	trim	nmin	ıg be	egin	ıs.
·		<15	:0>			CK_	_DIV	,		ck c s va					er c	of 32	kH	z clo	ock	cycl	es,	olus	1, p	oer I	HZ c	clock	ссус	cle.				

4.3.2.2 RTC Alarm Register (RTAR)

The RTAR, Table 4-38, is a 32-bit register. The processor can both read and write to this register. Following each rising edge of the HZ clock, this register is compared to the RCNR. If the two are equal and RTSR[ALE] is set, then RTSR[AL] is set.

Because of the asynchronous nature of the HZ clock relative to the processor clock, writes to this register are controlled by a hardware mechanism that delays the actual write to the register by two 32 kHz clock cycles after the processor store is performed.

The RTAR is initialized to 0x0 at reset.

Table 4-38. RTAR Bit Definitions

			PI		ical 1090		dres 004	S							RT	AR							S	yst	em	Inte	gra	tion	Un	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
																RT	MV															
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	escr	ipti	on										
	<31:0> RTMV								RTC Target Match Value. The value compared against the RTC counter.																							



4.3.2.3 RTC Counter Register (RCNR)

The RCNR, shown in Table 4-39, is a read/write register. The counter may be written by the processor at any time although it is recommended that the operating system prevent inadvertent writes to the RCNR through the use of the MMU protection mechanisms (refer to the *Intel*® *XScale*TM *Microarchitecture for the Intel*® *PXA255 Processor User's Manual* for details of MMU operation.)

Because of the asynchronous nature of the HZ clock relative to the processor clock, writes to this counter are controlled by a hardware mechanism that delays the actual write to the counter after the processor store is performed by approximately two 32 kHz clock cycles. In case of multiple writes to RCNR in quick succession, the final update to the RCNR counter may be delayed by up to two 32 kHz clock cycles.

The RCNR may be read at any time. Reads reflect the value in the counter after it increments or has been written and does not have the two 32 kHz clock cycle delay.

Table 4-39. RCNR Bit Definitions

						Add 0_ 0 0		SS	RCNR 4 23 22 21 20 19 18 17 16 15 14 13 12											System Integration Unit												
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
																R	CV															
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	its			Na	me												De	escr	ipti	on										
	<31:0>				RC	CV		RTC Count Value. The current value of the RTC counter.																								

4.3.2.4 RTC Status Register (RTSR)

The RTSR, shown in Table 4-40, is cleared to all zeroes at hardware reset. The ALE and HZE bits enable both the interrupt for the functions as well as the updating the AL and HZ bits. The AL and HZ bits are status bits and are set by the RTC logic if the ALE and HZE bits are set respectively. They are cleared by writing ones to the AL and HZ bits. The AL and HZ bits are routed to the interrupt controller where they may be enabled to cause a first level interrupt. Write zeros to all reserved bits and ignore all reads to the reserved bits.

In Sleep mode, only AL events set the status bit in the RTSR register. The HZ bit is not set in Sleep mode since it is a recurring event.

Table 4-40 shows the bitmap of the RTC Status Register.



Table 4-40. RTSR Bit Definitions

			P			Ad 0_00	dres	ss							RT	SR							S	Syst	em	Inte	gra	tion	Un	it		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													r	ese	rve	d													HZE	ALE	HZ	AL
Reset	?	?	?	?	?	?	?	?	?	?	?	?	1	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0
		Bi	ts			Na	me												De	esci	ipti	on										
		<31	:4>			_			res	erve	ed																					
		<3	3>			HZ	ZE		0	– T	rrup he l	ΗZ i	inter	rupt																		
		<2	2>			Al	LE		0	– Т	arm he f	RTC	ala	ırm	inte	rrup																
		<1	l>			Н	ΙZ		0	- N	ng-e No ris	sing	g ed	ge h	nas I					ıd H	ZE I	oit is	set	t.								
		<()>			Α	۱L				arm lo R				as l	oeer	ı de	tect	ed.													

4.3.3 Trim Procedure

The HZ clock driving the RTC is generated by dividing the output of the oscillator multiplexor. The inherent inaccuracies of crystals, aggravated by varying capacitance of the board connections, as well as other variables, may cause the time base to be somewhat inaccurate. This requires a slight adjustment in the desired clock period. The processor, through the RTTR, lets you adjust (or trim) the HZ time base to an error of less than 1ppm. Such that if the HZ clock is set to be 1 Hz, there would be an error of less than 5 seconds per month.

1 - An RTC alarm has been detected (RTNR matches RCAR).and ALE bit is set

The RTTR is reset to its default value of 0x0000_7FFF each time the nRESET signal is asserted. This yields approximately a 1 Hz clock.

When the clock divisor count (RTTR[15:0]) is set to 0x0, the HZ clock feeding the RTC maintains a high level signal - essentially disabling the RTC. For all non-zero values programmed into the clock divisor count, the HZ clock frequency will be the 32 kHz clock source divided by the clock divisor count plus 1.

4.3.3.1 Oscillator Frequency Calibration

To determine the value programmed into the RTTR, you must first measure the output frequency at the oscillator multiplexor (approximately 32 kHz) using an accurate time base, such as a frequency counter. This clock is externally visible by selecting the alternate function for GPIO[12] or GPIO[72]. To gain access to the clock, program this pin as an output and then switch to the alternate function. Refer to Section 4.1, for details on how to make the clock externally visible. To trim the clock, divide the output of the oscillator by an integer value and fractional adjust it by periodically deleting clocks from the stream driving this integer divider.



4.3.3.2 RTTR Value Calculations

After the true frequency of the oscillator is known, it must be divided by the desired HZ clock frequency and this value split into integer and fractional portions. The integer portion of the value (minus one) is loaded into the Clock Divider Count field of the RTTR. This value is compared against a 16-bit counter clocked by the output of the oscillator multiplexor at approximately 32 kHz. When the two values are equal, the counter resets and generates a pulse which constitutes the raw HZ clock signal.

The fractional part of the adjustment is done by periodically deleting clocks from the clock stream driving the integer counter. The trim interval period is hardwired to be 2^{10} -1 periods of the HZ clock. If the HZ clock is programed to be 1 Hz the trim interval would be approximately 17 minutes. The number of clocks deleted (the trim delete value) is a 10-bit programmable counter allowing from 0 to 2^{10} -1 32 kHz clocks to be deleted from the input clock stream once per trim interval. RTTR[25:16] represents the number of 32 kHz clocks deleted per trim operation.

In summary, every 2^{10} -1 HZ clock periods, the integer counter stops clocking for a period equal to the fractional error that has accumulated. If this fractional error is programmed to be zero, then no trim operations occur and the RTC is clocked with the raw 32 kHz clock. The relationship between the HZ clock frequency and the nominal 32 kHz clock (f1 and f32K, respectively) is shown in the following equation.

$$f1 = \frac{(2^{10-1})^{*}(RTTR[CK_DIV]+1) - RTTR[DEL]}{(2^{10-1})^{*}(RTTR[CK_DIV]+1)} \quad * \quad \frac{f32k}{(RTTR[CK_DIV]+1)}$$

f1 = HZ clock frequency

f32k=RTC internal clock - either the 32.678 kHz crystal output or the 3.68 MHz crystal output divided down to 32.914 kHz

RTTR[DEL] = RTTR(25:16) $RTTR[CK_DIV] = RTTR(15:0)$

4.3.3.2.1 Trim Example #1 – Measured Value Has No Fractional Component

In this example, the desired HZ clock frequency is 1 Hz. The oscillator output is measured as 36045.000 cycles/s (Hz). This output is exactly 3277 cycles over the nominal frequency of the crystal (32.768 kHz) and has no fractional component. As such, only the integer trim function is needed - no fractional trim is required. Accordingly, RTTR[15:0] is loaded with the binary equivalent of 36045-1, or 0x0000_8CCC. RTTR[25:16] is left at zero (power-up state) to disable fractional trimming. This trim exercise leaves an error of zero in trimming.

4.3.3.2.2 Trim Example #2 – Measured Value Has a Fractional Component

This example is more common in that the measured frequency of the oscillator has a fractional component. Again, the desired HZ clock output frequency is 1 Hz. If the oscillator output is measured as 32768.92 cycles/s (Hz), an integer trim is necessary so that the *average* number of cycles counted before generating one 1 Hz clock is 32768.92. Similar to the previous example, the integer field RTTR[15:0] is loaded with the hexadecimal equivalent of 32768-1 or 0x0000_7FFF (reset value).

Because the actual clock frequency is 0.92 cycles per second faster than the integer value, the HZ clock generated by just the integer trimming is slightly faster than needed and must be slowed down. Accordingly, program the fractional trim to delete 0.92 cycles per second on average to



bring the HZ output frequency down to the proper value. Since the trimming procedure is performed every 1023 (2^{10} -1) seconds, the trim must be set to delete 941.16 clocks every 1023 seconds ($.92 \times 1023 = 941.16$). Load the counter with the hexadecimal equivalent of 941, or 0x3AD. The fractional component of this value cannot be trimmed out and constitutes the error in trimming, described below.

This trim setting leaves an error of. 16 cycles per 1023 seconds. The error calculation yields (in parts-per-million or ppm):

$$Error = \frac{0.16 \text{ cycles}}{1023 \text{ sec}} X \frac{1 \text{ sec}}{32768 \text{ cycles}} = 0.002 \text{ ppm}$$

4.3.3.2.3 Maximum Error Calculation Versus RTC Accuracy

As seen from trim example #2, the maximum possible error approaches 1 clock per 2¹⁰-1 seconds. Calculating the ppm error for this scenario yields:

Error (maximum) =
$$\frac{1 \text{ cycle}}{1023 \text{ sec}} X \frac{1 \text{ sec}}{32768 \text{ cycles}} = 0.03 \text{ ppm}$$

To maintain an accuracy of +/- 5 seconds per month, the required accuracy is calculated to be:

Error =
$$\frac{5 \text{ sec}}{\text{month}} X \frac{1 \text{ month}}{2592000 \text{ sec}} = 1.9 \text{ ppm}$$

This calculation indicates that the HZ clock output can be made very accurate through the use of the trim procedure. Likewise, use the trim procedure to compensate for a range of factors that can affect crystal oscillators. Such factors can include, but are not limited to:

- · Manufacturing and supplier variance in the crystals
- Crystal aging effects
- System voltage differences
- System manufacturing variance

The trim procedure can counteract these factors by providing a highly accurate mechanism to remove the variance and shifts from the manufacturing and static environment variables on an individual system level. However, since this is a calibration solution, it is not a practical solution for dynamic changes in the system and environment and can most likely only be done in a factory setting due to the equipment required.

4.4 Operating System (OS) Timer

The processor contains a 32-bit OS timer that is clocked by the 3.6864 MHz oscillator. The Operating System Count register (OSCR) is a free running up-counter. The OS timer also contains four 32-bit match registers (OSMR3, OSMR2, OSMR1, OSMR0). Developers can read and write to each register. When the value in the OSCR is equal to the value within any of the match registers, and the interrupt enable bit is set, the corresponding bit in the OSSR is set. These bits are



also routed to the interrupt controller where they can be programmed to cause an interrupt. OSMR3 also serves as a watchdog match register that resets the processor when a match occurs provided the OS Timer Watchdog Match Enable Register (OWER) is set. You must initialize the OSCR and OSMR registers and clear any set status bits before the FIQ and IRQ interrupts are enabled within the CPU.

4.4.1 Watchdog Timer Operation

The OSMR3 can also be used as a watchdog compare register. This function is enabled by setting OWER[0]. When a compare against this register occurs and the watchdog is enabled, reset is applied to the processor and most internal states are cleared. Internal reset is asserted for 256 processor clocks and then removed, allowing the processor to boot. See Section 3.4.2, "Watchdog Reset" on page 3-7 for details on reset functionality.

The following procedure is suggested when using OSMR3 as a watchdog – each time the operating system services the register:

- 1. The current value of the counter is read.
- 2. An offset is then added to the read value. This offset corresponds to the amount of time before the next time-out (care must be taken to account for counter wraparound).
- 3. The updated value is written back to OSMR3.

The OS code must repeat this procedure periodically before each match occurs. If a match occurs, the OS timer asserts a reset to the processor.

4.4.2 OS Timer Register Definitions

4.4.2.1 OS Timer Match Register 0-3 (OSMRx)

These registers are 32-bits wide and are readable and writable by the processor. They are compared against the OSCR after every rising edge of the 3.6864 MHz clock. If any of these registers match the counter register, and the appropriate interrupt enable bit is set, then the corresponding status bit in the OSSR is set. The status bits are routed to the interrupt controller where they can be unmasked to cause a CPU interrupt. The OSMR3 can also be used as a watchdog timer.

Table 4-41 shows the bitmap of the OS Timer Match register. All four registers are identical, except for location. A single, generic OS Timer match register is described, but all information is common to all four OS Timer Match Registers.

<31:0>



Table 4-41. OSMR[x] Bit Definitions

OSMV

Physical Address OS Timer Match Register 0-3 (OSMR3, OSMR2, OSMR1, 0x40A0_0000 0x40A0_0004 **System Integration Unit** 0x40A0_0008 OSMR0) 0x40A0_000C Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 **OSMV Bits Description** Name OS Timer Match Value.

The value compared against the OS timer counter.

4.4.2.2 OS Timer Interrupt Enable Register (OIER)

The OIER, shown in Table 4-42, contains four enable bits that indicate whether a match between one of the match registers and the OS timer counter sets a status bit in the OSSR. Each match register has a corresponding enable bit. Clearing an enable bit does not clear the corresponding interrupt status bit if it is already set.

Table 4-42. OIER Bit Definitions

Physical Address OS Timer Interrupt Enable System Integration Unit Register (OIER) 0x40A0_001C 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 Rit reserved Ш **E**2 Ш E **Bits Description** Name <31:4> reserved Interrupt enable channel 3. <3> E3 0 - A match between OSMR3 and the OS Timer will NOT assert OSSR[M3]. 1 - A match between OSMR3 and the OS Timer asserts OSSR[M3]. Interrupt enable channel 2. <2> E2 0 - A match between OSMR2 and the OS Timer will NOT assert OSSR[M2]. 1 - A match between OSMR2 and the OS Timer asserts OSSR[M2]. Interrupt enable channel 1. <1> E1 0 - A match between OSMR1 and the OS Timer will NOT assert OSSR[M1]. 1 - A match between OSMR1 and the OS Timer asserts OSSR[M1].

0 – A match between OSMR0 and the OS Timer will NOT assert OSSR[M0].
1 – A match between OSMR0 and the OS Timer asserts OSSR[M0].

Interrupt enable channel 0.

<0>

FΩ



4.4.2.3 OS Timer Watchdog Match Enable Register (OWER)

The OWER, shown in Table 4-43, contains a single control bit (bit 0) that enables the watchdog function. This bit is set by writing a one to it and can only be cleared by one of the reset functions such as, hardware reset, sleep reset, watchdog reset, and GPIO reset.

Table 4-43. OWER Bit Definitions

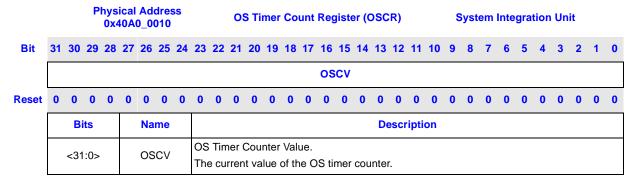
		Physical Address 0x40A0_0018							OS Timer Watchdog Match Enable Register (OWER)												System Integration Unit											
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
															re	serv	/ed															WME
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	esci	ripti	on										
		<31	:1>			-			res	erve	ed																					
	<0> WME			Watchdog Match Enable 0 – OSMR3 match will NOT cause a reset of the processor 1 – OSMR3 match causes a reset of the processor.																												

4.4.2.4 OS Timer Count Register (OSCR)

The OSCR, shown in Table 4-44, is a 32-bit counter that increments on rising edges of the 3.6864 MHz clock. This counter can be read or written at any time. It is recommended that the system write-protect this register through the MMU protection mechanisms.

After the OSCR is written, there is a delay before the register is actually updated. Software must make sure the register has changed to the new value before relying on the contents of the register.

Table 4-44. OSCR Bit Definitions



4.4.2.5 OS Timer Status Register (OSSR)

The OSSR, shown in Table 4-45, contains status bits that indicate a match has occurred between any of the four match registers and the OSCR. These bits are set when the match event occurs (following the rising edge of the 3.6864 MHz clock) and the corresponding interrupt enable bit is set in the OIER. The OSSR bits are cleared by writing a one to the proper bit position. Writing zeros to this register has no effect. Write all reserved bits as zeros and ignore all reads.

Bit



Table 4-45. OSSR Bit Definitions

Physical Address 0x40A0_0014	OS Timer Status Register (OSSR)	System Integration Unit
---------------------------------	---------------------------------	-------------------------

Bits	Name	Description
<31:4>	_	reserved
<3>	МЗ	Match status channel 3. If OIER[3] is set then: 0 - OSMR[3] has NOT matched the OS timer counter since last being cleared. 1 - OSMR[3] has matched the OS timer counter.
<2>	M2	Match status channel 2. If OIER[2] is set then: 0 – OSMR[2] has NOT matched the OS timer counter since last being cleared. 1 – OSMR[2] has matched the OS timer counter.
<1>	M1	Match status channel 1. If OIER[1] is set then: 0 - OSMR[1] has NOT matched the OS timer counter since last being cleared. 1 - OSMR[1] has matched the OS timer counter.
<0>	MO	Match status channel 0. If OIER[0] is set then: 0 - OSMR[0] has NOT matched the OS timer counter since last being cleared. 1 - OSMR[0] has matched the OS timer counter.

4.5 Pulse Width Modulator

Use the Pulse Width Modulator (PWM) to generate as many as two signals to be output from the processor. The signals are based on the 3.6864 MHz clock and must be a minimum of 2 clock cycles wide. These signals are output from the processor by configuring the GPIOs.

4.5.1 Pulse Width Modulator Operation

The processor contains two pulse width modulators: PWM0 and PWM1. Each PWM operates independently of the other, is controlled by its own set of registers. They provide a pulse width modulated signal on an external pin. Since each PWM contains identical circuitry, a generic PWMn, where n is 0 or 1, is described.

Each PWM contains:

- Two Pulse Width Modulator channels
- Enhanced Period control through 6-Bit Clock divider and 10-Bit Period counter
- 10-Bit Pulse control

A block diagram of one of the PWMs is shown in Figure 4-3.



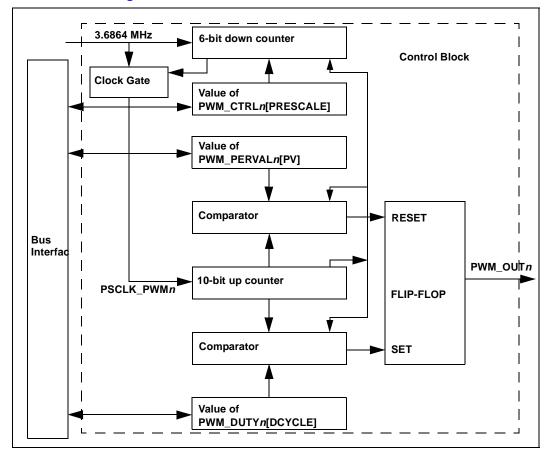


Figure 4-3. PWMn Block Diagram

4.5.1.1 Interdependencies

The PWM unit is clocked off the 3.6864 MHz oscillator output.

Each Pulse Width Modulator Unit (PWMn) is controlled by three registers:

- Pulse Width Control Register (PWM CTRL)
- Duty Cycle Control Register (PWM_DUTY)
- Period Control Register (PWM_PERVAL)

By setting the values in these registers the PWMn unit produces a pulse width modulated output signal. The registers contain the values for PWMn's counters and PWMn power management mode.

Each register contains one or more fields which determine an attribute of the PWM_OUTn waveform. PWM_CTRLn[PRESCALE] specifies the divisor for the PWM module clock. Note that the actual PWM module clock divisor used is 1 greater than the value programmed into PWM_CTRLn[PRESCALE]. This divided PWM module clock drives a 10 bit up-counter. This upcounter feeds 2 separate comparators. The first comparator contains the value of PWM_DUTYn[DCYCLE]. When the values match, the PWM_OUT signal is set high. The other



comparator contains PWM_PERVALn[PV] and clears the PWM_OUT signal low when PWM_PERVALn[PV] + 1 and the 10-bit up counter are equal. Both PWM_PERVALn[PV] and PWM_DUTYn[DCYCLE] are 10 bit fields.

Note: Take care to ensure that the value of the PWM_PERVAL*n* register remains larger than PWM_DUTY*n* register. In the case where PWM_PERVAL*n* is less than PWM_DUTY*n* the output maintains a high state.

4.5.1.2 Reset Sequence

A system reset results in no pulse width modulated signal. During system reset the PWM_CTRLn and PWM_DUTYn registers are reset to 0x0 and the PWM_PERVALn register is set to 0x004. This sets the PWM_OUTn pin low with a zero duty cycle. The six bit down-counter is reset to 0x0 and thus the 3.68 MHz input clock directly drives the 10 bit up-counter. The PWM_OUTn pin remains reset low until the PWM_DUTYn register is programmed with a non zero value.

A basic pulse width waveform is shown in Figure 4-4.

4.5.1.3 Power Management Requirements

Each PWM may be disabled through a pair of clock enable bits (see Section 3.6.2, "Clock Enable Register (CKEN)" on page 3-36). If the clock is disabled, the unit shuts down in one of two ways:

- Abrupt the PWM stops immediately.
- Graceful the PWM completes the current duty cycle before stopping.

Shutdown is selected by PWM_CTRL[PWM_SD] and described in Section 4.5.2.1.

4.5.2 Register Descriptions

The following paragraphs provide register descriptions for the Pulse Width Modulator.

4.5.2.1 PWM Control Registers (PWM CTRLn)

The PWM CTRLn, shown in Table 4-46, contains two fields:

• PRESCALE – The PRESCALE field contains the 6-bit prescale counter load value. This field allows the 3.6864 MHz input clock PSCLK_PWMn, to be divided by values between 1 (PWM_CTL[PRESCALE] = 0) and 64 (PWM_CTL[PRESCALE] = 63).

Note: The value of the divisor is one greater than the value programmed into the PRESCALE field.

• PWM_SD – PWMn can shut down in one of two ways, gracefully or abruptly, depending on the setting of PWM_CTRLn[PWM_SD]. If gracefully is chosen, then the duty cycle counter completes its count before PWMn is shut down. If abruptly is chosen, then the prescale counter and the duty cycle counter are reset to the reload values in their associated registers and PWMn is immediately shut down.

Note: During abrupt shut down the PWM_OUT*n* signal may be delayed by up to one PSCLK_PWM*n* clock period.



Table 4-46. PWM_CTRLn Bit Definitions

	Physical Address 0x40B0_0000 0x40C0_0000				PWM Control Registers (PWM_CTRL0, PWM_CTRL1)									System Integration Unit																		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							reserved								OS PRESCALE																	
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Naı	me												De	escr	ipti	on										
		<31	:7>			_	_		rese	erve	ed																					
	<6> PWM_SD)	PWMn Shutdown Method: 0 — Graceful shutdown of PWMn when the clock enable bit in the CKEN register is cleared. 1 — Abrupt shutdown of PWMn when the clock enable bit in the CKEN register is cleared.												ed.															
		<5:	0>		PF	RES	CAI	E.	Det	erm	ines	scal s the WMr	fre	que	ncy								`			f the	3.8	6 M	Hz (clock	()	

4.5.2.2 PWM Duty Cycle Registers (PWM_DUTY*n*)

The PWM_DUTY*n*, shown in Table 4-47, contains two fields:

- FDCYCLE
- DCYCLE

The FDCYCLE bit determines whether or not PWM_OUT*n* is a function of the DCYCLE bits in the PWM_DUTY*n* register or is set high. When the FDCYCLE bit is cleared low (normal operation), the output waveform of PWM_OUT*n* is cyclic, with PWM_OUT*n* being high for the number of PSCLK_PWM*n* periods equal to DCYCLE.

If FDCYCLE=0x0 and DCYCLE=0x0, PWM_OUTn is set low and does not toggle.

Note: If FDCYCLE is 0b1, PWM_OUT*n* is high for the entire period and is not influenced by the value programmed in the DCYCLE bits.



Table 4-47. PWM_DUTYn Bit Definitions

			P	0x4	10B	0_0	dres 004 004										•	egisters System Integration Unit														
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										res	serv	red										FDCYCLE				ı	DCY	CLI	=			
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Na	me												De	scr	ipti	on										
		<31	:11>			_			res	erve	ed																					
<10> FDCYCLE PWMn Full Duty Cycle 0 - PWM clock (PWM_OI 1 - PWM_OUTn is set high												ned	l by	DC'	YCL	_E fi	eld.															
		<9	:0>		I	DCY	′CLI	=	Du	ty cy	/cle	ty C	WN	1n c			. the	nur	mbe	er of	PS	CLK	_PV	۷M	cycl	les F	PWN	/In is	s as	serte	ed	

4.5.2.3 PWM Period Control Register (PWM_PERVAL*n*)

The PWM_PERVALn, shown in Table 4-48, contains a 10 bit field called PV. This field determines the period of the PWM_OUTn waveform in terms of the PSCLK_PWMn clock. If this field is cleared to zero PWMn is effectively turned off and PWM_OUTn remains in a high state. For any non-zero value written to the PV field, the output frequency of PWMn is the frequency of the PSCLK_OUTn divided by the value of (PV + 1). The range of the clock gate extends from a pass-through of the PSCLK_PWMn to a clock delay of 2^6 or 64 input clocks per output pulse.

When the value of the 10 bit up-counter equals the value of (PV +1), the up-counter and the flip-flop are reset and the values of PWM_CTRLn, PWM_PERVALn and PWM_DUTYn are loaded into the internal versions of these registers. Resetting this flip-flop causes PWM_OUTn to go low and the PWM cycle to start again.

Writing all zeroes to this register results in the output maintaining a high state unless FDCYCLE=0x0 and DCYCLE=0x0. If FDCYCLE=0x0 and DCYCLE=0x0, the output maintains a low state regardless of the value in the PV bit field.

Note: Due to internal timing requirements, all changes to any of the PWM registers must be complete a minimum of 4 core clock cycles before the start of end of a PWM clock cycle in order to guarantee that the following PWM cycle implements the new values.



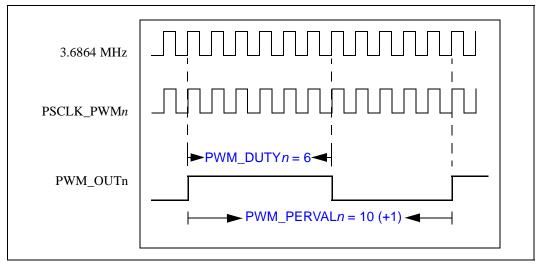
Table 4-48. PWM_PERVALn Bit Definitions

	Physical Address 0x40B0_0008 0x40C0_0008				PWM Period Control Registers (PWM_PERVAL0, PWM_PERVAL1)								\$	System Integration Unit																		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										r	ese	rve	d														P	V				
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0
		Bi	ts			Na	me												D	esci	ripti	on										
		<31:	10>			_	_		res	erve	ed																					
									PW	/Mn	Per	iod	Cor	trol:																		
	The number of PSCLK_PWMn cycles that comprise one PWM_OUTn cycle NOTE: If PV = 0x0, the PWMn clock (PWM_OUTn) is set high and does not togg FDCYCLE=0x0 and DCYCLE=0x0. In this case PWM_OUTn is set low a not toggle regardless of the value in PV.																															

4.5.3 Pulse Width Modulator Output Wave Example

Figure 4-4 is an example of the output of a Pulse Width Modulator for reference.

Figure 4-4. Basic Pulse Width Waveform



PWM_PERVAL[PV] = 0xA PWM_DUTY[FDCYCLE] = 0x0 PWM_DUTY[DCYCLE] = 0x6 PWM_CTRL[PRESCALE] = 0x0

The output waveform in Figure 4-4 is created by writing PWM_PERVALn[PV] with a decimal value of 10 (11 clocks) and PWM_DUTYn[DCYCLE] with 6. Figure 4-4 also shows that PWM_CTRLn[PRESCALE] is configured with a value of 0x0 loaded, which results in the PSCLK_PWMn having the same frequency as the 3.6864 MHz input clock.



4.6 System Integration Unit Register Summary

4.6.1 **GPIO Register Locations**

Table 4-49 shows the registers associated with the GPIO block and their physical addresses.

Table 4-49. GPIO Register Addresses (Sheet 1 of 2)

Address	Name	Description
0x40E0_0000	GPLR0	GPIO pin level register GPIO[31:0]
0x40E0_0004	GPLR1	GPIO pin level register GPIO[63:32]
0x40E0_0008	GPLR2	GPIO pin level register GPIO[80:64]
0x40E0_000C	GPDR0	GPIO pin direction register GPIO[31:0]
0x40E0_0010	GPDR1	GPIO pin direction register GPIO[63:32]
0x40E0_0014	GPDR2	GPIO pin direction register GPIO[80:64]
0x40E0_0018	GPSR0	GPIO pin output set register GPIO[31:0]
0x40E0_001C	GPSR1	GPIO pin output set register GPIO[63:32]
0x40E0_0020	GPSR2	GPIO pin output set register GPIO[80:64]
0x40E0_0024	GPCR0	GPIO pin output clear register GPIO[31:0]
0x40E0_0028	GPCR1	GPIO pin output clear register GPIO[63:32]
0x40E0_002C	GPCR2	GPIO pin output clear register GPIO[80:64]
0x40E0_0030	GRER0	GPIO rising-edge detect enable register GPIO[31:0]
0x40E0_0034	GRER1	GPIO rising-edge detect enable register GPIO[63:32]
0x40E0_0038	GRER2	GPIO rising-edge detect enable register GPIO[80:64]
0x40E0_003C	GFER0	GPIO falling-edge detect enable register GPIO[31:0]
0x40E0_0040	GFER1	GPIO falling-edge detect enable register GPIO[63:32]
0x40E0_0044	GFER2	GPIO falling-edge detect enable register GPIO[80:64]
0x40E0_0048	GEDR0	GPIO edge detect status register GPIO[31:0]
0x40E0_004C	GEDR1	GPIO edge detect status register GPIO[63:32]
0x40E0_0050	GEDR2	GPIO edge detect status register GPIO[80:64]
0x40E0_0054	GAFR0_L	GPIO alternate function select register GPIO[15:0]
0x40E0_0058	GAFR0_U	GPIO alternate function select register GPIO[31:16]
0x40E0_005C	GAFR1_L	GPIO alternate function select register GPIO[47:32]



Table 4-49. GPIO Register Addresses (Sheet 2 of 2)

0x40E0_0060	GAFR1_U	GPIO alternate function select register GPIO[63:48]
0x40E0_0064	GAFR2_L	GPIO alternate function select register GPIO[79:64]
0x40E0_0068	GAFR2_U	GPIO alternate function select register GPIO[80]

4.6.2 Interrupt Controller Register Locations

Table 4-50 shows the registers associated with the interrupt controller block and their physical addresses.

Table 4-50. Interrupt Controller Register Addresses

Address	Name	Description
0x40D0_0000	ICIP	Interrupt controller IRQ pending register
0x40D0_0004	ICMR	Interrupt controller mask register
0x40D0_0008	ICLR	Interrupt controller level register
0x40D0_000C	ICFP	Interrupt controller FIQ pending register
0x40D0_0010	ICPR	Interrupt controller pending register
0x40D0_0014	ICCR	Interrupt controller control register

4.6.3 Real-Time Clock Register Locations

Table 4-51 describes the location of the RTC registers.

Table 4-51. RTC Register Addresses

Address	Name	Description
0x4090_0000	RCNR	RTC count register
0x4090_0004	RTAR	RTC alarm register
0x4090_0008	RTSR	RTC status register
0x4090_000C	RTTR	RTC trim register

4.6.4 OS Timer Register Locations

Table 4-52 shows the registers associated with the OS timer and the physical addresses used to access them.

Table 4-52. OS Timer Register Addresses (Sheet 1 of 2)

Address	Name	Description
0x40A0_0000	OSMR0	OS timer match register 0
0x40A0_0004	OSMR1	OS timer match register 1
0x40A0_0008	OSMR2	OS timer match register 2



Table 4-52. OS Timer Register Addresses (Sheet 2 of 2)

0x40A0_000C	OSMR3	OS timer match register 3
0x40A0_0010	OSCR	OS timer counter register
0x40A0_0014	OSSR	OS timer status register
0x40A0_0018	OWER	OS timer watchdog enable register
0x40A0_001C	OIER	OS timer interrupt enable register

4.6.5 Pulse Width Modulator Register Locations

Table 4-53 shows the registers associated with the PWM and the physical addresses used to access them.

Table 4-53. Pulse Width Modulator Register Addresses

Address	Name	Description
0x40B0_0000	PWM_CTRL0	PWM0 Control Register
0x40B0_0004	PWM_PWDUTY0	PWM0 Duty Cycle Register
0x40B0_0008	PWM_PERVAL0	PWM0 Period Control Register
0x40C0_0000	PWM_CTRL1	PWM1 Control Register
0x40C0_0004	PWM_PWDUTY1	PWM1 Duty Cycle Register
0x40C0_0008	PWM_PERVAL1	PWM1 Period Control Register

DMA Controller

This chapter describes the on-chip DMA controller (DMAC) for the PXA255 processor. The DMAC transfers data to and from main memory in response to requests generated by internal and external peripherals. The peripherals do not directly supply addresses and commands to the memory system. The DMAC has 16 DMA channels, 0 through 15, and every DMA request from the peripheral generates at least one memory bus cycle.

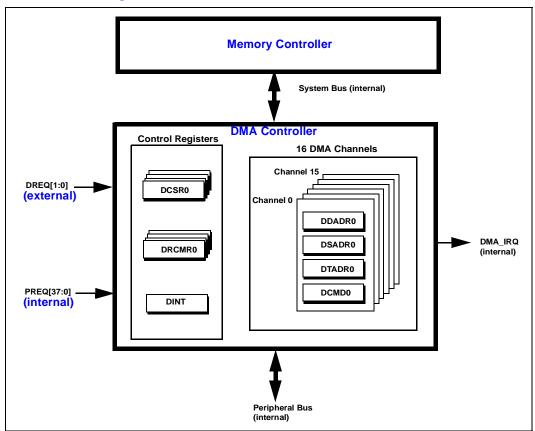
5.1 DMA Description

The DMAC supports only flow-through transfers.

Flow-through data passes through the DMAC before the data is latched by the destination in its buffers/memory. This DMAC can perform memory-to-memory moves with flow-through transfers.

Figure 5-1 provides an overview of the DMAC. Table 5-1 provides a list of the DMAC signals and descriptions.

Figure 5-1. DMAC Block Diagram





5.1.1 DMAC Channels

The DMAC has 16 channels, each controlled by four 32-bit registers. Each channel can be configured to service any internal peripheral or one of the external peripherals for flow-through transfers. Each channel is serviced in increments of the peripheral device's burst size and is delivered in the granularity appropriate to that device's port width. The burst size and port width for each device is programmed in the channel registers and is based on the device's FIFO depth and bandwidth needs. Due to performance issues, it is highly recommended that the user set the burst size equal to the FIFO DMA interrupt trigger level, also called the FIFO threshold level. When multiple channels are actively executing, the DMAC services each channel with a burst of data. After the data burst is sent, the DMAC may perform a context switch to another active channel. The DMAC performs context switches based on a channel's activity, whether its target device is currently requesting service, and where that channel lies in the priority scheme.

Channel information must be maintained on a per-channel basis and is contained in the DMAC registers see in Table 5-13. The DMAC supports two methods of loading the DMAC register, No-Descriptor and Descriptor Fetch Modes. The fetch modes are discussed in further detail in Section 5.1.4.

Software must ensure cache coherency when it configures the DMA channels. The DMAC does not check the cache so target and source addresses must be configured as non-cacheable in the Memory Management Unit.

Each demand for data that a peripheral generates results in a read or write to memory data. A peripheral must not request a DMA transfer unless it is prepared to read or write the full data block (8, 16, or 32 bytes) and it is equipped to handle reads and writes less than a full data block. Reads and writes less than a full data block can occur at the end of a DMA transfer.

5.1.2 Signal Descriptions

The DREQ[1:0], PREQ[37:0] and DMA_IRQ signals are controlled by the DMAC as indicated in Table 5-1.

Table 5-1. DMAC Signal List

Signal	Signal Type In/Out	To/From	Definition
DREQ[1:0]	Input	Pins	External companion chip request lines. DMA detects the positive edge of this signal as a request.
DMA_IRQ	Output	Interrupt Controller	Active high signal indicating an interrupt.
PREQ[37:0]	Input	On-chip peripherals	Internal peripheral DMA request lines. On chip peripherals send requests using the PREQ signals. The DMAC does not sample the PREQ signals until it completely finishes the data transfer from peripheral to the memory.

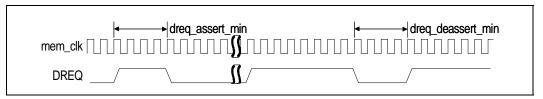
5.1.2.1 DREQ[1:0] and PREQ[37:0] Signals

The external companion chip asserts the positive edge triggered DREQ[1:0] signals when a DMA transfer request is needed. The DREQ[1:0] signal must remain asserted for four MEMCLKs to allow the DMA to recognize the 0 to 1 transition. When the DREQ[1:0] signals are deasserted, they



must remain deasserted for at least four MEMCLKs. The DMAC registers the transition from 0 to 1 to identify a new request. The external companion chip must not assert another DREQ until the previous DMA data transfer starts.

Figure 5-2. DREQ timing requirements



The PREQ[37:0] bits are the active high internal signals from the on-chip peripherals. Unlike DREQ[1:0], they are level sensitive. The DMAC does not sample the PREQ[37:0] signals until it completely finishes the current data transfer. For a write request to the on-chip peripheral, the DMAC begins to sample the PREQ[37:0] signals after it sends the last byte of the write request. For a read request, the DMAC begins to sample the PREQ[37:0] signals after it sends the last byte that pertains to the read on the internal bus.

The DCSR[REQPEND] bit indicates the status of the pending request for the channel.

If a DREQx assertion sets the DCSR[REQPEND] bit and software resets the DCSR[RUN] bit to stop the channel, the DCSR[REQPEND] bit and the internal registers that hold the DREQx assertion information may remain set even though the channel has stopped. To reset the DCSR[REQPEND] bit, software must send a dummy descriptor that transfers some data.

5.1.2.2 DMA_IRQ Signal

The processor has 16 IRQ signals, one for each DMA channel. Each DMA IRQ can be read in the DINT register that is shown in Table 5-6. The user can mask some bits that cause interrupts on a channel, such as ENDIRQEN, STARTIRQEN, and STOPIRQEN.

When DMA interrupt occurs, it is visible in Pending Interrupt Register Bit 25 (see Section 4.2.2.5, "Interrupt Controller Pending Register (ICPR)" on page 4-25). When a pending interrupt becomes active, it is sent to the CPU if its corresponding ICMR mask Bit 25 (see Section 4.2.2.1, "Interrupt Controller Mask Register (ICMR)" on page 4-22) is set to a one.

5.1.3 DMA Channel Priority Scheme

The DMA channel priority scheme allows peripherals that require high bandwidth to be serviced more often than those requiring less bandwidth. The DMA channels are internally divided into four sets. Each set contains four channels. The channels get a round-robin priority in each set. Set zero has the highest priority. Set 1 has higher priority than sets two and three. Sets two and three are low priority sets. Refer to Table 5-2 for details. High bandwidth peripherals must be programmed in set zero. Memory-to-memory moves and low bandwidth peripherals must be programmed in set two or three. When all channels are running concurrently, set zero is serviced four times out of any eight consecutive channel servicing instances. Set one is serviced twice and sets two and three are each serviced once.

If two or more channels are active and request a DMA, the priority scheme in Table 5-2 applies.

Request priority does not affect requests that have already started. The DMAC priority scheme is considered when the smaller dimension of the DCMDx[SIZE] or DCMDx[LENGTH] is complete.



If all channels request data transfers, the Sets are prioritized in following order:

- Set zero
- · Set one
- · Set zero
- Set two
- Set zero
- Set one
- Set zero
- Set three

The pattern repeats for the next eight channel services. In each set, the channels are given round-robin priority.

Table 5-2. Channel Priority (if all channels are running concurrently)

Set	Channels	Priority	Number of times served
0	0,1,2,3	Highest	4 / 8
1	4,5,6,7	Higher	2/8
2	8,9,10,11	Low	1/8
3	12,13,14,15	Low	1/8

The state machine used to determine the priority of the DMA channels is shown in Table 5-3. Use this table to determine the exact sequence the DMA controller gives to each channel when not all channels are running concurrently.

Table 5-3. Channel Priority

State Machine State	DMA Set Priority within each State Machine State
0	S0 > S1 > S2 > S3
1	S1 > S0 > S3 > S2
2	S0 > S1 > S2 > S3
3	S2 > S3 > S0 > S1
4	S0 > S1 > S2 > S3
5	S1 > S0 > S3 > S2
6	S0 > S1 > S2 > S3
7	S3 > S2 > S1 > S0

The channels get a round-robin priority in each set. Out of reset, the state machine state is zero. If a channel in set zero has a pending request, that channel is serviced. If a channel in set one has a pending request, that channel is serviced and so on. Once a request is serviced, the state machine



state is incremented, wrapping around from state machine state seven back to state machine state zero. If there is no pending request, the state machine stays in the current state machine state until there is a pending request. See Table 5-4 for priority scheme examples.

Table 5-4. Priority Schemes Examples

Channels Programmed	DMA Channel Priority
ch0, ch1	0,1,0,1,0,1,0,1,etc.
ch0, ch15	0,0,0,15,0,0,0,15,etc.
ch0, ch4, ch8, ch12	0,4,0,8,0,4,0,12,etc.
ch0, ch1, ch8, ch12	0,1,0,8,0,1,0,12,etc.
ch0, ch4	0,4,0,0,0,4,0,4,etc.
ch8, ch12	8,12,8,8,8,12,8,12,etc.

5.1.4 DMA Descriptors

The DMAC operates in two distinct modes: Descriptor Fetch Mode and No-Descriptor Fetch Mode. The mode used is determined by the DCSRx[NODESCFETCH] bit.

The Descriptor Fetch and No-Descriptor modes can be used simultaneously on different channels. This means that some DMA channels can be active in one mode while other channels are active in the other mode.

A channel must be stopped before it can be switched from one mode to the other.

If an error occurs in a channel, it returns to its stopped state and remains there until software clears the error condition and writes a 1 to the DCSR[RUN] register.

5.1.4.1 No-Descriptor Fetch Mode

In No-Descriptor Fetch Mode, the DDADRx is reserved. Software must not write to the DDADRx and must load the DSADRx, DTADRx, and DCMDx registers. When the Run bit is set, the DMAC immediately begins to transfer data. No-Descriptor fetches are performed at the beginning of the transfer. The channel stops when it finishes the transfer.

Ensure that the software does not program the channel's DDADx No-Descriptor Fetch Mode.

A typical No-Descriptor Fetch Mode (DCSR[NODESCFETCH] = 1) operation follows:

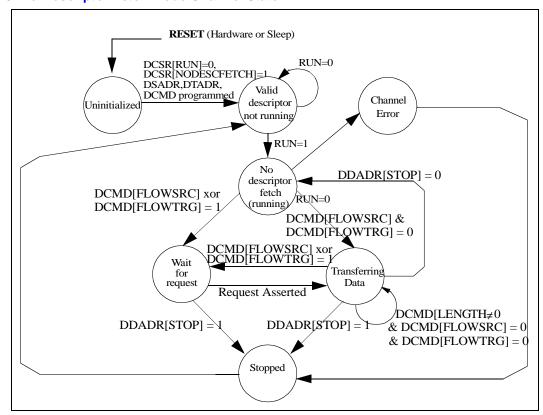
- 1. The channel is in an uninitialized state after reset.
- 2. The DCSR[RUN] bit is set to a 0 and the DCSR[NODESCFETCH] bit is set to a 1.
- 3. The software writes a source address to the DSADR register, a target address to the DTADR register, and a command to the DCMD register. The DDADR register is reserved in this No-Descriptor Fetch Mode and must not be written.
- 4. The software writes a 1 to the DCSR[RUN] bit and the No-Descriptor fetches are performed.
- 5. The channel waits for the request or starts the data transfer, as determined by the DCMD[FLOW] source and target bits.
- 6. The channel transmits a number of bytes equal to the smaller of DCMD[SIZE] and DCMD[LENGTH].



- 7. The channel waits for the next request or continues with the data transfer until the DCMD[LENGTH] reaches zero.
- 8. The DDADR[STOP] is set to a 1 and the channel stops.

Figure 5-3 summarizes typical No-Descriptor Fetch Mode operation.

Figure 5-3. No-Descriptor Fetch Mode Channel State



5.1.4.2 Descriptor Fetch Mode

In Descriptor Fetch Mode, the DMAC registers are loaded from DMA descriptors in main memory. Multiple DMA descriptors can be chained together in a list. This allows a DMA channel to transfer data to and from a number of locations that are not contiguous. The descriptor's protocol design allows descriptors to be added efficiently to the descriptor list of a running DMA stream.

A typical Descriptor Fetch Mode (DCSR[NODESCFETCH] = 0) operation follows:

- 1. The channel is in an uninitialized state after reset.
- The software writes a descriptor address (aligned to a 16-byte boundary) to the DDADR register.
- 3. The software writes a 1 to the DCSR[RUN] bit.
- 4. The DMAC fetches the four-word descriptor (assuming that the memory is already set up with the descriptor chain) from the memory indicated by DDADR.
- 5. The four-word DMA descriptor, aligned on a 16-byte boundary in main memory, loads the following registers:



- a. Word [0] -> DDADRx register and a single flag bit. Points to the next four-word descriptor.
- b. Word [1] -> DSADRx register for the current transfer.
- c. Word [2] -> DTADRx register for the current transfer.
- d. Word [3] -> DCMDx register for the current transfer.
- 6. The channel waits for the request or starts the data transfer, as determined by the DCMD[FLOW] source and target bits.
- 7. The channel transmits a number of bytes equal to the smaller of DCMD[SIZE] and DCMD[LENGTH].
- 8. The channel waits for the next request or continues with the data transfer until the DCMD[LENGTH] reaches zero.
- 9. The channel stops or continues with a new descriptor fetch from the memory, as determined by the DDADR[STOP] bit.

Bit [0] (STOP) of Word [0] in a DMA descriptor (the low bit of the DDADRx field) marks the descriptor at the end of a descriptor list. The value of the STOP bit does not affect the manner in which the channel's registers load the descriptor's fields. If a descriptor with its STOP bit set is loaded into a channel's registers, the channel stops after it completely transfers the data that pertains to that descriptor. Figure 5-4 summarizes this operation.

Software must set the DCSR[RUN] bit to 1 after it loads the DDADR. The channel descriptor fetch does not take place unless the DDADR register is loaded and the DCSR[RUN] bit is set to a 1.

The DMAC priority scheme does not affect DMA descriptor fetches. The next descriptor is fetched immediately after the previous descriptor is serviced.



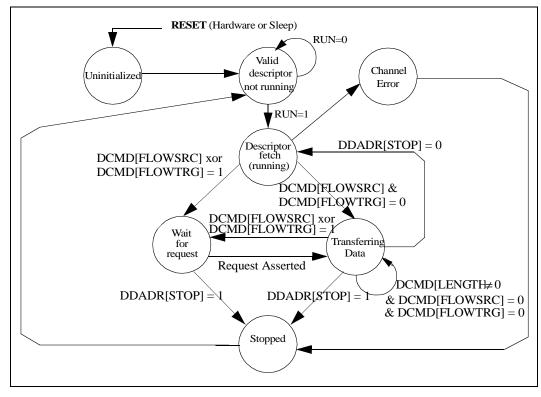


Figure 5-4. Descriptor Fetch Mode Channel State

5.1.4.3 Servicing an Interrupt

If software receives an interrupt caused by a successful descriptor fetch, i.e. DCSRx[STARTINTR] = 0b1, then software must write a 1 to this bit to reset the corresponding interrupt. Software normally accomplishes this by reading the DCSRx register, modifying the data value by setting the DCSRx[STARTINTR]=0b1 and leaving the DCSRx[RUN] bit set, and then writing this modified value back to the DCSRx. If the channel stops, DCSRx[RUN] = 0b0, before writing this value back to the DCSRx, then software can inadvertently set the DCSRx[RUN] bit before properly configuring other DMA registers. In order to avoid this problem, after writing the modified value back to the DCSRx, software must read the DCSRx and check to see if DCSRx[RUN] and DCSRx[STOPSTATE] are both set. If they are, then software must clear the DCSRx[RUN] bit and re-initialize the DMA channel.

5.1.5 Channel States

A DMA channel can go through any of the following states:

- Uninitialized: Channel is in an uninitialized state after reset.
- Valid Descriptor, Not Running: Software has loaded a descriptor in the DDADR of the channel, in the Descriptor Fetch Mode, or has programmed DSADR, DTADR and DCMD values, in No-Descriptor Fetch Mode, but the corresponding run bit in the DCSR[RUN] register is not set to a 1.
- Descriptor Fetch, Running: Fetching four words of descriptors from the memory.



- Wait for Request: Channel is waiting for a request before it starts to transfer the data.
- Transfer Data: Channel is transferring data.
- Channel Error: Channel has an error. It remains in the stopped state until software clears the error condition, re-initializes the channel, and writes a 1 to the DCSR[RUN] bit. See Section 5.3.1 and Section 5.3.2 for details.
- Stopped: Channel is stopped.

Figure 5-3 and Figure 5-4 show the progression from state to state.

5.1.6 Read and Write Order

The DMAC does not ensure the order of programmed I/O reads and writes made from the processor to the I/O devices (including the on-chip I/O devices). Software must ensure the order.

The DMAC ensures that all memory references made by a single DMA data stream are presented to main memory in the order in which they were made. The descriptor fetches occurs between the data blocks. This allows self-modifying DMA descriptor chains to function correctly (see Example 4 on page 5-27). It also allows schemes in which a DMA stream writes data blocks followed by status blocks and schemes in which another DMA stream (probably from the processor) polls the same field in the status block.

The DMAC ensures that data is not retained in per-channel buffers between descriptors. When a descriptor is completely processed, any read data that is buffered in the channel is discarded and any write data that is buffered in the channel is sent to memory (although it may not be there yet). The DMA interrupt is not posted until the descriptor is completely processed.

5.1.7 Byte Transfer Order

The DCMD[ENDIAN] bit indicates the byte ordering in a word when data is read from or written to memory. Refer to Figure 5-5 for details. The DCMD[ENDIAN] bit must be set to 0, which is little endian transfers.

Figure 5-5 shows the order which data is transferred as determined by the DCMD[ENDIAN] and DCMD[SIZE] bits.

If data is being transferred from an internal device to memory, DCMD[ENDIAN] is set to a 0, and DCMD[SIZE] is set to a 1, the memory receives the data in the following order:

- 1. Byte[0]
- 2. Byte[1]
- 3. Byte[2]
- 4. Byte[3]



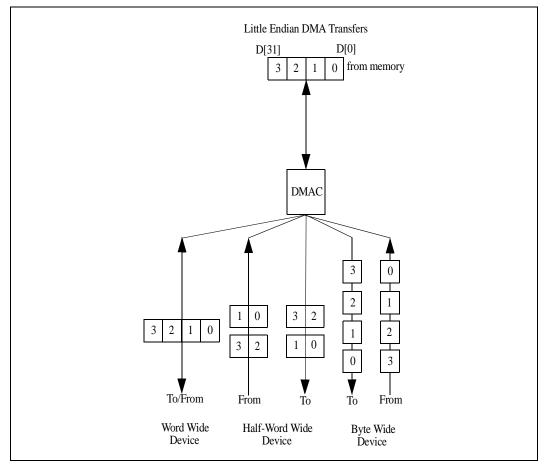


Figure 5-5. Little Endian Transfers

5.1.8 Trailing Bytes

The DMA normally transfers bytes equal to the transaction size specified by DCMD[SIZE]. As the descriptor nears the end its data, the number of trailing bytes in the DCMD[LENGTH] field may be smaller than the transfer size. The DMA can transfer the exact number of trailing bytes if the DCMD[FLOWSRC] and DCMD[FLOWTRG] bits are both 0 or if it receives a corresponding request from a peripheral or companion chip.

Trailing bytes must be considered in the following cases:

- Memory-to-Memory Moves: The DMA transfers a number bytes equal to the smaller of DCMD[LENGTH] or DCMD[SIZE].
- Companion-Chip Related Transfers: The companion-chip must assert the request if the DMAC must handle the trailing bytes. If the request is asserted, the DMA transfers a number of bytes equal to the smaller of DCMD[LENGTH] or DCMD[SIZE].
- Memory to Internal Peripheral Transfers: Most peripherals send a request for trailing bytes
 during memory to internal peripheral transfers. Refer to the appropriate section in this
 document for details of a peripheral's operation. The DMA transfers bytes equal to the smaller
 of DCMD[LENGTH] or DCMD[SIZE].



- Internal Peripheral to Memory Transfers: Most peripherals do not send a request for trailing bytes for on-chip peripheral to memory transfers. Refer to the appropriate section in this document for details of a peripheral's operation. If the peripheral sends out a request, the DMA transfers the number of bytes equal to the smaller of DCMD[LENGTH] or DCMD[SIZE]. If software must us programmed I/O to handle the trailing bytes, it must follow this sequence of operation:
- 1. Writing a 0 to the DCSR[RUN] bit to stop the DMA channel.
- 2. Wait until the channel to stops.
- 3. Make reads to the channel's registers to check the channel's status.
- 4. Perform the programmed I/O transfers to the peripheral.
- 5. Set the DCSR[RUN] bit to a 1 and reset the DMA channel for future data transfers.

5.2 Transferring Data

The internal peripherals are connected to the DMAC via the peripheral bus and use flow-through data transfers. The DMAC can also transfer data to and from any memory location with memory-to-memory moves in flow-through transfer mode. External devices, such as companion chips, that are directly connected to the external data pins must use flow-through data transfers.

Main memory includes any memory that the processor supports, except writes to flash. Writes to flash are not supported and cause a bus error.

In flow-through transfer mode, data passes through the DMAC before it is latched by the destination in its buffers/memory. The DMAC can also perform memory-to-memory moves in flow-through transfer mode.

5.2.1 Servicing Internal Peripherals

The DMAC provides the DMA Request to Channel Map Registers (DRCMRx) that contain four bits used to assign a channel number for each possible DMA request. An internal peripheral can be mapped to any of the 16 available channels. See Table 5-5 to configure the internal peripherals for DMA accesses. Internal peripherals assert the request bit through the peripheral request bus (PREQ). The signals from the PREQ are sampled on every peripheral clock (PCLK) and if any of the PREQ signals are not zeroes, a lookup is performed on the corresponding bits of the DRCMRx. This allows the request to be mapped to one of the channels.

If the internal peripheral address is in the DSADR, the DCMDx[FLOWSRC] bit must be set to a 1. This allows the processor to wait for the request before it initiates the transfer. If the internal peripheral address is in the DTADR, the DCMDx[FLOWTRG] bit must be set to a 1.

If DCMDx[IRQEN] is set to a 1, a DMA interrupt is requested at the end of the last cycle associated with the byte that caused DCMDx[LENGTH] to decrement to 0.



5.2.1.1 Using Flow-Through DMA Read Cycles to Service Internal Peripherals

A flow-through DMA read for an internal peripheral begins when the internal peripheral sends a request, via the PREQ bus, to a DMAC channel that is running and configured for a flow-through read. The number of bytes to be transferred is specified with DCMDx[SIZE]. When the request is the highest priority request, the following process begins:

- 1. The DMAC sends the memory controller a request to read the number of bytes addressed by DSADRx[31:0] into a 32-byte staging buffer in the DMAC.
- 2. The DMAC transfers the data to the I/O device addressed in DTADRx[31:0]. DCMD[WIDTH] specifies the width of the internal peripheral to which the data is transferred.
- 3. At the end of the transfer, DSADRx is increased by the smaller value of DCMDx[LENGTH] and DCMD[SIZE]. DCMDx[LENGTH] is decreased by the same value.

For a flow-through DMA read to an internal peripheral, use the following settings for the DMAC register bits:

- DSADR[SRCADDR] = external memory address
- DTADR[TRGADDR] = internal peripheral's address
- DCMD[INCSRCADDR] = 1
- DCMD[FLOWSRC] = 0
- DCMD[FLOWTRG] = 1

5.2.1.2 Using Flow-Through DMA Write Cycles to Service Internal Peripherals

A flow-through DMA write for an internal peripheral begins when the internal peripheral sends a request, via the PREQ bus, to a DMAC channel that is running and configured for a flow-through write. The number of bytes to be transferred are specified with DCMDx[SIZE]. When the request is the highest priority request, the following process begins:

- The DMAC transfers the required number of bytes from the I/O device addressed by DSADRx[31:0] to the DMAC write buffer.
- 2. The DMAC transfers the data to the memory controller via the internal bus. DCMD[WIDTH] specifies the width of the internal peripheral to which the transfer is being made.
- 3. At the end of the transfer, DTADRx is increased by the smaller value of DCMDx[LENGTH] and DCMD[SIZE]. DCMDx[LENGTH] is decreased by the same number.

For a flow-through DMA write to an internal peripheral, use the following settings for the DMAC register bits:

- DSADR[SRCADDR] = internal peripheral address
- DTADR[TRGADDR] = external memory address
- DCMD[INCTRGADDR] = 1
- DCMD[FLOWSRC] = 1
- DCMD[FLOWTRG] = 0



5.2.2 Quick Reference for DMA Programming

Use Table 5-5 as a quick reference sheet for programming the DMA.

Table 5-5. DMA Quick Reference for Internal Peripherals (Sheet 1 of 2)

Unit	Function	FIFO Address	Width (bytes)	DCMD. Width (binary)	Burst Size (bytes)	Source or Target	DRCMR
12S	receive	0x4040_0080	4	11	8, 16, 32	Source	0x4000_0108
123	transmit	0x4040_0080	4	11	8, 16, 32	Target	0x4000_010c
	receive	0x4020_0000	1	01	8, 16, 32	Source	0x4000_0110
BTUART	transmit	0x4020_0000	1	01	8, 16, 32 or trailing	Target	0x4000_0114
	receive	0x4010_0000	1	01	8, 16, 32	Source	0x4000_0118
FFUART	transmit	0x4010_0000	1	01	8, 16, 32 or trailing	Target	0x4000_011c
	microphone	0x4050_0060	4	11	8, 16, 32	Source	0x4000_0120
	modem receive	0x4050_0140	4	11	8, 16, 32	Source	0x4000_0124
AC97	modem transmit	0x4050_0140	4	11	8, 16, 32	Target	0x4000_0128
	audio receive	0x4050_0040	4	11	8, 16, 32	Source	0x4000_012c
	audio transmit	0x4050_0040	4	11	8, 16, 32	Target	0x4000_0130
SSP	receive	0x4100_0010	2	10	8, 16	Source	0x4000_0134
337	transmit	0x4100_0010	2	10	8, 16	Target	0x4000_0138
	receive	0x4080_000C	1	01	8, 16, 32	Source	0x4000_0144
FICP	transmit	0x4080_000C	1	01	8, 16, 32 or trailing	Target	0x4000_0148
	receive	0x4070_0000	1	01	8, 16, 32	Source	0x4000_014c
STUART	transmit	0x4070_0000	1	01	8, 16, 32, or trailing	Target	0x4000_0150
MMC	receive	0x4110_0040	1	01	32 or trailing	Source	0x4000_0154
IVIIVIC	transmit	0x4110_0044	1	01	32 or trailing	Target	0x4000_0158



Table 5-5. DMA Quick Reference for Internal Peripherals (Sheet 2 of 2)

Unit	Function	FIFO Address	Width (bytes)	DCMD. Width (binary)	Burst Size (bytes)	Source or Target	DRCMR
	endpoint 1 transmit	0x4060_0100	1	01	32	Target	0x4000_0164
	endpoint 2 receive	0x4060_0180	1	01	32	Source	0x4000_0168
	endpoint 3 transmit	0x4060_0200	1	01	32	Target	0x4000_016C
	endpoint 4 receive	0x4060_0400	1	01	32	Source	0x4000_0170
	endpoint 6 transmit	0x4060_0600	1	01	32	Target	0x4000_0178
USB	endpoint 7 receive	0x4060_0680	1	01	32	Source	0x4000_017C
ОЗБ	endpoint 8 transmit	0x4060_0700	1	01	32	Target	0x4000_0180
	endpoint 9 receive	0x4060_0900	1	01	32	Source	0x4000_0184
	endpoint 11 transmit	0x4060_0B00	1	01	32	Target	0x4000_018C
	endpoint 12 receive	0x4060_0B80	1	01	32	Source	0x4000_0190
	endpoint 13 transmit	0x4060_0C00	1	01	32	Target	0x4000_0194
	endpoint 14 receive	0x4060_0E00	1	01	32	Source	0x4000_0198

5.2.3 Servicing Companion Chips and External Peripherals

Companion chips and external peripherals can be serviced with flow-through transfers. The DMAC provides DMA Request to Channel Map Registers (DRCMRx) that contain four bits that assign a channel number for each of the possible DMA requests. The companion-chip requests are DREQ[1:0]. The DREQ signal can be mapped to one of the 16 available channels. The DREQ signals are sampled on every peripheral clock (PCLK) and if any of the DREQ signals are sampled non-zero, a lookup is performed on the corresponding bits in the DRCMRx. This allows requests to one of the channels to be mapped. If the external peripheral address is in the DSADR, the DCMDx[FLOWSRC] bit must be set to a 1. If the external peripheral address is in the DTADR, the DCMDx[FLOWTRG] bit must be set to a 1. This allows the processor to wait for the request before it initiates the transfer.

If DCMDx[IRQEN] is set to a 1, a DMA interrupt can be requested at the end of the last cycle associated with the byte that caused DCMDx[LENGTH] to decrease from a 1 to a 0.



5.2.3.1 Using Flow-Through DMA Read Cycles to Service External Peripherals

A flow-through DMA read for an external peripheral begins when the external peripheral sends a request, via the DREQ[1:0] bus, to a DMAC channel that is running and configured for a flow-through read. DCMDx[SIZE] specifies the number of bytes to be transferred. When the request is the highest priority request, the follow process begins.

- 1. The DMAC sends a request to the memory controller to read the number of bytes addressed by DSADRx[31:0] into a 32-byte staging buffer in the DMAC.
- 2. The DMAC transfers the data in the buffer to the external device addressed in DTADRx[31:0].
- 3. At the end of the transfer, DSADRx is increased by the smaller value of DCMDx[LENGTH] and DCMD[SIZE]. DCMDx[LENGTH] is decreased by the same value.

Note: The process shown for a flow-through DMA read to an external peripheral indicates that the external address increases. Some external peripherals, such as FIFOs, do not require an increment in the external address.

For a flow-through DMA read to an external peripheral, use the following settings for the DMAC register bits:

- DSADR[SRCADDR] = external memory address
- DTADR[TRGADDR] = companion chip's address
- DCMD[INCSRCADDR] = 1
- DCMD[INCTRGADDR] = 0
- DCMD[FLOWSRC] = 0
- DCMD[FLOWTRG] = 1

5.2.3.2 Using Flow-Through DMA Write Cycles to Service External Peripherals

A flow-through DMA write to an external peripheral begins when the external peripheral sends a request, via the DREQ bus, to a DMAC channel that is running and configured for a flow-through write. DCMDx[SIZE] specifies the number of bytes to be transferred. When the request is the highest priority request, the following process begins:

- 1. The DMAC transfers the required number of bytes from the I/O device addressed by DSADRx[31:0] to the DMAC write buffer.
- 2. The DMAC transfers the data to the memory controller via the internal bus.
- 3. At the end of the transfer, DTADRx is increased by the smaller value of DCMDx[LENGTH] and DCMD[SIZE]. DCMDx[LENGTH] is decreased by the same number.

Note: The process shown for a flow-through DMA write to an external peripheral indicates that the external address increases. Some external peripherals, such as FIFOs, do not require an increment in the external address.



For a flow-through DMA write to an external peripheral, use the following settings for the DMAC register bits:

- DSADR[SRCADDR] = companion chip address
- DTADR[TRGADDR] = external memory address.
- DCMD[INCSRCADDR] = 0
- DCMD[INCTRGADDR] = 1
- DCMD[FLOWSRC] = 1
- DCMD[FLOWTRG] = 0

5.2.4 Memory-to-Memory Moves

Memory-to-memory moves do not involve the DREQ and PREQ request signals. The processor writes to the DCSR[RUN] bit and a channel is configured for a memory-to-memory move. The DCMDx[FLOWSRC] and the DCMD[FLOWTRG] bits must be set to 0.

If DCMD[IRQEN] is set to a 1, a DMA interrupt is requested at the end of the last cycle associated with the byte that caused DCMDx[LENGTH] to decrease from 1 to 0.

A flow-through DMA memory-to-memory read or write goes through these steps:

- The processor writes to the DCSR[RUN] register bit and starts the memory-to-memory moves.
- 2. If the processor is in the Descriptor Fetch Mode, the channel configured for the move fetches the four-word descriptor. The channel transfers data without waiting for PREQ or DREQ to be asserted. The smaller value of DCMDx[SIZE] or DCMDx[LENGTH] specifies the number of bytes to be transferred.
- 3. The DMAC sends a request to the memory controller to read the number of bytes addressed by DSADRx[31:0] into a 32-byte staging buffer in the DMAC.
- 4. The DMAC generates a write cycle to the location addressed in DTADRx[31:0].
- 5. At the end of the transfer, DSADRx and DTADRx are increased by the smaller value of DCMD[SIZE] and DCMDx[LENGTH]. If DCMD[SIZE] is smaller than DCMDx[LENGTH], DCMDx[LENGTH] is decreased by DCMD[SIZE]. If DCMD[SIZE] is equal to or larger than DCMDx[LENGTH], DCMDx[LENGTH] is zero.

Note: The process shown for a memory-to-memory transfer indicates that the external address increases. Some external peripherals, such as FIFOs, do not require an increment in the external address.

For a memory-to-memory read or write, use these settings for the DMAC registers:

- DSADR[SRCADDR] = external memory address
- DTADR[TRGADDR] = external memory address
- DCMD[INCSRCADDR] = 1
- DCMD[INCTRGADDR] = 1
- DCMD[FLOWSRC] = 0
- DCMD[FLOWTRG] = 0
- DCSR[RUN] =1



5.3 DMAC Registers

The section describes the DMAC registers.

5.3.1 DMA Interrupt Register (DINT)

The DINT, shown in Table 5-6, logs the interrupts for each channel.

An interrupt is generated if any of these events occur:

- Any kind of transaction error on the internal bus that is associated with the relevant channel.
- The current transfer finishes successfully and the DCMD[ENDIRQEN] bit is set.
- The current descriptor loads successfully and the DCMD[STARTIRQEN] bit is set.
- The DCSR:STOPIRQEN is set to a 1 and the relevant channel is in the uninitialized or stopped state.

Software must set the corresponding DCSR register error bit to reset the interrupt.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 5-6. DINT Bit Definitions

			P	•		Add 0_00		SS				DM/	A Int	terri	upt	Reg	iste	r (D	INT)				0	MA	Co	ntro	oller	r			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							ı	ese	rve	d							Chlintr15	Chlintr14	Chlintr13	Chlintr12	Chlintr11	Chlintr10	Chlintr9	ChlIntr8	ChlIntr7	ChlIntr6	ChlIntr5	Chlintr4	ChlIntr3	Chlintr2	Chllntr1	Chlintr0
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Na	me												De	scr	iptio	on										
		31	:16			-			res	erve	ed																					
		15	5:0		С	HLII	NTF	₹x	0	_	No i	' Int inter rrup	rupt	pt (r	ead	-onl	y).															

5.3.2 DMA Channel Control/Status Register (DCSRx)

The DCSRx, shown in Table 5-7 contains the control and status bit for each channel. Read this register to find the source of an interrupt. Write the read value back to the register to clear the interrupt.



Table 5-7. DCSRx Bit Definitions (Sheet 1 of 2)

																•																
		0x4		•	ical / 00 -				030	C		DM	_			Cor (DC			itus	;				I	OM	A Co	ontr	olle	r			
Bit	31	30	29	28	27	26	25	24	2:	3 22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	RUN	NODESCFETCH	STOPIRGEN									r	ese	rve	d									REQPEND	•	rese	rve	d	STOPSTATE	ENDINTR	STARTINTR	BUSERRINTR
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0
		Bi	ts			Nar	ne												De	escr	ipti	on										
		3	1			RU	IN		Le bu So At D ar ch	un Bi	stop start oftwa vill core n ne cl [STO rrup el.	os the ts the are s omp nust hanr OPS ot afte ing the	e che ch tart lete writ nel s TAT er th	nann or s bef te to stops E] b ne c	top fore DD s, th bit to han bit to x[LE	the ADI ne Do det nel s	cha Rx b CSF erm stop	nnel efor R[ST ine t s. Se	is some it of the oftware.	stop sets STA cha vare	ped s thi TE] nne mu in e	l. s bit i l's s st w	for is se tatu rite	Des et to s or a 1 f	crip 1. S set to th	otor Softv the he b	Fetoware STO it to	th Me mu DPIF rest	ode ist p RQE art a	oll ti N to a sto	he ofor oppe	ce ed n if
		30	0			ODI FET			D ₁	o-De 0 – 1 – etern this t	Des No-I nine	cript Desc s if tl	or F cript he d	etc or F char 0, 1	h M etch nnel the	ode n Mo has char	ode a d	escr is in	•		pto	r Fet	tch l	Mod	e. S	See ·	Sec	tion	5.1.	4.2 ·	for	

If this bit is set to a 1, the channel is in No-Descriptor Fetch Mode. See Section 5.1.4.1 for



Table 5-7. DCSRx Bit Definitions (Sheet 2 of 2)

Physical Address 0x4000_0000 - 0x4000_003C DMA Channel Control/Status Register (DCSRx) **DMA Controller**

Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	RUN	NODESCFETCH	STOPIRGEN										rese	erve	d									REQPEND	r	ese	rve	d	STOPSTATE	ENDINTR	STARTINTR	BUSERRINTR
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0

Bits	Name	Description
		Stop State (read-only). 0 – channel is running 1 – channel is in uninitialized or stopped state.
3	STOPSTATE	If the channel is in the uninitialized or stopped state, this status bit is set. If the DCSR[STOPIRQEN] is set to 1, the DMAC generates an interrupt. Channel States are described in Section 5.1.5.
		Software must reprogram the DDADRx and write a 1 to the DCSR[RUN] bit to restart the channel and clear this bit.
		Software must write a 0 to the DCSR[STOPIRQEN] bit to reset the interrupt
		End Interrupt (read / write).
2	ENDINTR	 0 - no interrupt 1 - interrupt caused because the current transaction was successfully completed and DCMD[LENGTH] = 0.
		DCMD[ENDIRQEN] bit must be set for an interrupt to occur. Software must write a 1 to this bit to reset the corresponding interrupt. Writing a 0 to this bit has no effect.
		Start Interrupt (read / write).
1	STARTINTR	0 – no interrupt 1 – interrupt caused due to successful descriptor fetch
		DCMD[STARTIRQEN] bit must be set for an interrupt to occur. Software must write a 1 to this bit to reset the corresponding interrupt. Writing a 0 to this bit has no effect.
		Bus Error Interrupt (read / write).
		0 - no interrupt 1 - bus error caused interrupt
0	BUSERR INTR	Indicates that there was an error while transferring data. An error during data transfer occurs when the channel has a bad descriptor, source, or target address. An address is considered bad when it points to a non-busable location or reserved space. Software must write a 1 to this bit to reset the corresponding interrupt. Writing a 0 to this bit has no effect. Only one error incidence per channel is logged. The channel that caused the error is updated at the end of the transfer and is accessible after it logs an error until it is reprogrammed and the corresponding run bit is set.



5.3.3 DMA Request to Channel Map Registers (DRCMRx)

DRCMRx, shown in Table 5-8, map each DMA request to a channel. Refer to Table 5-13 for details.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 5-8. DRCMRx Bit Definitions

		0x40		hysi _01					9C			MA		que gis					Ma	p					DM.A	A Co	ontro	olle	r			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d											MAPVLD		reserved		C	HLI	NUI	VI
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bit	Bits Name 31:8 —																De	escr	ipti	on										
		31:							res	erve	ed																					
		7			N	ИΑР	'VL[)	0 1 Det	– I – F erm ues	ines t is r	uest uest wh nap	is is is rether	unm map er th	app ped ne ro a ch	l to a eque nann	est i el ir	s m	appe	ed to in E	a o DRC	han MR	nel x[3:	or n 0]. Ii	x[3: not. I f the	f the						
		6:4	4			-	-		res	erve	ed																					
				C	HLI	NUN	Л	Indi Do	cate	map	e c	han o ac	nel ctive	nun rec	nber ques	if C	o the		me (char	nnel	. It p	rod	a 1.	s un	pred	dicta	ıble	resu	ılts.		

5.3.4 DMA Descriptor Address Registers (DDADRx)

The DDADRx) (see Table 5-9) contain the memory address of the next descriptor for a specific channel. On power up, the bits in this register are undefined. The address must be aligned to a 16-byte boundary. This means that bits [3:1] of the address are reserved and must be read and written as zeroes. DDADR must not contain the address of any other internal peripheral register or DMA register.

DDADR is reserved if the channel is No-Descriptor Fetch Mode.



Table 5-9. DDADRx Bit Definitions

			0x	4000	_02	2x0			ľ	M	A De	esc			Add		s R	legis	ter				ı	DM <i>A</i>	A Co	ontro	olle	r			
Bit	31 30	29	28	27	26	25	24	23 2	2 2	1 2	20 ′	19	18	17	16	15	14	1 13	12	11	10	9	8	7	6	5	4	3	2	1	0
										DES	SCR	IP	TOR	RAI	DDR	ESS	S											res	serve	ed	STOP
Reset													U	nini	itiali	zed	ı														0
	Ві	ts			Naı	me												De	esc	ript	ion										
	31	:4				IPT RES		Addre	ss	of r	next	de	scri	ptor	r (rea	ad /	wr	ite).													
	3	:1			_	-		reserv	/ed																						
	()			ST	OP		Stop (0 - 1 - If this it fetc on the	Ru Sto de bit nes	un o op o scri is s the	char char iptor set, t e ne	nne nne r, i.e the ext o	l. l aft e., [cha desc	OCN inne cript	/ID[L el to tor. I	ENG stop f the	GT os a e D	DAD). it cc)Rx[mp ST(letel DP] i	y pro	oces	sses a ne	the	des escr	scrip ripto	otor	and b	efc	ore

5.3.5 DMA Source Address Registers

DSADRx, shown in Table 5-10, are read only in the Descriptor Fetch Mode and are read/write in the No-Descriptor Fetch Mode.

DSADRx contains the Source Address for the current descriptor of a specific channel. The Source Address is the address of the internal peripheral or a memory location. On power up, the bits in this register are undefined. If the Source Address is the address of a companion chip or external peripheral, the source address must be aligned to an 8-byte boundary. This allows bits [2:0] of the address to be reserved. If the source address is the address for an internal peripheral, the address must be 32-bit aligned, so bits [1:0] are reserved. DSADR cannot contain the address of any other internal DMA, LCD, or MEMC registers.



Table 5-10. DSADRx Bit Definitions

	0	x4000_02x4	DMA Source Addr Register (DSADRx) DMA Control	oller		
Bit	31 30 29 2	8 27 26 25 24	23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5	4 3	3 2	1 0
			SOURCE ADDRESS			reserved
Reset			Uninitialized			
	Bits	Name	Description			
	31:3	SRCADDR	Source Address (read / write). Address of the internal peripheral or address of a memory location. Address of a memory location for companion -chip transfer			
	31:3	SRCADDR	Address of the internal peripheral or address of a memory location.			

5.3.6 DMA Target Address Registers (DTADRx)

To software, DTADRx (Table 5-11) is read only in the Descriptor Fetch Mode and is read/write in the No-Descriptor Fetch Mode.

These registers contain the target address for the current descriptor in a channel. The target address is the address of an internal peripheral or a memory location. On power up, the bits in this register are undefined. If the target address is the address of a companion chip or external peripheral, the target address must be aligned to an 8-byte boundary. This allows bits [2:0] of the address to be reserved. If the target address is the address for an internal peripheral, the address must be 32-bit aligned so that bits [1:0] are reserved. DTADRx must not contain the address of any other internal DMA, LCD, or MEMC register.

The DTADRx must not contain a flash address because writes to flash from the DMAC are not supported.



Table 5-11. DTADRx Bit Definitions

	O	x4000_02x8	DMA Target Addr Register (DTADRx) DMA	Contr	ollei	r		
Bit	31 30 29 2	8 27 26 25 24	23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7	5	4	3	2	1 0
			TARGET ADDRESS					reserved
Reset			Uninitialized					
	Bits	Name	Description					
	31:3	TRGADDR	Target Address (read / write): Address of the on chip peripheral or the address of a memory local Address of a memory location for companion chip transfer	tion				
	2	TRGADDR	Target Address Bit 2 Reserved if DTADR.TrgAddr is an external memory location Not reserved if DTADR.trgAddr is an internal peripheral (read / writer)	te).				
	1:0	_	reserved					

5.3.7 DMA Command Registers (DCMDx)

For software, DCMDx (Table 5-12) is read only in Descriptor Fetch Mode and is read/write in No-Descriptor Fetch Mode.

These registers contain the channel's control bits and the length of the current transfer in that channel. On power up, the bits in this register are set to 0.



Table 5-12. DCMDx Bit Definitions (Sheet 1 of 2)

	0x4000_02xC						DMA Command Register (DCMDx)											DMA Controller														
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	INCSRCADDR	INCTRGADDR	FLOWSRC	FLOWTRG		re	serv	ed		STARTIRGEN	ENDIRGEN	pokrosor		ENDIAN	SIZE	1	WIDTH		reserved						LE	NG	тн					
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

· _	0 0 0 0	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0									
	Bits	Name	Description									
	31	INCSRCADDR	Source Address Increment Setting. (read / write). 0 - Do not increment Source Address 1 - Increment Source Address at the end of each internal bus transaction initiation by DCMD[SIZE] If the source address is an internal peripheral's FIFO address or external IO address, the address is not incremented on each successive access. In this case, this bit must be 0.									
	30	INCTRGADDR	Target Address Increment Setting. (read / write). 0 - Do not increment Target Address 1 - Increment Target Address at the end of each internal bus transaction initiated by DCMD[SIZE] If the target address is an internal peripheral's FIFO address or external IO address, the address is incremented on each successive access. In this cases the bit must be 0.									
	29	FLOWSRC	Flow Control by the source. (read / write). 0 — Start the data transfer immediately. 1 — Wait for a request signal before initiating the data transfer. Indicates the flow control of the source. This bit must be '1' if the source is an onchip or external peripheral. If either the DCMD[FLOWSRC] or DCMD[FLOWTRG] bit is set, the current DMA does not nitiate a transfer until it receives a request. Do not set both the DCMD[FLOWTRG] and DCMD[FLOWSRC] bit to1.									
	28	FLOWTRG	Flow Control by the target. (read / write). 0 — Start the data transfer immediately. 1 — Wait for a request signal before initiating the data transfer. Indicates the Flow Control of the target. This bit must be '1' if the target is an onchip or external peripheral. If either the DCMD[FLOWSRC] or DCMD[FLOWTRG] bit is set, the current DMA does not initiate a transfer until it receives a request. Do not set both the DCMD[FLOWTRG] and DCMD[FLOWSRC] bit to 1.									
Ī	27:23	_	reserved									
	22	Start Interrupt Enable (read / write), Reserved for the No-Descriptor Fetch Mode 0 - no interrupt is generated. 1 - Allow interrupt to pass when the descriptor (i.e., 4 words) for the channel at Sets DCSR[StartIntr] interrupt for the channel when this descriptor is loaded.										
	21	ENDIRQEN	End Interrupt Enable (read / write). 0 - no interrupt is generated. 1 - set DCSR[EndIntr] interrupt for the channel when DCMD[LENGTH] is decreased to zero. Indicates that the interrupt is enabled as soon as the data transfer is completed.									
ſ	20:19		reserved									
_												



Table 5-12. DCMDx Bit Definitions (Sheet 2 of 2)

Table 3-12. DOMDX bit Definitions (Sheet 2 of 2)																															
	0x4000_02xC						DMA Command Register (DCMDx)											DMA Controller													
Bit	31	30	29	28	27	26	25	24	23 2	2 2	1 2	0 19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	INCSRCADDR	INCTRGADDR	FLOWSRC	FLOWTRG		res	serve	ed		SIAKIIKGEN	ENDIRGEN	reserved	ENDIAN		SIZE	WIDTH		reserved						LE	ENG	тн					
Reset	0	0	0	0	0	0	0	0	0 (()	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	Bits Name					Description																									
	18					END	IAN		0 –	Ву	ndian-ness. (read / write). te ordering is little endian served																				
	17:16 SIZE								Maximum Burst Size of each data transferred (read / write). 00 – reserved 01 – 8 Bytes 10 – 16 Bytes 11 – 32 Bytes If DCMDx[LENGTH] is less than DCMDx[SIZE] the data transfer size equals DCMDx[LENGTH].																						
	15:14 WID					ЭТН		Width of the on-chip peripheral. (read / write/). 00 – reserved 01 – 1 byte 10 – HalfWord (2 bytes) 11 – Word (4 Bytes) Must be programmed 00 for memory-to-memory moves or companion-chip related operations.																							
	13 — reserved																														
	Length of transfer in bytes. (read / write). LENGTH LENGTH								Des	ny c	of																				



5.4 Examples

This section contains examples that show how to:

- Set up and start a channel
- Initialize a descriptor list for a channel that is running
- Add a descriptor to the end of a descriptor list for a channel that is running
- Initialize a channel that is going to be used by a direct DMA master

Example 1. How to set up and start a channel:

The following example shows how to set up a channel to transfer LENGTH words from the address DSADR to the I/O address DTADR. The example also shows how to start the transfer. The example sets the stop bit in the DDADR, so the DMA channel stops after it completely transfers LENGTH bytes of data.

```
// build real descriptor
desc[0].ddadr = STOP
desc[0].dsadr = DSADR
desc[0].dtadr = DTADR
desc[0].dcmd = DCMD

// start the channel
DMANEXT[CHAN] = &desc[0]
DRUN = 1
```

Example 2. How to initialize a descriptor list for a channel that is running:

```
// Allocate a new descriptor, and make it an end
// descriptor whose "ddadr" field points back at itself
newDesc = New Desc()
newDesc->ddadr = newDesc | STOP
// make it a zero length descriptor
newDesc->dcmd = ZERO
// Start the channel
DMANEXT[CHAN] = newDesc
DRUN = 1
```

The channel starts, loads the descriptor in its registers, and stops because the transfer length is 0 and the STOP bit is set. No data is transferred in this example. The channel can be restarted by writing to its DDADR and writing a 1 to the DCSR[RUN] bit.

Example 3. How to add a descriptor to the end of a descriptor list for a channel that is running:

The example in this section assumes that the Descriptor Fetch Mode is active.

DMA descriptor lists are used as queues of full buffers for network transmitters and as queues of empty buffers for network receivers. Because the buffers in a queue are often small (in particular, as small as an ATM cell), on-the-fly DMA descriptor lists manipulation must be efficient.

- 1. Write a 0 to DCSR[RUN].
- 2. Wait until the channel stops. The channel stop state is reflected in the DCSR:STOPSTATE bit.



- 3. In memory, create the descriptor to be added and set its stop bit to a 1.
- 4. In the memory, manipulate the DDADR of the current chain's last descriptor such that its DDADR points to the descriptor created in Step 3.
- 5. In the memory, create a new descriptor that has the same DDADR, DSADR, DTADR, and CMD as those of the stopped DMA channel. The new descriptor is the next descriptor for the list
- 6. Examine the DMA channel registers and determine if the channel stopped in the chain's last descriptor of the chain. If it did, manipulate the DDADR of the last descriptor in the memory so that its DDADR points to the descriptor created in Step 3. Otherwise, continue to Step 7.
- 7. Program the channel's DDADR with the descriptor created in Step 5.
- 8. Set the DCSR[RUN] to a 1.

Example 4. How to initialize a channel that is going to be used by a direct DMA master:

The most efficient way to move data between an I/O device and main memory is the processor's descriptor-based DMA system. Each application has different requirements, so a descriptor-based DMA may be best for some applications while a non-descriptor-based DMA is best for others. For applications that can not tolerate the time needed to fetch a descriptor before each DMA transfer, choose the non descriptor-based DMA method. For applications that can tolerate it, a descriptor-based DMA method can reduce the amount of core intervention.

Self-Modifying Descriptors: The descriptor-based DMA system can be used to provide true direct memory access to devices that require it.

In this example, a companion chip has these requirements:

- 1. When the companion chip asserts DREQ from 0 to 1, the DMA must fetch four words of the descriptor from one of the chip's ports.
- 2. Based on the information contained in the four descriptor words, the DMA must transfer data from the source address to the destination address without waiting for another request from the companion chip.
- 3. After it transfers the number of bytes in DCMD:LENGTH, the DMA returns to Step 1.

An external device with these requirements can use a constant descriptor in memory.

```
struct {longddadr;
       longdsadr;
       longdtadr;
       shortlength;
       shortdcmd;
} desc[2];
desc[0].ddadr = &desc[1]
desc[0].dsadr = I_ADR + I_DESC_OFFS
desc[0].dtadr = &desc[1].dsadr
desc[0].length = 8;
desc[0].dcmd = CMD_IncTrgAdr | CMD_FlowThru;
desc[1].ddadr = &desc[0]
desc[1].dtadr = I_ADR + I_DATA_OFFS
desc[1].dsadr = 0
desc[1].length = 0
desc[1].dcmd = 0
```



When the external device has data to transfer, it makes a DMA request in the standard way. The DMAC wakes up and reads four words from the device's I_DESC_OFFS address (the DMAC only transfers four words because the first descriptor has an 8-byte count.). The four words from the external device are written in the DSADR, DTADR, and DCMD fields of the next descriptor. The DMAC then steps into the next (dynamically modified) descriptor and, using the I_DATA_OFFS address on the external device, starts the transfer that the external device requested. When the transfer is finished, the DMAC steps back into the first descriptor and the process is repeated.

This example lends itself to any number of variations. For example, a DMA channel that is programmed in this way can be used to transfer messages from a network device directly into client buffers. Each block of data would be preceded by its final destination address and a count.

5.5 DMA Controller Register Summary

This section describes the DMAC Registers. Refer to Table 5-13.

Table 5-13. DMA Controller Register Summary (Sheet 1 of 5)

Address	Name	Description
0x4000_0000	DCSR0	DMA Control / Status Register for Channel 0
0x4000_0004	DCSR1	DMA Control / Status Register for Channel 1
0x4000_0008	DCSR2	DMA Control / Status Register for Channel 2
0x4000_000C	DCSR3	DMA Control / Status Register for Channel 3
0x4000_0010	DCSR4	DMA Control / Status Register for Channel 4
0x4000_0014	DCSR5	DMA Control / Status Register for Channel 5
0x4000_0018	DCSR6	DMA Control / Status Register for Channel 6
0x4000_001C	DCSR7	DMA Control / Status Register for Channel 7
0x4000_0020	DCSR8	DMA Control / Status Register for Channel 8
0x4000_0024	DCSR9	DMA Control / Status Register for Channel 9
0x4000_0028	DCSR10	DMA Control / Status Register for Channel 10
0x4000_002C	DCSR11	DMA Control / Status Register for Channel 11
0x4000_0030	DCSR12	DMA Control / Status Register for Channel 12
0x4000_0034	DCSR13	DMA Control / Status Register for Channel 13
0x4000_0038	DCSR14	DMA Control / Status Register for Channel 14
0x4000_003C	DCSR15	DMA Control / Status Register for Channel 15
0x4000_00F0	DINT	DMA Interrupt Register
0x4000_0100	DRCMR0	Request to Channel Map Register for DREQ 0 (companion chip request 0)
0x4000_0104	DRCMR1	Request to Channel Map Register for DREQ 1 (companion chip request 1)
0x4000_0108	DRCMR2	Request to Channel Map Register for I2S receive Request
0x4000_010C	DRCMR3	Request to Channel Map Register for I2S transmit Request



Table 5-13. DMA Controller Register Summary (Sheet 2 of 5)

Address	Name	Description									
0x4000_0110	DRCMR4	Request to Channel Map Register for BTUART receive Request									
0x4000_0114	DRCMR5	Request to Channel Map Register for BTUART transmit Request.									
0x4000_0118	DRCMR6	Request to Channel Map Register for FFUART receive Request									
0x4000_011C	DRCMR7	Request to Channel Map Register for FFUART transmit Request									
0x4000_0120	DRCMR8	Request to Channel Map Register for AC97 microphone receive Request									
0x4000_0124	DRCMR9	Request to Channel Map Register for AC97 modem receive Request									
0x4000_0128	DRCMR10	Request to Channel Map Register for AC97 modem transmit Request									
0x4000_012C	DRCMR11	Request to Channel Map Register for AC97 audio receive Request									
0x4000_0130	DRCMR12	Request to Channel Map Register for AC97 audio transmit Request									
0x4000_0134	DRCMR13	Request to Channel Map Register for SSP receive Request									
0x4000_0138	DRCMR14	Request to Channel Map Register for SSP transmit Request									
0x4000_013C	DRCMR15	reserved									
0x4000_0140	DRCMR16	reserved									
0x4000_0144	DRCMR17	Request to Channel Map Register for FICP receive Request									
0x4000_0148	DRCMR18	Request to Channel Map Register for FICP transmit Request									
0x4000_014C	DRCMR19	Request to Channel Map Register for STUART receive Request									
0x4000_0150	DRCMR20	Request to Channel Map Register for STUART transmit Request									
0x4000_0154	DRCMR21	Request to Channel Map Register for MMC receive Request									
0x4000_0158	DRCMR22	Request to Channel Map Register for MMC transmit Request									
0x4000_015C	DRCMR23	reserved									
0x4000_0160	DRCMR24	reserved									
0x4000_0164	DRCMR25	Request to Channel Map Register for USB endpoint 1 Request									
0x4000_0168	DRCMR26	Request to Channel Map Register for USB endpoint 2 Request									
0x4000_016C	DRCMR27	Request to Channel Map Register for USB endpoint 3 Request									
0x4000_0170	DRCMR28	Request to Channel Map Register for USB endpoint 4 Request									



Table 5-13. DMA Controller Register Summary (Sheet 3 of 5)

Address	Name	Description
0x4000_0174	DRCMR29	reserved
0x4000_0178	DRCMR30	Request to Channel Map Register for USB endpoint 6 Request
0x4000_017C	DRCMR31	Request to Channel Map Register for USB endpoint 7 Request
0x4000_0180	DRCMR32	Request to Channel Map Register for USB endpoint 8 Request
0x4000_0184	DRCMR33	Request to Channel Map Register for USB endpoint 9 Request
0x4000_0188	DRCMR34	reserved
0x4000_018C	DRCMR35	Request to Channel Map Register for USB endpoint 11 Request
0x4000_0190	DRCMR36	Request to Channel Map Register for USB endpoint 12 Request
0x4000_0194	DRCMR37	Request to Channel Map Register for USB endpoint 13 Request
0x4000_0198	DRCMR38	Request to Channel Map Register for USB endpoint 14 Request
0x4000_019C	DRCMR39	reserved
0x4000_0200	DDADR0	DMA Descriptor Address Register channel 0
0x4000_0204	DSADR0	DMA Source Address Register channel 0
0x4000_0208	DTADR0	DMA Target Address Register channel 0
0x4000_020C	DCMD0	DMA Command Address Register channel 0
0x4000_0210	DDADR1	DMA Descriptor Address Register channel 1
0x4000_0214	DSADR1	DMA Source Address Register channel 1
0x4000_0218	DTADR1	DMA Target Address Register channel 1
0x4000_021C	DCMD1	DMA Command Address Register channel 1
0x4000_0220	DDADR2	DMA Descriptor Address Register channel 2
0x4000_0224	DSADR2	DMA Source Address Register channel 2
0x4000_0228	DTADR2	DMA Target Address Register channel 2
0x4000_022C	DCMD2	DMA Command Address Register channel 2
0x4000_0230	DDADR3	DMA Descriptor Address Register channel 3
0x4000_0234	DSADR3	DMA Source Address Register channel 3
0x4000_0238	DTADR3	DMA Target Address Register channel 3
0x4000_023C	DCMD3	DMA Command Address Register channel 3
0x4000_0240	DDADR4	DMA Descriptor Address Register channel 4
0x4000_0244	DSADR4	DMA Source Address Register channel 4
0x4000_0248	DTADR4	DMA Target Address Register channel 4
0x4000_024C	DCMD4	DMA Command Address Register channel 4
0x4000_0250	DDADR5	DMA Descriptor Address Register channel 5
0x4000_0254	DSADR5	DMA Source Address Register channel 5



Table 5-13. DMA Controller Register Summary (Sheet 4 of 5)

0x4000_0258DTADR5DMA Target Address Register chann0x4000_025CDCMD5DMA Command Address Register chann0x4000_0260DDADR6DMA Descriptor Address Register chann0x4000_0264DSADR6DMA Source Address Register chann0x4000_0268DTADR6DMA Target Address Register chann0x4000_026CDCMD6DMA Command Address Register ch0x4000_0270DDADR7DMA Descriptor Address Register chann0x4000_0274DSADR7DMA Source Address Register chann0x4000_0278DTADR7DMA Target Address Register chann0x4000_027CDCMD7DMA Command Address Register ch0x4000_0280DDADR8DMA Descriptor Address Register chann0x4000_0284DSADR8DMA Source Address Register chann0x4000_0288DTADR8DMA Target Address Register chann0x4000_028CDCMD8DMA Command Address Register chann0x4000_0290DDADR9DMA Descriptor Address Register chann0x4000_0294DSADR9DMA Source Address Register chann0x4000_0298DTADR9DMA Target Address Register chann0x4000_029CDCMD9DMA Target Address Register chann0x4000_029CDCMD9DMA Command Address Register chann0x4000_02AdDDADR10DMA Descriptor Address Register chann0x4000_02AdDDADR10DMA Source Address Register chann0x4000_02AdDSADR10DMA Source Address Register chann	
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0x4000_0288DTADR8DMA Target Address Register chann0x4000_028CDCMD8DMA Command Address Register ch0x4000_0290DDADR9DMA Descriptor Address Register ch0x4000_0294DSADR9DMA Source Address Register chann0x4000_0298DTADR9DMA Target Address Register chann0x4000_029CDCMD9DMA Command Address Register ch0x4000_02A0DDADR10DMA Descriptor Address Register ch	nannel 8
0x4000_028C DCMD8 DMA Command Address Register ch 0x4000_0290 DDADR9 DMA Descriptor Address Register ch 0x4000_0294 DSADR9 DMA Source Address Register chann 0x4000_0298 DTADR9 DMA Target Address Register chann 0x4000_029C DCMD9 DMA Command Address Register ch 0x4000_02A0 DDADR10 DMA Descriptor Address Register ch	nel 8
0x4000_0290 DDADR9 DMA Descriptor Address Register ch 0x4000_0294 DSADR9 DMA Source Address Register chans 0x4000_0298 DTADR9 DMA Target Address Register chans 0x4000_029C DCMD9 DMA Command Address Register ch 0x4000_02A0 DDADR10 DMA Descriptor Address Register ch	nel 8
0x4000_0294 DSADR9 DMA Source Address Register chann 0x4000_0298 DTADR9 DMA Target Address Register chann 0x4000_029C DCMD9 DMA Command Address Register ch 0x4000_02A0 DDADR10 DMA Descriptor Address Register ch	nannel 8
0x4000_0298 DTADR9 DMA Target Address Register chann 0x4000_029C DCMD9 DMA Command Address Register ch 0x4000_02A0 DDADR10 DMA Descriptor Address Register ch	nannel 9
0x4000_029C DCMD9 DMA Command Address Register ch 0x4000_02A0 DDADR10 DMA Descriptor Address Register ch	nel 9
0x4000_02A0 DDADR10 DMA Descriptor Address Register ch	nel 9
	nannel 9
0x4000 02A4 DSADR10 DMA Source Address Register change	nannel 10
	nel 10
0x4000_02A8 DTADR10 DMA Target Address Register chann	nel 10
0x4000_02AC DCMD10 DMA Command Address Register ch	nannel 10
0x4000_02B0 DDADR11 DMA Descriptor Address Register ch	nannel 11
0x4000_02B4 DSADR11 DMA Source Address Register change	nel 11
0x4000_02B8 DTADR11 DMA Target Address Register chann	nel 11
0x4000_02BC DCMD11 DMA Command Address Register ch	nannel 11
0x4000_02C0 DDADR12 DMA Descriptor Address Register ch	nannel 12
0x4000_02C4 DSADR12 DMA Source Address Register change	nel 12
0x4000_02C8 DTADR12 DMA Target Address Register chann	nel 12
0x4000_02CC DCMD12 DMA Command Address Register ch	nannel 12
0x4000_02D0 DDADR13 DMA Descriptor Address Register ch	nannel 13
0x4000_02D4 DSADR13 DMA Source Address Register change	nel 13
0x4000_02D8 DTADR13 DMA Target Address Register chann	nel 13
0x4000_02DC DCMD13 DMA Command Address Register ch	nannel 13
0x4000_02E0 DDADR14 DMA Descriptor Address Register ch	nannel 14
0x4000_02E4 DSADR14 DMA Source Address Register change	nel 14
0x4000_02E8 DTADR14 DMA Target Address Register chann	nel 14



Table 5-13. DMA Controller Register Summary (Sheet 5 of 5)

Address	Name	Description
0x4000_02EC	DCMD14	DMA Command Address Register channel 14
0x4000_02F0	DDADR15	DMA Descriptor Address Register channel 15
0x4000_02F4	DSADR15	DMA Source Address Register channel 15
0x4000_02F8	DTADR15	DMA Target Address Register channel 15
0x4000_02FC	DCMD15	DMA Command Address Register channel 15

Memory Controller

6

This chapter describes the external memory interface structures and memory-related registers supported by the PXA255 processor.

6.1 Overview

The processor external memory bus interface supports Synchronous Dynamic Memory (SDRAM), synchronous and asynchronous burst modes, Page-mode flash, Synchronous Mask ROM (SMROM), Page Mode ROM, SRAM, SRAM-like Variable Latency I/O (VLIO), 16-bit PC Card expansion memory, and Compact Flash. Memory types can be programmed through the Memory Interface Configuration registers. Figure 6-1 is a block diagram of the maximum configuration of the memory controller.



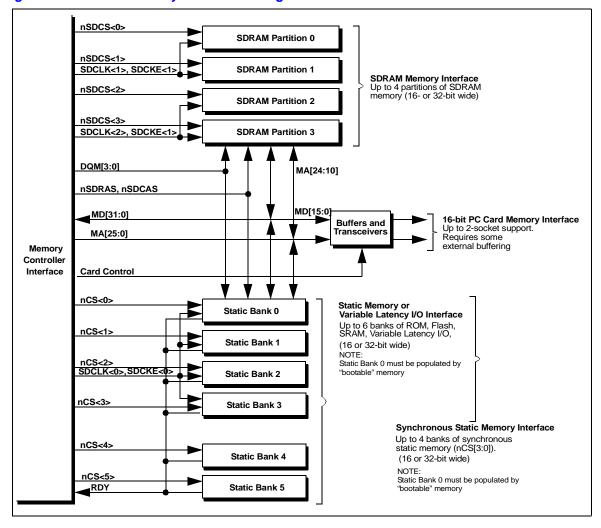


Figure 6-1. General Memory Interface Configuration

6.2 Functional Description

The processor has three different memory spaces: SDRAM, Static Memory, and Card Memory.

SDRAM has four partitions, Static Memory has six, and Card space has two. When memory access attempts to burst across the boundary between adjacent partitions, ensure that the configurations for the partitions are identical. The configurations must be identical in every aspect, including external bus width and burst length.

6.2.1 SDRAM Interface Overview

The processor supports the SDRAM interface, which supports four 16- and 32-bit-wide SDRAM partitions. Each partition is allocated 64 Mbytes of the internal memory map, but the actual size of each partition depends on the SDRAM configuration. The four partitions are divided into two



partition pairs: the 0/1 pair and the 2/3 pair. The partitions in a pair must be identical in size and configuration. The two pairs may be different (for example, the 0/1 pair can be 100 MHz SDRAM on a 32-bit data bus, while the 2/3 pair can be 50 MHz SDRAM on a 16-bit data bus).

The processor SDRAM Controller includes the following signals:

- 4 partition selects (nSDCS[3:0])
- 4 byte selects (DQM[3:0])
- 15 multiplexed bank/row/column address signals (MA[24:10])
- 1 write enable (nWE)
- 1 column-address strobe (nSDCAS)
- 1 row-address strobe (nSDRAS)
- 1 clock enable (SDCKE[1])
- 2 clocks (SDCLK[2:1])
- 32 data (MD[31:0])

The processor performs auto-refresh (CBR) during normal operation, and supports self-refreshing SDRAM during Sleep mode. An SDRAM *auto-power-down mode* bit can be set so that the clock and clock enable to SDRAM are automatically de-asserted whenever none of the corresponding partitions is being accessed.

The processor supports x8, x16, and x32 SDRAM chips.

Upon enabling an SDRAM partition, a mode register set command (MRS), see Section 6.5.6, is sent to the SDRAM devices by writing to the MDMRS register. The PXA255 processor adds support for low-power SDRAM by giving software access to the Extended Mode Register via the MDMRSLP register.

MRS commands always configure SDRAM internal mode registers for sequential burst type and a burst length of four.

The CAS latency is determined by the DTC0 or DTC2 field of MDCNFG.

6.2.2 Static Memory Interface / Variable Latency I/O Interface

The static memory and variable latency I/O interface has six chip selects (nCS[5:0]) and 26 bits of byte address (MA[25:0]) for accesses of up to 64 Mbytes of memory in each of six banks. Each chip select is individually programmed for selecting one of the supported static memory types:

- Non-burst ROM or Flash memory is supported on nCS[5:0]
- Burst ROM or Flash (with non-burst writes) is supported on nCS[5:0]
- Burst and non-burst SRAM is supported on nCS[5:0]
- Variable Latency I/O is supported on nCS[5:0]
- Synchronous static memory is supported on nCS[3:0]

The Variable Latency I/O interface differs from SRAM in that it allows the data-ready input signal, RDY, to insert a variable number of wait states. For all static memory types, each chip select can be individually configured to a 16-bit or 32-bit-wide data bus. nOE is asserted on all reads, nPWE is



asserted on writes to Variable Latency I/O devices, and nWE is asserted on writes to all other static devices, both synchronous and asynchronous. For SRAM and variable latency I/O, DQM[3:0] are byte selects for both reads and writes.

When the processor comes out of reset, it starts fetches and executes instructions at address 0x00, which corresponds to memory selected by nCS<0>. The boot ROM must be located at this address. The BOOT_SEL pins determine the type of boot memory (refer to Section 6.10.1).

6.2.3 16-Bit PC Card / Compact Flash Interface

The processor card interface is based on *The PC Card Standard - Volume 2 - Electrical Specification, Release 2.1*, and *CF+ and CompactFlash Specification Revision 1.4*. The 16-bit PC Card/Compact Flash interface provides control signals to support any combination of 16-bit PC Card/Compact Flash for two card sockets, using address line (MA[25:0]) and data lines (MD[15:0]).

The processor 16-bit PC Card / Compact Flash Controller provides the following signals.

- nPREG is muxed with MA[26] and selects register space (I/O or attribute) versus memory space
- nPOE and nPWE allow memory and attribute reads and writes
- nPIOR, nPIOW, and nIOIS16 control I/O reads and writes
- nPWAIT allows extended access times
- nPCE2 and nPCE1 are byte select high and low for a 16-bit data bus
- PSKTSEL selects between two card sockets

6.3 Memory System Examples

This section provides examples of memory configurations that are possible with the processor. Figure 6-2 shows a system that uses 1M x 16-bit x 4-bank SDRAM devices for a total of 48 Mbytes.



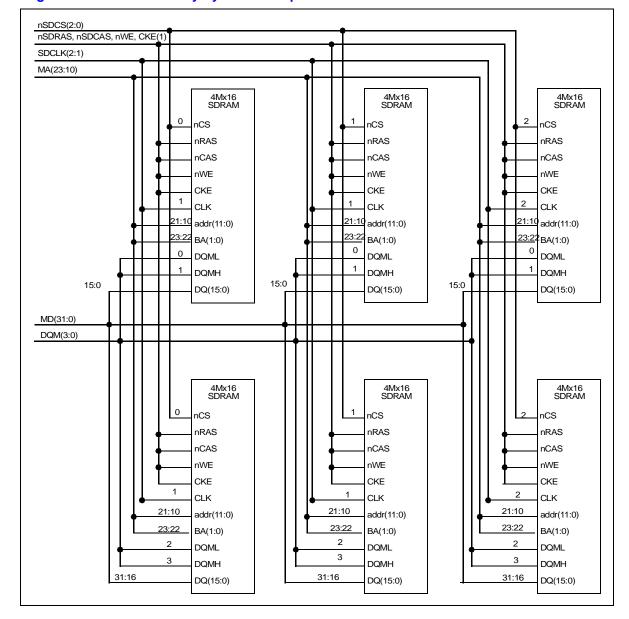
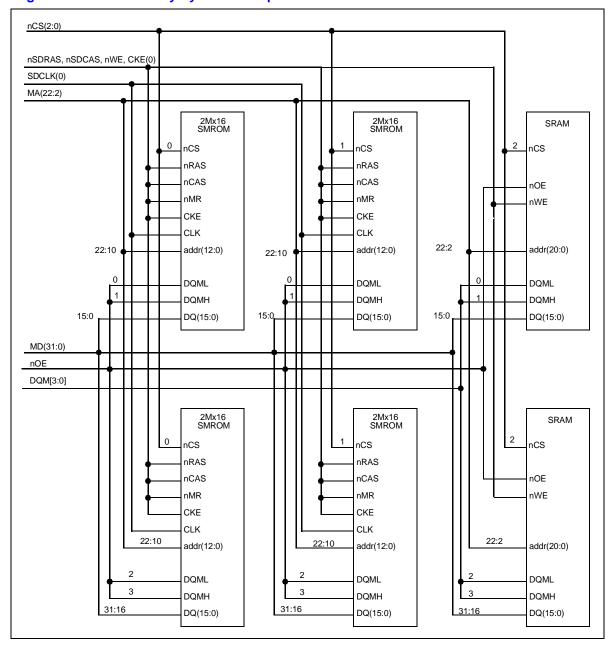


Figure 6-2. SDRAM Memory System Example



Figure 6-3 shows an alternate memory configuration. This system uses 2M x 16 SMROM devices in static banks 0 and 1, and RAM devices in static bank 2.

Figure 6-3. Static Memory System Example





6.4 Memory Accesses

If a memory access is followed by an idle bus period, the control signals return to their inactive state. The address and data signals remain at their previous values to avoid unnecessary bus transitions and eliminate the need for multiple pull-up resistors.

Table 6-1 lists all the transactions that the processor can generate. No burst can cross an aligned 32-byte boundary. On a 16-bit data bus, each full word access becomes a two half-word burst, with address bit 1 set to a 0. Each write access to Flash memory space must take place in one non-burst operation, regardless of the bus size.

Table 6-1. Device Transactions

Bus Operation	Burst Size (Words)	Start Address Bits [4:2]	Description
Read single	1	Any	Generated by core, DMA, or LCD request.
Read burst	4	0 4	Generated by DMA or LCD request.
Read burst	8	0	Generated by cache line fills.
Write single	1	Any	14 bytes are written as specified by the byte mask. Generated by DMA request.
Write burst	2	0,1,2 4,5,6	All 4 bytes of each word are written. Generated by DMA request.
Write burst	3	0,1 4,5	All 4 bytes of each word are written. Generated by DMA request.
Write burst	4	0 4	All 4 bytes of each word are written. Generated by DMA request.
Write burst	8	0	Cacheline copyback. All 32 bytes are written.



6.4.1 Reads and Writes

DQM[3:0] are data masking bits. When asserted (high), the corresponding bit masks the associated byte of data on the MD[31:0] bus. When deasserted (low), the corresponding bit does not mask the associated byte of data on the MD[31:0] bus.

- DQM[3] corresponds to MD[31:24]
- DQM[2] corresponds to MD[23:16]
- DQM[1] corresponds to MD[15:8]
- DQM[0] corresponds to MD[7:0]

For writes to SDRAM, SRAM, or Variable Latency I/O memory spaces, the DQM[3:0] lines enable the corresponding byte of the data bus. Flash memory space stores must be exactly the width of the Flash data bus, either 16- or 32-bits. See Section 6.7.6 for more information.

For reads to all memory types, the DQM[3:0] lines are deasserted (set low so data is not masked).

6.4.2 Aborts and Nonexistent Memory

Accessing reserved portions of the memory map results in a data abort exception.

Hardware does not detect reads and writes from or to enabled and nonexistent memory. If memory in an enabled partition is not present, a read returns indeterminate data.

If memory does not occupy all 64 MB of the partition, reads and writes from or to the unoccupied portion are processed as if the memory occupies the entire 64 MB of the memory partition.

A single word (or half-word if the data bus width is defined as 16-bits) access to a disabled SDRAM partition (MDCNFG:DEx=0) causes a CBR refresh cycle to all four partitions. This technique is used in the hardware initialization procedure. Read return data is indeterminate and writes are not executed on the external memory bus.

A burst read access to a disabled SDRAM partition results in a target-abort exception. Target aborts are also generated for burst writes to Flash/ROM space and bursts to configuration space. Attempted single beat writes to ROM are not aborted. Bursts to configuration space also result in target aborts. Target aborts can be either Data or Prefetch abort depending on the source of the attempted burst transaction.

6.5 Synchronous DRAM Memory Interface

Each possible SDRAM portion of the Memory Map is referred to as a partition, to distinguish them from banks internal to SDRAM devices.

The signals used to control the SDRAM memory are listed in Section 6.2.1.

6.5.1 SDRAM MDCNFG Register (MDCNFG

MDCNFG, shown in Table 6-2, is a read/write register and contains control bits for configuring the SDRAM. Both SDRAM partitions in a pair (0/1 or 2/3) must be implemented with the same type of SDRAM devices, but the two partition pairs may differ.



This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 6-2. MDCNFG Bit Definitions (Sheet 1 of 3)

Ia	DIE	U -	Z. I	VID	CIN	1 0	וט	טו	CIIII	IIII	Ulli	5 (5	,,,,	CL	. 0	. 3)																
				0x4	800	0_00)00							N	IDC	NF	3							Me	emo	ry (Con	trol	ler			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	res	serv	ed	DSA1111_2	DLATCH2	DADDR2	DT	C2	DNB2	C D V d U	7040	27470	7000	DWID2	DE3	DE2	res	serv	ed	DSA1111_0	DLATCH0	DADDR0	DT	C0	DNB0	00,490	2247	0000		DWID0	DE1	DE0
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Naı	me												De	escr	iptic	on										
		()			DE	Ξ0		For MD refre disa	ead CN esh able SI	ch S FG: cyc ed. DRA	DR/ DW/ cle to	ID0=	part ='1') part	ition acc titior	, the ess ns. \	(rea Whe	is an ad oi n all	r wri	ite) t	o a	disa	ble	JŻ b	DRA	Мр	ártit	ion t	rigg	ers a		
		1	I			DE			For MD refre disa	ead CN esh able SI	ch S FG: cyc ed. DRA	DR/ DW/ cle to	ID0=	part ='1') part	ition acc titior	, the ess ns. \	(rea Whe	is an ad oi n all	r wri	ite) t	o a	disa	ble	JŻ b	DRA	Мр	ártit	ion t	rigg	ers		
		2	2			DW	ID0)	SDF 0 - 1 -	32	2 bits	S	us \	widtl	h for	· pa	rtitic	n pa	air O	/1												
		4:	:3		DO	CAC	;0[1	:0]	00 - 01 - 10 -	8 c 9 c 10	colu colu col	mn a mn a umn	umr addi addi addi add	ress ress dres	bits bits s bit	ts	its fo	or pa	artiti	on p	air (0/1										
		6:	:5		DF	RAC	;0[1	:0]	00 - 01 - 10 -	- 11 - 12 - 13	l rov 2 rov	wad wad wad	ddre Idres Idres	ss b ss b	its its	oun	t for	part	titior	n pai	ir 0/	1										
		7	7			DN	В0		0 –	2 i	inte	rnal	ks i	RAN	1 ba	nks	itior	n pai	r													

1 - 4 internal SDRAM banks



Та	ble (6-2.	MD	CN	FG	Ві	it D	efin	itic	ns	s (S	She	et	2 (of 3)															
			0x4	1800)_0(000							N	۱D	CNF	G							M	emo	ry C	on	troll	er			
Bit	31 3	30 29	28	27	26	25	24	23	22	21	20	19	18	17	7 16	1	5 1	4 13	12	11	10	9	8	7	6	5	4	3	2	1	0
	rese	erved	DSA1111_2	DLATCH2	DADDR2	DI	ΓC2	DNB2	DRAC2		27470	7	DWID2	DE3	DE2	r	ese	erved	DSA1111_0	DLATCH0	DADDRO	Dī	rC0	DNB0	DRACO		DCACO		DWID0	DE1	DE0
Reset	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	() (0 0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bits			Na	me												D	escr	ipti	on										
		9:8		D	TC	0[1:	0]	00 - 01 - 10 - 11 - tWF Use refe equ See	tRF tRF tRF tRF d to rred ivale	P = P = P = P = P = P = P = P = P = P =	2 cll 2 cll 3 cll 3 cll reconfigu in th	ks,	CL = CL = CL = ry til the imin	= 2 = 3 = 3, me SD gs clo	, tRC , tRC , tRC) is f PRAM abov ocks	D D D tixe I tii	= 1 = 2 = 3 = 3 d at min are sed	clks, clks, clks, clks, t 2 clo gs to the n on th	tRAS tRAS tRAS cks. the S umb e MI	S(m S(m S(mi SDR er o	in) in) in) AM	= 5	clks, clks, clks, nufa Ks. xDB	tRC tRC tRC	C = 8 =10 C = 1 er's	clks clks 1 cl	s ks cific			Clod	cks
		10		See Figure 6-5 for a description of these timing numbers. DADDR0 reserved For an explanation on how the alternate addressing works, see Figu Return Data from SDRAM latching scheme for pair 0/1 0 - Latch return data using fixed delay from MEMCLK															igur	e 6	-4										
		11		D)LA	ГСН	10	0 – 1 – This	Lat Lat bit	tch tch mu	retu retu ıst a	irn d irn d Iwa	data data ys b	us wi	ing f th re writte	ixe tur n v	d d n cl vith	elay f	rom to er	MEN nabl	ИС e u	LK sing						CLI	< for		
		12		DS	SA1	111	_0	add	ress	sing	bit	pro	grar	nm	ed ir	١Ň	IDC	ode fo NFG:	DAE	DR	0.							le th	ne		
		15:13			_	_		rese	erve	d																					
		16			DI	E2		For MD refre disa	eac CNF esh bled	h S G:l cyc d.	DR/ DWI le to	AM ID2: o all	part ='1') par tion	itio ac titio	cess	ere s (r Wh	ead	an er I or wi all pa	ite) t	o a	dis	able	d ŠE	DRA	Мра	rtiti	on t	rigg	ers a		
		17			DI	E3		For MD refre disa	eac CNF esh bled	h S G:l cyc d.	DR/ DWI le to	AM ID2: o all	part ='1') par tion	itio ac titio	cess	ere s (r Wh	ead	an er I or wr all pa	ite) t	o a	dis	able	d Š[DRA	Мра	artiti	ion t	rigg	ers a		
								SDI	RAN	1 da	ata b	us	widt	h f	or pa	rtit	ion	pair 2	2/3												

18

DWID2

0 - 32 bits 1 - 16 bits



Table 6-2. MDCNFG Bit Definitions (Sheet 3 of 3)

Ia	DIC	0-2	"						•			ک , د		•		. 0,																
				0x4	800)_0	000							N	IDC	NF	3							Me	emo	ry C	ont	roll	er			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	res	erve	ed	DSA1111_2	DLATCH2	DADDR2	D.	TC2	DNB2	DRAC2		DCAC2	1	DWID2	DE3	DE2	res	erve	ed	DSA1111_0	DLATCH0	DADDR0	DT	C0	DNB0	DRACO		DCACO		DWID0	DE1	DE0
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bit	s			Na	ame												De	scr	iptio	on										
		20:	19		DC	CAC	C2[′	1:0]	00 - 01 - 10 -	mbe - 8 c - 9 c - 10 - 11	colu colu	ımn ımn lumr	add add add	ress ress dres	bits bits sb bits	s s its	ts fo	or pa	rtiti	on p	air 2	2/3										
		22::	21		DF	RAC	C2[′	1:0]	00 - 01 - 10 -	RAM - 11 - 12 - 13 - res	rov rov rov	v ad w ad w ad	dres	ss b	its its	ount	for	parti	itior	n pai	ir 2/:	3										
		23	3			10	NB2		0 –	mbe 2 ii 4 ii	nter	rnal	SDF	RAN	I ba	nks	titior	n pai	r													
		25::	24		D	TC	2[1	:0]	00 - 01 - 10 - tWI The spe ma	- tR - tR - tRI R (w ese becific y no	P = P = rite pits ation	2 cl 2 cl 3 cl 3 cl reco are	ks, ks, ks, ks, ver ver use Cloc lival	CL = CL = CL = Ty tir d to cks i lent	= 2, = 2, = 3, = 3, ne) cor refer to n	tRC tRC tRC tRC is fix infigured nem	D = D = D = ced a re th to ir	2/3 1 cli 2 cli 3 cli at 2 che Sin the clock	ks, ks, cloc DR tim ks b	tRA tRA tks. cks. AM nings	S(m S(m S(m timin s ab	nin) nin) in) : ngs ove n th	= 5 0 =7 0 = 7 0 per are e Mi	the the	tRC tRC tRC SDI num	C = 8 =10 C = 1 RAM	3 clk clks 11 cl 1 ma of S	ks s lks anuf SDC				Ks
		26	6		С)A[DDR	2		erve		lana	tion	on	how	, the	alte	ernat	e a	ddre	essi	na v	vork	s.s	ee F	iaur	e 6-	-4.				
		27	7		D	ıLA	тсн	H2	Ret 0 – 1 – Thi	urn La La s bit	Dat tch tch mu	retu retu retu	m S rn d rn d way	SDR lata lata /s b	AM usir with	lato ng fix reto ritter	hing ked urn d	g sch delag clock th a '	nem y fro	e foom I	r pa MEN	iir 2 MCL	/3 .K sing	the	retu	rn c	lock	SD		(for		
		28	3		DS	SA	1111	_2	ado	lress	sing	j bit į	orog	gran	nme	d in	МD	lode CNF 1111	G:I	DAC	DR:	2.		•					de th	ie		
		31:	29			-			res	erve	d																					



6.5.2 SDRAM Mode Register Set Configuration Register (MDMRS)

The MDMRS, shown in Table 6-3, issues an Mode Register Set (MRS) command to the SDRAM. The value written in this register is placed directly on address lines MA[24:17] during the MRS command. For MA[16:10], values which are fixed or derived from the MDCNFG register are placed on the address bus. When setting the values to be written out on the address lines, base the values on the addressing mode being used. Although writing to this register triggers an MRS command, the corresponding chip-select values are asserted only if the memory banks are enabled via the MDCNFG register. Therefore, to appropriately write a new MRS value to SDRAM, first enable the memory via the MDCNFG register, and then write the MDMRS register. This register is used solely for generating the MRS command.

All values in the MDCNFG register must be programmed correctly to ensure proper operation of the device.

The MDMRS[MDBLx] bits configure the SDRAM to a burst length of four. This value is fixed and cannot be changed. For transfer cycles that require more data than the set burst length of four, the controller can preform as many bursts as necessary to transfer the required amount of data. For example, during a cache line fill the controller can perform a four-beat burst followed immediately by another four-beat burst. This approach requires the controller to generate the first address for the second burst. During transfer cycles less than four beats, the controller ignores the data it does not need. For instance, if the SDRAM is configured as non-cacheable, single beat reads are seen on the bus as a four-beat read with only one beat that is used by the processor. This also applies to a single-beat write.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 6-3. MDMRS Bit Definitions (Sheet 1 of 2)

			0)x4	800	_00)40								MDI	/IRS	;							Me	emo	ry (on	trol	ler			
Bit	31	30 29	2	8	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	reserved			M	DM	RS	2			N	DC	L2	MDADD2		MDBL2		reserved			N	IDN	IRS	0			M	DCI	-0	MDADD0	M	DBI	-0
Reset	0	0 0	(0	0	0	0	0	0	0	1	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	1	0
		31 — reserved 30:23 MDMRS2 MRS value to be written to SDRAM for partition pair 2.																														
·																																
•		22:20			ľ	MD	CL2		SD Re			artiti y.	on p	air 2	2 C <i>F</i>	AS L	ater	су -	- dei	rive	d fro	om N	ИDC	CNF	G:D	TC2	2. W	rites	are	ign	ore	d.
•		19			M	IDA	DD	2	l			artiti ads (air 2	2 Bu	ırst -	Гуре	e. Fiz	x to	sec	uer	itial	add	ress	sing.	. Wr	ites	are	igno	ored		
•		18:16			ľ	MDI	BL2					artiti ads (2 bu	rst l	engt	h. F	ixed	d to	a bi	urst	lenç	gth c	of fo	ur. V	Vrite	es a	re ig	nore	ed.	
•		15				_	_		res	erv	ed																					
•		14:7			М	IDN	1RS	0	MR	S v	alu	e to	be w	ritte	en to	SD	RAI	√l fo	r Pa	rtiti	on F	Pair (0.									



Table 6-3. MDMRS Bit Definitions (Sheet 2 of 2)

				0x4	1800_	00	40							ı	MDI	VIRS	6							Me	emo	ry (Con	troll	er			
Bit	31	30	29	28	27 2	6	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	reserved			N	IDMR	S2	2			M	DCI	_2	MDADD2		MDBL2		reserved			N	ИDN	IRS	0			M	DCI	_0	MDADD0	ME	BL	_0
Reset	0	0	0	0	0 ()	0	0	0	0	1	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	1	0
		6:	4		М	DC	CLO			RAN s fie						AS L	.ateı	псу	- de	rive	d fro	om N	ИDC	NF	G:D	TCC). W	rites	are	igno	ore	d.
		3	3		МЕ	Α	DD			RAN ⁄ays				air () Bu	ırst [·]	Туре	e. Fi	x to	sec	quen	itial	add	ress	ing.	Wr	ites	are	igno	ored.		
		2:	0		М	DE	BL0			RAN ays					0 bu	rst I	eng	th. F	ixe	d to	a bı	urst	lenç	gth c	of fo	ur. V	Vrite	s a	re ig	nore	d.	

6.5.2.1 Low-Power SDRAM Mode Register Set Configuration Register

The Low-Power SDRAM Mode Register Set Configuration register (MDMRSLP) is used to issue a special low-power MRS command to SDRAM. Writing this register will trigger a two-stage MRS command to be issued to external SDRAM. The first stage will write the low-power MRS value to SDRAM partitions 0 and 1; the second stage will write the low-power MRS value to SDRAM partitions 2 and 3. The value written in this register will be placed directly on address lines MA[24:10] during the MRS command. When setting the values to be written out on the address lines, they must be written out properly based on the addressing mode which is being used. Although writing to this register will trigger an MRS command, the corresponding chip select values will be asserted only if the memory banks are enabled via the MDCNFG register and if the corresponding MDMRSLP[MDLPENx] bit is set. To write a new low-power MRS value to SDRAM, first enable the memory via the MDCNFG register, and then write the MDMRSLP register with the enable bits set.

This register is not used with in the processor except to write the value during the MRS command. All values in the MDCNFG register must be programmed correctly to ensure proper operation of the SDRAM. The register is used by a low-power SDRAM to control the Partial Array Self-Refresh (PASR) and Temperature Compensated Self-Refresh (TCSR) settings.



Table 6-4. MDMRSLP Register Bit Definitions

				0 X	480	0 00	58							M	DM	RSL	.P															
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	MDLPEN2								MDMRSLP2								MDLPENO								MDMRSLP0							
Reset	0	Name Description															0	0														
·		ENABLE BIT FOR LOW POWER MRS VALUE FOR PARTITION F																														
·		3	1		MD	LPE	EN2							Di	sabl	le lo	w po	owe	r MF	RS f	or F	Parti	ion	Pai	r 2/3	_	PAF	RTIT	ION	I PA	IR 2	2/3:
·		30:	16		MD	MR	SLF	2					_		_		MR:	-	٩LU	ΕТ	ОВ	ΕW	'RIT	TEN	N TC) SE	DRA	M F	OR			
·		1	5		MD	LPE	EN0						Ena 0 – 1 –	Di	sabl	le lo	w po		r Mi	RS f	or F	Partit	ion	Pai	ion F r 0/1 0/1		0/1.	•				
·		14	:0		MD	MR	SLF	90									MR:		٩LU	ΕТ	ОВ	ΕW	'RIT	TEN	N TC) SE	DRA	ΜF	OR			

6.5.3 SDRAM MDREFR Register (MDREFR)

MDREFR, shown in Table 6-5, is a read/write register and contains control bits that refresh both SDRAM partition pairs. MDREFR also contains control and status bits for SDRAM self-refresh, SDRAM/SMROM clock divisors, SDRAM/SMROM clocks running, and SDRAM/SMROM clock-enable pin states. Independent control/status is provided for each of the clock pins (SDCLK[2:0]) and clock-enable pins (SDCKE[1:0]).

The clock-run bits (K0RUN, K1RUN, and K2RUN) and clock-enable bits (E0PIN and E1PIN) provide software control of SDRAM and Synchronous Static Memory low-power modes. When the clock-run bits and clock-enable bits are cleared, the corresponding memory is inaccessible.

Automatic power-down, enabled by the APD bit, is an mechanism for minimizing power consumption in the processor SDCLK pin drivers and/or the SDRAM/Synchronous Static Memory devices. A latency penalty of one memory cycle is incurred when SDCLK and SDCKE are restarted between non-consecutive SDRAM/Synchronous Static Memory transfers.

The following conditions determine whether SDRAM refreshes occur.

- No refreshes are sent to SDRAM when the refresh counter is cleared to zero.
- If a single transaction to a disabled SDRAM partition is requested, a refresh to all four partitions is performed.
- If all four SDRAM partitions are disabled, the refresh counter is disabled.
- If the clock frequency is changed, the register must be rewritten, even if the value has not changed. This results in a refresh and the refresh counter being reset to the next refresh interval.



This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 6-5. MDREFR Bit Definitions (Sheet 1 of 3)

Ta	ble	6-	5. N	IDI	REF	R Bit	De	efini	tior	ıs (S	She	et 1	of	3)																
				0x4	800_	0004						N	IDR	EFF	2							Me	emo	ory (Con	troll	er			
Bit	31	30	29	28	27 2	6 25	24	23 2	2 2	1 20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		r	ese	rvec	ı	K2FREE	K1FREE	KOFREE	SLFRSH	APD	K2DB2	K2RUN	K1DB2	K1RUN	E1PIN	K0DB2	KORUN	EOPIN						D	RI					
Reset	0	0	0	0	0 0	1	1	1	1 (0	1	0	1	0	0	1	*	*	1	1	1	1	1	1	1	1	1	1	1	1
		Bi	its		N	lame											De	escr	iptio	n										
		31:	:26			_		resei	ved																					
		2	5		K2	FREE	.	SDR 0 – 1 – Prov circu	SDC SDC ides	LK2 LK2	is no	ot fre	e-ru unni	ınnii ng (i	ng igno														nal	
		2	4		K1	FREE	•	SDR 0 – 1 – Prov	SDC	LK1 LK1	is no	ot fre	e-ru unnii	ınniı ng (i	ng igno				-	-				-		-			itry.	
		2	3		K0	FREE	Ē	SDR 0 – 1 – Prov	SDC	LK0 LK0	is no	ot fre	e-ru unnii	ınnii ng (i	ng igno				-	_				-		-		,	itry.	
		2	2		SL	FRSH	ł	SDR Cont hard 0 – 1 – SLFF be cl SDR SLFF prohi	rol/s ware Self Self RSH eare AM SSH	tatus refre refre can d. Th nas e). Thi autor	bit for eep sh dish e so e so e ap nteros cap matic	or energy reservations in the control of the contro	nteri et. led ed / sof oriat elf-r lity r	twa e cle efre nusi	re to ock sh a t be for a	o for run and use any	ce a bits mus ed w com	a sel (K1 st be ith e	f-ref RUI set extre	resh N an prio me o	co d/o r to cau	mm r K2 exi tion	and RU ting be	d. E1 IN) n self	PIN nust refi se th	I do t ren resh ne re	es n nain (cle sult	ot h set earir	ave unt	e to iil
		2	1			_		resei	ved																					
		2	0		,	APD		SDR If AP corre If no mode If one not b If no these Section	D=1 sportsports SDF e and e SC eing Syn e par	, the nding RAM d the RAM used chror	clock partic clock par d are	k en titions ks a tition turi	ablens and of the control of the con	s ar re be beiock clock cein off.	nd cleing ing a cen g us	lock acce able sed a	pinessesse pines	s are sed u ed, th ns ar the	e au unles ne S e tu othe	tomass the DRA rned er is eing	etica e K AM l off not	ally (xFF chip f. t, the	dea REE os a e cle	asse E bits Ire p ock t	are ut in to th	e se nto F ne pa d cla	t. Powe artiti ock	er-D on t ena	owr hat ble	is to
		1	9		K	2DB2		SDR 0 – 1 –	SDC	LK2	is sa	me	freq	uen	су а	s M	IEM	CLK				S								



Ta	ble	6	-5.	MD	RE	FR	Bit	t De	efin	itic	ons	(S	he	et 2	2 0	3)																
				0x	480	0_0	004							N	/IDR	EFF	2							Me	emo	ry (Con	tro	ller			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			res	erve	d		K2FREE	K1FREE	KOFREE	SLFRSH	reserved	APD	K2DB2	K2RUN	K1DB2	K1RUN	E1PIN	K0DB2	KORUN	EOPIN						D	RI					
Reset	0	0	0	0	0	0	1	1	1	1	0	0	1	0	1	0	0	1	*	*	1	1	1	1	1	1	1	1	1	1	1	1
		E	Bits			Na	ame												D	esc	riptio	on										
			18			K2	RUN	l	0 – 1 – K2F prol	SI SI RUN nibit	ts au K1F	K2 o K2 o o ca utom	disal enak an b natio	bled bled e cl tra	l eare nsiti	ed by ons	y pro	ogra any	ım. con	Use nma	wit ands	h ca	utio	n be						•		
			17			K1	DB2		0 –	SI	M CI DCL DCL	K1 i	s sa	me	frec	luen	cy a	as M	IEM	CLŁ	<			IS								
			16			K1	RUN		0 – 1 – K1F auto	SI SI RUN oma	atic t K1F	K1 o K1 e n be rans	disal enal e cle sitio	bled oled ared	l d by or a	soft	war	e. L	Jse ds.	with	cau	utior	ı be	caus and				·		·		bits
			15			E1	PIN		0 – 1 – E1F K2F auto be s	SI SI PIN RUN oma set I SL	l=1, atic t by p FRS	E1 i be and rand rand rogr	is di is er clea I SL sition am 0).	sab nabl red FRS ns fo to c	led ed by p SH= or m aus	orog 0). l ode e a	ram Jse reg pow	to o with iste	caus r ca r se low	se a ution t, re	pov n be ad, tit co	wer- ecau writ omn	dow ise t e, a nand	n co he r nd re d (if I	esul efres K1R	ting sh c UN	sta comr =1 a	te p nar ind	orohi nds. /or K	ibits E1F (2R	NI (can
			14			K0	DB2		0 – 1 –	SI SI	ono DCL DCL	K0 r K0 r	uns	at t	he r	nem half	ory the	cloo me	ck fr mor	equ y cl	enc ock	y. freq	uen	•	by	2 C	ontr	ol/S	Statu	S		

This bit is automatically set upon hardware or sleep reset.



Table 6-5. MDREFR Bit Definitions (Sheet 3 of 3)

	D .C	0-5	• ••			• • •						,,				٠٠,																
				0x4	1800	0_00	04							N	IDF	REF	?							Me	emo	ry (Con	troll	er			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		re	se	rve	d		K2FREE	K1FREE	KOFREE	SLFRSH	reserved	APD	K2DB2	K2RUN	K1DB2	K1RUN	E1PIN	K0DB2	KORUN	EOPIN						D	RI					
Reset	0	0	0	0	0	0	1	1	1	1	0	0	1	0	1	0	0	1	*	*	1	1	1	1	1	1	1	1	1	1	1	1
		Bit	S			Na	me												De	esci	iptic	on										
		13				KOR	RUN		0 – 1 – Set synd K0R	SD SD on e chro	CLI CLI exit nou car	K0 i K0 i fron is m	s dis s en n ha nemo	able able rdw ory t	ed ed are type	and e and e. / the	sle	ep r	eset	t if ti	he E	300 :apa	T_S bility	SEL y mı	sigr ust b	nals be u	are sed	conf	ca			a
		12				EOF	PIN		0 - 1 - This conf E0F Use Reg	SD Sbit figur PIN o with	CK is s ed can h ca	E0 i E0 i et o for a be o autio	s dis s er on ex a sy clea on bo omm	sabl lable dit fr nchi red ecau	ed om ron by use I ar	hardous the period the indicate	dwai men orog resu ad c	re an nory ram ilting	nd s typ to c g sta man	sleep e. caus ate p ate,	o res	set i	f the	e BC	OOT on co	_SE omm tran	EL si nanc	ignal I (if I	ls a K0R or N	re UN: 1ode	,	
		112	0			DI	RI		The row appropries of the DRI (Rei Must prog Must Wheel If all only If the chair	nur is re is	mbe effre ble t nme ue t Num h tir e pro e les lear r SI forr pock	r of she to all the the to all th	mer d in II SE umb mus r of I row ann 1 to 0, AM p I who quen resu	mory eac PRA er, S et be men rs) x ed t no parti- en a cy i	/ cle h S M ii 31 i 31 i e loa ncll Me co b RA refi tior a sii s cl	hall pa oock of DRA in the is added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added added	eycle AM to founded linto les- y cle arec nax) aes a e distran ed, esh	es di pank r pa to t to thi 31) ock I by for re s sable sact this	c du nrtition his r s re- y 32 freq both the ed, 1 tion reg the	ring ons. num giste = juen h pa SDF to th the to a iste refi	each To be read to the series of the series	ch C calconal calcona	BR ulat cula airs airs coursed S erecurrent	refree the second ted. The access. Interpolation of the period of the control of	esh e re ultip as f ne sr is d AM tten, rese	cyclifres fres fres follor mall ed. lisat part , eve	le. The inh	terva 32.	refi equ	rval om 1 mus resh este	is his est be	e are

6.5.4 Fixed-Delay or Return-Clock Data Latching

The Return-clock data latching works in all situations. Fixed-delay latching is not guaranteed to work in all situations and must not be used. Fixed-delay latching does not work all the time because the delays through the output pads can vary widely over a range of conditions. Because this delay can be greater than a clock cycle, it is impossible to determine the clock to latch the data from the SDRAM. Program the MDCNFG:DLATCHx and SXCNDF:SXLATCHx fields to a 1 to enable latching using the return clock SDCLK.



6.5.5 SDRAM Memory Options

The Dynamic Memory interface supports up to four partitions, organized as two pairs. Both partitions in a pair must have the same SDRAM size, configuration, timing category, and data bus width. Initialization software must set up the memory interface configuration register with:

- SDRAM timing category
- · Data-bus width
- Number of row, column, and bank address bits
- Addressing scheme
- · Data-latching scheme

Table 6-6 shows a sample of the supported SDRAM configurations.

Table 6-6. Sample SDRAM Memory Size Options

SDRAM Configuration	Oldin Olma		r Chips/ ition	Bank Bits x		on Size Partition)
(Words x Bits)	Chip Size	16-Bit Bus	32-Bit bus	Row Bits x Column Bits	16-Bit Bus	32-Bit Bus
1M x 16	16 Mbit	1	2	1 x 11 x 8	2 Mbyte	4 Mbyte
2 M x 8	16 Mbit	2	4	1 x 11 x 9	4 Mbyte	8 Mbyte
2 M x 32	64 Mbit	N/A	1	2 x 11 x 8	N/A	8 Mbyte
4 M x 16	64 Mbit	1	2	1 x 13 x 8 2 x 12 x 8	8 Mbyte	16 Mbyte
8 M x 8	64 Mbit	2	4	1 x 13 x 9 2 x 12 x 9	16 Mbyte	32 Mbyte
8 M x 16	128 Mbit	1	2	2 x 12 x 9	16 Mbyte	32 Mbyte
16 M x 8	128 Mbit	2	4	2 x 12 x 10	32 Mbyte	64 Mbyte
16 M x 16	256 Mbit	1	2	2 x 13 x 9	32 Mbyte	64 Mbyte
32 M x 8	256 Mbit	2	4	2 x 13 x 10	64 Mbyte	128 Mbyte - exceeds partition size

Figure 6-4 shows the bank/row/column address multiplexing using a 2x13x9 32-bit SDRAM as an example for the normal bank-addressing scheme. All unused address bits during RAS and CAS time - including MA[9:0] bits not shown here - are not guaranteed, and will be either driven to zero or one.

6.5.5.1 SDRAM Addressing Modes

The processor supports two addressing modes: Normal Bank Address mode and SA-1111 Address mode. The addressing mode alters the order of the address bits that are driven on the individual memory address pins and control the SDRAM components.

Refer to Table 6-7 through Table 6-9 for a listing of address mapping options.



Table 6-4 shows how the SDRAM row and column addresses are mapped to the internal SDRAM address. The SDRAM row and column addresses are muxed. The SDRAM row is sent during an Active command and is followed by the column address during the read or write command. MA<20> is driven with 0 during column addressing. BA[1:0] is used to tell the SDRAM which bank is being read from and remains stable during column addressing. During SDRAM configuration, all the address pins are used to transfer the MRS command.

Figure 6-4. External to Internal Address Mapping Options

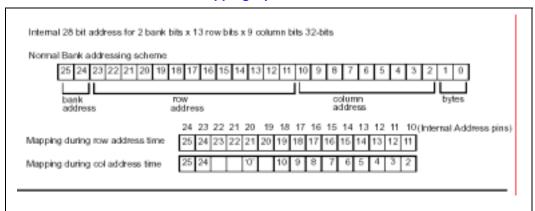


Table 6-7. External to Internal Address Mapping for Normal Bank Addressing (Sheet 1 of 3)

# Bits Bank x		Ex	teri	nal <i>i</i>	Add		s pi MA				ΑM	RAS	S Tii	me			Ex	teri	nal <i>I</i>	Add		s pi				AM (13 12 11 10 5 4 3 2 4 3 2 1 5 4 3 2 4 3 2 1 5 4 3 2 4 3 2 1 5 4 3 2 1 3 2 1 bination)											
Row x Col x Data	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10								
1x11x8x32				21	20	19	18	17	16	15	14	13	12	11	10				21	'0'			9	8	7	6	5	4	3	2								
1x11x8x16				20	19	18	17	16	15	14	13	12	11	10	9				20	'0'			8	7	6	5	4	3	2	1								
1x11x9x32				22	21	20	19	18	17	16	15	14	13	12	11				22	'0'		10	9	8	7	6	5	4	3	2								
1x11x9x16				21	20	19	18	17	16	15	14	13	12	11	10				21	'0'		9	8	7	6	5	4	3	2	1								
1x11x10x32				23	22	21	20	19	18	17	16	15	14	13	12				23	'0'	11	10	9	8	7	6	5	4	3	2								
1x11x10x16				22	21	20	19	18	17	16	15	14	13	12	11				22	'0'	10	9	8	7	6	5	4	3	2	1								
1x11x11x32		NC)T V	ALI	D (i	lleg	al a	ddr	ess	ing	cor	nbir	natio	on)			NC)T V	ALI	D (i	lleg	al a	ddr	ess	ing	6 5 4 3 2 7 6 5 4 3 2 7 6 5 4 3 2 7 6 5 4 3 2 7 6 5 4 3 2 ng combination) 7 6 5 4 3 6 5 4 3 2 7 6 5 4 3 7 6 5 4 3												
1x11x11x16		NC)T V	'ALI	D (i	lleg	al a	ddr	ess	ing	cor	nbir	natio	on)			NC)T V	'ALI	D (i	lleg	al a	ddr	ess	ing	con	nbir	natio										
1x12x8x32			22	21	20	19	18	17	16	15	14	13	12	11	10			22		'0'			9	8	7	6	5	4	3	2								
1x12x8x16			21	20	19	18	17	16	15	14	13	12	11	10	9			21		'0'			8	7	6	5	4	3	2	1								
1x12x9x32			23	22	21	20	19	18	17	16	15	14	13	12	11			23		'0'		10	9	8	7	6	5	4	3	2								
1x12x9x16			22	21	20	19	18	17	16	15	14	13	12	11	10			22		'0'		9	8	7	6	5	4	3	2	1								
1x12x10x32			24	23	22	21	20	19	18	17	16	15	14	13	12			24		'0'	11	10	9	8	7	6	5	4	3	2								



Table 6-7. External to Internal Address Mapping for Normal Bank Addressing (Sheet 2 of 3)

# Bits		Ex	teri	nal <i>i</i>	Add		s pi MA•				AM I	RAS	Tir	ne			Ex	teri	nal <i>i</i>	Add		s pi				AM (CAS	S Tiı	ne	
Bank x Row x Col x Data	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10
1x12x10x16			23	22	21	20	19	18	17	16	15	14	13	12	11			23		'0'	10	9	8	7	6	5	4	3	2	1
1x12x11x32			25	24	23	22	21	20	19	18	17	16	15	14	13			25	12	'0'	11	10	9	8	7	6	5	4	3	2
1x12x11x16			24	23	22	21	20	19	18	17	16	15	14	13	12			24	11	'0'	10	9	8	7	6	5	4	3	2	1
1x13x8x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23			'0'			9	8	7	6	5	4	3	2
1x13x8x16		22	21	20	19	18	17	16	15	14	13	12	11	10	9		22			'0'			8	7	6	5	4	3	2	1
1x13x9x32		24	23	22	21	20	19	18	17	16	15	14	13	12	11		24			'0'		10	9	8	7	6	5	4	3	2
1x13x9x16		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23			'0'		9	8	7	6	5	4	3	2	1
1x13x10x32		25	24	23	22	21	20	19	18	17	16	15	14	13	12		25			'0'	11	10	9	8	7	6	5	4	3	2
1x13x10x16		24	23	22	21	20	19	18	17	16	15	14	13	12	11		24			'0'	10	9	8	7	6	5	4	3	2	1
1x13x11x32		13	26	25	24	23	22	21	20	19	18	17	16	15	14		13		12	'0'	11	10	9	8	7	6	5	4	3	2
1x13x11x16		25	24	23	22	21	20	19	18	17	16	15	14	13	12		25		11	'0'	10	9	8	7	6	5	4	3	2	1
2x11x8x32			22	21	20	19	18	17	16	15	14	13	12	11	10			22	21	'0'			9	8	7	6	5	4	3	2
2x11x8x16			21	20	19	18	17	16	15	14	13	12	11	10	9			21	20	'0'			8	7	6	5	4	3	2	1
2x11x9x32			23	22	21	20	19	18	17	16	15	14	13	12	11			23	22	'0'		10	9	8	7	6	5	4	3	2
2x11x9x16			22	21	20	19	18	17	16	15	14	13	12	11	10			22	21	'0'		9	8	7	6	5	4	3	2	1
2x11x10x32			24	23	22	21	20	19	18	17	16	15	14	13	12			24	23	'0'	11	10	9	8	7	6	5	4	3	2
2x11x10x16			23	22	21	20	19	18	17	16	15	14	13	12	11			23	22	'0'	10	9	8	7	6	5	4	3	2	1
2x11x11x32		NC	T V	'ALI	D (i	lleg	al a	ddr	ess	ing	con	nbir	atio	on)			NC)T V	'ALI	D (i	lleg	al a	ddr	ess	ing	con	nbir	natio	on)	
2x11x11x16		NC	T V	'ALI	D (i	lleg	al a	ddr	ess	ing	con	nbir	atio	on)			NC)T V	'ALI	D (i	lleg	al a	ddr	ess	ing	con	nbir	natio	on)	
2x12x8x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'			9	8	7	6	5	4	3	2
2x12x8x16		22	21	20	19	18	17	16	15	14	13	12	11	10	9		22	21		'0'			8	7	6	5	4	3	2	1
2x12x9x32		24	23	22	21	20	19	18	17	16	15	14	13	12	11		24	23		'0'		10	9	8	7	6	5	4	3	2
2x12x9x16		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'		9	8	7	6	5	4	3	2	1
2x12x10x32		25	24	23	22	21	20	19	18	17	16	15	14	13	12		25	24		'0'	11	10	9	8	7	6	5	4	3	2
2x12x10x16		24	23	22	21	20	19	18	17	16	15	14	13	12	11		24	23		'0'	10	9	8	7	6	5	4	3	2	1
2x12x11x32		26	25	24	23	22	21	20	19	18	17	16	15	14	13		26	25	12	'0'	11	10	9	8	7	6	5	4	3	2
2x12x11x16		25	24	23	22	21	20	19	18	17	16	15	14	13	12		25	24	11	'0'	10	9	8	7	6	5	4	3	2	1



Table 6-7. External to Internal Address Mapping for Normal Bank Addressing (Sheet 3 of 3)

# Bits Bank x		Ex	teri	nal A	Add				at S :10>		AM I	RAS	S Tiı	ne			Ex	teri	nal <i>i</i>	Add		s pi MA•				AM (CAS	S Tiı	ne	
Row x	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10
2x13x8x32	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23			'0'			9	8	7	6	5	4	3	2
2x13x8x16	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	23	22			'0'			8	7	6	5	4	3	2	1
2x13x9x32	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	25	24			'0'		10	9	8	7	6	5	4	3	2
2x13x9x16	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23			'0'		9	8	7	6	5	4	3	2	1
2x13x10x32		•		•	N	ЮТ	VAI	LID	(toc	big	g)	•	•							N	ОТ	VAI	LID	(toc	big	g)				
2x13x10x16	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	25	24			'0'	10	9	8	7	6	5	4	3	2	1
2x13x11x32					N	ЮТ	VAI	LID	(toc	big	g)									N	ОТ	VAI	LID	(toc	big	g)				
2x13x11x16					N	ЮТ	VAI	LID	(toc	big	g)									N	ОТ	VAI	LID	(toc	big	g)				

Table 6-8. External to Internal Address Mapping for SA-1111 Addressing (Sheet 1 of 3)

# Bits		Ext	tern	al A	Add		s pii MA<				AM	RA	S Ti	me			Ex	tern	al A	Add		•		at S :10:		AM	CA	S Ti	ime	!
Row x Col x Data	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10
1x11x8x32				21	20	19	18	17	16	15	14	13	12	11	10				21	'0 '			9	8	7	6	5	4	3	2
1x11x8x16		NO	OT VALID (illegal addressing combination) 21 20 19 18 17 16 15 14 13 12 11 21 20 19 18 17 16 15 14 13 12 11														NO	T V	ALI	D (i	lleg	al a	ddr	ess	ing	COI	mbi	nati	ion))
1x11x9x32				21	20	19	18	17	16	15	14	13	12	11	10				21	'0'		22	9	8	7	6	5	4	3	2
1x11x9x16				21	20	19	18	17	16	15	14	13	12	11	10				21	'0'		9	8	7	6	5	4	3	2	1
1x11x10x32				21	20	19	18	17	16	15	14	13	12	11	10				21	'0'	23	22	9	8	7	6	5	4	3	2
1x11x10x16				21	20	19	18	17	16	15	14	13	12	11	10				21	'0'	22	9	8	7	6	5	4	3	2	1
1x11x11x32		NO	T V	ALII	D (il	lleg	al a	ddr	ess	ing	COI	mbi	nat	ion))		NO	T V	ALI	D (i	lleg	al a	ddr	ess	ing	COI	mbi	nati	ion))
1x11x11x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nat	ion))		NO	T V	ALI	D (i	lleg	al a	ddr	ess	ing	COI	mbi	nati	ion))
1x12x8x32			22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0'			9	8	7	6	5	4	3	2
1x12x8x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)			NO	T V	ALI	D (i	leg	al a	ddr	ess	ing	COI	mbi	nati	ion))
1x12x9x32			22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0'		23	9	8	7	6	5	4	3	2
1x12x9x16			22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0 '		9	8	7	6	5	4	3	2	1
1x12x10x32			22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0'	24	23	9	8	7	6	5	4	3	2



Table 6-8. External to Internal Address Mapping for SA-1111 Addressing (Sheet 2 of 3)

# Bits		Ext	tern	al A	Add		s pi				ΑM	RAS	S Ti	me			Ext	tern	al A	\dd		s pi				ΑM	CA	S Ti	me	
Bank x Row x Col x Data	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10
1x12x10x16			22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0'	23	9	8	7	6	5	4	3	2	1
1x12x11x32		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
1x12x11x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
1x13x8x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10				21	°°			9	8	7	6	5	4	3	2
1x13x8x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
1x13x9x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0'		24	9	8	7	6	5	4	3	2
1x13x9x16		23	22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0'		9	8	7	6	5	4	3	2	1
1x13x10x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0'	25	24	9	8	7	6	5	4	3	2
1x13x10x16		23	22	21	20	19	18	17	16	15	14	13	12	11	10				21	'0'	24	9	8	7	6	5	4	3	2	1
1x13x11x32		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
1x13x11x16		23 22 21 20 19 18 17 16 15 14 13 12 11 1 2 3 22 21 20 19 18 17 16 15 14 13 12 11 1 2 3 22 21 20 19 18 17 16 15 14 13 12 11 1 1 2 3 22 21 20 19 18 17 16 15 14 13 12 11 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1															NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
2x11x8x32			22	21	20	19	18	17	16	15	14	13	12	11	10			22	21	'0'			9	8	7	6	5	4	3	2
2x11x8x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
2x11x9x32			22	21	20	19	18	17	16	15	14	13	12	11	10			22	21	'0'		23	9	8	7	6	5	4	3	2
2x11x9x16			22	21	20	19	18	17	16	15	14	13	12	11	10			22	21	'0'		9	8	7	6	5	4	3	2	1
2x11x10x32			22	21	20	19	18	17	16	15	14	13	12	11	10			22	21	'0'	24	23	9	8	7	6	5	4	3	2
2x11x10x16			22	21	20	19	18	17	16	15	14	13	12	11	10			22	21	'0'	23	9	8	7	6	5	4	3	2	1
2x11x11x32		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
2x11x11x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
2x12x8x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'			9	8	7	6	5	4	3	2
2x12x8x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
2x12x9x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'		24	9	8	7	6	5	4	3	2
2x12x9x16		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0 '		9	8	7	6	5	4	3	2	1
2x12x10x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'	25	24	9	8	7	6	5	4	3	2
2x12x10x16		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'	24	9	8	7	6	5	4	3	2	1
2x12x11x32		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)			NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	
2x12x11x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)	1		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	nbi	nati	on)	



Table 6-8. External to Internal Address Mapping for SA-1111 Addressing (Sheet 3 of 3)

# Bits		Ext	ern	al A	Add		s pii MA<				ΑM	RA	S Ti	me			Ex	tern	al A	Add		s pi MA				AM	CA	S Ti	me	
Bank x Row x Col x Data	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10
2x13x8x32	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'			9	8	7	6	5	4	3	2
2x13x8x16		NO	T V	ALII	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)			NO	T V	ALI	D (il	leg	al a	ddr	ess	ing	COI	mbi	nati	on)	1
2x13x9x32	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'		25	9	8	7	6	5	4	3	2
2x13x9x16	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'		9	8	7	6	5	4	3	2	1
2x13x10x32					N	от	VAL	.ID	(too	bi	g)	•	•							N	ОТ	VAI	LID	(to	bi	g)				
2x13x10x16	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		'0'	25	9	8	7	6	5	4	3	2	1
2x13x11x32					N	от	VAL	.ID	(too	bi	g)									N	ОТ	VAI	LID	(to	bi	g)				
2x13x11x16					N	от	VAL	.ID	(too	bi	g)									N	от	VAI	LID	(to	o bi	g)				

Use the information below to connect the processor to the SDRAM devices. Some of the addressing combinations may not apply in SA1111 addressing mode. See Table 6-10 for a complete listing of supported addressing combinations and how to connect the PXA255 processor to the SA1111.

Table 6-9. Pin Mapping to SDRAM Devices with Normal Bank Addressing (Sheet 1 of 3)

# Bits Bank x			MA			ng to SI sent the						_	sor.		
Row x Col x Data	MA24	MA23	MA22	MA21	MA20	MA19	MA18	MA17	MA16	MA15	MA14	MA13	MA12	MA11	MA10
1x11x8x32				BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	A0
1x11x8x16				BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	A0
1x11x9x32				BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x11x9x16				BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x11x10x32				BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	A0
1x11x10x16				BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x11x11x32					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
1x11x11x16					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
1x12x8x32			BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x12x8x16			BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	A0
1x12x9x32			BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	A0



Table 6-9. Pin Mapping to SDRAM Devices with Normal Bank Addressing (Sheet 2 of 3)

# Bits			MA			_			s for No			_	sor.		
Bank x Row x Col x Data	MA24	MA23	MA22	MA21	MA20	MA19	MA18	MA17	MA16	MA15	MA14	MA13	MA12	MA11	MA10
1x12x9x16			BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	Α0
1x12x10x32			BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x12x10x16			BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x12x11x32			BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	Α0
1x12x11x16			BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	Α0
1x13x8x32		BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	A0
1x13x8x16		BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x13x9x32		BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x13x9x16		BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x13x10x32		BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x13x10x16		BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x13x11x32		BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x13x11x16		BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
2x11x8x32			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
2x11x8x16			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
2x11x9x32			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	Α0
2x11x9x16			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
2x11x10x32			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
2x11x10x16			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	A0
2x11x11x32					NOT	VALID	(illegal	addres	ssing c	ombin	ation)				
2x11x11x16					NOT	VALID	(illegal	addres	ssing c	ombin	ation)				
2x12x8x32		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	A0
2x12x8x16		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	A0
2x12x9x32		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	A0
2x12x9x16		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
2x12x10x32		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	A0
2x12x10x16		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0



Table 6-9. Pin Mapping to SDRAM Devices with Normal Bank Addressing (Sheet 3 of 3)

# Bits			MA			•				ormal A		_	sor.		
Bank x Row x Col x Data	MA24	MA23	MA22	MA21	MA20	MA19	MA18	MA17	MA16	MA15	MA14	MA13	MA12	MA11	MA10
2x12x11x32		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	A0
2x12x11x16		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	A0
2x13x8x32	BA1	BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
2x13x8x16	BA1	BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	A0
2x13x9x32	BA1	BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	A0
2x13x9x16	BA1	BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	A0
2x13x10x32		•		•			NOT V	ALID (t	oo big)	•			•	
2x13x10x16	BA1	BA0	A12	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	A0
2x13x11x32		ı		ı			NOT V	ALID (t	oo big)	ı			ı	
2x13x11x16							NOT V	ALID (t	oo big)					

Table 6-10. Pin Mapping to SDRAM Devices with SA1111 Addressing (Sheet 1 of 3)

# Bits		N			ping to resent t						_	•		r.	
Bank x Row x Col x Data	MA24	MA23	MA22	MA21	MA20	MA19	MA18	MA17	MA16	MA15	MA14	MA13	MA12	MA11	MA10
1x11x8x32				BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x11x8x16		I.	l	l	NOT	VALID	(illegal	addres	ssing c	ombin	ation)	I.	l		
1x11x9x32				BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x11x9x16				BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	A0
1x11x10x32				BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	Α0
1x11x10x16				BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	Α0
1x11x11x32		I.	l		NOT	VALID	(illegal	addres	ssing o	ombin	ation)	I.	l		
1x11x11x16					NOT	VALID	(illegal	addre	ssing c	ombin	ation)				
1x12x8x32			A11	BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x12x8x16		ı	ı	ı	NOT	VALID	(illegal	addres	ssing c	ombin	ation)	ı	ı	ı	
1x12x9x32			A11	BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	A0



Table 6-10. Pin Mapping to SDRAM Devices with SA1111 Addressing (Sheet 2 of 3)

# Bits		N	Pi 1A[24:1		oing to esent t						_			r.	
Bank x Row x Col x Data	MA24	MA23	MA22	MA21	MA20	MA19	MA18	MA17	MA16	MA15	MA14	MA13	MA12	MA11	MA10
1x12x9x16			A11	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	A0
1x12x10x32			A11	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	A0
1x12x10x16			A11	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
1x12x11x32			•		NOT	VALID	(illegal	addres	ssing c	ombina	ation)		•	•	
1x12x11x16					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
1x13x8x32		A12	A11	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	A0
1x13x8x16					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
1x13x9x32		A12	A11	BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x13x9x16		A12	A11	BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x13x10x32		A12	A11	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	Α0
1x13x10x16		A12	A11	BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
1x13x11x32					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
1x13x11x16					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
2x11x8x32			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
2x11x8x16					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
2x11x9x32			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	A0
2x11x9x16			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
2x11x10x32			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	Α0
2x11x10x16			BA1	BA0	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	A0
2x11x11x32					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
2x11x11x16					NOT	VALID	(illegal	addres	ssing c	ombina	ation)				
2x12x8x32		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	A0
2x12x8x16					NOT	VALID	(illegal	addres	ssing c	ombin	ation)				
2x12x9x32		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A 1	A0
2x12x9x16		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A1	Α0
2x12x10x32		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	A3	A2	A1	A0
2x12x10x16		BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	Α0



Table 6-10. Pin Mapping to SDRAM Devices with SA1111 Addressing (Sheet 3 of 3)

# Bits		Pin mapping to SDRAM devices for SA1111 Addressing Options. MA[24:10] represent the address signals driven from the PXA255 processor.														
Bank x Row x Col x Data	MA24	MA23	MA22	MA21	MA20	MA19	MA18	MA17	MA16	MA15	MA14	MA13	MA12	MA11	MA10	
2x12x11x32	NOT VALID (illegal addressing combination)															
2x12x11x16	NOT VALID (illegal addressing combination)															
2x13x8x32	A12 BA1 BA0 A11 A10 A9 A8 A7 A6 A5 A						A4	А3	A2	A 1	A0					
2x13x8x16	NOT VALID (illegal addressing combination)															
2x13x9x32	A12	BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	A0	
2x13x9x16	A12	BA1	BA0	A11	A10	A9	A8	A7 A6		A5 A4		А3	A2	A 1	A0	
2x13x10x32							NOT V	ALID (t	oo big)						
2x13x10x16	A12	BA1	BA0	A11	A10	A9	A8	A7	A6	A5	A4	А3	A2	A 1	A0	
2x13x11x32		•					NOT V	ALID (t	oo big)		•	•	•		
2x13x11x16							NOT V	ALID (t	oo big)						

6.5.6 SDRAM Command Overview

The processor accesses SDRAM with the following subset of standard interface commands:

- Mode Register Set (MRS)
- Bank Activate (ACT)
- Read (READ)
- Write (WRITE)
- Pre-charge All Banks (PALL)
- Pre-charge One Bank (PRE)
- Auto-Refresh (CBR)
- Power-Down (PWRDN)
- Enter Self-Refresh (SLFRSH)
- Exit Power-Down (PWRDNX)
- No Operation (NOP)

Table 6-11 shows the SDRAM interface commands. The table assumes the bank bits for the SDRAM are sent out on external address lines MA<24:23>.



Table 6-11. SDRAM Command Encoding

		Pins														
Command		SDCKE	SDCKE	nSDCS	-CDDAC	nSDCAS	nWE	DQM		MA <24:10>						
		(at clk n-1)	(at clk n)	3:0	nSDRAS	HSDCAS	HWE	3:0	24:23	22:21	20	19:10				
PWF	RDN	1	0	1	1	1	1	1	х							
PWR	DNX	0	1	1	1	1	1	1	х							
SLF	RSH	1	0	0	0	0	1	0	х							
CE	BR	1	1	0	0	0	1	х	x							
MF	RS	1	х	0	0	0	0	0	OP code							
AC	СТ	1	х	0	0	1	1	х	bank row							
RE	AD	1	х	0	1	0	1	0	bank	col	0	col				
WR	ITE	1	х	0	1	0	0	mask	bank	col	0	col				
PALL	All	1	x	0	0	1	0		х	x	1	v				
PRE	Bank	1	X	U		1		х	bank	X	0 x					
NC	NOD		Х	1	х	х	х	х	х							
NOP		1	^	0	1	1	1									

The programmable opcode for address bits MA<24:17> used during the mode-register set command (MRS) is exactly what is programmed in the MDMRS register.

Table 6-12. SDRAM Mode Register Opcode Table

Address Bits	Option	Value
MA<24:17>	reserved	MDMRSx
MA[16:14]	CAS Latency = 2	010
WA[10.14]	CAS Latency = 3	011
MA[13]	Sequential Burst	0
MA[12:10]	Burst Length = 4	010

6.5.7 SDRAM Waveforms

6-28

Normal operation of the SDRAM controller is shown in Figure 6-5 through Figure 6-11.



Figure 6-5. Basic SDRAM Timing Parameters

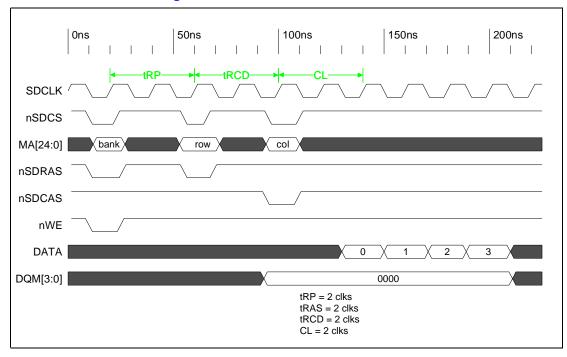


Figure 6-6. SDRAM_Read_diffbank_diffrow

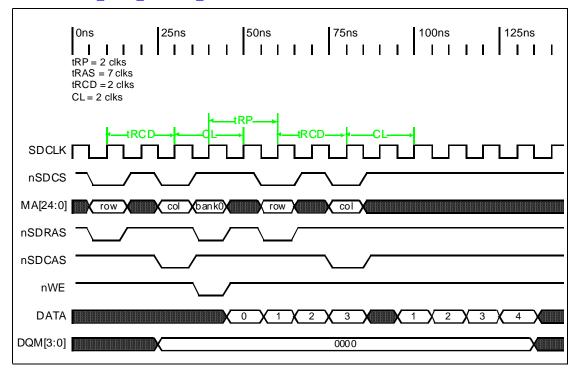




Figure 6-7. SDRAM_read_samebank_diffrow

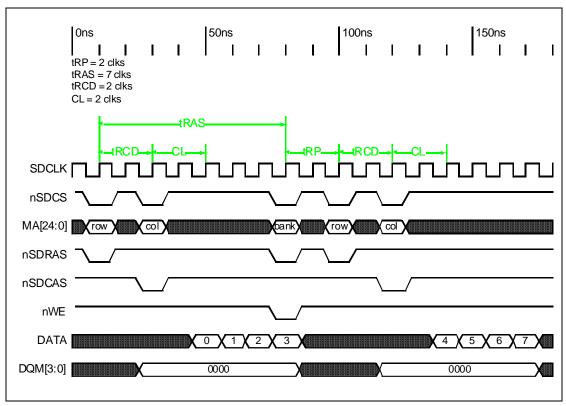


Figure 6-8. SDRAM_read_samebank_samerow

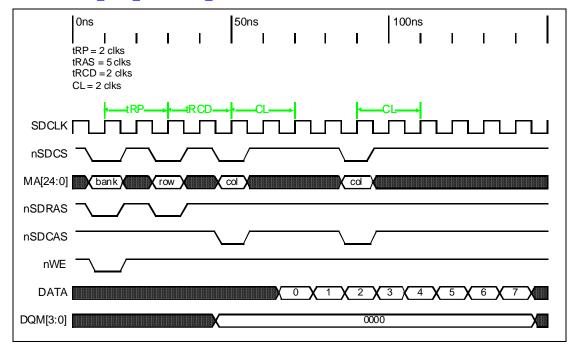




Figure 6-9. SDRAM_write

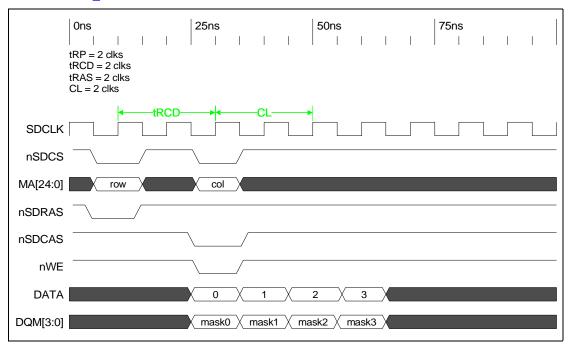
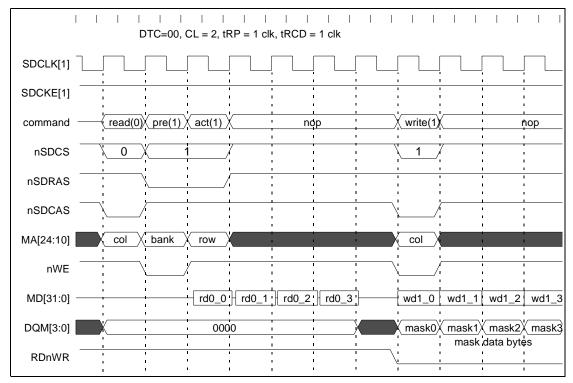


Figure 6-10. SDRAM 4-Beat Read/ 4-Beat Write To Different Partitions





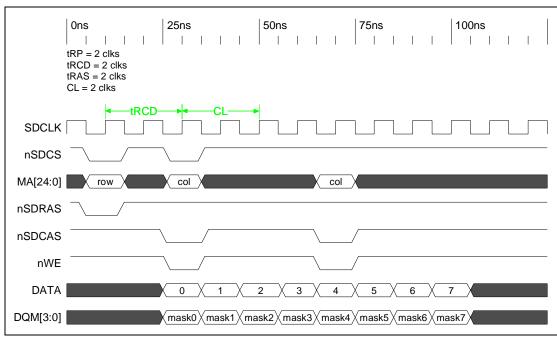


Figure 6-11. SDRAM 4-Beat Write / 4-Write Same Bank, Same Row

6.6 Synchronous Static Memory Interface

The synchronous static memory interface supports SMROM and non-SDRAM-like Flash memories. The synchronous static memory can be configured for any of the nCS[3:0] signals. Chip Select 0 must be used for boot memory. Synchronous static memories in bank pairs 1/0 or 3/2 must be set to the same timing.

If any of the nCS[3:0] banks are configured for Synchronous Static Memory via SXCNFG[SXENx], the corresponding half-words of MSC0 (see Section 6.7.2) and MSC1, except the data width in MSCx[RBWx], are ignored.

6.6.1 Synchronous Static Memory Configuration Register (SXCNFG)

SXCNFG controls all synchronous static memory. SXCNFG[15:0] configures chip select signals 0 and 1. SXCNFG[31:16] configures chip select signals 2 and 3.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 6-13. SXCNFG Bit Definitions (Sheet 1 of 4)

	0x4800_001C										SXCNFG									Memory Controller											
Bit	31	30	29 28	27	26	25	24	23	22	21	20	19	18	17	7 16	15	14	13 1	2	11	10	9	8	7	6	5	4	3	2	1	0
	reserved	SXLATCH2	SXTP2		SXCA2		SXRA2	S	XRL	2	SXCL2 NH X			reserved	SXLATCH0	SXTP0		SXCA0				SXRL0			SXCL0		ONEXS	CALIN			
Reset	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	*	*	*	*	*	*	*	*	*	*	*	*	*	0	*
		Bit	ts		Na	ame)											Des	cr	iptic	on										
		3′	1		-	_		reserved																							
								SXMEM latching scheme for pair 2/3																							
	30				30 SXLATCH2				O – Latch return data with fixed delay on MEMCLK 1 – Latch return data with return clock																						
									Must be set to a 1 to enable the return clock SDCLK for latching data. For more details on this return data latching, see Section 6.5.4.														on								
							SX Memory type for partition pair 2/3																								
	20.20			29:28 SXTP2			00 – Synchronous Mask ROM (SMROM) 01 – reserved																								
					0,1		-	10 – non-SDRAM-like Synchronous Flash																							
								11 - reserved																							
								SX Memory column address bit count for partition pair 2/3																							
		27:26 SXCA2			00 – 7 column address bits 01 – 8 column address bits																										
					•		_	10 – 9 column address bits																							
								11 -	- 10	col	umı	n ad	dres	ss I	bits																
								SX Memory row address bit count for partition pair 2/3 00 – 12 row address bits																							
		25:	24		SX	RA	2		– 12 – 13																						
								10	– res	serv	ed																				
								11 -	– res	serv	ed																				
											•				•			pair 2					۸.							- ^ -	
																		tweer the e								nano	ano	d th	e KI	=AD	,
									SXTI			0" (8	SMF	ROI	M):																
			000) – 1	clo	ck																									

IF SXTP2 = 10 (non-SDRAM timing Fast Flash), this field is not used and must be

SXRL2

23:21

001 – 2 clocks 010 – 3 clocks

011 – 4 clocks 100 – 5 clocks 101 – 6 clocks 110 – 7 clocks 111 – 8 clocks

programmed to 111.



Table 6-13. SXCNFG Bit Definitions (Sheet 2 of 4)

Ta	Table 6-13. SXCNFG Bit Definitions (Sheet 2 of 4)															
	0x4	1800_001C		8	SXCNF		Memory Controller									
Bit	31 30 29 28	27 26 25 24	23 22 21	20 19 18	17 16	15	14	13 12	11 10	9 8	7 6 5	4	3 2	1 0		
	SXLATCH2 SXTP2	SXCA2 SXRA2	SXRL2	SXCL2	SXEN2	reserved	SXLATCH0	SXTP0	SXCA0	SXRA0	SXRL0	SXCL0		SXENO		
Reset	0 0 0 0	0 0 0 0	0 0 0	0 0 0	0 0	0	*	* *	* *	* *	* * *	*	* *	0 *		
•	Bits	Name						Descr	iption							
	20:18	SXCL2	Number of data. The uthe memory as to which IF SXTP2 = 000 - resel 010 - 3 clo 011 - 4 clo 100 - 7 clo 110 - 3 clo 011 - 4 clo 100 - 7 clo 110 - 5 clo 110 - 3 clo 011 - 4 clo 100 - 5 clo 111 - 7 clo	rved cks cks cks rved rved = 10 (non-S rved cks cks cks cks cks cks	DCLK cy SXCL2 uency (M ncy to us DM):	cles is th MDR se, u	bet e e: EFI se t	ween rexternal services and the nexternal services are recorded as the nextended as the nexternal services are recorded as the nexternal services are recorded as the nextended as the ne	eceiving SDCLK 32 = 1), t t larger.	cycle. V	When SX M y is 2*mem	emory	runs	at half		
	17:16	SXEN2	Enable Bits for SX Memory Partition 2 (bit 16) and Partition 3 (bit 17) 0 - Partition is not Enabled as SX Memory 1 - Partition is Enabled as SX Memory													
	15	— reserved														
	14	SXLATCH0	0 – Latch 1 – Latch Must alway	tching sche return data return data /s be writter I on this retu	with fixe with reton with a	ed do urn o 1 to	elay cloc ena	k able the	return o		OCLK for lat	tching	data.	For		
	13:12	SXTP0	SX Memory type for partition pair 0/1 00 – Synchronous Mask ROM (SMROM) 01 – reserved 10 – non-SDRAM-like Synchronous Flash													

11 - reserved



Table 6-13. SXCNFG Bit Definitions (Sheet 3 of 4)

				0x4	180	0_0	010	;						,	SXCI	NFC	3					Me	emo	ory (Con	troll	ler		
Bit	31	30	29	28	27	26	25	5 24	23	22	21	20	19	18	17	16	15	14	13 12	11 10	9	8	7	6	5	4	3	2	1 0
	reserved	SXLATCH2	SXTD2		0,000	SACAZ		SXRA2	•	SXR	L2	8	XCI	_2	SXEN2		reserved	SXLATCH0	SXTP0	SXCA0	CVDAD	SARAU	00	SXRI	L 0	S	XCL	0	SXENO
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	*	* *	* *	*	*	*	*	*	*	*	*	0 *
		Bi	ts			Na	ıme)											Desci	ription									
		11:	10			SX	CA	0	00 01 10	- 7 - 8 - 9	colu colu	, ımn ımn	add add	lres: lres: lres:	ddres s bits s bits s bits	; ;	oit co	ount	for par	tition pa	ir O/	1							
												•	w a			it co	ount	for	partitio	n pair 0/	′1								



Table 6-13. SXCNFG Bit Definitions (Sheet 4 of 4)

Ia	DIE	0-1	J. 0	λΟ.		ט כ		eiiiiii		3 (0	лес	7	01 4	')													
			0 x	480	0_00)1C					:	SXC	NF	3						Me	emor	/ Coi	ntro	ller			
Bit	31	30	29 28	27	26	25	24	23 22	21	20	19 18	17	16	15	14	13 12	11	10	9	8	7	5	4	3	2	1	0
	reserved	SXLATCH2	SXTP2		SXCA2	CVDVO	24446	SXR	L 2	SX	CL2		SXENZ	reserved	SXLATCH0	SXTP0	SXCA0		UVAXS	ONING TOTAL	SX	RL0	•	SXC	LO	CVEND	SAEINU
Reset	0	0	0 0	0	0	0	0	0 0	0	0	0 0	0	0	0	*	* *	*	*	*	*	*	* *	*	*	*	0	*
		Bit	s		Na	me										Descr	riptio	n									
								CASI	aten	cv fo	r SX N	1em	orv	narti	tion	nair 0/	1										
	Bits Name Description CAS Latency for SX Memory partition pair 0/1 Number of external SDCLK cycles between reception of the READ command and I of the data. The unit size for SXCL0 is the external SDCLK cycle. When SX Memor at half the memory clock frequency (MDREFR:KODB2 = 1), the delay is 2*MEMCLI in doubt as to which CAS Latency to use, the next larger must be used. IF SXTP0 = 00 (SMROM): 000 - reserved 001 - reserved 010 - 3 clocks 011 - 4 clocks 100 - 5 clocks 110 - reserved 111 - reserved IF SXTP0 = 10 (non-SDRAM timing Fast Flash) 000 - reserved 001 - reserved 101 - 3 clocks 011 - 4 clocks 100 - 5 clocks 110 - 7 clocks														emory	/ is ı	run										
		1:0	0		SXI	EN0		Enable 0 - Pa 1 - Pa For res	artiti artiti	on is on is	not er enable	able ed a	ed a	s SX K Me	(Me emo		and I	Par	titio	n 1 ((bit 1)						

6.6.1.1 SMROM Memory Options

Table 6-15 shows the possible external-to-internal address multiplexing options. For SMROM, there are no bank-address bits, but the corresponding bits are put on the external address bus. The number of banks per device always defaults to four.



			_									_													-	_		_		
# Bits Bank x		Ext	tern	al A	Add				at S :10>		EM	RA	S Ti	me			Ext	terr	al A	Add		s pi MA-				EM	CA	S Ti	me	
Row x Col x Data	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10
2x12x7x32		22	21	20	19	18	17	16	15	14	13	12	11	10	9		22	21		0				8	7	6	5	4	3	2
2x12x7x16		21	20	19	18	17	16	15	14	13	12	11	10	9	8		21	20		0				7	6	5	4	3	2	1
2x12x8x32		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		0			9	8	7	6	5	4	3	2
2x12x8x16		22	21	20	19	18	17	16	15	14	13	12	11	10	9		22	21		0			8	7	6	5	4	3	2	1
2x12x9x32		24	23	22	21	20	19	18	17	16	15	14	13	12	11		24	23		0		10	9	8	7	6	5	4	3	2
2x12x9x16		23	22	21	20	19	18	17	16	15	14	13	12	11	10		23	22		0		9	8	7	6	5	4	3	2	1
2x12x10x32		25	24	23	22	21	20	19	18	17	16	15	14	13	12		25	24		0	11	10	9	8	7	6	5	4	3	2
2x12x10x16		24	23	22	21	20	19	18	17	16	15	14	13	12	11		24	23		0	10	9	8	7	6	5	4	3	2	1
2x13x7x32	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	23	22			0				8	7	6	5	4	3	2
2x13x7x16	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	22	21			0				7	6	5	4	3	2	1
2x13x8x32	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23			0			9	8	7	6	5	4	3	2
2x13x8x16	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	23	22			0			8	7	6	5	4	3	2	1
2x13x9x32	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	25	24			0		10	9	8	7	6	5	4	3	2
2x13x9x16	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	24	23			0		9	8	7	6	5	4	3	2	1
2x13x10x16	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	25	24			0	10	9	8	7	6	5	4	3	2	1

Table 6-15. Synchronous Static Memory External to Internal Address Mapping Options

6.6.2 Synchronous Static Memory Mode Register Set Configuration Register (SXMRS)

On power up, a MRS command that contains the default boot-up value is written to the external memory if the system is configured to boot out of SMROM (see Section 6.10.2). Otherwise, SXMRS, shown in Table 6-16, is used to issue an MRS command to SMROM. The value written in this register is placed directly on address lines MA<24:10> during the MRS command. Writing to this register triggers a two-stage MRS command that is sent to the external synchronous static memory. The first state issues an MRS command to banks 0 and 1. The second stage issues an MRS command to banks 2 and 3. The corresponding chip select values are only asserted if the memory banks are enabled via the SXCNFG register and the memory type is configured as SMROM.

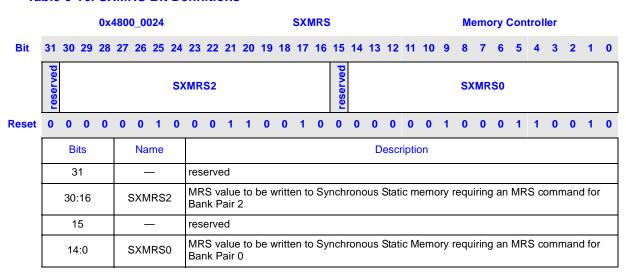
To write a new MRS value to a synchronous static memory, first enable and configure the memory via the SXCNFG register, then write the SXMRS register. This register is only used for the value written during the MRS command. All values in the SXCNFG register must be programmed correctly to ensure proper device operation (refer to the external memory chip product documentation for proper MRS encoding). Information programmed in the SXCNFG[CL] and



SXCNFG[RL] fields must match any CAS latencies and RAS latencies programmed in this SXMRS register. Software must ensure that fields match the latencies. In some cases, duplicate information must be programmed.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 6-16. SXMRS Bit Definitions



6.6.3 Synchronous Static Memory Timing Diagrams

Figure 6-12 shows a three-beat read cycle for SMROM.



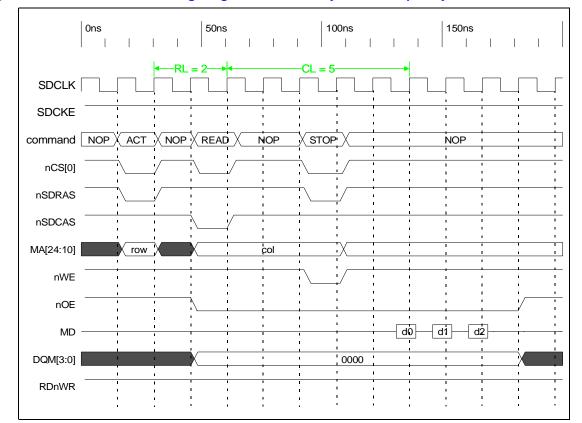


Figure 6-12. SMROM Read Timing Diagram Half-Memory Clock Frequency

6.6.4 Non-SDRAM Timing SXMEM Operation

Non-SDRAM Timing Synchronous Flash operation resets to asynchronous mode (page-mode for reads and asynchronous single word writes). Software can change the Read Configuration Register (RCR) to synchronous mode (burst-timing synchronous reads and asynchronous single word writes). At boot time, the Non-SDRAM Timing Synchronous Flash operates similarly to an asynchronous boot ROM (See Section 6.8).

Consult the documentation for the memory being used. Table 6-17 is provided only as a reference. The frequency configuration must be determined based on the CLK-to-output delay, the CLK period, and the nADV-to-output delay timing parameters of the Flash device.

The values for this part number are shown as an example. For Intel part number 28F800F3, programming values for this register to ensure proper operation with the processors are shown in Table 6-17.

Software must ensure that the CLK-to-output delay is less than 1 SDCLK period for non-SDRAM Timing Synchronous Flash.



Table 6-17. Read Configuration Register Programming Values

Bits	Field Name	Value to Program
2:0	BURST LENGTH	010 8 Word Burst
5:3	reserved	000
6	CLOCK CONFIGURATION	1 Use rising edge of clock
7	BURST SEQUENCE	1 Linear burst Order (INTEL BURST ORDER IS NOT SUPPORTED)
8	WAIT CONFIGURATION	N/A nWAIT from the Flash device is ignored by the processor.
9	DATA OUTPUT CONFIGURATION	0 Hold data for one clock
10	reserved	0
13:11	FREQUENCY CONFIGURATION	010 -> CAS Latency 3 011 -> CAS Latency 4 100 -> CAS Latency 5 101 -> CAS Latency 6 110 -> CAS Latency 7 Chosen based on the AC Characteristics - Read only
		Operation section of the Flash device data sheet
14	reserved	0
15	READ MODE	0 - Synchronous Operation 1 - Asynchronous Operation

Table 6-18 shows sample frequency configurations for programming non-SDRAM Timing Fast Flash. When in doubt, the higher frequency configuration and corresponding CAS latency must be used.

Table 6-18. Frequency Code Configuration Values Based on Clock Speed (Sheet 1 of 2)

MEMCLK Frequency	SDCLK0 Frequency	MDREFR: K0DB2	Valid Frequency Configurations	Corresponding CAS Latencies
20	20	0	2/3/4/5/6	3/4/5/6/7
33	33	0	3/4/5/6	4/5/6/7
50	50	0	4/5/6	5/6/7
30	25	1	2/3/4/5/6	3/4/5/6/7
66	66	0	5/6	6/7
00	33	1	3/4/5/6	4/5/6/7
100	50	1	4/5/6	5/6/7
118	59	1	5/6	6/7



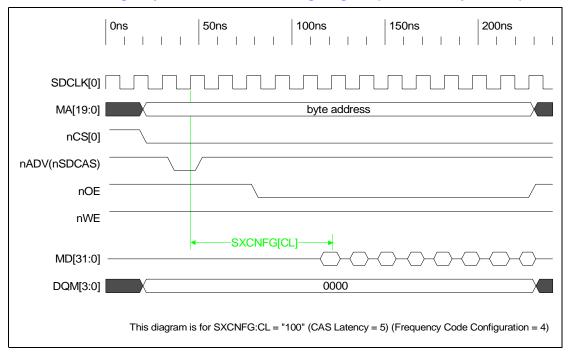
Table 6-18. Frequency Code Configuration Values Based on Clock Speed (Sheet 2 of 2)

MEMCLK Frequency	SDCLK0 Frequency	MDREFR: K0DB2	Valid Frequency Configurations	Corresponding CAS Latencies
133	66	1	5/6	6/7
147	Not supported			
166	Not supported			

6.6.4.1 Non-SDRAM Timing Flash Read Timing Diagram

Figure 6-12 shows the burst-of-eight read timing diagram.

Figure 6-13. Burst-of-Eight Synchronous Flash Timing Diagram (non-divide-by-2 mode)



In Figure 6-13, the following timing parameters apply:

- nADV asserted time = 1 MEMCLK
- MA, nCS setup to nADV asserted = 1 MEMCLK
- nADV deasserted to nOE asserted = Code 2 MEMCLKs



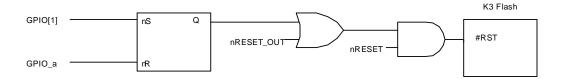
For divide-by-two mode, the following timing parameters apply:

- nADV assert time = 3 MEMCLKs
- MA, nCS setup to nADV asserted = 1 MEMCLK
- nADV deasserted to nOE asserted = (Code * 2) 4 MEMCLKs

6.6.4.2 K3 Synchronous StrataFlash Reset

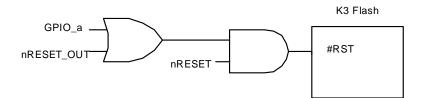
The PXA255 processor nRESET_OUT pin must be connected to the K3 #RST pin for Hardware reset, Watchdog reset and sleep mode to work properly. GPIO reset however does not reset the contents of the memory controller configuration register. If GPIO reset operation is required, a state machine is necessary between nRESET, nRESET_OUT, GPIO[1], and #RST to ensure that #RST is asserted during Hardware reset, Watchdog reset, and sleep mode, and not asserted during GPIO reset. Figure 6-14 shows the required logic. GPIO_a is an unused GPIO that is driven low by software during the initialization sequence and left high during normal operation. After this is completed, then enable GPIO reset.

Figure 6-14. Flash Memory Reset Using State Machine



If Watchdog reset is not necessary, a secondary GPIO can control nRESET_OUT using the equation #RST = nRESET & (nRESET_OUT | GPIO_a). This allows sleep mode entry to reset the flash memory while keeping it in synchronous mode during a GPIO reset. Figure 6-15 shows the required logic. GPIO_a is an unused GPIO that is kept high during normal operation and driven low before sleep mode entry.

Figure 6-15. Flash Memory Reset Logic if Watchdog Reset is Not Necessary



6.7 Asynchronous Static Memory

6.7.1 Static Memory Interface

The Static Memory interface is made up of six chip selects: nCS[5:0]. The chip selects can be configured as the following:



- Non-burst ROM or Flash memory
- · Burst ROM or Flash
- SRAM
- SRAM-like variable latency I/O devices

The Variable Latency I/O interface differs from SRAM in that it allows the use of the data-ready input signal, RDY, to insert a variable number of memory-cycle wait states. The data bus width for each chip-select region can be programmed as 16- or 32-bit. nCS[3:0] can also be configured for Synchronous Static Memory (refer to Section 6.6). During Variable Latency I/O writes, nPWE is used instead of nWE so SDRAM refreshes can be executed while performing the VLIO transfers.

The use of the signals nOE, nWE, and nPWE is summarized below:

- nOE is asserted for all reads
- nWE is asserted for Flash and SRAM writes
- nPWE is asserted for Variable Latency I/O writes

For SRAM and Variable Latency I/O implementations, DQM[3:0] signals are used for the write byte enables, where DQM[3] corresponds to the MSB. The processor supplies 26-bits of byte address for access of up to 64 Mbytes per chip select. This byte address is sent out on the 26 external address pins. Do not connect MA[1:0] for 32-bit systems. Do not connect MA[0] for 16-bit systems (the PXA255 processor operating in 16-bit mode). For all reads on a 32 bit system DQM[3:0] and MA[1:0] are 0. For all reads on a 16 bit system DQM[1:0] and MA[0] are 0. In the timing diagrams, these byte addresses are shown and referred to as "addr".



Table 6-19. 32-Bit Bus Write Access

Data Size	MA[1:0]	DQM[3:0]
8-bit	00	1110
8-bit	01	1101
8-bit	10	1011
8-bit	11	0111
16-bit	00	1100
16-bit	10	0011
32-bit	00	0000

Table 6-20, 16-Bit Bus Write Access

Data Size	MA[0]	DQM[1:0]
8-bits	0	10
8-bits	1	01
16-bits	0	00

The RT fields in the MSCx registers specify the type of memory:

- Non-burst ROM or Flash
- SRAM
- Variable Latency I/O
- · Burst-of-four ROM or Flash
- Burst-of-eight ROM or Flash

The RBW fields specify the bus width for the memory space selected by nCS[5:0]. For a 16-bit bus width transactions occur on MD[15:0]. The BOOT_SEL pins and/or SXCNFG register must be used to configure nCS[3:0] for SMROM or some other type of Synchronous Static Memory.

6.7.2 Asynchronous Static Memory Control Registers (MSCx)

The MSCx, shown in Table 6-21, are read/write registers and contain control bits for configuring Static Memory (or Variable Latency I/O) that correspond to chip-select pairs nCS(1:0), nCS(3:2), and nCS(5:4), respectively. Timing fields are specified as numbers of memory clock cycles. Each of the three registers contain two identical CNFG fields One for each chip select in the pair.

When programming a different memory type in an MSC register, ensure that the new value has been accepted and programmed before issuing a command to that memory. To do this, the MSC register must be read after it is written and before an access to the memory is attempted. This is especially important when changing from ROM/Flash to an unconstrained writable memory type (such as SRAM).

If any of the nCS[3:0] banks is configured for Synchronous Static Memory via SXCNFG[SXENx], the corresponding half-words of MSC0 and/or MSC1 are ignored, except MSCx:RBWx, the data width. Another exception is non-SDRAM timing Synchronous Flash, which writes asynchronously and requires these programmed values.



This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 6-21. MSC0/1/2 Bit Definitions (Sheet 1 of 3)

			0x4	1800	0_0 0_0 0_0	00C						•		MS MS MS	C1								M	lem	ory	y Co	onti	roll	er			
Bit	31	30 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	, (6 5	5	4	3	2	1	0
	RBUFF1/3/5	RRR1/	/3/5	F	RDN	1/3/	5	R	DF1	1/3/	5	RBW1/3/5	R	Γ1/3	/5	RBUFF0/2/4		RRR0/2/4		F	RDN	10/2	2/4		RD)F0/	2/4	,	RBW0/2/4	R	T0/2	2/4
Reset	0	1 1	1	1	1	1	1	1	1	1	1	0	0	0	0	0	1	1	1	1	1	1	1	1	•	1 1	1	1	*	0	0	0
		Bits			Acc	ess	i		Nar	ne										De	esc	ript	ion									
		15			R	/W		RBI	JFF	×		Wh slow rem the bus RBI tran tran Wh ban dev spe 0 - 1 -	en s v SI lain sys ma JFF lsac en Si k, tl ice) cific SI Fa	Data RAM idle idle tem by ha bit ction Sync his v cally owe aster	er n M/RC whi is a alt w doe beg sta chro value e re bee	nem DM) ile a allow while is no gins rts. nou e wil egist en p vice vice	ory , this ll da red t all o at aff on t s St ll de er b rogr	dev s bit ta is to p data fect he r atic faul it wi ami etur	rices t mu s be roce a is i the men me lt to iill sti med n Da ning	s are st b ing ess of return beh nory mor Stre ill re I to a	e be e re rea othe rne avi bu y d eam ad a 1.	eing d from or of s, it evices ing as (Th	use to a om to forn om to f the muse ses h beh	allovenation he construction has a construction he construction has a con	w th dev dev terr e co e be or (ne syrice. Wherice. nal romp	yste By en Th ner lete ena um Bu	em res set ne v mor ed b able ing	to n setti t, the alue y bu pefo ed fo a fa) un	ot hing the interest of the in	his tern the Onc inot	bit, al e a her en
		14:12			R	/W		RRI	₹x<:	2:0>	>	Chi san mei This	p se ne s mcll s fie	SRAM elect station ks. eld m ulse	dea me	asse emoi be ¡	erted ry ba	d aft ank) gran	er a) or nme	nSE d w	ith t	as the	serte max	ed i imu	is e	qual of t _C	l to OFF	(RÌ	RRx vide	: * 2 d b	y 2)	,
		11:8			R	/W		RDI	Nx<	3:0>	^	Add to (I nW The Var	lres RDI E as nC	elay s to Nx + sser E (r e La s nu	data 1) i tion PW tend	a va men for /E) o	lid fonctk write deas O is	or si s. e ac sser eqi	ces rt tim ual t	ses ne b o (F	to setw	SR <i>A</i> /eer	M is n ead · 2) r	s ec ch b ner	qual bea mcll	I to (t of ks. F	(RE	DFx ad/w	+ 1 vrite) me	emc	lks.



Table 6-21. MSC0/1/2 Bit Definitions (Sheet 2 of 3)

 0x4800_0008
 MSC0

 0x4800_000C
 MSC1
 Memory Controller

 0x4800_0010
 MSC2

Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

RBUFF0/2/4 BUFF0/2/4 BUFF0	RDN0/2/4 RDF0/2/4 PT0/2/4
--------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------	---------------------------

Reset 0 1 1 1 1 1 1 1 1 1 1 0 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 1 * 0 0 0

Bits	Access	Name	Description
7:4	R/W	RDFx<3:0>	ROM delay first access. RDF programmed RDF value interpreted 0-11 0-11 12 13 13 15 14 18 15 23 Address to data valid for the first read access from all devices except VLIO is equal to (RDFx + 2) memclks. Address to data valid for subsequent read accesses to non-burst devices is equal to (RDFx + 1) memclks. nWE assertion for write accesses (which are non-burst) to all Flash is equal to (RDFx + 1) memclks. nOE (nPWE) assert time for each beat of read (write) is equal to (RDFx + 1) memclks for Variable Latency I/O (nCS[5:0]). For Variable Latency I/O, RDFx must be greater than or equal to 3.
3	R/W	RBWx	ROM bus width 0 - 32 bits 1 - 16 bits For reset value for RBW0, see Section 6.8. This value must be programmed with all memory types including Synchronous Static Memory. This value must not change during normal operation.



Table 6-21. MSC0/1/2 Bit Definitions (Sheet 3 of 3)

				0x4	1800	0_0	008 00C 010								MS	SC0 SC1 SC2								M	emo	ory (Con	trol	ler			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	(
	3/5												5				4/		4										4			

	4																Ľ															
Reset	0	1	1	1	1	1	1	1	1	1	1	1	0	0	0	0	0	1	1	1	1	1	1	1	1	1	1	1	*	0	0	0

Bits	Access	Name	Description
2:0	R/W	RTx<2:0>	ROM type 000 - Nonburst ROM or Flash Memory 001 - SRAM 010 - Burst-of-four ROM or Flash (with non-burst writes) 011 - Burst-of-eight ROM or Flash (with non-burst writes) 100 - Variable Latency I/O (VLIO) 101 - reserved 110 - reserved 111 - reserved Burst refers to the device's timing. When the subsequent reads from the device take less time than the first read from a device, it is referred to as burst timing. The address bits must also be taken into account for burst timing devices. For example, in a burst-of-four device, only the lower two non-byte address bits can change for burst timing. For 32-bit devices, this is MA[3:2]. The address order can go 00, 01, 10, 11 where the reads from 01, 10, and 11, take less time to come out of the device. For burst-of-eight devices, the lower three non-byte address bits can change. Writes to these devices are non-burst.



Table 6-22 provides a comparison of supported Asynchronous Static Memory types.

Table 6-22. Asynchronous Static Memory and Variable Latency I/O Capabilities

			1	Γiming (Memo	ry Clocks)		
MSCx[RTx]	Device Type	Burst Read Address Assert	nOE Assert	Burst nOE Deassert	Burst Write Address Assert	nWE Assert	Burst nWE De- assert
000	Non-burst ROM or Flash	RDF+1	RDF+1	0	N/A	RDF+1	N/A
001	SRAM	RDF+1	RDF+1	0	RDN+2	RDN+1	1
010	Burst-of-4 ROM or Flash (non- burst writes)	RDF+1 (0,4) RDN+1 (1:3,5:7)	RDF+1 (0,4) RDN+1 (1:3,5:7)	0	N/A	RDF+1	N/A
011	Burst-of-8 ROM or Flash (non-burst writes)	RDF+1 (0) RDN+1 (1:7)	RDF+1 (0) RDN+1 (1:7)	0	N/A	RDF+1	N/A
100	Variable Latency I/O	RDF+ RDN+2+waits	RDF+1+ waits	RDN+2	RDF+ RDN+2+waits	RDF+1+ waits	RDN+2

6.7.3 ROM Interface

The processor provides programmable timing for both burst and non-burst ROMs. The RDF field in MSCx is the latency (in memory clock cycles) for the first, and all subsequent, data beats from non-burst ROMs, and the first data beat from a burst ROM. RDN is the latency for the burst data beats after the first for burst ROMs. RRR delays the following access to a different memory space to allow time for the current ROM to three-state the data bus.

RRR must be programmed with the maximum t_{OFF} value, as specified by the ROM manufacturer.

For hardware reset initialization values, refer to Section 6.8. MSC0[15:0] is selected when the address space corresponding to nCS0 is accessed. The processor supports a ROM burst size of 1, 4, or 8 by configuring the MSCx[RTx] register bits to 0, 2 or 3 respectfully.



6.7.3.1 ROM Timing Diagrams and Parameters

Figure 6-17, Figure 6-18, and Figure 6-19 show the timings for burst and non-burst ROMs.

Figure 6-17. 32-Bit Burst-of-Eight ROM or Flash Read Timing Diagram (MSC0[RDF] = 4, MSC0[RDN] = 1, MSC0[RRR] = 1)

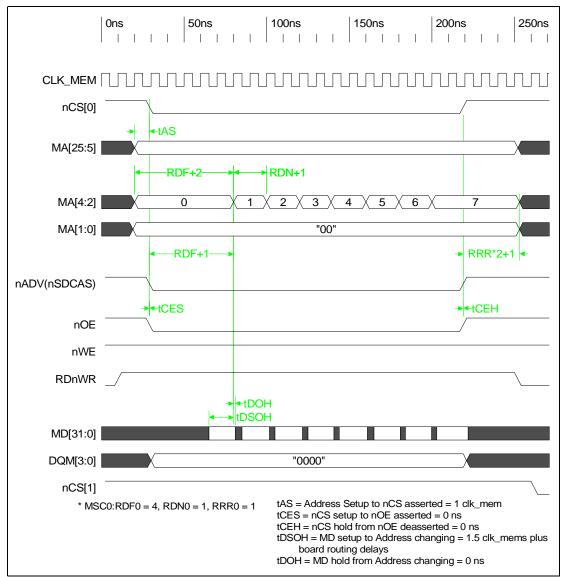
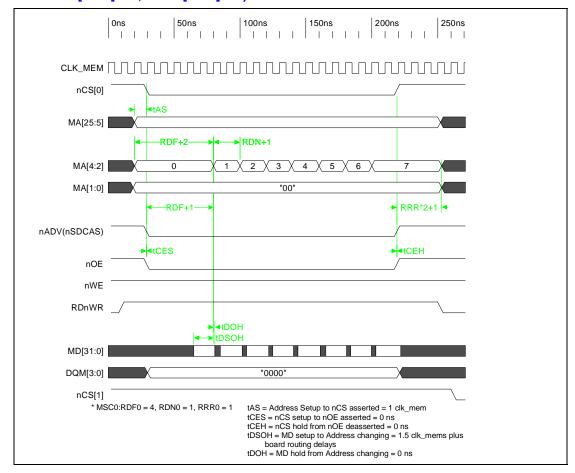




Figure 6-18. Eight-Beat Burst Read from 16-Bit Burst-of-Four ROM or Flash (MSC0[RDF] = 4, MSC0[RDN] = 1, MSC0[RRR] = 0)





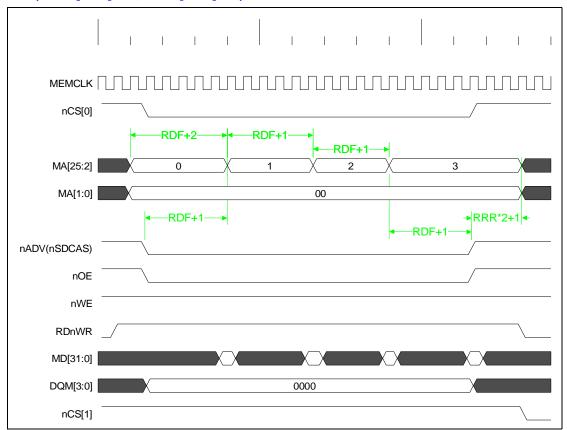


Figure 6-19. 32-Bit Non-burst ROM, SRAM, or Flash Read Timing Diagram - Four Data Beats (MSC0[RDF] = 4, MSC0[RRR] = 1)

6.7.4 SRAM Interface Overview

The processor provides a 16-bit or 32-bit asynchronous SRAM interface that uses the DQM pins for byte selects on writes. nCS[5:0] select the SRAM bank. nOE is asserted on reads and nWE is asserted on writes. Address bits MA[25:0] allow up to 64 Mbytes of SRAM per bank to be addressed.

The timing for a read access is identical to that for a non-burst ROM (see Section 6.7.3.1). The RDF fields in the MSCx registers select the latency for a read access. The MSCx[RDN] field controls the nWE low time during a write cycle. MSCx[RRR] is the time from nCS deassertion after a memory access to the start of another memory access. MSCx[RTx] must be configured to 0b001 to select SRAM.

6.7.4.1 SRAM Timing Diagrams and Parameters

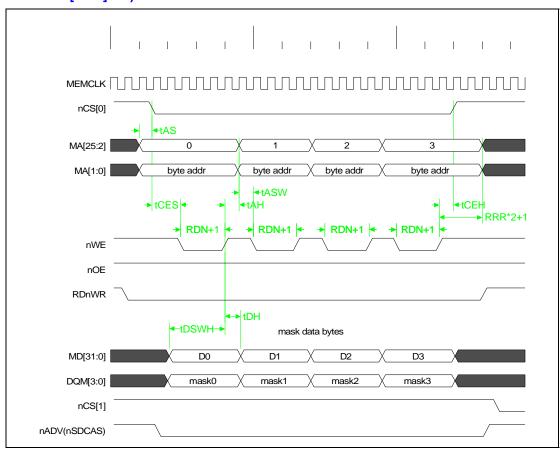
As shown in Figure 6-17, SRAM reads have the same timing as non-burst ROMs, except DQM[3:0] are used as byte selects. For all reads, DQM[3:0] are 0b0000. During writes, all 32 data pins are actively driven by the processor regardless of the state of the individual DQM pins.



For writes to SRAM, if all byte enables are turned off (masking out the data, DQM = 1111), then the write enable are 1 (nWE = 1) for this write beat. This can result in a period when nCS is asserted, but neither nOE nor nWE is asserted. This happens with a write of 1 beat to SRAM, but all byte enables are turned off.

Figure 6-20 shows the timing for SRAM writes.

Figure 6-20. 32-Bit SRAM Write Timing Diagram (4-beat Burst (MSC0[RDN] = 2, MSC0[RRR] = 1)



In Figure 6-20, the parameters are defined as follows:

- tAS = Address setup to nCS = 1 MEMCLK
- tCES = nCS setup to nWE = 2 MEMCLKs
- tASW = Address setup to nWE low (asserted) = 1 MEMCLK
- tDSWH = Write data setup, DQM to nWE high (deasserted) = (RDN+2) = 4 MEMCLKs
- tDH = Data, DQM hold after nWE high (deasserted) = 1 MEMCLK
- tCEH = nCS held asserted after nWE deasserted = 1 MEMCLK
- tAH = Address hold after nWE deasserted = 1 MEMCLK
- nWE high time between burst beats = 2 MEMCLKs



6.7.5 Variable Latency I/O (VLIO) Interface Overview

Variable Latency I/O read accesses differ from SRAM read accesses in that the nOE toggles for each beat of a burst. The first nOE assertion occurs two memory cycles after the assertion of the chip select nCS<x>. Also, for Variable Latency I/O writes, nPWE is used instead of nWE so SDRAM refreshes can be executed while performing the VLIO transfers. Variable Latency I/O is selected by programming the MSCx[RTx] bits as 0b100.

Both reads and writes for VLIO differ from SRAM in that the processor samples the data-ready input, RDY. The RDY signal is level sensitive and goes through a two-stage synchronizer on input. When the internal RDY signal is high, the I/O device is ready for data transfer. This means that for a transaction to complete at the minimum assertion time for either nOE or nPWE (RDF+1), the RDY signal must be high two clocks prior to the minimum assertion time for either nOE or nPWE (RDF-1). Data will be latched on the rising edge of MEMCLK once the internal RDY signal is high and the minimum assertion time of RDF+1 has been reached. Once the data has been latched, the address may change on the next rising edge of MEMCLK or any cycles thereafter. The nOE or nPWE signal will de-assert one MEMCLK after data is latched. Before a subsequent data beat, nOE or nPWE remains deasserted for RDN+2 memory cycles. The chip select and byte selects, DQM[3:0], remain asserted for one memory cycle after the burst's final nOE or nPWE deassertion.

For both reads and writes from/to VLIO, a DMA mode exists that does not increment the address to the VLIO, which will allow port-type VLIO chips to interface to the processor. See DCMDx[INCSRCADDR] and DCMDx[INCTRGADDR] in Table 5-12, "DCMDx Bit Definitions" on page 5-24.

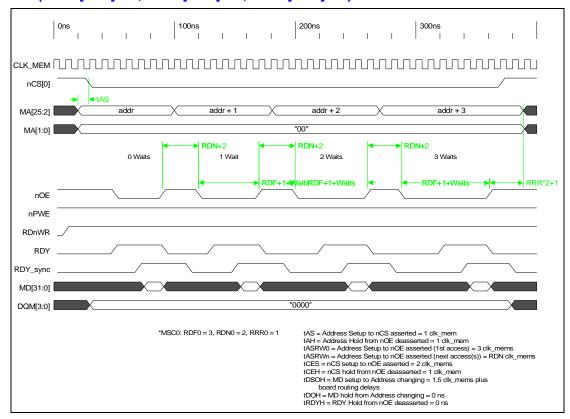
For writes to VLIO, if all byte enables are turned off, masking out the data, DQM = 1111, the write enable is suppressed (nPWE = 1) for this write beat to VLIO. This can result in a period when nCS is asserted, but neither nOE nor nPWE is asserted (this happens when there is a write of 1 beat to VLIO, but all byte enables are turned off).



6.7.5.1 Variable Latency I/O Timing Diagrams and Parameters

Figure 6-21 shows the timing for Variable Latency I/O reads and Figure 6-22 shows the timing for Variable Latency I/O writes.

Figure 6-21. 32-Bit Variable Latency I/O Read Timing (Burst-of-Four, One Wait Cycle Per Beat)
(MSC0[RDF] = 2, MSC0[RDN] = 2, MSC0[RRR] = 1)





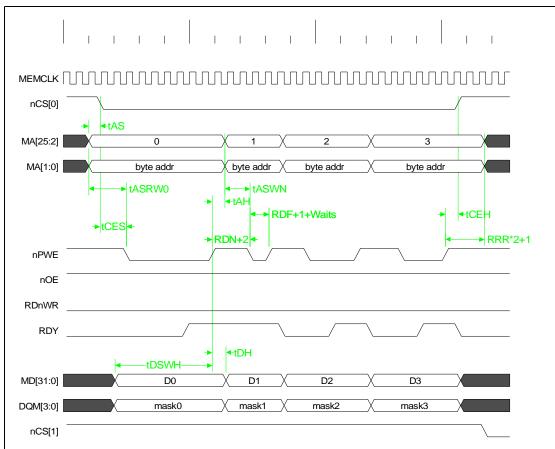


Figure 6-22. 32-Bit Variable Latency I/O Write Timing (Burst-of-Four, Variable Wait Cycles Per Beat)

In Figure 6-21 and Figure 6-22, some of the parameters are defined as follows:

- tAS = Address setup to nCS = 1 MEMCLK
- tCES = nCS setup to nOE or nPWE = 2 MEMCLKs
- tASRW0 = Address setup to nOE or nPWE low (asserted) = 3 MEMCLKs
- tASRWn = Address setup to nOE or nPWE low (asserted) = RDN MEMCLKs
- tDSWH,min = Minimum write data, DQM setup to nPWE high (deasserted) = (RDF+2) MEMCLKs
- tDHW = Data, DQM hold after nPWE high (deasserted) = 1 MEMCLK
- tDHR = Data hold required after nOE deasserted = 0 ns
- tCEH = nCS held asserted after nOE or nPWE deasserted = 1 MEMCLK
- tAH = Address hold after nOE or nPWE deasserted = 1 MEMCLK
- nOE or nPWE high time between burst beats = (RDN+2) MEMCLKs



6.7.6 FLASH Memory Interface

The processor provides an SRAM-like interface for access of Flash memory. The RDF fields in the MSCx registers are the latency for each read access to non-burst Flash, or the first read access to burst Flash. The RDF fields also control the nWE low time during a write cycle to Flash. The RDN field controls subsequent read access times to burst Flash and the nWE low time during a write cycle to non-burst Flash. RRR is the time from nCS deassertion after a read to the start of a read from a different memory, or after a write to another memory access.

Reads from Flash memory have the following requirements:

- Because Flash defaults to Read-Array mode, burst reads are permitted out of Flash, which allows instruction caching and DMA reads from Flash.
- Software partitions commands and data and writes the commands to Flash before the read. The Memory controller does not insert any commands before Flash reads.

Writes to Flash memory have the following requirements:

- Flash memory space must be uncacheable and unbuffered.
- Burst writes to Flash are not supported. Writes to Flash must be exactly the width of the populated Flash devices on the data bus and must be a burst length of one write, for example no byte writes to a 32-bit bus. The allowable writes are: 2 bytes written to a 16-bit bus, and 4-bytes written to a 32-bit bus.
- For asynchronous writes to Flash, the command and data must be given in separate write instructions to the Memory controller, the first carries the command, the next carries the data.
- The Memory controller does not insert any commands before Flash writes. Software must
 write the commands and data in the correct order.
- No Flash writes can be bursts. DMA must never write to Flash.

For writes to Flash, if all byte enables are turned off (masking out the data, DQM = 1111), the write enable is suppressed (nWE = 1) for the write beat, which can result in a period when nCS is asserted, but neither nOE nor nWE is asserted. This happens when there is a 1-beat write to Flash, but all byte enables are turned off.

6.7.6.1 FLASH Memory Timing Diagrams and Parameters

Non-burst Flash reads have the same timing as non-burst ROMs reads. Figure 6-23 shows the timing for writes to non-burst asynchronous Flash.



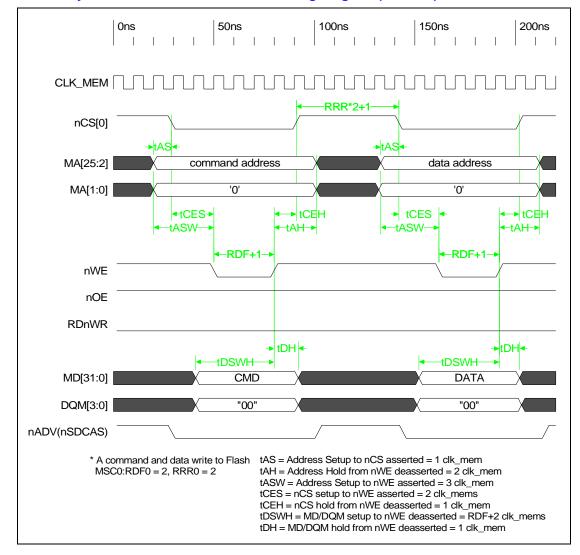


Figure 6-23. Asynchronous 32-Bit Flash Write Timing Diagram (2 Writes)

In Figure 6-23 some of the parameters are defined as follows:

- tAS = Address setup to nCS = 1 MEMCLK
- tCES = nCS setup to nWE = 2 MEMCLKs
- tASW = Address setup time to nWE asserted = 3 MEMCLKs
- tDSWH = Write data, DQM setup to nWE deasserted = (RDF+2) MEMCLKs
- tDH = Data, DQM hold after nWE deasserted = 1 MEMCLKs
- tCEH = nCS held asserted after nWE deasserted = 1 MEMCLK
- tAH = Address hold after nWE deasserted = 1 MEMCLKs



6.8 16-Bit PC Card/Compact Flash Interface

The following sections provide information on the card interface based on the *PC Card Standard - Volume 2 - Electrical Specification, Release 2.1*, and *CF+ and CompactFlash Specification Revision 1.4*. Only 8- and 16-bit data transfers are supported.

6.8.1 Expansion Memory Timing Configuration Register

MCMEM0, MCMEM1, MCATT0, MCATT1, MCIO0, and MCIO1 are read/write registers that contain control bits for configuring the timing of the 16-bit PC Card/Compact Flash interface.

The programming of each of the four fields in each of the six registers lets software to individually select the duration of accesses to I/O, common memory, and attribute space for each of two 16-bit PC Card/Compact Flash card slots.

Refer to Table 6-23 through Table 6-25 for bitmaps of the MCMEMx registers. Also refer to Table 6-26. Refer to Figure 6-29 and Figure 6-30 for a 16-bit PC Card/Compact Flash timing diagram.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 6-23. MCMEM0/1 Bit Definitions

)_00)_00										EM	•							Me	emo	ry (Con	troll	er			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					r	ese	rved	t						ME	Mx_	HO	LD		reserved		N	IEM	x_A	SS	т		l	MEN	/lx_	SET	•	
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
·		Ві	its			Na	me												De	scr	ipti	on										
·		31:	20			_	_		res	erve	ed																					
·		19:	:14		MC	OL	_												set IEMx					fore	e co	mm	and	ass	ertic	n fo	r	
·		13:	:12			_	_		res	erve	ed																					
		11	:7		MC	CME SS	-	_A			or th								. See	e Ta	able	6-20	5 fo	rac	lesc	ripti	on c	of thi	s co	ode a	and	its
,		6	:0		MC	CME E	-	_											set EMx				s be	fore	со	mm	and	ass	ertic	n fo	r	

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 6-24. MCATT0/1 Bit Definitions

)_00)_00									_	ATTO								Me	emo	ry (Con	trol	ler			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13 1	12	11	10	9	8	7	6	5	4	3	2	1	0
					r	ese	rved	d						AT	Tx_	HOI	LD		reserved		A	ιΤτ	_A	SST	Γ			ΑT	Гх_9	SET		
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Naı	me												Des	scri	iptio	n										
		31:	20			_			rese	erve	ed																					
		19:	:14		MC	CATT LI	_	Ю											set Tx_H				s be	efore	e co	mm	and	ass	ertic	on fo	r	
		13:	:12			_	_		rese	erve	ed																					
		11	:7		MC	CAT S	_	AS								serti sser			. See	· Ta	ble (6-26	fo	rac	desc	ripti	on c	of th	is co	ode	and	its
		6:	:0		МС	TAC	_	SE											set Tx_S				s be	fore	co	mm	and	ass	ertic	on fo	r	

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 6-25. MCIO0/1 Bit Definitions

					1800 1800	_									MC MC									Me	emo	ory (Con	troll	er			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					r	ese	rvec	ł						IC)x_H	HOL	D		reserved			IOx.	_AS	ST				IO	c_S	ET		
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	scr	iptio	on										
		31:	20			_	_		res	erve	ed																					
		19:	:14		MC	XOIS	(_H(set + 2.		add	ress	bef	ore	con	nma	nd a	asse	rtior	n for	МС	Ю
		13:	12			_	_		res	erve	ed																					
		11	:7		МС	KOIC T	_				or th								. See	e Ta	able	6-26	o for	a d	esc	ripti	on c	of thi	s cc	de a	and	its
		6:	:0		МС	CIO	_SE				m N ket x								set 2.	up	add	ress	bef	ore	con	nma	nd a	asse	rtior	n for	МС	Ю



Table 6-26. Card Interface Command Assertion Code Table

MCMEMx_ MCATTx_/		x_ASST_WAIT	x_ASST	r_HOLD	x_ASST x_ASST	
MCIOx_A	SST		(nPIOW asserted)	(nPIOR asserted)	(nPIOW asserted)	(nPIOR asserted)
Programmed Bit Value	Code decimal value	# MEMCLKs (minimum) to wait before checking for nPWAIT='1'	# MEMCLKs (minimum) to assert command (nPIOW) after nPWAIT='1'	# MEMCLKs (minimum) to assert command (nPIOR) after nPWAIT='1'	# MEMCLKs (minimum) command assertion time	# MEMCLKs (minimum) command assertion time
(Code)	(Code)	(Code + 2)	(2*Code + 3)	(2*Code + 4)	(3*Code + 5)	(3*Code + 6)
00000	0	2	3	4	5	6
00001	1	3	5	6	8	9
00010	2	4	7	8	11	12
00011	3	5	9	10	14	15
00100	4	6	11	12	17	18
00101	5	7	13	14	20	21
00110	6	8	15	16	23	24
00111	7	9	17	18	26	27
01000	8	10	19	20	29	30
01001	9	11	21	22	32	33
01010	10	12	23	24	35	36
01011	11	13	25	26	38	39
01100	12	14	27	28	41	42
01101	13	15	29	30	44	45
01110	14	16	31	32	47	48
01111	15	17	33	34	50	51
10000	16	18	35	36	53	54
10001	17	19	37	38	56	57
10010	18	20	39	40	59	60
10011	19	21	41	42	62	63
10100	20	22	43	44	65	66
10101	21	23	45	46	68	69
10110	22	24	47	48	71	72
10111	23	25	49	50	74	75
11000	24	26	51	52	77	78
11001	25	27	53	54	80	81
11010	26	28	55	56	83	84
11011	27	29	57	58	86	87
11100	28	30	59	60	89	90



Table 6-26. Card Interface Command Assertion Code Table

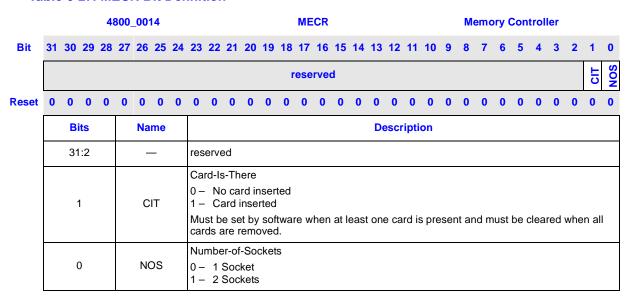
MCMEMx_ MCATTx_/		x_ASST_WAIT	x_ASST	HOLD	x_ASST_ x_ASST	
MCIOx_A	SST		(nPIOW asserted)	(nPIOR asserted)	(nPIOW asserted)	(nPIOR asserted)
Programmed Bit Value	Code decimal value	# MEMCLKs (minimum) to wait before checking for nPWAIT='1'	# MEMCLKs (minimum) to assert command (nPIOW) after nPWAIT='1'	# MEMCLKs (minimum) to assert command (nPIOR) after nPWAIT='1'	# MEMCLKs (minimum) command assertion time	# MEMCLKs (minimum) command assertion time
(Code)	(Code)	(Code + 2)	(2*Code + 3)	(2*Code + 4)	(3*Code + 5)	(3*Code + 6)
11101	29	31	61	62	92	93
11110	30	32	63	64	95	96
11111	31	33	65	66	98	99

6.8.2 Expansion Memory Configuration Register (MECR)

To eliminate external hardware, the two bits in MECR, shown in Table 6-27, are used to signal the memory controller when a card (16-Bit PC Card/Compact Flash) is inserted in the socket and the number of cards supported in the system. The number-of-sockets bit is required because the PSKTSEL pin is used as the nOE for the data transceivers in single socket mode. The card-is-there bit is used to reduce external hardware by ignoring nIOIS16 and nPWAIT when there is no card inserted in the socket.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 6-27. MECR Bit Definition

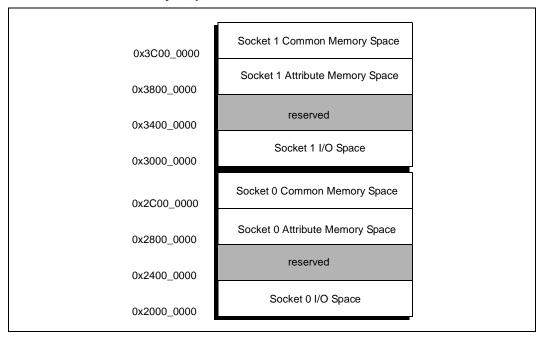




6.8.3 16-Bit PC Card Overview

The PXA255 processor 16-bit PC Card interface provides control for one 16-bit PC Card card slot with a PSKTSEL pin for support of a second slot. The PXA255 processor interface supports 8- and 16-bit peripherals and handles common memory, I/O, and attribute-memory accesses. The duration of each access is based on the values programmed in the fields in the MCMEMx, MCATTx, and MCIOx registers. Figure 6-26 shows the memory map for the 16-bit PC Card space.

Figure 6-26. 16-Bit PC Card Memory Map



The 16-bit PC Card Memory Map space is divided into eight partitions, four for each card slot. The four partitions for each card slot are: common memory, I/O, attribute memory, and a reserved space. Each partition starts on a 64-Mbyte boundary.

During an access, pins MA[25:0], nPREG, and PSKTSEL are driven at the same time. nPCE1 and nPCE2 are driven concurrently with the address signals for common memory and attribute-memory accesses. For I/O accesses, their value depends on the value of nIOIS16 and is valid a fixed amount of time after nIOIS16 is valid.

Common memory and attribute memory accesses assert the nPOE or nPWE control signals. I/O accesses assert the nIOR or nIOW control signals and use the nIOIS16 input signal to determine the bus width of the transfer (8 or 16 bits). The PXA255 processor uses nPCE2 to indicate to the expansion device that the upper half of the data bus (MD[15:8]) are used for the transfer, and nPCE1 to indicate that the lower half of the data bus (MD[7:0]) are used. nPCE1 and nPCE2 are asserted for 16-bit accesses.

Refer to Table 6-28 through Table 6-35 for signal combinations during common memory, I/O, and attribute accesses.



When writes goes to a card sockets and a byte has been masked via an internal byte enable, the write does not occur on the external bus. For reads, one half-word is always read from the socket, even if only 1 byte is requested. In some cases, based on internal address alignment, one word is read, even if only 1 byte is requested.

All DMA modes are supported in the Card interface increment the address.

Table 6-28. Common Memory Space Write Commands

nPCE2	nPCE1	MA<0>	nPOE	nPWE	MD[15:8]	MD[7:0]
0	0	0	1	0	Odd Byte	Even Byte
1	0	0	1	0	Unimportant	Even Byte
1	0	1	1	0	Unimportant	Odd Byte

Table 6-29. Common Memory Space Read Commands

nPCE2	nPCE1	MA<0>	nPOE	nPWE	MD[15:8]	MD[7:0]
0	0	0	0	1	Odd Byte	Even Byte

Table 6-30. Attribute Memory Space Write Commands

nPCE2	nPCE1	MA<0>	nPOE	nPWE	MD[15:8]	MD[7:0]
0	0	0	1	0	Unimportant	Even Byte
1	0	0	1	0	Unimportant	Even Byte
1	0	1	1	0	Unimportant	Unimportant

Table 6-31. Attribute Memory Space Read Commands

nPCE2	nPCE1	MA<0>	nPOE	nPWE	MD[15:8]	MD[7:0]
0	0	0	0	1	Unimportant	Even Byte

Table 6-32. 16-Bit I/O Space Write Commands (nIOIS16 = 0)

nPCE2	nPCE1	MA<0>	nPIOR	nPIOW	MD[15:8]	MD[7:0]
0	0	0	1	0	Odd Byte	Even Byte
1	0	0	1	0	Unimportant	Even Byte
1	0	1	1	0	Unimportant	Odd Byte

Table 6-33. 16-Bit I/O Space Read Commands (nIOIS16 = 0)

nPCE2	nPCE1	MA<0>	nPIOR	nPIOW	MD[15:8]	MD[7:0]
0	0	0	0	1	Odd Byte	Even Byte



Table 6-34. 8-Bit I/O Space Write Commands (nIOIS16 = 1)

nPCE2	nPCE1	MA<0>	nPIOR	nPIOW	MD[15:8]	MD[7:0]
1	0	0	1	0	Unimportant	Even Byte
1	0	1	1	0	Unimportant	Odd Byte

Table 6-35. 8-Bit I/O Space Read Commands (nIOIS16 = 1)

nPCE2	nPCE1	MA<0>	nPIOR	nPIOW	MD[15:8]	MD[7:0]
1	0	0	0	1	Unimportant	Even Byte
1	0	1	0	1	Unimportant	Odd Byte

6.8.4 External Logic for 16-Bit PC Card Implementation

The PXA255 processor requires external glue logic to complete the 16-bit PC Card socket interface that allows either 1-socket or 2-socket solutions.

Figure 6-27 and Figure 6-28 show general solutions for a one- and two-socket configuration. The pull-ups shown are included as specified in the *PC Card Standard - Volume 2 - Electrical Specification*. Low-power systems must remove power from the pull-ups during sleep to avoid unnecessary power consumption.

GPIO or memory-mapped external registers can be used to control the reset of the 16-bit PC Card interface, power selection (V_{CC} and V_{PP}), and drive enables. The INPACK# signal is not used.

Figure 6-27 and Figure 6-28 provide the logical connections necessary to support hot insertion capability. For dual-voltage support, level shifting buffers are required for all PXA255 processor input signals. Hot insertion capability requires that each socket be electrically isolated from the other and from the remainder of the memory system. If one or both of these features is not required, then some of the logic shown in the following diagrams can be eliminated.

Software is responsible for setting the MECR[NOS] and MECR[CIT] bits. NOS indicates the number of sockets that the system support while CIT is written when the Card is in place. Input pins nPWAIT and nIOIS16 are three stated until card detect (CD) signal is asserted. To achieve this, software programs the MECR[CIT] bit when a card is detected. If the MECR[CIT] is 0, the nPWAIT and nIOIS16 inputs are ignored.

Figure 6-27 shows the minimal glue logic needed for a 1-socket system, including: data transceivers, address buffers, and level shifting buffers. The transceivers are enabled by the PSKTSEL signal. The DIR pin of the transceiver is driven by the RD/nWR pin. A GPIO is used for the three-state signal of the address and nPWE lines. These signals must be three-stated because they are used for memories other than the card interface. The Card Detect[1:0] signals are driven by the signal device.



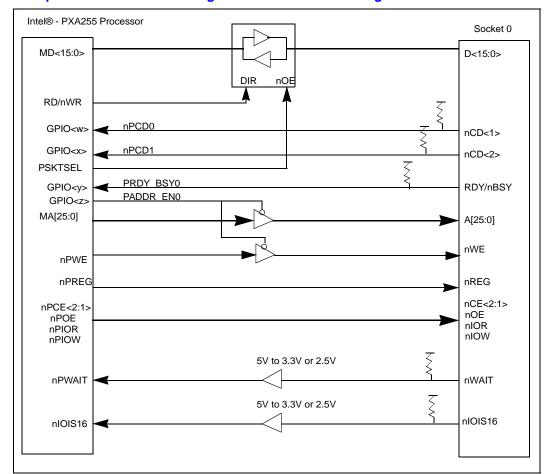


Figure 6-27. Expansion Card External Logic for a One-Socket Configuration

Figure 6-28 shows the glue logic need for a 2-socket system. RDY/nBSY signals are routed through a buffer to two separate GPIO pins. In the data bus transceiver control logic, nPCE1 controls the enable for the low byte lane and nPCE2 controls the enable for the high byte lane.\



PXA255 **Processor** Socket 0 D(15:0) D(15:0) DIR OE# nPCEx Socket 1 D(15:0) DIR OE# nPOE nPIOR nPCEx CD1# GPIO(w) CD2# CD1# GPIO(x) CD2# GPIO(y) RDY/BSY# GPIO(z) RDY/BSY# PSKTSEL MA(25:0) nPREG A(25:0) REG# A(25:0) nPCE(1:2) CE(1:2)# OE# WE# IOR# IOW# REG# nPOE, nPWE CE(1:2)# OE# WE# nPIOW, nPIOR ₹ IOR# IOW# WAIT# ₹ nPWAIT WAIT# ₹ IOIS1616# ₹ nPIOIS16 IOIS1616#

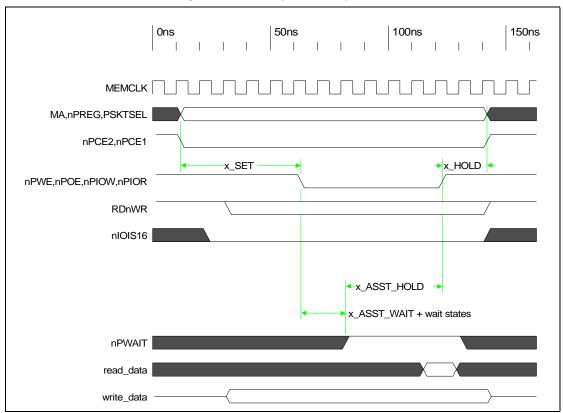
Figure 6-28. Expansion Card External Logic for a Two-Socket Configuration



6.8.5 Expansion Card Interface Timing Diagrams and Parameters

Figure 6-29 shows a 16-bit access to a 16-bit memory or I/O device. When common memory is accessed, the MCMEM0 and MCMEM1 registers are used, depending on whether card socket 0 or 1 is addressed. MCIO0 and MCIO1 are used for I/O accesses and MCATT0 and MCATT1 are used for access to attribute memory.

Figure 6-29. 16-Bit PC Card Memory or I/O 16-Bit (Half-word) Access





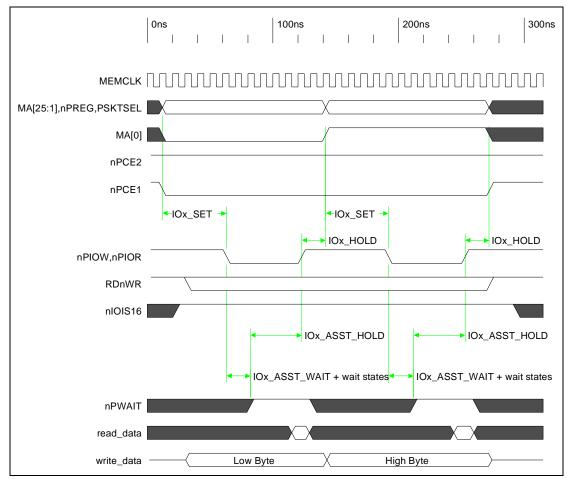


Figure 6-30. 16-Bit PC Card I/O 16-Bit Access to 8-Bit Device

The interface waits the smallest possible amount of time (x_ASST_WAIT) before it checks the value of the nPWAIT signal. If the nPWAIT signal is asserted (active low), the interface continues to wait (for a variable number of wait states) until nPWAIT is deasserted. When the nPWAIT signal is deasserted, the command continues to be asserted for a fixed amount of time (x_ASST_HOLD).

6.9 Companion Chip Interface

The processor can be connected to a companion chip in two different ways:

- Alternate Bus Master Mode
- Variable Latency I/O (See Section 6.7.5)

The connection methods are illustrated in Figure 6-31 and Figure 6-32.



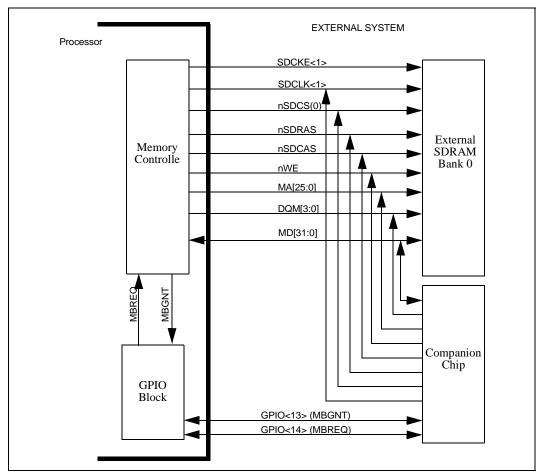
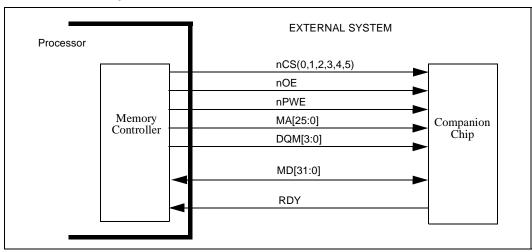


Figure 6-31. Alternate Bus Master Mode

Figure 6-32. Variable Latency IO





6.9.1 Alternate Bus Master Mode

The processor supports the presence of an alternate master on the SDRAM memory bus. The alternate master is given control of the bus with a hardware handshake that is performed through MBREQ and MBGNT, which are invoked through the alternate functions on GPIO[14] and GPIO[13], respectively. The Memory Controller performs an SDRAM refresh if SDRAM clocks and clock enable are turned on. When the alternate master must take control of the memory bus, it asserts MBREQ. It then deasserts SDCKE<1> and three-states all memory bus pins used with SDRAM bank 0 (nSDCS<0>, MA[25:0], nOE, nWE, nSDRAS, nSDCAS, SDCLK<1>, MD[31:0], DQM[3:0]). All other memory and 16-bit PC Card pins remain driven. RD/nWR remain low. Then the processor asserts MBGNT, the alternate master starts to drive all pins including SDCLK<1>, and the processor reasserts SDCKE<1>.

The grant sequence and timing follow:

- 1. The Alternate master asserts MBREO.
- 2. The Memory Controller performs an SDRAM refresh if SDRAM clocks and clock enable are turned on
- 3. If the MDCNFG:SA1111x bit is enabled, the Memory Controller sends the SDRAMs an MRS command to change the SDRAM burst length to one. The burst length is changed to one for SA-1111 compatibility.
- 4. The processor deasserts SDCKE<1> at time (t).
- 5. The processor three-states SDRAM outputs at time (t + 1 MEMCLK).
- 6. The processor asserts MBGNT at time (t + 2 MEMCLKS).
- 7. The Alternate master drives SDRAM outputs before time (t + 3 MEMCLKS).
- 8. The processor asserts SDCKE<1> at time (t + 4 MEMCLKS).

During the three-state period, both MBREQ and MBGNT remain high and an external device can take control of the three-stated pins. The external device must drive all the three-stated pins. Floating inputs can cause excessive crossover current and erroneous SDRAM commands. During the three-state period, the processor can not perform SDRAM refresh cycles.

The alternate master must assume the responsibility for SDRAM integrity during the three-state period. The system must be designed to ensure that the period of alternate mastership is limited to less than the refresh period or that the alternate master implements a refresh counter to perform refreshes at the proper intervals.

To surrender the bus, the alternate master deasserts MBREQ. The processor deasserts SDCKE<1> and MBGNT. The alternate master stops driving the SDRAM pins. The processor drives all SDRAM pins and then re-asserts SDCKE<1>.

The release sequence and timing follows:

- 1. The Alternate master deasserts MBREQ.
- 2. The processor deasserts SDCKE<1> at time (t).
- 3. The processor deasserts MBGNT at time (t + 1 MEMCLK).
- 4. The Alternate master three-states SDRAM outputs prior to time (t + 2 MEMCLKS).
- 5. The processor drives SDRAM outputs at time (t + 3 MEMCLKS).
- 6. The processor asserts SDCKE<1> at time (t + 4 MEMCLKS).



- 7. The Memory Controller performs an SDRAM refresh if SDRAM clocks and clock enable are turned on.
- 8. The Memory Controller sends an MRS command to the SDRAMs if the MDCNFG:SA1111x bit is enabled. This changes the SDRAM burst length back to four.

If the refresh counter for the processor requested a refresh cycle during the alternate master's tenure, a refresh cycle runs first, followed by any other bus transactions that stalled during that period.

To enable alternate bus master, the set up the signals by writing the following registers:

- Write the GPIO Pin Direction register GPDR0 to set bit 13 (make GPIO[13] an output) and clear bit 14 (make GPIO[14] an input)
- Write the GPIO Alternate Function register GAFR0_L to set bits 27 and 26 to 0b11 (enable the MBGNT alternate function 3) and set bits 29 and 28 to 0b01 (enable the MBREQ alternate function 1).

6.9.1.1 **GPIO** Reset

During GPIO reset, the GPIOs, including MBREQ and MBGNT, are set to their reset state. The MBREQ and MBGNT pins become general purpose inputs. The system must have external putdowns on these pins to prevent the pins from floating.

If a transaction is in progress when GPIO reset is asserted, the alternate master loses ownership of the bus. The alternate master must immediately give up the bus when MBGNT is deasserted. Because the memory controller is not reset, an SDRAM refresh can occur immediately after the GPIO reset assertion.

6.9.1.2 nVDD_FAULT/nBATT_FAULT with PMCR[IDAE] Disabled

If an nVDD_FAULT or nBATT_FAULT occurs, the processor places the GPIOs into their sleep states. MBGNT must be programmed to go low during sleep.

The memory controller prevents the processor from entering sleep until all outstanding transactions have completed. This includes waiting for the MBREQ signal from the alternate master to deassert. For best sleep performance, the alternate master must immediately give up the bus when MBGNT is deasserted. If necessary, the alternate master can hold the bus until its transaction is completed. After the memory controller has completed all outstanding transactions, it places SDRAM into self-refresh and allows the processor to complete the sleep entry sequence.

Note: The alternate bus master must de-assert MBREQ when nVDD_FAULT or nBATT_FAULT is asserted.

6.9.1.3 nVDD_FAULT/nBATT_FAULT with PMCR[IDAE] Enabled

If an nVDD_FAULT or nBATT_FAULT occurs with PMCR[IDAE] enabled, the processor causes an Imprecise Data Abort Exception. This allows the processor to do any required actions before sleep entry. Sleep entry with PMCR[IDAE] enabled is similar to normal sleep entry. The processor places the GPIOs into their sleep states. MBGNT must be programmed to go low during sleep.

The memory controller prevents the processor from entering sleep until all outstanding transactions have completed. This includes waiting for the MBREQ signal from the alternate master to deassert. For best sleep performance, the alternate master must immediately give up the bus when MBGNT



is deasserted or, as part of the sleep entry routine, the alternate master can be disabled. If necessary, the alternate master can hold the bus until its transaction is completed. After the memory controller has completed all outstanding transactions, it places SDRAM into self-refresh and allows the processor to complete the sleep entry sequence.

Note: The alternate bus master must de-assert MBREQ when nVDD_FAULT or nBATT_FAULT is asserted.

6.10 Options and Settings for Boot Memory

This section explains the settings that control Boot Memory configurations.

6.10.1 Alternate Booting

The processor allows six boot configurations. These configurations are determined by the three BOOT_SEL(2:0) pins and are described in Table 6-36. A description of the effect of these input pins on the Configuration registers at boot time is included in Section 6.8.

Table 6-36. BOOT SEL Definitions

	BOOT_SEL		Doct From
2	1	0	Boot From
0	0	0	Asynchronous 32-bit ROM
0	0	1	Asynchronous 16-bit ROM
0	1	0	reserved
0	1	1	reserved
1	0	0	1- 32-bit Synchronous Mask ROM (64 Mbit) 2- 16-bit Synchronous Mask ROMs = 32 bits (32 Mbit each)
1	0	1	1- 16-bit Synchronous Mask ROM (64 Mbit)
1	1	0	2- 16-bit Synchronous Mask ROMs = 32 bits (64 Mbit each)
1	1	1	1- 16-bit Synchronous Mask ROM (32 Mbit)

6.10.2 Boot Time Defaults

The following sections provide information on boot time default parameters.

6.10.2.1 BOOT_DEF Read-Only Register (BOOT_DEF)

BOOT_DEF, shown in Table 6-37, contains the boot-up values for the three BOOT_SEL pins and the single package-type bit.

This is a read-only register. Ignore reads from reserved bits.



Table 6-37. BOOT_DEF Bitmap

				0x4	1800	0_00	044							В	тос	_DI	ĒF							Me	emo	ry (Con	troll	ler			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													r	ese	rve	t													PKG_TYPE		BOOT_SEL	
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	*	*	*	*
		В	its			Na	me												D	esci	riptio	on										
		31	1:4			_			rese	rve	d																					
									PRO	OCE	SS	OR	TYI	PE (read	d on	ly):															
		;	3		Pł	(G_	TYF	PΕ	0 – 1 –					esso	or																	
									ВОС				,		•	,																
		2	:0		В	TOC	_SE	ΞL	Con											-	-		•									
									See Time						id bo	oot o	conf	igur	atio	ns. S	See	Sec	tion	6.1	0.2.	.2 fc	r de	scri	ptio	ns o	f Bo	oot

Table 6-38. Valid Boot Configurations Based on Processor Type

Processor Type	Boot_Sel Signals Valid Booting Configurations
	000
	001
(PXA255 processor)	100
(FAA255 processor)	101
	110
	111



6.10.2.2 Boot-Time Configurations

The boot time configurations are shown in Figure 6-33 - Figure 6-35. A boot from a single 32-Mbit SMROM with nWORD = 1 is not supported.

Three Configuration registers are affected at reset - MSC0:RBW0, MDREFR:E0PIN/K0RUN, and SXCNFG.

Figure 6-33. Asynchronous Boot Time Configurations and Register Defaults

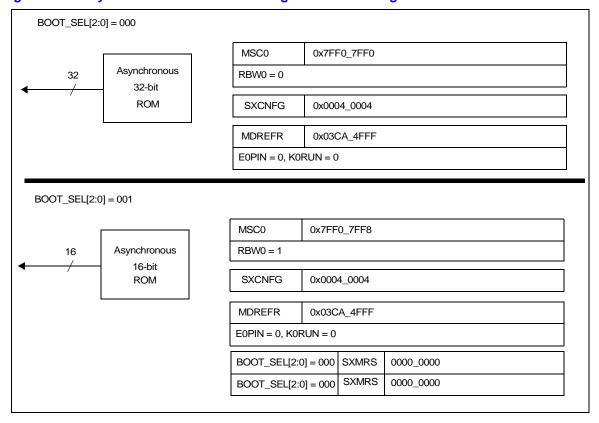




Figure 6-34. SMROM Boot Time Configurations and Register Defaults

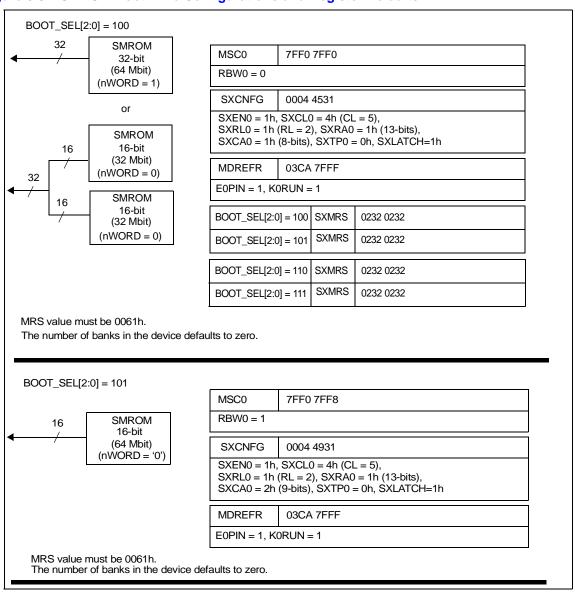
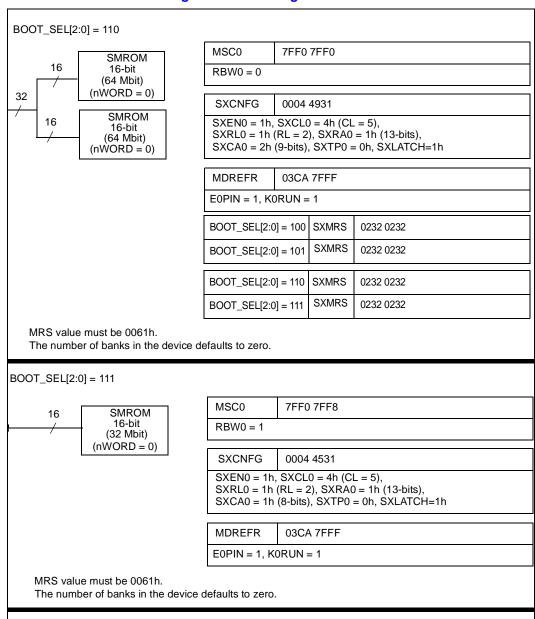




Figure 6-35. SMROM Boot Time Configurations and Register Defaults



6.10.3 Memory Interface Reset and Initialization

On reset, the SDRAM Interface is disabled. Reset values for the Boot ROM are determined by BOOT_SEL. The memory pins and controller are in the state shown in Table 6-39.



Table 6-39. Memory Controller Pin Reset Values

Pin Name	PXA255 Processor Reset Value
SDCLK [2:0]	000
SDCKE <1>	00
SDCKE <0>	1 if BOOT_SEL = Synchronous Memory
DQM [3:0]	0000
nSDCS [3:0]	1111
nWE	1
nSDRAS	1
nSDCAS	1
nOE	1
MA [25:0]	0x000000h
RDnWR	0
MD [31:0]	0x0000000h
nCS <0>	1
nCS <5:1>	GPIO Input
nPIOIR	GPIO Input
nPIOIW	GPIO Input
nPOE	GPIO Input
nPWE	GPIO Input

In sleep mode, the memory pins and controller are in the same state as they are after a hardware reset, except that the GPIO signals are driven high. If SDRAMs are in self-refresh, they are held there by setting SDCKE<1> to a 0.

6.11 Hardware, Watchdog, or Sleep Reset Operation

Software performs the following procedures when the processor comes out of a reset:

- 1. After hardware reset, complete a power-on wait period of 200 µs, which allows the internal clocks that generate SDCLK to stabilize. Enable MDREFR:K0RUN and E0PIN for Synchronous Static memory. When MDREFR is written, a refresh interval value (MDREFR:DRI) must also be written. The following writes are allowed:
 - a. Write MSC0, MSC1, MSC2
 - b. Write MECR, MCMEM0, MCMEM1, MCATT0, MCATT1, MCIO0, MCIO1
 - c. Write MDREFR:K0RUN and MDREFR:E0PIN. Configure MDREFR:K0DB2. Retain the current values of MDREFR:APD and MDREFR:SLFRSH. MDREFR:DRI must contain a valid value. Deassert MDREFR:KxFREE.
- 2. In systems that contain Synchronous Static memory, write to the SXCNFG to configure all appropriate bits, including the enable bits. Software must perform a sequence that involves a subsequent write to SXCNFG to change the RAS latencies. While any SMROM banks are



being configured, the SDRAM banks must be disabled and MDREFR:APD must be deasserted (auto-power-down disabled).

- a. Write SXCNFG (with enable bits asserted).
- b. Write to SXMRS to trigger an MRS command to all enabled banks of synchronous static memory.
- c. SXLCR must only be written when it is required by the SDRAM-like synchronous flash device for command encoding.
- 3. In systems that contain SDRAM, transition the SDRAM controller through the following state sequence:
 - a. self-refresh and clock-stop
 - b. self-refresh
 - c. power-down
 - d. PWRDNX
 - e. NOP
- 4. The SDRAM clock run and enable bits (MDREFR:K1RUN, K2RUN, and E1PIN), described in Section 6.5.3. MDREFR:SLFRSH must not be asserted.
 - a. Write MDREFR:K1RUN, K2RUN (self-refresh and clock-stop -> self-refresh). Configure MDREFR:K1DB2,K2DB2.
 - b. Write MDREFR:SLFRSH (self-refresh -> power-down).
 - c. Write MDREFR:E1PIN (power-down -> PWRDNX).
 - d. a write is not required for this state transition (PWRDNX -> NOP).
 - e. Configure, but do not enable, each SDRAM partition pair.
 - f. Write MDCNFG (with enable bits deasserted), MDCNFG:DE3:2,1:0 set to '0'.
- 5. For systems that contain SDRAM, wait a specified NOP power-up waiting period required by the SDRAMs to ensure the SDRAMs receive a stable clock with a NOP condition
- 6. Ensure the Data Cache bit (DCACHE) is disabled. If this bit is enabled, the refreshes triggered by the next step may not pass through to the Memory Controller properly.
- 7. On a hardware reset in systems that contain SDRAM, trigger the specified number (typically eight) of refresh cycles by attempting non-burst read or write accesses to any disabled SDRAM bank. Each such access causes a simultaneous CBR refresh cycles for all four banks, which causes a pass through the CBR state and back to NOP. On the first pass, the PALL state occurs before the CBR state.
- 8. Re-enable the DCACHE bit if it is disabled.
- 9. In systems that contain SDRAM, enable SDRAM partitions by setting MDCNFG:DE3:2,DE1:0.
- 10. In systems containing SDRAM, write the MDMRS register to trigger an MRS command to all enabled banks of SDRAM. For each SDRAM partition pair that has one or both partitions enabled, this forces a pass through the MRS state and back to NOP. The CAS latency must be the only variable option and is derived from the value programmed in the MDCNFG:MDTC0,2 fields. The burst type is programmed to sequential and the length is set to four.



11. Optionally, in systems that contain SDRAM or Synchronous Static memory, enable auto-power-down by setting MDREFR[APD].

6.12 GPIO Reset Procedure

On a GPIO Reset, the Memory Controller registers keep the values they had before the reset. No new configuration programming is required. However, SDRAM refreshes do not occur during the reset time. After nRESET_OUT is deasserted, the memory controller will continue refreshing. By ensuring a refresh time for SDRAM that is smaller than the default, it is possible to preserve the SDRAM contents. To do this, follow this procedure:

The SDRAM refresh time is chosen by taking the specified refresh time, typically 64 ms, and subtracting the GPIO Reset time (found in the *Intel*® *PXA255 Applications Processors Electrical, Mechanical, and Thermal Specification*). For example, the GPIO Reset time is ~360 microseconds, leaving an SDRAM refresh time of (64 ms - .360 ms) = 63.64 ms. Use this time to program the MDREFR[DRI].

In the boot code, determine the type of reset. If the reset was a GPIO reset, then refresh all the SDRAM rows. Refreshing all the SDRAM rows preserves their value in case GPIO reset occurs again.

After all the SDRAM rows have been refreshed, enable GPIO reset

6.13 Memory Controller Register Summary

Table 6-40 shows the registers associated with the memory interface and the physical addresses used to access them. These registers must be mapped as non-cacheable and non-bufferable and can only be a single word access. They are grouped together in one page and all have the same memory protections.

Table 6-40. Memory Controller Register Summary (Sheet 1 of 2)

Physical Address	Symbol	Register Name
0x4800_0000	MDCNFG	SDRAM Configuration Register
0x4800_0004	MDREFR	SDRAM Refresh Control Register
0x4800_0008	MSC0	Static Memory Control Register 0
0x4800_000C	MSC1	Static Memory Control Register 1
0x4800_0010	MSC2	Static Memory Control Register 2
0x4800_0014	MECR	Expansion Memory (16-bit PC Card / Compact Flash) Bus Configuration register
0x4800_001C	SXCNFG	Synchronous Static Memory Control Register
0x4800_0024	SXMRS	MRS value to be written to SMROM
0x4800_0028	МСМЕМО	Card interface Common Memory Space Socket 0 Timing Configuration
0x4800_002C	MCMEM1	Card interface Common Memory Space Socket 1 Timing Configuration
0x4800_0030	MCATT0	Card interface Attribute Space Socket 0 Timing Configuration
0x4800_0034	MCATT1	Card interface Attribute Space Socket 1 Timing Configuration
0x4800_0038	MCIO0	Card interface I/O Space Socket 0 Timing Configuration



Table 6-40. Memory Controller Register Summary (Sheet 2 of 2)

Physical Address	Symbol	Register Name
0x4800_003C	MCIO1	Card interface I/O Space Socket 1 Timing Configuration
0x4800_0040	MDMRS	MRS value to be written to SDRAM
0x4800_0044	BOOT_DEF	Read-Only Boot-time register. Contains BOOT_SEL and PKG_SEL values.
0x4800 0058	MDMRSLP	Low-Power SDRAM Mode Register Set Configuration Register



LCD Controller

The LCD controller provides an interface from the PXA255 processor to a passive (DSTN) or active (TFT) flat panel display. Monochrome and several color pixel formats are supported.

7.1 Overview

The processor LCD controller supports single- or dual-panel displays. Encoded pixel data created by the core is stored in external memory in a frame buffer in 1, 2, 4, 8, or 16-bit increments. The data is fetched from external memory and loaded into a first-in first-out (FIFO) buffer on a demand basis, using the LCD controller's dedicated dual-channel DMA controller (DMAC). One channel is used for single-panel displays and two are used for dual-panel displays.

Frame buffer data contains encoded pixel values that are used by the LCD controller as pointers to index a 256-entry x 16-bit-wide palette. For 16 bit per pixel frame buffer entries, the palette RAM is bypassed. Monochrome palette entries are eight bits wide, and color palette entries are 16 bits wide. The encoded pixel data determines the number of possible colors within the palette as follows:

- 1-bit-wide pixels address the top 2 locations of the palette
- 2-bit-wide pixels address the top 4 locations of the palette
- 4-bit-wide pixels address the top 16 locations of the palette
- 8-bit-wide pixels address any of the 256 entries within the palette
- 16-bit-wide pixels bypass the palette

When passive color 16-bit pixel mode is enabled, the color pixel values bypass the palette and are fed directly to the LCD controller's Frame Rate Control logic. When active color 16-bit pixel mode is enabled, the pixel value bypasses the palette and the Frame Rate Control logic and is sent directly to the LCD controller's data pins. Optionally, the palette RAM is loaded for each frame by the LCD controller's DMAC.

Once the encoded pixel value is used to select a palette entry, the value programmed within the entry is transferred to the Frame Rate Control logic, which uses the Temporal Modulated Energy Distribution (TMED) algorithm to produce the pixel data that is sent to the screen. Frame Rate Control is a technique used to create additional color shades by rapidly turning on and off a pixel on the LCD screen. This is also known as temporal dithering. The data output from the dither logic is grouped into the selected format (e.g., 8-bit color, dual panel, 16-bit color., etc.) and placed in a FIFO buffer before being sent out on the LCD controller's pins and driven to the display using the pixel clock.

Depending on the type of panel used, the LCD controller is programmed to use either 4-, 8-, or 16-pixel data output pins. Single-panel monochrome displays use either four or eight data pins to send 4 or 8 pixels for each pixel clock. Single-panel color displays use eight pins to send 2-2/3 pixels each pixel clock (8 pins / 3 colors/pixel = 2 2/3 pixels per clock). The LCD controller also supports dual-panel mode, in which the LCD controller's data lines are split into two groups, one to drive the top half and one to drive the bottom half of the screen. For dual-panel displays, the number of pixel data output pins is doubled, allowing twice as many pixels to be sent each pixel clock to the two halves of the screen.



In active color display mode, the LCD controller can drive TFT displays. When using 1-, 2-, 4-, or 8-bit modes, the LCD's dither logic is bypassed, and the pixel value is sent from the palette buffer directly to the LCD's data output pins. 16-bit pixel mode bypasses both the palette and the dither logic.

7.1.1 Features

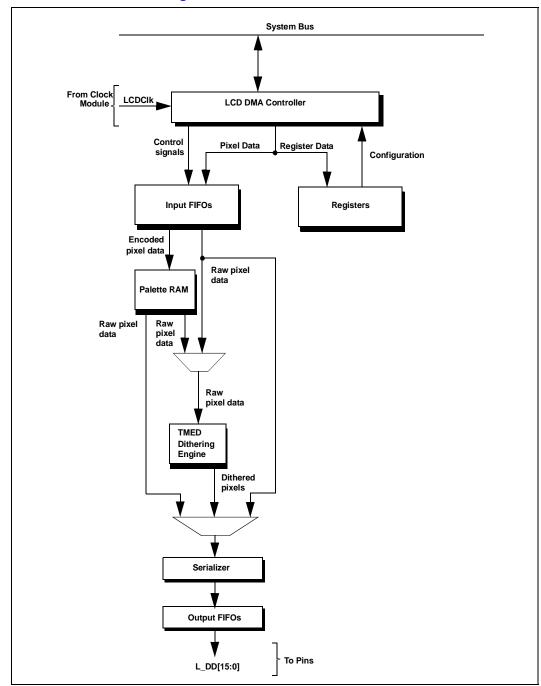
The processor LCD controller supports the following features:

- Display modes:
 - single- or dual-panel displays
 - up to 256 gray-scale levels (8 bits) in Passive Monochrome Mode
 - a total of 65536 possible colors in Passive Color Mode (using the 16-bit TMED dithering algorithm)
 - up to 65536 colors in Active Color Mode (16 bits, bypasses palette)
 - passive 8-bit color single-panel displays
 - passive 8-bit (per panel) color dual-panel displays
- Display sizes up to 1024x1024 pixels, recommended maximum of 640x480
- Internal color palette RAM 256 entry by 16 bits (can be loaded automatically at the beginning of each frame)
- Encoded pixel data of 1, 2, 4, 8, or 16 bits
- Programmable toggle of AC bias pin output (toggled by line count)
- Programmable pixel clock from 195 kHz to 83 MHz (100 MHz/512 to 166 MHz/2)
- Integrated 2-channel DMA (one channel for palette and single panel, the other channel for second panel in dual-panel mode).
- Programmable wait-state insertion at the beginning and end of each line
- Programmable polarity for output enable, frame clock, and line clock
- Programmable interrupts for input and output FIFO underrun
- Programmable frame and line clock polarity, pulse width, and wait counts



Figure 7-1 illustrates a simplified, top-level block diagram for the processor LCD Controller.

Figure 7-1. LCD Controller Block Diagram





7.1.2 Pin Descriptions

When the LCD controller is enabled, all of the LCD pins are outputs only. When the LCD controller is disabled, its pins can be use for general-purpose input/output (GPIO). Refer to the System Integration Unit chapter for details.

Table 7-1 describes the LCD controller's pins. For more detailed information, see Section 7.3.5. All of the LCD pins are outputs only.

Table 7-1. Pin Descriptions

Pin	Definition
L_DD[7:0]	These data lines transmit either four or eight data values at a time to the LCD display. For monochrome displays, each pin value represents a single pixel. For passive color, groupings of three pin values represent one pixel (red, green, and blue subpixel data values). In single-panel monochrome mode, L_DD<3:0> pins are used. For double-pixel data, single-panel monochrome, dual-panel monochrome, single-panel color, and active color modes, L_DD[7:0] are used.
L_DD[15:8]	When dual-panel color or TFT (active color mode) operation is programmed, these data outputs are also required to send pixel data to the screen.
L_PCLK	The Pixel Clock is used by the LCD display to clock the pixel data into the line shift register. In passive mode, the pixel clock toggles only when valid data is available on the data pins. In active mode, the pixel clock toggles continuously, and L_BIAS serves as an output to signal when data is valid on the LCD's data pins.
L_LCLK	The Line Clock is used by the LCD display to signal the end of a line of pixels. The display transfers the line data from the shift register to the screen and increments the line pointer. In active mode, it is the horizontal synchronization signal.
L_FCLK	The Frame Clock is used by the LCD display to signal the start of a new frame of pixels. The display resets the line pointer to the top of the screen. In active mode, it is the vertical synchronization signal.
L_BIAS	AC Bias is used to signal the LCD display to switch the polarity of the power supplies to the row and column drivers of the screen to counteract DC offset. In active mode, it serves as the output enable to signal when data is latched from the data pins using the Pixel Clock.

7.2 LCD Controller Operation

7.2.1 Enabling the Controller

If the LCD controller is being enabled for the first time after system reset or sleep reset, all of the LCD registers must be programmed as follows:

- 1. Configure the General Purpose I/O (GPIO) pins for LCD controller functionality. See Chapter 4, "System Integration Unit" for details.
- 2. Write the frame descriptors and, if needed, the palette descriptor to memory.
- 3. Program all of the LCD configuration registers *except* the Frame Descriptor Address Registers (FDADRx) and the LCD Controller Configuration Register 0 (LCCR0). See Section 7.6 for details of all registers.
- 4. Program FDADRx with the memory address of the palette/frame descriptor, as described in Section 7.6.5.2.
- 5. Enable the LCD controller by writing to LCCR0, as described in Section 7.6.1.



If the LCD controller is being re-enabled, there has *not* been a reset since the last programming, and the GPIO pins are still configured for LCD Controller functionality, only the registers FDADRx and LCCR0 need to be reprogrammed. The LCD Controller Status Register (LCSR) must also be written to clear any old status flags before re-enabling the LCD controller. See Section 7.6.7 for details.

7.2.2 Disabling the Controller

The LCD controller can be disabled in two ways: regular and quick.

Regular disabling, the recommended method for stopping the LCD controller, is accomplished by setting the disable bit, LCCR0[DIS]. The other bits in LCCR0 must not be changed — read the register, set the DIS bit, and rewrite the register. This method causes the LCD controller to stop cleanly at the end of the frame currently being fetched from memory. If the LCD DMAC is fetching palette data when DIS is set, the palette RAM load is completed, and the next frame is displayed before the LCD is disabled. The LCD Disable Done bit, LCSR[LDD], is set when the LCD controller finishes displaying the last frame fetched, and the enable bit, LCCR0[ENB], is cleared automatically by hardware.

Quick disabling is accomplished by clearing the enable bit, LCCR0[ENB]. The LCD controller will finish any current DMA transfer, stop driving the panel, and shut down immediately, setting the quick-disable bit, LCSR[QD]. This method is intended for situations such as a battery fault, where system bus traffic has to be minimized immediately so the processor can have enough time to store critical data to memory before the loss of power. The LCD controller must not be reenabled until the QD bit is set, indicating that the quick shutdown is complete.

Once disabled, the LCD Controller automatically disables its clocks to conserve power.

7.2.3 Resetting the Controller

At reset, the LCD Controller is disabled, and the output pins are configured as GPIO pins. All LCD Controller Registers are reset to the conditions shown in the register descriptions.

7.3 Detailed Module Descriptions

This section describes the functions of the modules in the LCD Controller:

7.3.1 Input FIFOs

Data fetched from external memory by the dedicated DMAC is placed in one of two input FIFO buffers. Each input FIFO comprises 128 bytes, organized as 16 entries by 8 bytes. In single-panel mode, one FIFO is used to queue both encoded pixel data and data for writing to the internal palette RAM. In dual-panel mode, this FIFO queues data for the internal palette RAM and the upper half of the LCD display, while the second FIFO buffer holds data for the lower half of the LCD display.

The FIFO signals a service request to the DMAC whenever four FIFO entries are empty. In turn, the DMAC automatically fills the FIFO with a 32-byte burst. Pixel data from the frame buffer remains packed within individual 8-byte entries when it is loaded into the FIFO. If the pixel size is



1, 2, 4, or 8-bits, the FIFO entries are unpacked and used to index the palette RAM to read the color value. In 16-bit passive mode, the entries bypass the palette and go directly to the TMED dither logic. In 16-bit active mode, the pixels are sent directly to the pins.

7.3.2 Lookup Palette

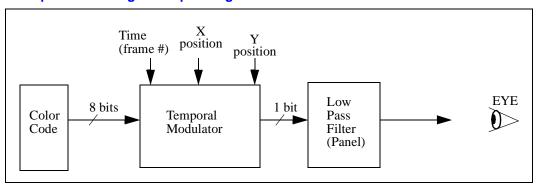
The internal palette RAM holds up to 256 16-bit color values. Color palette RAM entries are 16 bits wide, with 5 bits of red, 6 bits of green, and 5 bits of blue. Monochrome entries are 8 bits wide. Encoded pixel values from the input FIFO are used as an address to index and select individual palette locations. 1-bit pixel encodings address the first 2 entries, 2-bit pixel encodings address the first 4 entries, 4-bit pixel encodings address 16 locations, and 8-bit pixel encodings select any of the 256 palette entries. In 16-bit pixel mode, the palette RAM is not used and must not be loaded.

7.3.3 Temporal Modulated Energy Distribution (TMED) Dithering

For passive displays, entries selected from the lookup palette (or directly from memory for 16-bit pixels) are sent to the TMED dithering algorithm. TMED is a form of temporal dithering, also know as frame rate control. The algorithm determines whether a pixel is on or off.

Understanding how the TMED dithering algorithm works is not necessary to use the processor LCD controller. However, certain characteristics of the algorithm can be controlled through the use of the TMEDRGB Seed Register (Table 7-14) and the TMED Control Register (TCR, Table 7-15). If these registers are to be modified from their default values, refer to this section. Figure 7-2 illustrates the TMED concept.

Figure 7-2. Temporal Dithering Concept - Single Color



This dithering concept is applied separately to each color displayed. Each color has zeros added to make the data for each color 8 bits. If a monochrome display is used, only a single matrix (blue) is used.

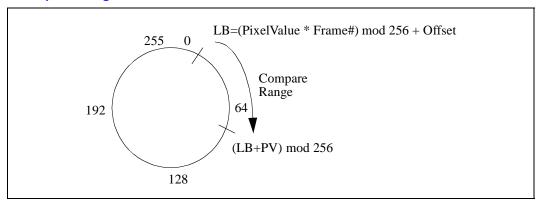
The processor LCD Controller implements the following algorithm, which is used by TMED to determine an upper and lower boundary:

```
LowerBoundary = [(PixelValue * FrameNumber) mod 256] + Offset UpperBoundary = [(PixelValue + LowerBoundary) mod 256]
```

A 16x16 matrix uses the row (line), column (pixel number), and frame number (which wraps back to 0 from 255) to select a matrix value. When the matrix value is between the lower and upper boundaries from the algorithm, the LCD controller sends a "1" to the LCD panel. The boundaries created by the algorithm are circular, wrapping from 255 back to 0, as shown in Figure 7-3.



Figure 7-3. Compare Range for TMED



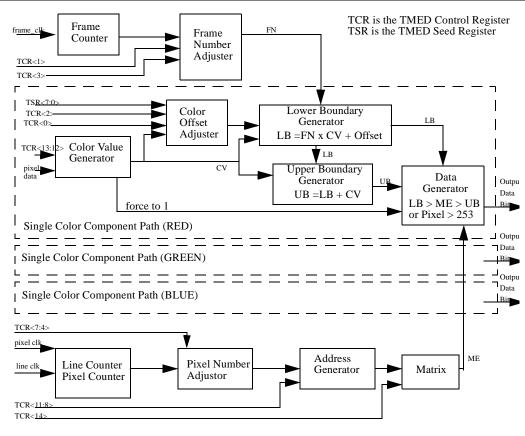
Either of two matrices may be used for each color, chosen by bits 0, 1, and 14 of the TMED Control Register (TCR, Section 7.6.10). Offsets may be selected for the shading of each color to avoid gray color problems. Although these offset values are panel dependent, the recommended values are listed in Section 7.6.9, with the blue data path being used for monochrome modes. Offsets may also be chosen in the TMED Control Register for shifting the row (horizontal) value, line (vertical) value, and frame number.

Figure 7-4 shows the block diagram for TMED. Pixel data (up to 8 bits) enters the module and is sent through the Color Value (CV) generator. Depending on the value of the TCR[TSCS] field, the CV generator rounds off between 0 and 3 of the least-significant bits, creating a new CV. If the original pixel value is 254 or 255, the final data output is set to one. Otherwise, the following occurs:

- 1. The new CV is sent through the Color Offset Adjuster, where it is used as a lookup into the matrix selected by TCR[COAM].
- 2. Either the 8-bit output of the chosen matrix or 00h, as selected by TCR[COAE], is added to the appropriate color's seed register value in register TRGBR to form an offset.
- 3. This offset is added to the result of the multiplication of the Frame Number and the CV to form the algorithm's lower boundary (only the lower 8 bits are used).
- 4. The CV is added to the lower boundary to obtain the upper boundary.
- 5. Row (line) and column (pixel) counters are combined with beat suppression (offset) values in the Pixel Number Adjuster and Address Generator to form yet another address for a matrix lookup.
- 6. The output of the chosen matrix is compared to the lower and upper boundaries in the Data Generator.
- 7. If the matrix output is between these boundaries or the original pixel value is 254 or 255, then the data output to the panel is one. In all other cases, it is zero.



Figure 7-4. TMED Block Diagram



7.3.4 Output FIFOs

The LCD controller has two output FIFOs to queue pixel data before it is sent out to the pins. Each output FIFO is 16 bytes, organized as 16 entries by 8 bits wide. Pixel values are accumulated in a serial shifter and written to the FIFO buffers in 4-, 8-, or 16-bit quantities. Four pins are used for single-panel monochrome screens, 8 pins are used for single- and dual-panel monochrome screens and single-panel color displays, and 16 pins are used for dual-panel color and active displays. Each time a value is taken from the bottom of the FIFO, the entry is invalidated, and all data in the FIFO moves down one position.

7.3.5 LCD Controller Pin Usage

The timing of the line (L_LCLK) and frame (L_FCLK) clocks is programmable to support both passive display and active display modes. Programming options include: wait state insertion at the beginning and end of each line and frame, pixel clock (L_PCLK), line clock, frame clock (L_FCLK), output enable signal polarity, and frame clock pulse width.

See Section 7.5 for pin timing diagrams. When the LCD controller is disabled, all of its pins can be used for GPIO. See Chapter 4, "System Integration Unit" for further details. See also Table 7-1.



7.3.5.1 Passive Display Timing

In passive display mode (LCCR0[PAS] = 0), L_PCLK toggles only when data is being written to the panel. When an entire line of pixels has been sent to the display, L_LCLK is asserted. When an entire frame of pixels has been sent to the display, L_FCLK is asserted.

If an output FIFO underrun occurs (i.e., the LCD controller runs out of data), L_PCLK stalls until valid data is available. This results in a slower pixel clock, but data sent to the display is always valid.

To prevent a D.C. charge from building within a passive display, its power and ground supplies must be switched periodically. Many modern panels do this automatically. If not, the LCD controller can toggle the AC bias pin (L_BIAS) to signal the display to switch the polarity. The frequency of the L_BIAS toggle is controlled by programming the number of line clock transitions between each toggle (LCCR3[ACB]).

7.3.5.2 Active Display Timing

In active display mode (LCCR0[PAS] = 1), L_PCLK toggles continuously as long as the LCD controller is enabled. The other pins function as follows:

L_BIAS — output enable. When asserted, the LCD latches L_DD data using L_PCLK.

L_LCLK — horizontal synchronization signal (HSYNC)

L_FCLK — vertical synchronization signal (VSYNC)

If an output FIFO underrun occurs, the data on the L_DD pins is repeated, L_BIAS stays asserted, and L_PCLK keeps running. As valid data enters the output FIFO, it is sent to the display. Additional pixel clocks are inserted at the end of the line to drain the remaining valid pixels from the output FIFO before HSYNC is asserted. This mechanism allows an underrun to corrupt only a single line rather than an entire frame.

7.3.5.3 Pixel Data Pins (L_DDx)

Pixel data is removed from the bottom of the output FIFO and driven in parallel onto the LCD data lines on the edge of the pixel clock selected by Pixel Clock Polarity (LCCR3[PCP]). For a 4-bit wide bus, data goes out on the LCD data lines L_DD[3:0]. For an 8-bit wide bus, data goes out on L_DD[7:0]. For a 16-bit bus, data goes out on L_DD[15:0]. In monochrome dual-panel mode, the pixels for the upper half of the screen go out on L_DD[3:0] and those for the lower half on L_DD[7:4]. In color dual-panel mode, the upper panel pixels go out on L_DD[7:0] and the lower panel pixels on L_DD[15:8]. The LCD data pins are driven at their last value during the inactive portion of the LCD frame.

7.3.6 DMA

Values for palette RAM entries and encoded pixel data are stored in off-chip memory and are transferred to the LCD controller's input FIFO buffers, on a demand basis, using the LCD controller's dedicated DMA controller (DMAC). The LCD's descriptor-based DMAC contains two channels that transfer data from external memory to the input FIFOs. One channel is used for single-panel displays and two are used for dual-panel displays.

The LCD controller issues a service request to the DMAC after it has been initialized and enabled. The DMAC automatically performs eight word transfers, filling four entries of the input FIFO. Values are fetched from the bottom of the FIFO, one entry at a time, and each 64-bit value is



unpacked into individual pixel encodings of 1, 2, 4, 8, or 16 bits each. After the value is removed from the bottom of the FIFO, the entry is invalidated, and all data in the FIFO is shifted down one entry. When four of the entries are empty, a service request is issued to the DMAC. If the DMAC is not able to keep the FIFO filled with enough pixel data (due to insufficient external memory access speed) and the FIFO is emptied, the appropriate FIFO underrun status bit is set (bit IUL or IUU in register LCSR), and an interrupt request is issued (unless it is masked).

7.4 LCD External Palette and Frame Buffers

The LCD controller supports a variety of user-programmable options, including display type and size, frame buffer location, encoded pixel size, and output data width. Although all programmable combinations are possible, the displays available on the market dictate which combinations of these programmable options are practical. The processor external system memory limits the throughput of the LCD controller's DMAC, which, in turn, limits the size and type of display that can be controlled. The user must also determine the maximum bandwidth of the processor external bus that the LCD controller is allowed to use without negatively affecting all other functions that the processor must perform.

7.4.1 External Palette Buffer

The external palette buffer is an off-chip memory area containing up to 256 16-bit entries to be loaded into the internal palette RAM. The palette buffer data does not have to be at the beginning of the external frame buffer, it can also be in a separate memory location. Palette data is 8 bytes (4 entries) for 1 and 2-bit pixels (the last 2 entries are loaded but not used with 1-bit pixels), 32 bytes (16 entries) for 4-bit pixels, and 512 bytes (256 entries) for 8-bit pixels. The palette RAM is not used and must not be loaded when using 16 bits per pixel.

After enabling the LCD controller, the user must first load the palette RAM before processing any frame data. After the initial load, the palette can be reloaded optionally on a frame-by-frame basis. This would be done when the color selection changes frame to frame. The palette RAM is always loaded with DMA channel 0.

Figure 7-5 shows the format of the palette entries in little endian. "Endian" does *not* imply endianness with respect to bytes and half-words within memory. It refers strictly to the ordering of the palette entries; i.e., whether palette entry 0 is at the MSB or the LSB of a word boundary. The ordering of RGB values within the 16-bit entry is fixed for little endian. In Figure 7-5, "Base" is the palette buffer base programmed in register FSADR.

Palette entry 16

Palette entry 254



Figure 7-5. Palette Buffer Format

							Indiv	idual P	alet	te E	ntry						
Bit	15	14	13	12	11	10	9	8	7	•	6	5	4	3	2	1	0
Color		F	Red (R)				Gree	n (G	3)	3)						
Bit	15	14	13	12	11	10	9	8	7	,	6	5	4	3	2	1	0
Mono			unused Monochrome (M)														
,																	
			Little Endian Palette Entry Ordering														
								or 256			•		Ū				
Bit		31						1	6	15							0
Base +	0x0			Pa	alette (entry 1							Palet	te entr	y 0		
Base +	0x4			Pa	alette (entry 3							Palet	te entr	y 2		
			ies 4 th pixel.	rough	255 d	o not e	xist fo	r 1 and	2								,
Base +	0x1C			Pa	lette e	ntry 15	5						Palett	e entry	/ 14		

Base + 0x20 Palette entry 17

Entries 16 through 255 do not exist for 1, 2, and 4 bits/pixel.

Palette entry 255

Base + 0x1FC

7.4.2 External Frame Buffer

The external frame buffer is an off-chip memory area used to supply enough encoded pixel values to fill the entire screen one or more times. The number of pixel data values depends on the size of the screen (for example, $640 \times 480 = 307,200$ encoded pixel values). Figure 7-6 through Figure 7-18 show the memory organization within the frame buffer for each size pixel encoding.

In the following figures, "Base" refers to the initial address programmed in the FSADR register, "Palette Buffer Index" refers to the data that specifies the location in the palette buffer, and "Raw Pixel Data" refers to the actual 16-bit RGB data when the palette RAM is bypassed.

Figure 7-6. 1 Bit Per Pixel Data Memory Organization

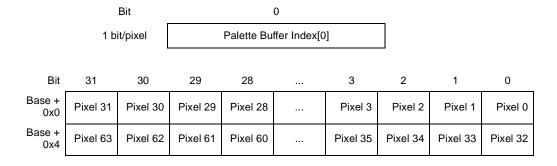




Figure 7-7. 2 Bits Per Pixel Data Memory Organization

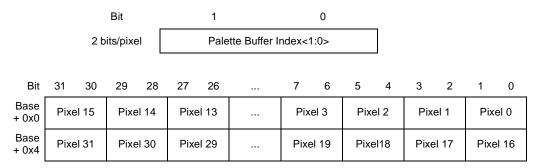


Figure 7-8. 4 Bits Per Pixel Data Memory Organization

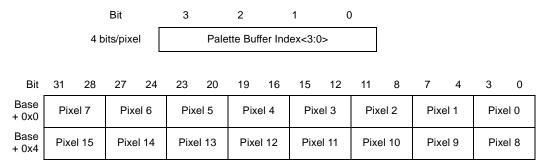


Figure 7-9. 8 Bits Per Pixel Data Memory Organization

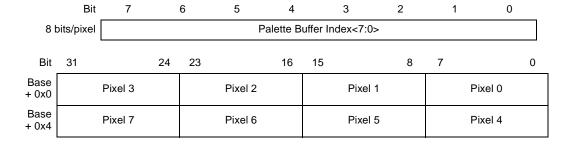
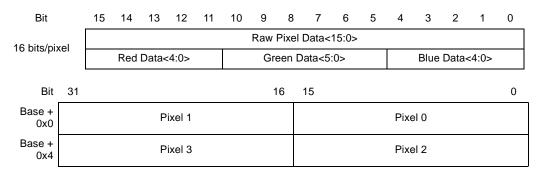


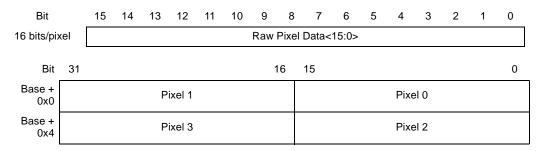


Figure 7-10. 16 Bits Per Pixel Data Memory Organization - Passive Mode



Note: For passive 16 bits per pixel operation, the Raw Pixel Data must be organized as shown above.

Figure 7-11. 16 Bits Per Pixel Data Memory Organization - Active Mode



Note: For active 16 bits per pixel operation, the Raw Pixel Data is sent directly to the LCD panel pins and must be in the same format as required by the LCD panel.

In dual-panel mode, pixels are presented to two halves of the screen at the same time (upper and lower). A second DMA channel, input FIFO, and output FIFO exist to support dual-panel operation. The palette buffer is implemented in DMA channel 0, but not channel 1. The frame source address for the lower half always points to the top of the encoded pixel values for channel 1. The frame source address of both DMA channels must be configured such that the least significant three address bits are all zero (address bits 2 through 0 must be zero). This requires that the source address of the frame buffer start at 8-byte boundaries.

Each line in the frame buffer must start at a word boundary. For the various pixel sizes, this requires each line of the display to have pixels in multiples of: 32 pixels for 1-bit pixels, 16 pixels for 2-bit pixels, 8 pixels for 4-bit pixels, 4 pixels for 8-bit pixels, and 2 pixels for 16-bit pixels. If the LCD screen does not naturally have the correct multiple of pixels per line, the user must adjust the start address for each line by adding dummy pixel values to the end of the previous line.

Note: There are two special conditions: 8 bits/pixel monochrome screens with double-pixel-data mode and 8 or 16 bits/pixel passive color screens require a multiple of 8 pixels for each line.

For example, if the screen being driven is 107 pixels wide, and 4 bits/pixel mode is used, each line is 107 pixels or nibbles in length (53.5 bytes). The next nearest 8-pixel boundary (for 4-bit pixels) occurs at 112 pixels or nibbles (56 bytes). Each new line in the frame buffer must start at multiples of 56 bytes by adding an extra 5 dummy pixels per line (2.5 bytes) to LCCR1[PPL].

If dummy pixels are to be inserted, the panel must ignore the extra pixel clocks at the end of each line that correspond to the dummy pixels.



Use the following equation to calculate the total size of the frame buffer (in bytes). This calculation is used to encode the length of the frame buffer in the DMA descriptors (Section 7.6.5.4). The first term is the size required for the encoded pixel values. "Lines" is the number of lines for the display. "Pixels" is the number of pixels per line. Use the actual line/pixel count, not minus 1 as in the LCCR registers. "n" = the number of extra dummy pixels required per line, as described above. For dual-panel mode, the frame buffer size is equally distributed between the two DMA channels. Therefore, "Lines" in this equation are divided in half for dual-panel mode.

FrameBufferSize =
$$\frac{BitsPerPixel \bullet Lines \bullet (Pixels + n)}{8}$$

The bandwidth required for the LCD Controller can be calculated using the following equations. FrameBufferSize is the result of the previous equation. Bandwidth is always an important part of any system analysis. Systems with large panels and high bits per pixel must ensure that the panel is not starved for data.

```
BusBandwidth = (FrameBufferSize + PaletteSize) \bullet RefreshRate \\ BusBandwidth(DualPanel) = (FrameBufferSize \bullet 2 + PaletteSize) \bullet RefreshRate
```

Sample calculations for an 640x480 panel, 16 bits per pixel, 60 Hz refresh rate:

```
FrameBufferSize = 16*640*480/8 = 614,400 bytes

Bus Bandwidth = 614,400 * 60 = 36.9 MB/sec
```

7.5 Functional Timing

Figure 7-12 through Figure 7-14 illustrate LCD controller timing in passive display mode. The example used is a 320x240 panel. Figure 7-15 and Figure 7-16 illustrate the LCD controller timing in active display mode. For precise timing relationships, see the *Intel® PXA255 Processor Electrical, Mechanical, and Thermal Specification*.



Figure 7-12. Passive Mode Start-of-Frame Timing

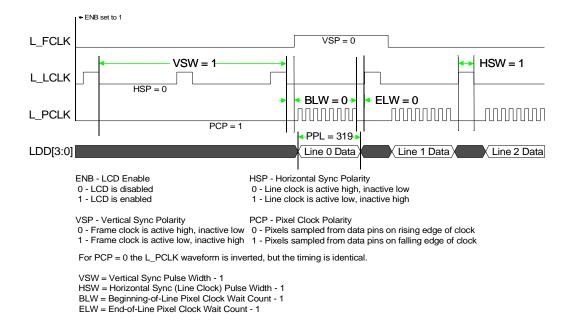


Figure 7-13. Passive Mode End-of-Frame Timing

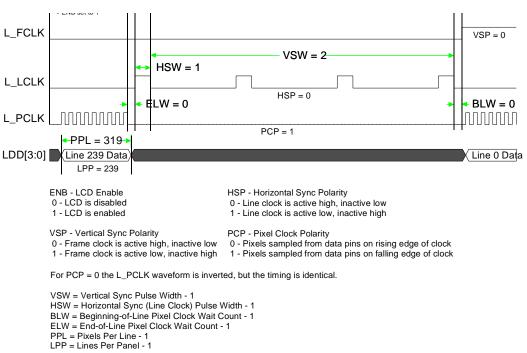
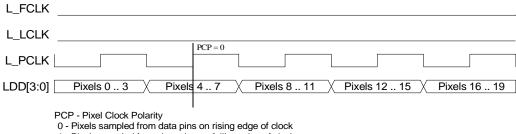




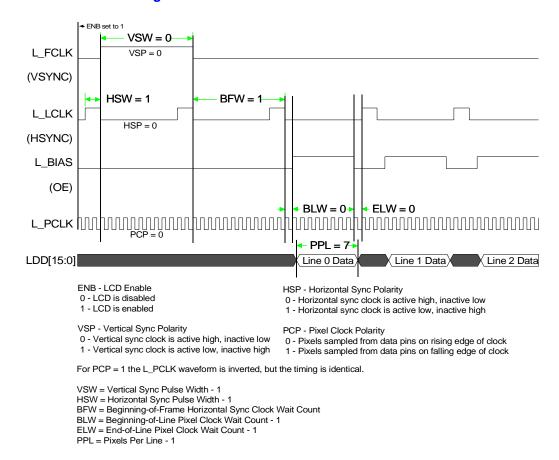
Figure 7-14. Passive Mode Pixel Clock and Data Pin Timing



1 - Pixels sampled from data pins on falling edge of clock

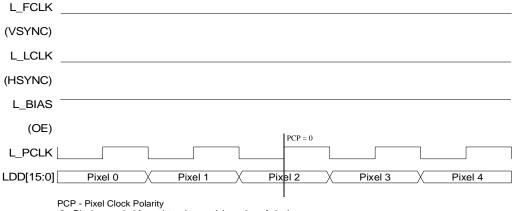
For PCP = 1 the L_PCLK waveform is inverted, but the timing is identical.

Figure 7-15. Active Mode Timing









^{0 -} Pixels sampled from data pins on rising edge of clock.

For PCP = 1 the L_PCLK waveform is inverted, but the timing is identical.

7.6 Register Descriptions

The LCD controller contains four control registers, ten DMA registers, one status register, and a 256-entry palette RAM. Table 7-16 lists their locations in physical memory. All of the LCD registers must be accessed as 32-bit values. Reads and writes to undefined addresses in the LCD register space yield unpredictable results and must not be attempted.

The control registers contain bit fields to do the following:

- Enable and disable the LCD controller.
- Define the height and width of the screen being controlled.
- Indicate single- or dual-panel display mode.
- Indicate color versus monochrome mode.
- Indicate passive versus active display.
- Set the polarity of the control pins.
- Set the pulse width of the line and frame clocks, pixel clock, and ac bias pin frequency.
- Set the AC bias pin toggles per interrupt.
- Set the number of wait states to insert before and after each line and after each frame.
- Enable various interrupt masks.

An additional control field exists to tune the DMAC's performance based on the type of memory system with which the processor is used. This field controls the placement of a minimum delay between each LCD DMA request during palette loads to insure enough bus bandwidth is given to other bus masters for accesses.

^{1 -} Pixels sampled from data pins on falling edge of clock.



The DMA descriptor addresses are initially programmed by software. After that, the other DMA registers are programmed by the hardware. Section 7.6.5 provides a complete description of how the DMA is programmed.

The status registers contain bits that signal:

- Input and output FIFO overrun and underrun errors
- DMA bus errors
- When the DMAC starts and ends a frame
- When the last active frame has completed after the LCD is disabled
- Each time the L_BIAS pin has toggled a programmed number of times

Each of these hardware-detected events can signal an interrupt request to the interrupt controller.

7.6.1 LCD Controller Control Register 0 (LCCR0)

The control bits in LCCR0, shown in Table 7-3, within all other control registers must be programmed before setting ENB=1 (a word write can be used to configure LCCR0 while setting ENB after all other control registers have been programmed). The LCD controller must be disabled when changing the state of any control bit within the LCD controller.

LCD Output Fifo Underrun Mask (OUM) — used to mask interrupt requests that are asserted whenever an output FIFO underrun error occurs. When OUM=0, underrun interrupts are enabled, and whenever an output FIFO underrun (OU) status bit within the LCD status register (LCSR) is set (one), an interrupt request is made to the interrupt controller. When OUM=1, underrun interrupts are masked and the state of the underrun status bit (OU) is ignored by the interrupt controller. Setting OUM does not affect the current state of the status bit or the LCD controller's ability to set and clear it, it only blocks the generation of the interrupt request. Output FIFO underruns are more critical than Input FIFO underruns, since Output FIFO underruns will affect the display.

Branch Mask (BM) — used to mask interrupt requests that are asserted after the LCD Controller has branched to a new set of frame descriptors. See Section 7.6.6 for details.

Palette DMA Request Delay (PDD) — used to select the minimum number of internal bus clock cycles to wait between the servicing of each DMA request issued while the on-chip palette is loaded. When the palette is optionally loaded at the beginning of a frame, up to 512 bytes of data may be accessed by the LCD's DMAC. Using PDD allows other bus masters to gain access to shared memory in between palette DMA loads. PDD must be used carefully, as it can severely degrade LCD controller performance if not used properly. It is recommended to leave PDD zero and add delay only when necessary. PDD does not apply to the normal input FIFO DMA requests for frame buffer information, since these DMA requests do not occur back-to-back. The input FIFO DMA request rate is a function of the rate at which pixels are displayed on the screen.

After a word of palette data is written to the input FIFO, the value contained within PDD is loaded to a down counter that disables the palette from issuing another DMA request until the counter decrements to zero. This counter ensures that the LCD's DMAC does not tie up the full bandwidth of the processor system bus. Once the counter reaches zero, any pending or future DMA requests by the palette cause the DMAC to arbitrate for the bus. Once the DMA burst cycle has completed, the process starts over, and the value in PDD is loaded to the counter to create another wait state period, which disables the palette from issuing a DMA request. PDD can be programmed with a



value that causes the FIFO to wait from 0 to 255 clock cycles after the completion of one DMA request to the start of the next request. When PDD=0x00, the FIFO DMA request delay function is disabled.

LCD Quick Disable Interrupt Mask (QDM) — used to mask interrupt requests that are asserted after the LCD Enable bit (ENB) is cleared and the DMAC finishes the current burst transfer. The LCD controller immediately stops requesting new data and the current frame is not completed. This shutdown is for Sleep shutdown. When QDM=0, the quick disable interrupt is enabled, and whenever the LCD quick disable (QD) status bit in the LCD Status Register (LCSR) is set, an interrupt request is made to the interrupt controller. When QDM=1, the quick disable interrupt is masked and the state of the QD status bit is ignored by the interrupt controller. Setting QDM does not affect the current state of QD or the LCD controller's ability to set and clear QD, it only blocks the generation of the interrupt request.

LCD Disable (DIS) — During LCD Controller operation, setting DIS=1 causes the LCD Controller to finish fetching the current frame from memory and then shut down cleanly. If the LCD DMAC is loading the palette RAM when DIS is set, the load will complete followed by the next frame, and then the LCD controller is disabled. Completion of the current frame is signalled by the LCD when it sets the LCD disable done flag (LDD) in register LCSR. Use a read-modify-write procedure to set this bit, since the other bit fields within LCCR0 continue to be used until the current frame is completed. The LCD Enable bit (ENB) is cleared when the disable is completed. Section 7.2.2 for more information.

Double-Pixel Data (DPD) Pin Mode — selects whether four or eight data pins are used for pixel data output to the LCD screen in single-panel monochrome mode. When DPD=0, L_DD[3:0] pins are used to send 4 pixel values each pixel clock transition. When DPD=1, L_DD[7:0] pins are used to send 8 pixel values each pixel clock. See Table 7-3 for a comparison of how the LCD's data pins are used in each of its display modes.

Note: DPD does not affect dual-panel monochrome mode, any of the color modes, or active mode. Clear DPD in these modes.

Passive/Active Display Select (PAS) — selects whether the LCD controller operates in passive (STN) or active (TFT) display control mode. When PAS=0, passive mode is selected. All LCD data flow operates normally (including the LCD's dither logic), and all LCD controller pin timing operates as described in Table 7.5.

When PAS=1, active mode is selected. 1- and 2-bit pixel modes are not supported in active mode. For 4- and 8-bit pixel modes, pixel data is transferred via DMA from off-chip memory to the input FIFO, unpacked, and used to select an entry from the palette, just as in passive mode. However, the value read from the palette bypasses the LCD controller's dither logic and is sent directly to the output FIFO to be driven onto the LCD's data pins. This 16-bit output to the pins represents 5 bits of red, 6 bits of green, and 5 bits of blue data. For 16-bit pixel encoding mode, two 16-bit values are packed into each word in the frame buffer. Each 16-bit value is transferred via DMA from off-chip memory to the input FIFO. Unlike 4 and 8 bit per pixel modes, the 16-bit value bypasses both the palette and the dither logic and is placed directly in the output FIFO to be sent to the LCD's data pins. Using the 16-bit pixel encoding mode allows a total of 64K colors to be generated.

The 16-bit output from either the palette or frame buffer to the pins can be organized in any fashion necessary to correctly interface with the LCD panel. Typically, the output is configured into one of three user-specified RGB color formats: 6 bits of red, 5 bits of green, and 5 bits of blue data; 5 bits of red, 6 bits of green, and 5 bits of blue data; and 5 bits of red, 5 bits of green, and 6 bits of blue data. The RGB format 5:6:5 is normally used, since the human eye can distinguish more shades of green than of red or blue.



The LCD pin timing changes when active mode is selected. Timing of each pin is described in subsequent bit-field sections for both passive and active mode.

The LCD controller can be configured in active color display mode and used with an external DAC and optionally an external palette to drive a video monitor. Only monitors that implement the RGB data format can be used. The LCD controller does not support the NTSC standard. However, the 2X pixel clock mode allows the LCD controller to easily interface with an NTSC encoder, such as the Analog Devices 7171 encoder.

Figure 7-17 shows which bits are sent to the individual LCD data pins for both a frame buffer entry (for 16-bit/pixel mode) and a selected palette entry (for 1, 2, 4 and 8 bit/pixel mode). The pixel bits corresponding to L_DD pins when using an RGB format of 5:6:5 are also shown. In active mode, L_DD pins [15:8] are also used. The user must configure the proper GPIO pins for LCD operation to enable LCD Controller operation. See Chapter 4, "System Integration Unit" for GPIO configuration information.

The processor LCD controller may be used with active panels having more than 16 data pins, but the panel will be limited to a total of 64K colors. There are three options:

- 1. To maintain the panel's full range of colors and increase the granularity of the spectrum, connect the LCD controller's 16 data pins to the panel's most significant R, G, and B pixel data input pins and tie the panel's least significant R, G, and B data pins either high or low.
- 2. To maintain the granularity of the spectrum and limit the overall range of colors possible, connect the LCD controller's 16 data pins to the panel's least significant R, G, and B pixel data input pins and tie the panel's most significant data pins either high or low.
- 3. Sometimes, better results can be obtained by replicating the upper bits on the lower bits.

Figure 7-17. Frame Buffer/Palette Output to LCD Data Pins in Active Mode

4/8/16 Bits/Pixel Mode, Frame Buffer or Palette Entry

Plxel	R4	R3	R2	R1	R0	G5	G4	G3	G2	G1	G0	B4	В3	B2	B1	В0
Bit	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
L_DD Pin	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0

End of Frame Mask (EFM) — used to mask interrupt requests that are asserted at the end of each frame (when the DMA length of transfer counter decrements to zero). When EFM=0, the interrupt is enabled, and whenever the EOF status bit in the LCD status register (LCSR) is set to one, an interrupt request is made to the interrupt controller. When EFM=1, the interrupt is masked, and the state of the EOF status bit is ignored by the interrupt controller. Setting EFM does not affect the current state of EOF or the LCD controller's ability to set and clear EOF, it only blocks the generation of the interrupt request.

Input Fifo Underrun Mask (IUM) — used to mask interrupt requests that are asserted whenever an input FIFO underrun error occurs. When IUM=0, underrun interrupts are enabled, and whenever an input FIFO underrun (IUL, IUU) status bit in the LCD status register (LCSR) is set to one, an interrupt request is made to the interrupt controller. When IUM=1, underrun interrupts are masked and the state of the underrun status bits (IUL, IUU) is ignored by the interrupt controller. Setting IUM does not affect the current state of these status bits or the LCD controller's ability to set and clear them, it only blocks the generation of the interrupt requests.

Start Of Frame Mask (SFM) — used to mask interrupt requests that are asserted at the beginning of each frame when the LCD's frame descriptor has been loaded into the internal DMA registers. When SFM=0, the interrupt is enabled, and whenever the start of frame (SOF) status bit in the LCD



status register (LCSR) is set, an interrupt request is made to the interrupt controller. When SFM=1, the interrupt is masked and the state of the SOF status bit is ignored by the interrupt controller. Setting SFM does not affect the current state of SOF or the LCD controller's ability to set and clear SOF, it only blocks the generation of the interrupt request.

LCD Disable Done Interrupt Mask (LDM) — used to mask interrupt requests that are asserted after the LCD is disabled and the frame currently being sent to the output pins has completed. When LDM=0, the interrupt is not masked, and whenever the LCD disable done (LDD) status bit in the LCD status register (LCSR) is set to one, an interrupt request is made to the interrupt controller. When LDM=1, the interrupt is masked, and the state of the LDD status bit is ignored by the interrupt controller. Setting LDM does not affect the current state of LDD or the LCD controller's ability to set and clear LDD, it only blocks the generation of the interrupt request. This interrupt is used when the LCD must be disabled after the current frame being sent to the output pins has completed. Clearing LCD Enable (ENB) is a quick disable, and LDD is not set.

Note: This mask bit applies only to regular shutdowns using the LCD Disable (DIS) bit.

Single-/Dual-Panel Select (SDS) — In passive mode (PAS=0), SDS is used to select the type of display control implemented by the LCD screen. When SDS=0, single-panel operation is selected (pixels presented to screen a line at a time). When SDS=1, dual-panel operation is selected (pixels presented to screen two lines at a time). Single-panel LCD drivers have one line/row shifter and driver for pixels and one line pointer. Dual-panel LCD controller drivers have two line/row shifters (one for the top half of the screen, one for the bottom) and two line pointers (one for the top half of the screen, one for the bottom).

When dual-panel mode is programmed, both of the LCD controller's DMA channels are used. DMA channel 0 is used to load the palette RAM from the frame buffer and to drive the upper half of the display, and DMA channel 1 drives the lower half. The two channels alternate when fetching data for both halves of the screen, placing encoded pixel values in the two separate input FIFOs.

When dual-panel operation is enabled, the LCD controller doubles its pin usage. For monochrome screens, eight pins are used and for color screens, 16 pins are used.

Note: SDS must be set to 0 in active mode (PAS=1).

Table 7-3 shows the LCD data pins and GPIO pins used for each mode of operation and the ordering of pixels delivered to a screen for each mode of operation.

Note: In passive color mode, the data pin ordering switches. Figure 7-18 shows the LCD data pin pixel ordering.

Table 7-2. LCD Controller Data Pin Utilization (Sheet 1 of 2)

Color/Monochrome Panel	Single/ Dual Panel	Passive/ Active Panel	Screen Portion	Pins
Monochrome	Single	Passive	Whole	L_DD[3:0]
Monochrome	Single	Passive	Whole	L_DD[7:0] [†]
Monochrome	Dual	Passive	Тор	L_DD[3:0]
Worldchionie	Duai	rassive	Bottom	L_DD[7:4]
Color	Single	Passive	Whole	L_DD[7:0]



Table 7-2. LCD Controller Data Pin Utilization (Sheet 2 of 2)

Color/Monochrome Panel	Single/ Dual Panel	Passive/ Active Panel	Screen Portion	Pins
Color	Dual	Passive	Тор	L_DD[7:0]
Color	Duai	i assive	Bottom	L_DD[15:8]
Color	Single	Active	Whole	L_DD[15:0]

[†] Double-pixel data mode (DPD) = 1.

Figure 7-18. LCD Data-Pin Pixel Ordering

	orner of So		Column 2	Column 2	Column 4	Column 5	Column 6	Column 7	Column 9
Row 0	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[0]
Row 1	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[0]
Row 2	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[0]
Row 3	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[0]
	Pass	ive Mo	nochro	me Sing	gle-Pan	el Disp	lay Pixe	el Orde	ring
Top Left Co			Column 2	Column 3	Column 4	Column 5	Column 6	Column 7	Column 8
Row 0	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[4]	LDD[5]	LDD[6]	LDD[7]	LDD[0]
Row 1	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[4]	LDD[5]	LDD[6]	LDD[7]	LDD[0]
Row 2	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[4]	LDD[5]	LDD[6]	LDD[7]	LDD[0]
Row 3	LDD[0]	LDD[1]	LDD[2]	LDD[3]	LDD[4]	LDD[5]	LDD[6]	LDD[7]	LDD[0]
	Manaa	bromo	Cinala	Donal C)aubla	Divol Di	onlov F	Divol Or	dorina
Passive	MOHOC	nrome	Single-	Panei L	ouble-	Pixei Di	ispiay r	rixei Or	dering
Top Left Co	orner of So	reen							
								Column 7	
Row 0	LDD<0>	LDD<1>	LDD<2>	LDD<3>	LDD<0>	LDD<1>	LDD<2>	LDD<3>	LDD<0>
Row 1	LDD<0>	LDD<1>	LDD<2>	LDD<3>	LDD<0>	LDD<1>	LDD<2>	LDD<3>	LDD<0>
					. :				
Row n/2	LDD<4>	LDD<5>	LDD<6>	LDD<7>	LDD<4>	LDD<5>	LDD<6>	LDD<7>	LDD<4>
Row n/2+1	LDD<4>	LDD<5>	LDD<6>	LDD<7>	LDD<4>	LDD<5>	LDD<6>	LDD<7>	LDD<4>
n = # of row	s Pas	sive Mo	nochro	me Du	al-Pane	d Displa	ay Pixe	l Orderi	ng
Top Left Co	orner of So	reen							
	Column 0 Red	Column 0 Green	Column 0 Blue	Column 1 Red	Column 1 Green	Column 1 Blue	Column 2 Red	Column 2 Green	Column 2 Blue
Row 0	LDD<7>	LDD<6>	LDD<5>	LDD<4>	LDD<3>	LDD<2>	LDD<1>	LDD<0>	LDD<7>
Row 1	LDD<7>	LDD<6>	LDD<5>	LDD<4>	LDD<3>	LDD<2>	LDD<1>	LDD<0>	LDD<7>
Row 2	LDD<7>	LDD<6>	LDD<5>	LDD<4>	LDD<3>	LDD<2>	LDD<1>	LDD<0>	LDD<7>
Row 3	LDD<7>	LDD<6>	LDD<5>	LDD<4>	LDD<3>	LDD<2>	LDD<1>	LDD<0>	LDD<7>
	P	assive	Color S	ingle-P	anel Di	snlav F	ixel Or	derina	1
	•	400.70	00.0. 0	inigio i	u	opiay i	1,01 01	uoiiig	
Top Left Co				Caluman 0	Caluman 2		Caluman 4	Caluman F	Caluman F
	Red	Column 0 Green		Green	Column 2 Blue	,	Column 4 Blue	Column 5 Red	Column 5 Green
Row 0	LDD[7]	LDD[6]	T T	LDD[0]	LDD[7]	Γ	LDD[1]	LDD[0]	LDD[7]
Row 1	LDD[7]	LDD[6]	T 7	LDD[0]	LDD[7]	Γ -	LDD[1]	LDD[0]	LDD[7]
			-			-			
Row n/2	LDD[15]	LDD[14]	Τ -	LDD[8]	LDD[15]	Γ -	LDD[9]	LDD[8]	LDD[15]
Row n/2+1	LDD[15]	LDD[14]	Τ -	LDD[8]	LDD[15]	Γ -	LDD[9]	LDD[8]	LDD[15]



Color/Monochrome Select (CMS) — selects whether the LCD controller operates in color or monochrome mode. When CMS=0, color mode is selected. Palette entries are 16 bits wide (5-bits red, 6-bits green, 5-bits blue), 8 data pins are enabled for single-panel mode, 16 data pins are enabled for dual-panel mode, and all three dither blocks are used, one each for the red, green, and blue pixel components. When CMS=1, monochrome mode is selected, palette entries are 8 bits wide, 4 or 8 data pins are enabled for single-panel mode, 8 data pins are enabled for dual-panel mode, and the blue dither block is used.

LCD Enable (ENB) — used to enable and quickly disable all LCD controller operation. When ENB=0, the LCD controller is either disabled or in the process of quickly disabling, and all of the LCD pins can be used for GPIO. When ENB=1, the LCD controller is enabled.

All other control registers must be initialized before setting ENB. LCCR0 can be programmed last, and all bit fields can be configured at the same time with a word write to the register. If ENB is cleared while the LCD controller is enabled, the LCD controller will immediately stop requesting data from the LCD DMAC, and the current frame will not complete. The LCD controller must not be re-enabled until the QD status flag is set in register LCSR, indicating the quick disable is complete. Quick disable is for sleep shutdown. Regular shutdown of the LCD controller at the end of the frame can be done via the LCD Disable bit, LCCR0[DIS]. There are separate maskable interrupts for quick disable and regular disable. See Section 7.2.1 for more information.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 7-3. LCCR0 Bit Definitions (Sheet 1 of 2)

	Physical Address 0x4400_0000											LCCR0											LCD Controller											
Bit	31 30 29 28 27 26 25 24							23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0			
					OUM	BM				PE	DD				QDM	DIS	DPD	reserved	PAS	EFM	IUM	SFM	LDM	SDS	CMS	ENB								
Reset	x x x x		X	X	X	X	X	X	X	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0				
		Bits			Na	me		Description																										
		31:22			_	_		reserved																										
		21			Ol	JM		Out 0 = 1 =	FIF	=0 ı	und	erru	n er	rors	ger						erru	pt.												
		20			В	М			Ge	ener	ates										new te a			ıpt.										
		19:12			ΡI	DD		Palette DMA Request Delay: Value (0–255) specifies the number of internal bus clocks to wait before requesting another burst of palette data. The counter starts decrementing when the first word is written to the input FIFO buffer. PDD = 0x00 disables this function.																										
		11			QI	OM		0 =	Ge	ener	ate		nter	rupt	afte tatus						e an	inte	errup	ot.										
		10			D	IS		LC[0 = 1 =	LC	D C	Cont									in th	ie pr	oce:	ss o	f dis	sabli	ing.								



Table 7-3. LCCR0 Bit Definitions (Sheet 2 of 2)

Ta	Table 7-3. LCCR0 Bit Definitions (Sheet 2 of 2) Physical Address																														
				Addr 0_000						L	CCR	0				LCD Controller															
Bit	31	30	29	28	27	26 2	5 2	24	23 22	21	2	0 19	18	1	7 10	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				r	ese	rved				MNO	200	NA NA			F	DD				QDM	SIQ	DPD	reserved	PAS	EFM	IOM	SFM	LDM	SDS	CMS	ENB
Reset	X	X	X	X	X	X 2	Κ	X	x x	X)	K 0	0		0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Nam	е		Description																						
									Double-Pixel Data (DPD) pin mode: In passive monochrome single panel mode,																						
		ç	1			DPD	,		•							•															
			,			טו נ	,		$0 = L_{-}$ $1 = L_{-}$																		ISITIO	on.			
									In any	othe	er i	mode	, DF	P) mus	t be	0.														
		8	3			_			reserved Passive/Active Display select:																						
		7	,			PAS	2	Passive/Active Display select: 0 = Passive (or STN) display operation enabled.																							
		,				IAC	,		0 = Pa	assive	ve e (d	or TF	T) d	a ISI	iispia play (y ope opera	erati	on e ena	nac able	nea. d.											
									End of Frame Mask: 0 = Generates an interrupt at the end of a frame.																						
		6	6			EFM	1		0 = G 1 = E						•							an	inte	rrup	t.						
								Input FIFO Underrun Mask: 0 = FIFO underrun errors generate an interrupt.																							
		5)			IUM	l		0 = FI 1 = FI	FO FO	un un	nderru nderru	ın er	rc	ors ge	nera not	ite a gen	n in erat	terrı e ar	upt. n int	erru	pt.									
							_		Start of	f Fra	am	ne Ma	sk:																		
		2	ļ			SFM	1		0 = St 1 = St																	an i	nter	rupt			
								LCD Disable Done Mask: 0 = LCD disable done condition generates an interrupt (state of LDD status sent to the																							
		3	3			LDM	1		in	terru	ıpt	t cont	rolle	r)								. `									
						1 = LCD disable done condition does not generate an interrupt (LDD status bit ignored). Single-/Dual-panel Display Select:																									
		2	2			SDS	3		0 = Si 1 = Di																						
									SDS m								AS=	1).													
									Color/N																						
		1	l			CMS	s .		0 = Co 1 = M								led.														
									LCD C	ontr	oll	ler Er	nable	:																	

7.6.2 LCD Controller Control Register 1 (LCCR1)

1 = LCD controller enabled.

ENB

LCCR1, shown in Table 7-4, contains four bit fields that are used as modulus values for a collection of down counters, each of which performs a different function to control the timing of several of the LCD pins. These values must be programmed before enabling the LCD Controller.

0 = LCD controller disabled or in the process of quickly disabling.



Beginning-of-Line Pixel Clock Wait Count (BLW) — used to specify the number of dummy pixel clocks to insert at the beginning of each line or row of pixels. After the line clock for the previous line has been negated, the value in BLW is used to count the number of pixel clocks to wait before starting to output the first set of pixels in the next line. BLW generates a wait period ranging from 1 to 256 pixel clock cycles. BLW must be programmed with the desired number of pixel clocks minus one. L_PCLK does not toggle during these "dummy" pixel clock cycles in passive display mode. It does toggle continuously in active display mode.

End-of-Line Pixel Clock Wait Count (ELW) — used to specify the number of "dummy" pixel clocks to insert at the end of each line or row of pixels before pulsing the line clock pin. Once a complete line of pixels is transmitted to the LCD driver, the value in ELW is used to count the number of pixel clocks to wait before pulsing the line clock. ELW generates a wait period ranging from 1 to 256 pixel clock cycles. ELW must be programmed with the desired number of pixel clocks minus one. L_PCLK does not toggle during these dummy pixel clock cycles in passive display mode. It does toggle continuously in active display mode.

Horizontal Sync Pulse Width (HSW) — specifies the pulse width (minus 1) of the line clock in passive mode or the horizontal synchronization pulse in active mode. L_LCLK is asserted each time a line is sent to the display and a programmable number of pixel clock wait states have elapsed. When L_LCLK is asserted, the value in HSW is transferred to a 6-bit down counter, which decrements at the programmed pixel clock frequency. When the counter reaches zero, L_LCLK is negated. HSW can be programmed to generate a line clock pulse width ranging from 1 to 64 pixel clock periods.

The pixel clock does not toggle during the line clock pulse in passive display mode but does toggle in active display mode. The polarity (active and inactive state) of the line clock pin is programmed using the horizontal sync polarity (HSP) bit in LCCR3.

HSW must be programmed with the desired number of pixel clocks minus one.

Note: The term "pulse width" refers to the time which L_LCLK is asserted, rather than the time for a cycle of the line clock to occur.

Pixels Per Line (PPL) — used to specify the number of pixels in each line or row on the screen (minus one). PPL is a 10-bit value that represents between 1 and 1024 pixels per line. It is recommended not to exceed 640 pixels. It is used to count the number of pixel clocks that must occur before the line clock can be asserted. As discussed in Section 7.4.2, pixels per line must be multiples of: 32 pixels for 1-bit pixels, 16 pixels for 2-bit pixels, 8 pixels for 4-bit pixels, 4 pixels for 8-bit pixels, and 2 pixels for 16-bit pixels. The two special conditions are: 8-bits/pixel monochrome screens with double-pixel-data mode and 8 or 16 bits/pixel passive color screens require a multiple of 8 pixels for each line.

If the display used is not naturally a multiple of the above, "dummy" pixels must be added to each line to keep the frame buffer aligned in memory. For example, if the display being controlled is 250 pixels wide and the pixel-size is 8-bits, the nearest greater multiple of 8 is 256. Pixels per line must be set to 255. 6 extra "dummy" pixel values must be added to the end of each line in the frame buffer. The display being controlled must ignore the dummy pixel clocks at the end of each line.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 7-4. LCCR1 Bit Definitions

			P		ical 440(SS		LCD Controller Control Register 1													LCD Controller											
Bit	31	31 30 29 28		27 26 2		25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0			
				BL	LW				ELW									н	SW						PPL									
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
		Bi	ts			Na	me												De	escr	iptic	on												
		31:	24			BL	_W		Beginning-of-Line pixel clock Wait count: Value (0–255) specifies the number of pixel clock periods to add to the beginning of a line transmission before the first set of pixels is sent to the display. BOL wait = (BLW + 1).															ne										
		23:	16			EL	_W		End-of-line pixel clock wait count: Value (0–255) specifies the number of pixel clock periods to add to the end of a line transmission before line clock is asserted. EOL = (ELW+1). In passive display mode, Pixel Clock is held in its inactive state during the end-of-line wait period. In active display mode, it toggles.															ait										
		15:	10			HS	SW		Horizontal sync pulse width: Value (0–63) specifies the number of pixel clock periods to pulse the line clock at the end of each line. HSYNC pulse width = (HSW+1). In passive display mode, Pixel Clock is held in its inactive state during the generation of the line clock. In active display mode, it toggles.																									
		9:	0			PF	٦L			els p				er o	of pi	xels	cor	ıtair	ned v	with	in ea	ach	line	on t	the L	_CD) dis	play	/. Ac	tual	pix	els		

7.6.3 LCD Controller Control Register 2 (LCCR2)

per line = (PPL+1).

LCCR2, shown in Table 7-5, contains four bit fields that are used as values for a collection of down counters, each of which performs a different function to control the timing of several of the LCD's pins.

Beginning-of-Frame Line Clock Wait Count (BFW) — used in active mode (LCCR0[PAS] = 1) to specify the number of line clocks to insert at the beginning of each frame. The BFW count starts when the VSYNC signal for the previous frame is negated. The value in BFW is used to count the number of line clock periods to insert before starting pixel output in the next frame. BFW generates a wait period ranging from 0 to 255 extra L_LCLK cycles (BFW=0x00 disables the wait count). L_LCLK does toggle during the generation of the BFW line clock wait periods.

In passive mode, BFW must be set to zero so that no beginning-of-frame wait states are generated. Use VSW exclusively in passive mode to insert line clock wait states, which allow the LCD controller's DMAC to fill the palette and insert additional pixels before the start of the next frame.

End-of-Frame Line Clock Wait Count (EFW) — used in active mode (LCCR0[PAS] = 1) to specify the number of line clocks to insert at the end of each frame. Once a complete frame of pixels is transmitted to the LCD display, the value in EFW is used to count the number of line clock periods to wait. After the count has elapsed, the VSYNC (L_FCLK) signal is pulsed. EFW generates a wait period ranging from 0 to 255 line clock cycles (EFW = 0x00 disables the EOF wait count). L_LCLK does not toggle during the generation of the EFW line clock periods.



In passive mode, EFW must be set to zero so that no EOF wait states are generated. Use VSW exclusively in passive mode to insert line clock wait states, which allow the LCD controller's DMAC to fill the palette and insert additional pixels before the start of the next frame.

Vertical Sync Pulse Width (VSW) — used to specify the pulse width of the vertical synchronization pulse in active mode or to add extra "dummy" line clock wait states between the end and beginning of frame in passive mode.

In active mode (LCCR0[PAS]=1), L_FCLK is used to generate the vertical sync signal and is asserted each time the last line or row of pixels for a frame is sent to the display and a programmable number of line clock wait states as specified by LCCR1[BLW] have elapsed. When L_FCLK is asserted, the value in VSW is transferred to a 6-bit down counter, which uses the line clock frequency to decrement. When the counter reaches zero, L_FCLK is negated. VSW can be programmed to generate a vertical sync pulse width ranging from 1 to 64 line clock periods. VSW must be programmed with the desired number of line clocks minus one. The polarity (active and inactive state) of the L_FCLK pin is programmed using the vertical sync polarity (VSP) bit in LCCR3.

In passive mode (LCCR0[PAS]=0), VSW does not affect the timing of the L_FCLK pin, but rather can be used to add extra line clock wait states between the end of each frame and the beginning of the next frame. When the last line clock of a frame is negated, the value in VSW is transferred to a 6-bit down counter that uses the line clock frequency to decrement. When the counter reaches zero, the next frame begins. VSW can be programmed to generate from 1 to 64 dummy line clock periods between each frame in passive mode. VSW must be programmed to allow:

- enough wait states to occur between frames such that the LCD's DMAC is able to fully load the on-chip palette (if applicable)
- a sufficient number of encoded pixel values to be fetched from the frame buffer, to be
 processed by the dither logic and placed in the output FIFO, ready to be sent to the LCD data
 pins.

The number of wait states required is system dependent, depending on such factors as:

- palette buffer size (none; 8, 32 or 512 bytes)
- memory system speed (number of wait states, burst speed, number of beats)
- Palette DMA request delay, LCCR0[PDD].

The line clock pin does toggle during the insertion of the line clock wait state periods.

VSW does not affect generation of the frame clock signal in passive mode. Passive LCD displays require that the frame clock be active on the rising edge of the first line clock pulse of each frame, with adequate setup and hold time. To meet this requirement, the LCD controller's frame clock pin is asserted on the rising edge of the first pixel clock for each frame. The frame clock remains asserted for the remainder of the first line as pixels are sent to the display. It is then negated on the rising edge of the first pixel clock of the second line of each frame.

Lines Per Panel (LPP) — specifies the number of lines or rows present on the LCD panel being controlled. In single-panel mode, it represents the total number of lines for the entire LCD display. LPP is used to count the correct number of line clocks that must occur before the frame clock can be pulsed. In dual-panel mode, it represents half the number of lines of the entire LCD display, which is split into two panels. LPP is a 10-bit value that represents between 1 and 1024 lines per screen. It must be programmed with the actual height of the display minus one. It is recommended not to exceed 480 pixels. For portrait mode panels, more than 480 pixels can be used as long as total pixels do not exceed 480,000. For example, a 480x640 portrait mode panel can be used.



This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 7-5. LCCR2 Bit Definitions

		•	_					• • • •		•																						
					ical 4400			S			LC	D C	ont	roll	er C	ont	rol	Reg	iste	r 2					LCE	Co	ontro	olle	r			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				BF	w							EF	w						VS	w							LF	P				
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bit	s			Na	me												De	SCI	ipti	on										
									Bed	inn	ing-	of-fı	rame	e lin	e cl	ock	wait	COL	ınt:													
									`		•									-25	55) s	spec	ifies	s the	nu	mbe	er of	line	e clo	ck p	eric	ods
		31:2	24			BF	W		to a	dd	to th	ne b	egir	nin	g of	a fr	ame	bef	orè	the	first	set	of p	oixel	s is	sen	t to	the				
														00		U									lock	per	iods					
									BF\	N n	nust	be	clea	red	to z	ero	(dis	able	ii (be	n pa	ıssiv	e m	ode) .								
									End	l-of	-frar	ne I	ine (cloc	k wa	ait c	ount	::														
		23:1	16			EF	W		to a	dd		e e	nd o														er of e ins					
														red	to z	ero	(dis	able	d) ir	n pa	ssiv	e m	ode) .								
									Ver	tica	l syı	пс р	ulse	wic	dth:																	
		15:1	10			VS	SW		to p ela _l	uls se	e the	e L_ ram	.FCL e clo	K p	in a use	t the	en VS	d of YNC	eac	h fra Inal	ame	afte	r th	e er	nd-o	f-fra	er of ame ae cl	wai	t (El	FW)	per	iod
									per affe	ods cte	s to i	inse VS	rt a	ter t n pa	he ssiv	end- e m	of-f	rám . Th	e. T	hè t ie c	ime lock	for doe	whi	ch L	_FC	LK	ber of the	ssei	rted	is n	ot	
									Line	es p	er p	ane	el:																			
		9:0)			LF	PP															e-pa					repr					

7.6.4 LCD Controller Control Register 3 (LCCR3)

LCCR3, shown in Table 7-6, contains bits and bit fields used to control various functions within the LCD controller.

the whole LCD display. Lines per panel = (LPP+1).

number of lines on the LCD display. For dual-panel mode, it is half the number of lines on

Double Pixel Clock (DPC) — doubles the rate of the pixel clock on the L_PCLK pin. This allows direct connection to an NTSC encoder (e.g., the Analog Devices 7171). All of the LCD controller settings are still specified in terms of the "original" pixel clock and this mode affects *only* the L_PCLK output pin. If DPC is set to 1, the pixel clock divisor (PCD) must be greater than or equal to 1.

Bits Per Pixel (BPP) — BPP specifies the size of encoded pixel values in memory. Pixel sizes of 1, 2, 4, and 8 bits require that the internal palette RAM be loaded before pixels can be displayed on the screen. See Section 7.6.5 for details on programming the DMAC to load the palette RAM. BPP is programmed as follows:

0b000 = 1-bit pixels

0b001 = 2-bit pixels

0b010 = 4-bit pixels



0b011 = 8-bit pixels 0b100 = 16-bit pixels 0b101-0b111 = reserved

Output Enable Polarity (OEP) — In active display mode (LCCR0[PAS] = 1), the OEP bit selects the active and inactive states of the output enable signal (L_BIAS). In this mode, the AC bias pin serves as an enable that signals the off-chip device when data is actively being driven using the pixel clock, which continuously toggles in active mode. When OEP = 0, L_BIAS is active high and inactive low. When OEP = 1, L_BIAS is active low and inactive high. When L_BIAS is in its active state, data is driven onto the LCD data pins on the programmed edge of the pixel clock.

In passive display mode, OEP does not affect L_BIAS.

Pixel Clock Polarity (PCP) — selects the edge of the pixel clock (L_PCLK) on which data is sampled at the LCD pins. When PCP = 0, sampling occurs on the rising edge of L_PCLK. When PCP = 1, sampling occurs on the falling edge. PCP does not affect the timing of data being driven, it simply inverts L_PCLK.

Horizontal Sync Polarity (HSP) — selects the active and inactive states of the L_LCLK pin. When HSP = 0, L_LCLK is active high and inactive low. When HSP = 1, it is active low and inactive high. In active display mode, L_LCLK serves as the horizontal sync signal and in passive display mode, it is the line clock.

In **both** active and passive display modes, the L_FCLK pin is forced to its inactive state whenever pixels are transmitted. After the end of each line and a programmable number of pixel clocks occur (controlled by LCCR1[ELW]), the L_FCLK pin is forced to its active state for a programmable number of line clocks (controlled by LCCR1[HSW]), and is then again forced to its inactive state.

Vertical Sync Polarity (VSP) — selects the active and inactive states of the L_FCLK pin. When VSP = 0, L_FCLK is active high and inactive low. When VSP = 1, L_FCLK is active low and inactive high.

In **active display** mode (LCCR0[PAS] = 1), L_FCLK serves as the vertical sync signal. It is forced to its inactive state while pixels are transmitted during the frame. After the end of the frame and a programmable number of line clocks occur (controlled by LCCR2[EFW]), it is forced to its active state for a programmable number of line clocks (controlled by LCCR2[VSW]), and is then again forced to its inactive state.

In **passive display** mode, L_FCLK serves as the frame clock. It is forced to its active state on the rising edge of the first pixel clock of each frame. It remains active during the transmission of the entire first line of pixels in the frame and is then forced back to its inactive state on the rising edge of the first pixel clock of the second line of the frame. It remains at this state through the end of the frame.

AC Bias Pin Transitions Per Interrupt (API) — specifies the number of AC bias pin (L_BIAS) transitions to count before setting the AC bias count status (ACS) bit in the LCD Controller Status Register (LCSR), which signals an interrupt request. After the LCD controller is enabled, the value in API is loaded to a 4-bit down counter, and the counter decrements each time L_BIAS is inverted. When the counter reaches zero, it stops, and the AC bias count bit, LCSR[ABC], is set. Once ABC is set, the 4-bit down counter is reloaded with the value in API and is disabled until ABC is cleared. When ABC is cleared by the CPU, the down counter is enabled and again decrements each time the AC bias pin is inverted. The number of AC bias pin transitions between each interrupt request ranges from 1 to 15. Setting API to 0x0 disables the API function.



In **active display** mode (LCCR0[PAS] = 1), L_BIAS is the output enable signal. However, signalling of the API interrupt may still occur. The ACB bit field can be used to count line clock pulses in active mode. When the programmed number of line clock pulses occurs, an internal signal is toggled that is used to decrement the 4-bit counter used by the API interrupt logic. Once this internal signal toggles the programmed number of times, as specified by API, an interrupt is generated. The user must program API to zero if the API interrupt function is not required in active mode.

AC Bias Pin Frequency (ACB) — In passive display mode (LCCR0[PAS] = 1), the 8-bit ACB field specifies the number of line clocks to count between each toggle of the AC bias pin (L_BIAS). After the LCD controller is enabled, the value in ACB is loaded to an 8-bit down counter, which begins to decrement using the line clock (L_LCLK). When the counter reaches zero, it stops, L_BIAS is toggled, and the whole procedure starts again. The number of line clocks between each bias pin transition ranges from 1 to 256, corresponding to ACB values from 0 to 255. Thus, the value to program into ACB is the desired number of line clocks minus 1.

AC bias is used by a passive LCD display to periodically reverse the polarity of the power supplied to the screen in order to eliminate D.C. offset. If the LCD display being controlled has its own internal means of switching its power supply, set ACB to its maximum value (0xFF) to reduce power consumption. ACB must be programmed conservatively in a system with bandwidth problems that result in output FIFO underruns in the LCD Controller. In these cases, the pixel clock is stalled for passive displays, which can result in more time between line clocks than expected. See Section 7.3.5 for more information on how output FIFO underruns are handled.

In **active display** mode, the ACB bit field has no effect on the L_BIAS pin. Because the pixel clock toggles continuously in active mode, the AC bias pin is used as an output enable signal. It is asserted automatically by the LCD controller in active mode whenever pixel data is driven out to the data pins to signal the display when it may latch pixels using the pixel clock. ACB can be used in active mode to count line clocks and generate API interrupts.

Pixel Clock Divider (PCD) — selects the frequency of the pixel clock (L_CLK). PCD can be any value from 0 to 255. It generates a range of pixel clock frequencies from LCLK/2 to LCLK/512, where LCLK is the programmed frequency of the LCD/Memory Controller clock. LCLK can vary from 100MHz to 166 MHz.

The pixel clock frequency must be adjusted to meet the required screen refresh rate, which depends on:

- number of pixels for the target display
- number of panels (single or dual)
- display type (monochrome or color)
- number of pixel clock wait states programmed at the beginning and end of each line
- number of line clocks inserted at the beginning and end of each frame
- width of the VSYNC signal in active mode or VSW line clocks inserted in passive mode
- width of the frame clock or HSYNC signal.

All of these factors alter the time duration from one frame transmission to the next. Different display manufacturers require different frame refresh rates, depending on the physical characteristics of the display. PCD is used to alter the pixel clock frequency in order to meet these requirements. The frequency of the pixel clock for a set PCD value or the required PCD value to yield a target pixel clock frequency can be calculated using the two following equations. If double pixel clock mode (DPC) is enabled, PCD must be set greater than or equal to 1.



$$PixelClock = \frac{LCLK}{2(PCD+1)}$$

$$PCD = \frac{LCLK}{2(PixelClock)} - 1$$

where

LCLK = LCD/Memory Clock

PCD = LCCR3[7:0]

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 7-6. LCCR3 Bit Definitions (Sheet 1 of 2)

Physical Address 0x4400_000C

LCD Controller Control Register 3

LCD Controller

Bits	Name	Description
31:28	-	reserved
		Double Pixel Clock mode:
27	DPC	0 = The L_PCLK pin is driven at the frequency specified by PCD. 1 = The L_PCLK pin is driven at double the frequency specified by PCD.
		Bits Per Pixel:
26:24	ВРР	000 – 1-bits/pixel [4 entry, 8 byte palette buffer (only first 2 entries are used)] 001 – 2-bits/pixel [4 entry, 8 byte palette buffer] 010 – 4-bits/pixel [16 entry, 32 byte palette buffer] 011 – 8-bits/pixel [256 entry, 512 byte palette buffer] 100 – 16-bits/pixel [no palette buffer] 101, 110, 111 – reserved
		Output Enable Polarity:
23	OEP	0 = L_BIAS pin is active high and inactive low in active display mode. 1 = L_BIAS pin is active low and inactive high in active display mode.
		In active display mode, data is driven out to the LCD's data pins on the programmed pixel clock edge when the L_BIAS pin is active. OEP is ignored in passive display mode.
		Pixel Clock Polarity:
22	PCP	0 = Data is sampled on the LCD data pins on the rising edge of L_PCLK. 1 = Data is sampled on the LCD data pins on the falling edge of L_PCLK.
		Horizontal Sync Polarity:
21	HSP	0 = L_LCLK pin is active high and inactive low. 1 = L_LCLK pin is active low and inactive high.
		Vertical Sync Polarity:
20	VSP	0 = L_FCLK pin is active high and inactive low. 1 = L_FCLK pin is active low and inactive high.



Table 7-6. LCCR3 Bit Definitions (Sheet 2 of 2)

14	DIC	. 0.		.			U		011	٠, د	J	,,,,	_	,, <u>,</u>	,																
		ı		ical 4400		dres 00C	S			LC	D C	ont	rolle	er C	ont	rol	Re	giste	er 3				ı	LCI	D Co	ontro	oller				
Bit	31	30 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	r	eserve	ed	DPC		BPP	_	OEP	PCP	HSP	VSP		A	PI					A	СВ							PC	D			
Reset	X	х х	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bits			Na	me												D	esc	ript	ion										
								AC	bia	s Pir	n tra	nsit	tions	ре	r Int	erru	pt.														
	AC bias Pin transitions per Interrupt. Value (0–15) is used to specify the number of AC bias pin transitions to count before setting the line count status (ABC) bit, signalling an interrupt request. The counter is frozen when ABC is set and is restarted when ABC is cleared by software. This function is disabled when API=0x0. AC Bias pin frequency:														_																
		15:8			A	СВ		In p cloc pola disp (0x) acti	eass cks arity clay FF) ve r	ive- to co of t bein to co	dispount he p ng c ons e (L	blay bow conti erve	mod fore er su rolle e pov R0[F	de (l tog uppl d do wer PAS	gling y to bes AC] = 1	g the pre not B ca 1).	e A ver nee an l	S]=0 C bia nt D.0 ed to be us	as p C. c use sed	in hare E L_ in c	Γĥis ge bι BIAS onju	pin uildu S, pi incti	is us up w rogra ion v	sed ithir am	to p n the ACE	erio e dis 3 to i	dica play ts m	lly ir . If tl axir	nver he p num	t the ass	ive ue
		7:0			P	CD		Valu Cor LCI	ue (ntrol _K/5	ller o 512.	55) cloc	is us k (L	sed t	() fre	eque	ency	. T	requ he P CD+	ixel												

7.6.5 LCD Controller DMA

The LCD controller has two fully independent DMA channels used to transfer both the palette buffer and the frame buffer from external memory to the LCD controller. The LCD DMA controller (DMAC) behaves much like the processor DMAC in descriptor fetch mode. DMA channel 0 is used for single-panel display mode and the upper screen in dual-panel mode. DMA channel 1 is used exclusively for the lower screen in dual-panel mode. The palette RAM is always loaded through DMA channel 0. All of the information for the DMA transfers is maintained in registers within the LCD DMAC. These registers are loaded from *frame descriptors* located in main memory. One descriptor must be used per frame buffer in memory. A separate descriptor is also used when the palette RAM is loaded. Multiple descriptors can be chained together in a list, making it possible for the DMAC to transfer data from an essentially infinite number of discontiguous locations. The four DMA register types are numbered according to the associated DMA channel: See Section 7.6.5.2, Section 7.6.5.3, Section 7.6.5.4, and Section 7.6.5.5 for more information.

PCD must be programmed with a value of 1 or greater if Double Pixel Clock Mode is

7.6.5.1 Frame Descriptors

Although the FDADRx registers are loaded by software, the other DMA registers can only be loaded indirectly from DMA frame descriptors. A frame descriptor is a four-word block, aligned on a 16-byte boundary, in main memory:

word[0] contains the value for FDADRx

enabled.



word[1] contains the value for FSADRx word[2] contains the value for FIDRx word[3] contains the value for LDCMDx

Software must write the location of the first descriptor to FDADRx before enabling the LCD controller. Once the controller is enabled, the first descriptor is read, and all four registers are written by the DMAC. The next frame descriptor pointed to by FDADRx is loaded into the registers for the associated DMA channel after all data for the current descriptor has been

transferred.

The address in FDADRx is not used when the BRA bit in the Frame Branch Register (FBRx) is set.

In this case, the Frame Branch Address is used to fetch the descriptor for the next frame. Branches can be used to load a new palette or to process a regular frame, as detailed in Section 7.6.6.

Note: If only one frame buffer is used in external memory, the FDADRx field (word[0] of the frame descriptor) must point back to itself.

7.6.5.2 LCD DMA Frame Descriptor Address Registers (FDADRx)

FDADR0 and FDADR1, shown in Table 7-7, correspond to DMA channels 0 and 1 and contain the memory address of the next descriptor for the DMA channel. The DMAC fetches the descriptor at this location after finishing the current descriptor. On reset, the bits in this register are undefined. The target address must be aligned to a 16-byte boundary. Bits [2:0] of the address must be zero.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 7-7. FDADRx Bit Definitions

Physical Address FDADR0 channel 0: 0x4400_0200 **LCD Controller** FDADR1 channel 1: 0x4400 0210 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 **Descriptor Address** Reset ? ? ? ? ? ? ? ? **Bits** Name **Description** Descriptor Address of next descriptor. 31:0 Address Bits [2:0] must be zero for proper memory alignment.

7.6.5.3 LCD DMA Frame Source Address Registers (FSADRx)

FSADR0 and FSADR1, shown in Table 7-8, correspond to DMA channels 0 and 1 and contain the source address of the current descriptor for the DMA channel. The address must be aligned on an 8-byte boundary. Bits [2:0] must be zero. If this descriptor is a palette load, FSADRx points to the memory location at the beginning of the palette data. The size of the palette data must be four 16-bit entries for 1- and 2-bit pixels, sixteen 16-bit entries for 4-bit pixels, or 256 16-bit entries for 8-bit pixels. If this descriptor is for pixel data, FSADRx points to the beginning of the frame buffer in memory. This address is incremented as the DMAC fetches from memory. If desired, the DMA Frame ID Register can be used to hold the initial frame source address.

These read-only registers are loaded indirectly via the frame descriptors, as described in Section 7.6.5.1.



These are read-only registers. Ignore reads from reserved bits.

Table 7-8. FSADRx Bit Definitions

Physical Address FSADR0 LCD Controller channel 1: 0x4400_0214 FSADR1

Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Frame Source Address

Bits	Name	Description
31:0	Source	Address of the palette or pixel frame data in memory. Bits [2:0] must be zero for proper memory alignment.

7.6.5.4 LCD DMA Frame ID Registers (FIDRx)

FIDRx, shown in Table 7-9, correspond to DMA channels 0 and 1 and contain an ID field that describes the current frame. The particular use of this field is up to the user. This ID register is copied to the LCD Controller Interrupt ID Register when an interrupt occurs.

These read-only registers are loaded indirectly via the frame descriptors, as described in Section 7.6.5.1.

These are read-only registers. Ignore reads from reserved bits.

Table 7-9. FIDRx Bit Definitions

Physical Address FIDR0 LCD Controller channel 1: 0x4400_0218 FIDR1

Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Bits	Name	Description
31:3	Frame ID	Frame ID.
2:0	_	reserved



7.6.5.5 LCD DMA Command Registers (LDCMDx)

LDCMDx, shown in Table 7-10, correspond to DMA channels 0 and 1 and contain configuration fields and the length of the current descriptor for the DMA channel. On reset, the bits in these register are initialized to zero. Reserved bits must be written with zeros and reads from reserved bits must be ignored.

These read-only registers are loaded indirectly via the frame descriptors, as described in Section 7.6.5.1.

Load Palette (PAL) — indicates that data being fetched will be loaded into the palette RAM. If PAL is set to one, the palette RAM is loaded with the first 8, 32, or 512-bytes of data as follows:

8 bytes for 1 and 2-bit pixels

32 bytes for 4-bit pixels

512 bytes for 8-bit pixels.

Software must load the palette at least once after enabling the LCD. Otherwise, the palette entries will not be initialized, and the frame data will not have a valid frame palette to reference.

The palette must not be loaded if the LCD is operating in 16-bit pixel mode.

Note: The PAL bit must never be set in LDCMD1, since the palette is always loaded with Channel 0.

Start Of Frame Interrupt (SOFINT) — when set, the DMAC sets the start of frame bit (LCSR[SOF]) when starting a new frame. The SOF bit is set after a new descriptor is loaded from memory and before the palette/frame data is fetched.

In dual-panel mode, LCSR[SOF] is set only when both channels reach the start of frame and both frame descriptors have SOFINT set. SOFINT must not be set for palette descriptors in dual-panel mode, since only one channel is ever used to load the palette RAM.

End Of Frame Interrupt (EOFINT) — when set, the DMAC sets the end of frame bit (LCSR[EOF]) after fetching the last word in the frame buffer.

In dual-panel mode, LCSR[EOF] is set only when both channels reach the end of frame and both frame descriptors have EOFINT set. EOFINT must not be set for palette descriptors in dual-panel mode, since only one channel is ever used to load the palette RAM.

Transfer Length (LEN) — determines the number of bytes fetched by the DMAC. LEN = 0 is not valid. If PAL is set to one, LEN must be programmed with the size of the palette RAM. This corresponds to:

8 bytes for 1 and 2-bit pixels (only the top 2 entries are actually used for 1-bit pixels)

32 bytes for 4-bit pixels

512 bytes for 8-bit pixels.

Note: A separate descriptor must be used to fetch the frame data.

The value of LEN for frame data is a function of the screen size and the pixel size and it must be consistent with the values used for LCCR1[PPL], LCCR2[LPP], and LCCR3[BPP]. See Section 7.4.2 for instructions on calculating length. The LCD DMAC decrements LEN as it fetches data, allowing the user to read the number of bytes remaining for the current descriptor.

These are read-only registers. Ignore reads from reserved bits.



Table 7-10. LDCMDx Bit Definitions

Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

		res	serv	ed		PAL	res	serv	ed	SOFINT	EOFINT											LEN										
Reset	X	X	X	X	X	0	X	X	X	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Bits	Name	Description
31:27	_	reserved
26	PAL	Load Palette: 0 = DMA in progress is <i>not</i> the palette buffer. 1 = DMA in progress is the palette buffer. PAL must not be set in LDCMD1.
25:23	_	reserved
22	SOFINT	Start of Frame Interrupt: 0 = Do not set the SOF interrupt bit in the LCD status register when starting a new frame. 1 = Set the start of frame (SOF) interrupt bit in the LCD status register when starting a new frame (after loading the frame descriptor).
21	EOFINT	 End of Frame Interrupt: 0 = Do not set the EOF interrupt bit in the LCD status register when finished fetching the last word of this frame. 1 = Set the end of frame (EOF) interrupt bit in the LCD status register when finished fetching the last word of this frame.
20:0	LEN	Length of transfer in bytes: The two lowest bits [1:0] are part of the length calculation but must always be zero for proper memory alignment. LEN = 0 is illegal.



7.6.6 LCD DMA Frame Branch Registers (FBRx)

FBRx, one for each DMA channel, shown in Table 7-11, contain the addresses, aligned on a 4-byte boundary, of the descriptors to branch to.

When BRA is set, the Frame Descriptor Address Register is ignored. The next descriptor is fetched from the address in FBRx[31:4], regardless of whether frame data or palette RAM data is being processed. Setting BINT to one forces the DMAC to set the Branch Status interrupt bit (BS) in the LCD Controller Status Register after fetching the branched-to descriptor. BRA is automatically cleared by hardware when the branch is taken.

Note: In dual-panel mode, both FBR0 and FBR1 must be written in order to branch properly.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 7-11. FBRx Bit Definitions

Physical Address channel 0: 0x4400_0020 channel 1: 0x4400_0024

LCD DMA Frame Branch Registers

LCD Controller

Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

Frame Branch Address

Bits	Name	Description
31:4	Frame Branch Address	Frame Branch Address: Address of the descriptor for the branched-to frame.
3:2	_	reserved
1	BINT	Branch Interrupt: 0 = Do not set the BS interrupt bit in register LCSR after the branched-to descriptor is loaded. 1 = Set the BS interrupt bit in register LCSR after the branched-to descriptor is loaded.
0	BRA	Branch: 0 = Do not branch after finishing the current frame. 1 = Branch after finishing the current frame. The next descriptor will be fetched from the Frame Branch Address. BRA is automatically cleared after loading the new descriptor.



7.6.7 LCD Controller Status Register (LCSR)

LCSR, shown in Table 7-12, contains bits that signal:

- Underrun errors for both the input and output FIFOs
- AC bias pin transition count
- LCD disable and quick disable
- DMA start/end frame and branch status
- DMA transfer bus error conditions.

Unless masked, each of these hardware-detected events signals an interrupt request to the interrupt controller. Two bits, BER and ABC, generate nonmaskable interrupts.

Each of the LCD's status bits continues to signal an interrupt request for as long as the bit is set. Once the bit is cleared, the interrupt is cleared. Status bits are referred to as sticky (once set by hardware, they must be cleared by software). Writing one to a sticky status bit clears it. Writing zero has no effect. All LCD interrupts can be masked by programming the Interrupt Controller Mask Register (ICMR). See Section 4.2, "Interrupt Controller" on page 4-20 for more details.

Subsequent Interrupt Status (SINT) — set when an unmasked interrupt occurs and there is already a pending interrupt. The frame ID of the first interrupt is saved in the LCD controller interrupt ID register (LIIDR). SINT is set only for bus error, start of frame, end of frame, and branch status interrupts.

Note: If a branched-to descriptor has SOF set, both the SOF and branch interrupts are signalled at the same time, and SINT is not set.

Branch Status (BS) — set after the DMA controller has branched and loaded the descriptor from the frame branch address in the frame branch register, and the branch interrupt (BINT) bit in the frame branch register is set. When BS is set, an interrupt request is made to the interrupt controller if it is unmasked (LCCR0[BM] = 0).

In dual-panel mode (LCCR0[SDS = 1]), both DMA channels are enabled, and BS is set only after both channels' frames have been fetched. BS remains set until cleared by software.

End Of Frame Status (EOF) — set after the DMA controller has finished fetching a frame from memory and that frame's descriptor has the end-of-frame interrupt bit set (LDCMDx[EOFINT] = 1). When EOF is set, an interrupt request is made to the interrupt controller if it is unmasked (LCCR0[EFM] = 0).

When dual-panel mode is enabled (LCCR0[SDS] = 1), both DMA channels are enabled, and SOF is set only after both channels' frames have been fetched. EOF remains set until cleared by software.

LCD Quick Disable Status (QD) — set when LCD Enable (LCCR0[ENB]) is cleared and the DMA controller finishes any current data burst. When QD is set, an interrupt request is made to the interrupt controller if it is unmasked (LCCR0[QDM] = 0). This forces the LCD controller to stop immediately and quit driving the LCD pins. Quick disable is intended for use with Sleep shutdown.

Output FIFO Underrun Status (OU) — set when an output FIFO is completely empty and the LCD's data pin driver logic attempts to fetch data from the FIFO. It is cleared by writing one to the bit. OU is used for single- and dual-panel displays. In dual-panel mode (LCCR0[SDS] = 1), both FIFOs are filled and emptied at the same time, so that underrun occurs at the same time for both



panels. When OU is set, an interrupt request is made to the interrupt controller if it is unmasked (LCCR0[OUM] = 0). Output FIFO underruns are more important that Input FIFO underruns, because they affect the panel.

Input FIFO Underrun Upper Panel Status (IUU) — set when the upper panel's input FIFO is completely empty and the LCD controller's pixel unpacking logic attempts to fetch data from the FIFO. It is cleared by writing one to the bit. IUU is used in both single-panel (LCCR0[SDS] = 0) and dual-panel (SDS = 1) modes. When IUU is set, an interrupt request is made to the interrupt controller if it is unmasked (LCCR0[IUM] = 0).

Input FIFO Underrun Lower Panel Status (IUL) — used only in dual-panel mode (LCCR0[SDS] = 1) and is set when the lower panel's input FIFO is completely empty and the LCD controller's pixel unpacking logic attempts to fetch data from the FIFO. It is cleared by writing one to the bit. When IUL is set, an interrupt request is made to the interrupt controller if it is unmasked (LCCR0[IUM]=0).

AC Bias Count Status (ABC) — set each time the AC bias pin (L_BIAS) toggles the number of times specified in the AC bias pin transitions per interrupt (API) field in LCCR3. If API is programmed with a non-zero value, a counter is loaded with the value in API and is decremented each time L_BIAS toggles. When the counter reaches zero, ABC is set, which signals an interrupt request to the interrupt controller. The counter reloads using the value in API but does not start to decrement again until ABC is cleared by software.

Bus Error Status (BER) — set when a DMA transfer causes a system bus error. The error is signalled when the DMA controller attempts to access a reserved or nonexistent memory space. When this occurs, the DMA controller stops and remains halted until software installs a valid memory address into the FDADRx register. In dual-channel mode, both channels are stopped. FDADR0 and FDADR1 must be rewritten to continue LCD operation. BER remains set until cleared by software.

Start Of Frame Status (SOF) — set after the DMA controller has loaded a new descriptor and that descriptor has the start of frame interrupt bit set (LDCMDx[SOFINT] = 1). When SOF is set, an interrupt request is made to the interrupt controller if it is unmasked (LCCR0[SFM] = 0). In dual-panel mode (LCCR0[SDS] = 1), both DMA channels are enabled, and SOF is set only after both channels' descriptors have been loaded. SOF remains set until cleared by software.

LCD Disable Done Status (LDD) — set by hardware after the LCD has been disabled and the frame that is active has been sent to the LCD data pins. When the LCD controller is disabled by setting the LCD disable bit in LCCR0, the current frame is completed before the controller is disabled. After the last set of pixels is clocked out onto the LCD data pins by the pixel clock, the LCD controller is disabled, LDD is set, and an interrupt request is made to the interrupt controller if it is unmasked (LCCR0[LDM] = 0). LDD remains set until cleared by software.

Performing a quick disable by clearing LCCR0[ENB] does not set LDD.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 7-12. LCSR Bit Definitions (Sheet 1 of 2)

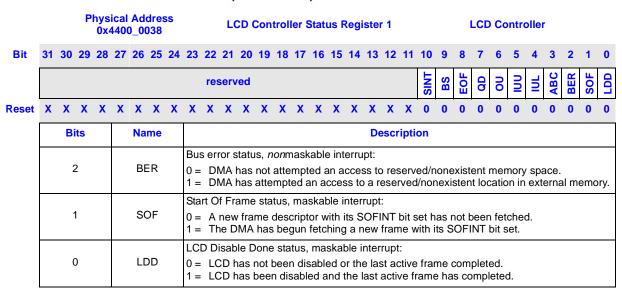
Physical Address 0x4400_0038 LCD Controller Status Register 1 LCD Controller

Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										res	serv	ed										SINT	BS	EOF	۵D	OU	IUU	IUL	ABC	BER	SOF	LDD
Reset	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	Y	0	0	0	0	0	0	0	0	0	0	0

Bits	Name	Description
31:11	_	reserved
		Subsequent Interrupt status, maskable interrupt:
10	SINT	 0 = A second unmasked branch, start of frame, end of frame, or bus error interrupt has NOT occurred before a previous interrupt has completed. 1 = A second unmasked branch, start of frame, end of frame, or bus error interrupt has occurred before the previous interrupt has been cleared. The value in the Interrupt Frame ID Register is not replaced with the value from the second interrupt.
		Branch Status, maskable interrupt:
9	BS	 The DMA has not loaded a branched-to descriptor, or the DMA has loaded a branched-to descriptor, but the branch interrupt (BINT) bit is not set in the Frame Branch Register. The DMA has loaded a branched-to descriptor, and the BINT bit is set.
		End Of Frame status, maskable interrupt:
8	EOF	 0 = A new frame with the EOFINT bit set in its descriptor has not been processed. 1 = The DMA has finished fetching a frame with the EOFINT bit set in its descriptor.
		LCD Quick Disable status, maskable interrupt:
7	QD	0 = LCD has not been quickly disabled by clearing LCCCR0[ENB].1 = LCD has been quickly disabled.
		Output FIFO Underrun status, maskable interrupt:
6	OU	 0 = Output FIFOs have not underrun. 1 = LCD dither logic is not supplying data to output FIFOs for the panel at a sufficient rate. The output FIFOs have completely emptied.
		Input FIFO Underrun Upper panel status, maskable interrupt:
5	Ιυυ	 0 = The input FIFO for the upper (dual-panel mode) or whole panel (single-panel mode) display has not underrun. 1 = DMA is not supplying data to the input FIFO for the upper or whole panel at a sufficient rate. The FIFO has completely emptied, and the pixel unpacking logic has attempted to take data from the FIFO.
		Input FIFO Underrun Lower panel status, dual-panel mode only, maskable interrupt:
4	IUL	 0 = The input FIFO for the lower panel display has not underrun. 1 = DMA is not supplying data to the input FIFO for the lower panel at a sufficient rate. The FIFO has completely emptied, and the pixel unpacking logic has attempted to take data from the FIFO.
		AC Bias Count status, <i>non</i> maskable interrupt:
3	ABC	 0 = The AC bias transition counter has not decremented to zero. 1 = The AC bias transition counter has decremented to zero, indicating that the L_BIAS pin has toggled the number of times specified by the LCCR3[API] control-bit field. The counter is reloaded with the value in API but is disabled until the user clears ABC.



Table 7-12. LCSR Bit Definitions (Sheet 2 of 2)



7.6.8 LCD Controller Interrupt ID Register (LIIDR)

LIIDR, shown in Table 7-13, contains a copy of the Frame ID Register (FIDR) from the descriptor currently being processed when a start of frame (SOF), end of frame (EOF), branch (BS), or bus error (BER) interrupt is signalled. LIIDR is written to only when an unmasked interrupt of the above type is signalled and there are no other unmasked interrupts in the LCD controller pending. As such, the register is considered to be sticky and will be overwritten only when the signalled interrupt is cleared by writing the LCD controller status register. Except for a bus error, in dual panel mode LIIDR is written only when both channels have reached a given state. LIIDR is written with the last channel to reach that state. (i.e. FIDR of the last channel to reach SOF would be written in LIIDR if SOF interrupts are enabled).

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 7-13. LIICR Bit Definitions

						Ad 0_00		SS				LC	D C		rolle Reg			rupt	ID					I	LCD	Co	ntro	olle	r			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														IFR	AM	EID														res	erv	ed
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	X	X	X
	Bits				Na	me												De	escr	ipti	on											
31:3				IF	RA	MEI	D	Inte	errup	ot Fr	rame	e ID																				
					-	-		res	erve	ed																						



7.6.9 TMED RGB Seed Register (TRGBR)

TRGBR, shown in Table 7-14 contains the three (red, green, blue) eight-bit seed values used by the TMED algorithm. This value is added into the modified pixel data value as an offset in creating the lower boundary for the algorithm. These values are used during the dithering process for passive (DSTN) displays. The default, recommended setting is 0x00AA5500. This setting provides superior display results in most cases.

This is a write-only register. Write zeros to reserved bits.

Table 7-14. TRGBR Bit Definitions

·		•		• • • •		٠.,	٥.,		,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,	,,,,																		
			P	hys 0x4		Ad 0_0		SS			TMEC	RG	B Se	ed Re	gister					I	LCE	Co	ontro	oll	ler			
Bit	31	30	29	28	27	26	25	24	23 22	21	20 19	18	17 1	6 15	14 1	3 12	11	10	9	8	7	6	5	4	4 3	2	1	0
			r	ese	rve	d					TBS					T	GS							1	TRS			
Reset	x x x x x x x							X			0xAA					0:	(55							0	0x00			
		В	its			Na	me								I)esc	ript	ion										
		31	:24			-	_		reserve	ed																		
	23:16 TBS								TME B	lue S	Seed va	alue																
	15:8 TGS								TME G	reen	Seed	valu	ie															
	7:0 TRS								TME R	ed S	Seed Va	alue																



7.6.10 TMED Control Register (TCR)

TCR, shown in Table 7-15, selects various options available in the TMED dither algorithm. There are two Temporal Modulated Energy Distribution algorithms that can be used. The default, recommended setting is 0x0000754F. This setting provides superior display results in most cases.

For more details on the effects of the individual fields within this register, please refer to Section 7.3.3.

TMED Energy Distribution Select (TED) — selects which matrix is used in the final step of TMED algorithm. TMED = 1 selects the (preferred) TMED2 matrix. TMED = 0 selects the older TMED matrix. After the pixel value has gone through the algorithm to determine a lower and upper boundary, the row and column counters are combined and run through one of the matrices to obtain a number that will be compared to the 2 boundaries. If that number is between the 2 boundaries, then the data out for this pixel in this frame is a 1, otherwise it is a 0.

TMED Horizontal Beat Suppression (THBS) — is the column shift value used as an offset that is combined with the row (line) counter and the pixel counter to create an address to lookup in the matrix. The matrix output is compared to the upper and lower boundaries defined in Section 7.3.3.

TMED Vertical Beat Suppression (TVBS) — is the block shift value used as an offset that is combined with the pixel counter.

TMED Frame Number Adjuster Enable (FNAME) — allows the frame number adjuster to add an offset to the current frame number before the value is sent through the algorithm. Setting this bit enables the addition of the current frame number to a value composed from the row and column counters. This value comes from one of the two look up matrices which is selected by TMED[FNAM].

TMED Color Offset Adjuster Enable (COAE) — enables the color offset adjuster for each color. The color offset adjuster creates the offset in the lower boundary in the TMED algorithm (refer to Section 7.3.3). The Offset is created by adding either the output of the lookup matrix (input was the Color Value) or '00' to the Seed value in the TSR for that color. The color offset adjuster for each color can be disabled by clearing this bit. When cleared, this bit allows only the Seed Register value to go through the algorithm.

TMED Frame Number Adjuster Matrix (FNAM) — selects which matrix is used when using the frame number adjuster. A 1 will select the (recommended) TMED2 matrix, and a 0 will select the older TMED matrix.

TMED Color Offset Adjuster Matrix (COAM) — selects which matrix is used when using the color offset adjuster. A 1 will select the (recommended) TMED2 matrix, and a 0 will select the older TMED matrix.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 7-15. TCR Bit Definitions

		P				ddres 0044	SS	TMED Control Register 4 23 22 21 20 19 18 17 16 15 14 13 1														LCE	Co	ontro	ollei	r					
Bit	31	30 29	28	27	26	25	24	23	22	21	20	19 1	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•							res	serv	ed								TED	reserved			TH	BS			TV	BS		FNAME	COAE	FNAM	COAM
Reset	X	х х	X	X	X	X	X	X	X	X	X	X	X	X	X	X	1	1	1	0	1	0	1	0	1	0	0	1	1	1	1
•		Bits			Na	ame												De	escr	ipti	on										
		31:15						res	erve	ed																					
•		14 TED							Se	elect	s Ma	Distri atrix atrix	1	tion	Mat	trix :	Sele	ct													
•		13:12	_		res	erve	ed																								
•		11:8 THBS										al Be			•		n														
•		7:4			T۱	/BS		TMED Vertical Beat Suppression Specifies the block shift value.																							
•		3			FN	AME		TMED Frame Number Adjuster Enable 0 = Disable frame number adjuster. 1 = Enable frame number adjuster.																							
•		2	C	DAE		TMED Color Offset Adjuster Enable 0 = Disable color offset adjuster. 1 = Enable color offset adjuster.																									
•	1 FNAM							TMED Frame Number Adjuster Matrix 0 = Selects Matrix 1 for frame number adjuster. 1 = Selects Matrix 2 for frame number adjuster.																							
•		0 COAM						TMED Color Offset Adjuster Matrix 0 = Selects Matrix 1 for color offset adjuster. 1 = Selects Matrix 2 for color offset adjuster.																							

7.7 LCD Controller Register Summary

Table 7-16 shows the registers associated with the LCD Controller and the physical addresses used to access them. All of the LCD registers must be accessed as 32-bit values.

Table 7-16. LCD Controller Register Summary (Sheet 1 of 2)

Address	Name	Description
0x4400_0000	LCCR0	LCD controller control register 0
0x4400_0004	LCCR1	LCD controller control register 1
0x4400_0008	LCCR2	LCD controller control register 2
0x4400_000C	LCCR3	LCD controller control register 3
0x4400_0020	FBR0	DMA channel 0 frame branch register



Table 7-16. LCD Controller Register Summary (Sheet 2 of 2)

Address	Name	Description
0x4400_0024	FBR1	DMA channel 1 frame branch register
0x4400_0038	LCSR	LCD controller status register
0x4400_003C	LIIDR	LCD controller interrupt ID register
0x4400_0040	TRGBR	TMED RGB Seed Register
0x4400_0044	TCR	TMED Control Register
0x4400_0200	FDADR0	DMA channel 0 frame descriptor address register
0x4400_0204	FSADR0	DMA channel 0 frame source address register
0x4400_0208	FIDR0	DMA channel 0 frame ID register
0x4400_020C	LDCMD0	DMA channel 0 command register
0x4400_0210	FDADR1	DMA channel 1 frame descriptor address register
0x4400_0214	FSADR1	DMA channel 1 frame source address register
0x4400_0218	FIDR1	DMA channel 1 frame ID register
0x4400_021C	LDCMD1	DMA channel 1 command register



intel® Synchronous Serial Port Controller

This chapter describes the Synchronous Serial Port Controller's (SSPC) signal definitions and operation for the PXA255 processor usage.

8.1 **Overview**

The SSPC is a full-duplex synchronous serial interface and can connect to a variety of external analog-to-digital (A/D) converters, audio and telecom codecs, and other devices that use serial protocols for transferring data. The SSPC supports National's Microwire*, Texas Instruments' Synchronous Serial Protocol* (SSP), and Motorola's Serial Peripheral Interface* (SPI) protocol.

The SSPC operates in master mode (the attached peripheral functions as a slave) and supports serial bit rates from 7.2 KHz to 1.84 MHz. Serial data formats may range from 4 to 16 bits in length. The SSPC provides 16 entries deep x 16 bits wide transmit and receive data FIFOs.

The FIFOs may be loaded or emptied by the Central Processor Unit (CPU) using programmed I/O, or DMA burst transfers of 4 or 8 half-words per transfer while receiving or transmitting.

8.2 **Signal Description**

This section describes the SSPC signals.

8.2.1 **External Interface to Synchronous Serial Peripherals**

Table 8-1 lists the external signals that connect the SSP to an external peripheral.

Table 8-1, External Interface to Codec

Name	Direction	Description
SSPSCLK	Output	Serial bit-rate clock
SSPSFRM	Output	Frame indicator
SSPTXD	Output	Transmit Data (serial data out)
SSPRXD	Input	Receive Data (serial data in)
SSPEXTCLK	Input	External clock which can be selected to drive the serial clock (SSPSCLK)

SSPSCLK is the bit-rate clock driven from the SSPC to the peripheral. SSPSCLK is toggled only when data is actively being transmitted and received.

SSPSFRM is the framing signal, indicating the beginning and the end of a serialized data word.

SSPTXD and SSPRXD are the Transmit and Receive serial data lines.



SSPEXTCLK is an external clock (input through GPIO27) that replaces the standard 3.6864 MHz clock used to generate the serial bit-rate clock (SSPSCLK). The external clock is internally divided by 2 and then further divided by the value in SSCR0[SCR].

If SSP operation is disabled, the five SSP pins are available for GPIO use. See Chapter 4, "System Integration Unit" for details on configuring pin direction and interrupt capabilities.

8.3 Functional Description

Serial data is transferred between the processor and an external peripheral through FIFO buffers in the SSPC. Data transfers to system memory are handled by either the CPU (using programmed I/O) or by DMA. Operation is full duplex - separate buffers and serial data paths permit simultaneous transfers to and from the external peripheral.

Programmed I/O transmits and receives data directly between the CPU and the transmit/receive FIFO's. The DMA controller transfers data during transmit and receive operations between memory and the FIFO's. DMA programming guidelines are found in Chapter 5, "DMA Controller".

8.3.1 Data Transfer

Transmit data is written by the CPU or DMA to the SSPC's transmit FIFO. The write takes the form of a programmed I/O or a DMA burst, with 4 or 8 half-words being transferred per burst. The SSPC then takes the data from the FIFO, serializes it, and transmits it via the SSPTXD signal to the peripheral. Data from the peripheral is received via the SSPRXD signal, converted to parallel words and is stored in the Receive FIFO. Read operations automatically target the receive FIFO, while write operations write data to the transmit FIFO. Both the transmit and receive FIFO buffers are 16 entries deep by 16 bits wide.

As the received data fills the receive FIFO, a programmable threshold triggers an interrupt to the Interrupt Controller. If enabled, an interrupt service routine responds by identifying the source of the interrupt and then performs one or several read operations from the inbound (receive) FIFO buffer.

8.4 Data Formats

The SSPC uses serial data formats to transfer and store data. This section describes the data formats.

8.4.1 Serial Data Formats for Transfer to/from Peripherals

Four signals are used to transfer data between the processor and external codecs or modems. Although there are three formats for serial data, they have the same basic structure:

- SSPSCLK—Defines the bit rate at which serial data is transmitted and received from the port.
- SSPSFRM–Depending on the transmission format selected, defines the boundaries of a data frame, or marks the beginning of a data frame.
- SSPTXD-Transmit signal for outbound data, from system to peripheral.



• SSPRXD–Receive signal for inbound data, from peripheral to system.

A data frame can be configured to contain from 4 to 16 bits. Serial data is transmitted most significant bit first.

The SSPC supports three formats: Motorola SPI, Texas Instruments SSP, and National Microwire. The three formats have significant differences, as described below.

SSPSFRM varies for each protocol as follows:

- For SPI and Microwire formats, SSPSFRM functions as a chip select to enable the external device (target of the transfer), and is held active-low during the data transfer.
- For SSP format, SSPSFRM is pulsed high for one serial bit-clock period at the start of each frame.

SSPSCLK varies for each protocol as follows:

- For Microwire, both transmit and receive data sources switch data on the falling edge of SSPSCLK, and sample incoming data on the rising edge.
- For SSP, transmit and receive data sources switch data on the rising edge of SSPSCLK, and sample incoming data on the falling edge.
- For SPI, the user has the choice of which edge of SSPSCLK to use for switching outgoing data, and for sampling incoming data. In addition, the user can move the phase of SSPSCLK, shifting its active state one-half period earlier or later at the start and end of a frame.

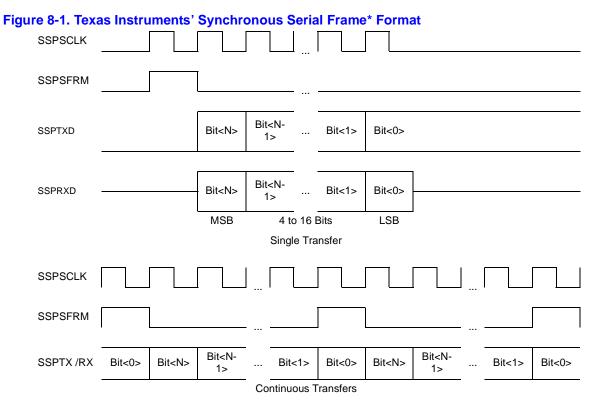
While SSP and SPI are full-duplex protocols, Microwire uses a half-duplex master-slave messaging protocol. At the start of a frame, a 1 or 2-byte control message is transmitted from the controller to the peripheral. The peripheral does not send any data. The peripheral interprets the message and, if it is a READ request, responds with requested data, one clock after the last bit of the requesting message. Return data (part of the same frame) can be from 4 to 16 bits in length. Total frame length is 13 to 33 bits.

The serial clock (SSPSCLK) only toggles during an active frame. At other times it is held in an inactive or idle state, as defined by its specified protocol.

8.4.1.1 SSP Format Details

When outgoing data in the SSP controller is ready to be transmitted, SSPSFRM is asserted for one clock period. On the following clock period, data to be transmitted is driven on SSPTXD one bit at a time, most significant bit first. Similarly, the peripheral drives data on the SSPRXD pin. Word length is from 4 to 16 bits. All transitions take place on the SSPSCLK rising edge and data is sampled on the falling edge. At the end of the transfer, SSPTXD retains the value of the last bit sent (bit 0) through the next idle period. If the SSP Port is disabled or reset, SSPTXD is forced to zero. Figure 8-1 shows the Texas Instruments' Synchronous Serial Frame* format for a single transmitted frame and when back-to-back frames are transmitted. When the bottom entry of the transmit FIFO contains data, SSPSFRM is pulsed high for one SSPSCLK period and the value to be transmitted is transferred from the transmit FIFO to the transmit logic's serial shift register. On the next rising edge of SSPSCLK, the most significant bit of the 4 to 16-bit data frame is shifted to the SSPTXD pin. Likewise, the most significant bit of the received data is shifted onto the SSPRXD pin by the off-chip serial slave device. Both the SSP and the off-chip serial slave device then latch each data bit into their serial shifter on the falling edge of each SSPSCLK. The received data is transferred from the serial shifter to the receive FIFO on the first rising edge of SSPSCLK after the last bit has been latched.





8.4.1.2 SPI Format Details

The SPI format has four sub-modes. The sub-mode used depends on the SSPSCLK edge selected for driving and sampling data and on the phase mode of SSPSCLK selected (see Section 8.7.2 for complete description of each mode).

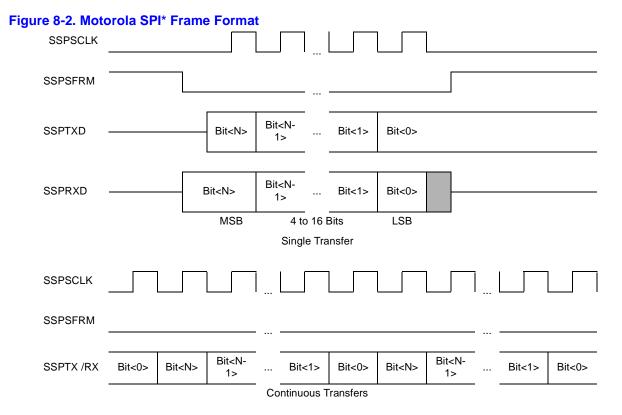
In idle mode or when the SSP is disabled, SSPSCLK and SSPTXD are low and SSPSFRM is high. When transmit (outgoing) data is ready, SSPSFRM goes low and stays low for the remainder of the frame. The most significant serial data bit is driven onto SSPTXD a half-cycle later, and halfway into the first bit period SSPSCLK asserts high and continues toggling for the remaining data bits. Data transitions on the configured SSPSCLK edge. From 4 to 16 bits may be transferred per frame.

When SSPSFRM is asserted, receive data is simultaneously driven from the peripheral on SSPRXD, most significant bit first. Data transitions on the configured SSPSCLK edge and is sampled by the controller on opposite edge. At the end of the frame, SSPSFRM is deasserted high one clock period after the last bit is latched at its destination and the completed incoming word is shifted into the incoming FIFO. The peripheral can tristate SSPRXD after sending the last bit of the frame. SSPTXD retains the last value transmitted when the controller goes into idle mode, unless the SSP port is disabled or reset (which forces SSPTXD to zero).

For back-to-back transfers, start and completion are similar to those of a single transfer but SSPSFRM does not deassert between words. Both transmitter and receiver know the word length and internally keep track of the start and end of words (frames). There are no dead bits. One frame's least significant bit is followed immediately by the next frame's most significant bit.



Figure 8-2 shows one of the four configurations for the Motorola SPI frame format for single and back-to-back frame transmissions.



Note: SSPSCLK's phase and polarity can be configured for four modes. This example shows one of those modes.

8.4.1.3 Microwire Format Details

Microwire format is similar to SPI, but it uses half-duplex transmissions with master-slave message passing rather than full-duplex. In idle state or when the SSP is disabled, SSPSCLK is low, SSPSFRM is high, and SSPTXD is low.

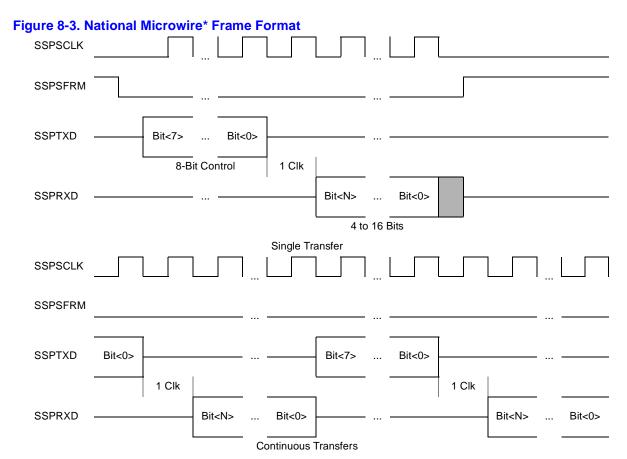
Each Microwire transmission begins with SSPSFRM assertion (low), followed by an 8 or 16-bit command word sent from controller to peripheral on SSPTXD. The command word data size is selected by the Microwire Transmit Data Size (MWDS) bit in SSP Control Register 1. SSPRXD is controlled by the peripheral and remains tristated. SSPSCLK goes high midway through the command's most significant bit and continues to toggle at the bit rate.

One bit-period after the last command bit, the peripheral must return the serial data requested, most significant bit first, on SSPRXD. Data transitions on SSPSCLK's falling edge and is sampled on the rising edge. SSPSCLK's last falling edge coincides with the end of the last data bit on SSPRXD and it remains low if it is the only or last word of the transfer. SSPSFRM deasserts high one-half clock period later.

The start and end of a series of back-to-back transfers are similar to those of a single transfer. However, SSPSFRM remains asserted (low) throughout the transfer. The end of a data word on SSPRXD is immediately followed by the start of the next command byte on SSPTXD.



Figure 8-3 shows the National Microwire frame format with 8-bit command words for single and back-to-back frame transmissions.



8.4.2 Parallel Data Formats for FIFO Storage

Data in the FIFOs is stored with one 16-bit value per data sample with no regard to the format's data word length. In each 16-bit field, the stored data sample is right-justified, the word's least significant bit is stored in bit 0, and unused bits are packed as zeroes above the most significant bit. Logic in the SSPC automatically left-justifies data in the Transmit FIFO so the sample is properly transmitted on SSPTXD in the selected frame format.



8.5 FIFO Operation and Data Transfers

Transmit and receive serial data use independent FIFOs. FIFOs are filled or emptied by programmed I/O or DMA bursts that the DMAC initiates. Bursts may be 4 or 8 half-words in length during transmission or reception.

8.5.1 Using Programmed I/O Data Transfers

Data words are 32 bits wide, but only 16-bit samples are transferred. Only the lower 2 bytes of a 32-bit word have valid data. The upper 2 bytes are not used and include invalid data that must be discarded.

The processor can fill or empty FIFOs in response to an interrupt from the FIFO logic. Each FIFO has a programmable interrupt threshold. When the threshold value is exceeded and an interrupt is enabled, an interrupt that signals the CPU to empty the receive FIFO or refill the transmit FIFO is generated.

The user can also poll the SSP Status Register (see Section 8.7.4) to determine how many samples are in a FIFO or whether the FIFO is full or empty.

8.5.2 Using DMA Data Transfers

The DMA controller can also be programmed to transfer data to and from the SSP's FIFO's. Refer to Chapter 5, "DMA Controller" for instructions on programming the DMA channels.

The steps for the DMA programming model are:

- 1. Program the transmit/receive byte count (buffer length) and burst size.
- 2. Program the DMA request to channel map register for SSP.
- 3. Set the run bit in the DMA control register.
- 4. Set the desired values in the SSP control registers.
- 5. Enable the SSP by setting the SSE bit in the SSP Control Register 0 (see Section 8.7.1).
- 6. Wait for both the DMA transmit and receive interrupt requests.

Note: If the transmit/receive byte count is not a multiple of the transfer burst size, the user must check the SSP Status Register (see Section 8.7.4) to determine if any data remains in the Receive FIFO.

8.6 Baud-Rate Generation

The baud (or bit-rate clock) is generated internally by dividing the internal clock (3.6864 MHz). The internal clock is first divided by 2 and this divided clock feeds a programmable divider to generate baud rates from 7.2 kbps to 1.8432 Mbps. Setting the External Clock Select (ECS) bit to 1 enables an external clock (SSPEXTCLK) to replace the 3.6864 MHz standard internal clock. The external clock is also divided by 2 before it is fed to the programmable divider.



8.7 SSP Serial Port Registers

The SSPC has five registers: two control, one data, one status, and one "reserved" register:

- The SSPC Control Registers (SSCR0 and SSCR1) are used to program the baud rate, data length, frame format, data transfer mechanism, and port enabling. They also control the FIFO "fullness" threshold that triggers an interrupt. These registers must be written before the SSP is enabled after reset and must only be changed when SSP is disabled.
- The SSPC Data Register (SSDR) is mapped as one 32-bit location that consists of two 16-bit registers. One register is for write operations and transfers data to the transmit FIFO. The other is for read operations and takes data from the receive FIFO. A write cycle, or burst write, loads successive half-words into the transmit FIFO. The write data occupies the lower 2 bytes of the 32-bit word. A read cycle, or burst read, similarly transfers data from the receive FIFO. The FIFOs are independent buffers that allow full duplex operation.
- The SSPC Status Register (SSSR) indicates the state of the FIFO buffers, whether the programmable threshold has been passed, and whether a transmit or receive FIFO service request is active. It also shows how many entries are occupied in the FIFO. Flag bits are set when the SSPC is actively transmitting or receiving data, when the transmit FIFO is not full, and when the receive FIFO is not empty. An error bit signals an overrun of the receive FIFO.

When the registers are programmed, reserved bits must be written as zeros and are read as undefined.

8.7.1 SSP Control Register 0 (SSCR0)

SSCR0, shown in Table 8-2, contains five bit fields that control SSP data size, frame format, external clock selection, clock divisor, and SSP enable. The SSE bit is reset to a zero state to ensure the SSP is disabled. The reset states for the other control bits are shown in the table, but each reset state must be set to the desired value before the SSPC is enabled.



Table 8-2. SSCR0 Bit Definitions

Та	ble	8-2	2. SS	CI	R0	Bit	De	efir	nition	S																					
			0:	x4 1	100	_000	00			S	SP (Con	trol	Re	egiste	er O	(S	SC	R0)		S	/nch	rone	ous	Seri	al P	ort (Con	trol	ler	
Bit	31	30	29 2	8 :	27	26	25	24	23 22	21	20	19	18	17	7 16	15	14	1	3 12	11	1	9	8	7	6	5	4	3	2	1	0
							re	esei	ved										S	CR				SSE	ECS	FR			DS	s	
Reset								Х											0	x00				0	0	0			0		
	Bi	its			Nar	ne												ſ	Desc	ript	ion										
	33	3:0 DSS							Data Size Select 0000 - reserved, undefined operation 0001 - reserved, undefined operation 0010 - reserved, undefined operation 0011 - 4-bit data 0100 - 5-bit data 0101 - 6-bit data 0110 - 7-bit data 0111 - 8-bit data 1000 - 9-bit data 1001 - 10-bit data 1011 - 12-bit data 1110 - 13-bit data 1111 - 16-bit data																						
		5:4 FRF							Frame Format 00 - Motorola's Serial Peripheral Interface (SPI) 01 - Texas Instruments' Synchronous Serial Protocol (SSP) 10 - National Microwire 11 - reserved, undefined operation External clock select:																						
		6 ECS							0 = On-chip clock used to produce the SSP's serial clock (SSPSCLK).1 = SSPEXTCLK is used to create the SSPSCLK.																						
	7	7 SSE						Synchronous Serial Port Enable: 0 = 0 - SSP operation disabled (pins may function as GPIOs) 1 = 1 - SSP operation enabled																							
	15	5:8			SC	R			Serial Value Bit rate	(0 tc	255	5) us			-									x (S0	CR +	- 1))					

8.7.1.1 Data Size Select (DSS)

reserved

31:16

The 4-bit data size select (DSS) field is used to determine the size of the data that the SSPC transmits and receives. The data can range from 4 to 16 bits in length. When data is programmed to be less than 16 bits, received data is automatically right-justified and the upper bits in the receive FIFO are zero-filled by receive logic. Do not left-justify transmit data before placing it in the



transmit FIFO. The transmit logic in the SSPC left-justifies the data sample according to the DSS bits before the sample is transmitted. Data sizes of 1, 2, and 3 bits are reserved and produce unpredictable results in the SSPC.

In National Microwire frame format, this bit field selects the size of the received data. The size of the transmitted command data is either 8-bit or 16-bit as selected by the MWDS bit in SSCR1.

8.7.1.2 Frame Format (FRF)

The 2-bit frame format (FRF) field is used to select Motorola SPI (FRF=00), Texas Instruments Synchronous Serial (FRF=01), or National Microwire (FRF=10) frame format.

FRF=11 is reserved and the SSPC produces unpredictable results if this value is used.

8.7.1.3 External Clock Select (ECS)

The external clock select (ECS) bit determines whether the SSPC uses the on-chip 3.6864-MHz clock or an off-chip clock supplied via SSPEXTCLK. When ECS=0, the SSPC uses the on-chip 3.6864-MHz clock to produce a range of serial transmission rates from 7.2 Kbps to 1.8432 Mbps. When ECS=1, the SSP uses SSPEXTCLK to access an off-chip clock. The off-chip clock's frequency can be any value up to 3.6864 MHz. The off-chip clock is useful when a serial transmission rate not evenly divisible from 3.6864 MHz is required for synchronization with the target off-chip slave device.

If the off-chip clock is used, the user must set the appropriate bits in the GPIO alternate function and pin direction registers that correspond to the pin. See Chapter 4, "System Integration Unit" for more details on configuring GPIO pins for alternate functions.

Note: Disable the SSPC by setting the SSPC Enable (SSE) to a 0 before setting the ECS bit to a 1. The ECS bit may be set to one either before the SSE is set to one or at the same time.

8.7.1.4 Synchronous Serial Port Enable (SSE)

The SSCR0[SSE] bit is used to enable and disable all SSP operations. When SSCR0[SSE]=0, the SSP is disabled. When SSCR0[SSE]=1, the SSP is enabled. When the SSP is disabled, all of its clocks are powered down to minimize power consumption. The SSP is disabled following a reset.

When the SSCR0[SSE] bit is cleared during active operation, the SSP is immediately disabled and the frame being transmitted is terminated. Clearing SSCR0[SSE] resets the SSP's FIFOs and the SSP status bits. The SSP's control registers are not reset when SSCR0[SSE] is cleared.

Note: After reset or after the SSCR0[SSE] is cleared, ensure that the SSCR1 and SSSR registers are properly reconfigured or reset before re-enabling the SSP with the SSCR0[SSE]. Other control bits in SSCR0 may be written at the same time as the SSCR0[SSE].

When the SSPC is disabled, its five pins may be used as GPIOs. They are configured as inputs or outputs with the control registers described in Chapter 4, "System Integration Unit". In Sleep mode, the pins' states are controlled by the GPIO sleep register. SSPC register settings have no effect on the pins in Sleep mode.



8.7.1.5 Serial Clock Rate (SCR)

The 8-bit serial clock rate (SCR) bit-field is used to select the SSPC bit rate. The SSPC has 256 bit rates, from 7.2 Kbps to 1.8432 Mbps. The serial clock generator uses the internal 3.6864 MHz clock or an external clock provided through SSPEXTCLK. The clock is divided by 2, then divided by the programmable SCR value (0 to 255) plus 1 to generate the serial clock (SSPSCLK). The resultant clock is driven on the SSPSCLK pin and is used by the SSP's transmit logic to drive data on the SSPTXD pin and to latch data on the SSPRXD pin. Depending on the frame format selected, each transmitted bit is driven on either SSPSCLK's rising or falling edge and sampled on the opposite clock edge.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

8.7.2 SSP Control Register 1 (SSCR1)

SSCR1, shown in Table 8-3, contains bit fields that control SSP functions.

Table 8-3. SSCR1 Bit Definitions (Sheet 1 of 2)

		0x4100_0004	SSP Control Register 1 (SS	SCR1)	Syr	nchron	ous	Seri	al P	ort	Cor	itro	ller					
Bit	31 30	29 28 27 26 25 24	23 22 21 20 19 18 17 16 15 14	13 12 11	10	9 8	7	6	5	4	3	2	1	0				
		r	eserved	RFT		Т	FT		MWDS	SPH	SPO	LBM	TIE	RIE				
Reset		2	X	0x0		0	x0		0	0	0	0	0	0				
	Bits	Name		Descripti	ion													
	0	RIE	Receive FIFO Interrupt Enable 0 = Receive FIFO interrupt is disable 1 = Receive FIFO interrupt is enable	errupt is disabled														
	1	TIE	Transmit FIFO Interrupt Enable 0 = Transmit FIFO interrupt is disable 1 = 1 - Transmit FIFO interrupt is ena	rupt is disabled														
	2	LBM	Loop-Back Mode 0 = Normal serial port operation enal 1 = Output of transmit serial shifter ir	ort operation enabled hit serial shifter internally connected to input of receive serial shifter														
	3	SPO	Motorola SPI SSPSCLK polarity settin 0 = The inactive or idle state of SSP3 1 = The inactive or idle state of SSP3	SCLK is lov														
	Motorola SPI SSPSCLK phase setting: 0 = SSPSCLK is inactive one cycle at the start of a frame and 1/2 cycle at the end of a frame. 1 = SSPSCLK is inactive 1/2 cycle at the start of a frame and one cycle at the end of a frame.																	
	5	MWDS	Microwire Transmit Data Size: 0 = 8-bit command words are transm 1 = 16-bit command words are transm															



Table 8-3. SSCR1 Bit Definitions (Sheet 2 of 2)

		0x4100_0004	SSP Control Register 1 (SS	SCR1) Syı	nchronous Ser	ial P	ort	Cor	itrol	ler						
Bit	31 30	29 28 27 26 25 24	23 22 21 20 19 18 17 16 15 14	13 12 11 10	9 8 7 6	5	4	3	2	1	0					
		r	eserved	RFT	TFT	MWDS	SPH	SPO	LBM	JIE	RIE					
Reset)	(0x0	0x0	0	0	0	0	0	0					
	Bits	Name	Description													
	9:6	TFT (Transmit FIFO Threshold)		reshold level at which Transmit FIFO generates an interrupt or DMA request. This ust be set to the desired threshold value minus 1.												
	13:10 RFT (Receive FIFO Threshold) Sets threshold level at which Receive FIFO generates an interrupt or DMA request. This level must be set to the desired threshold value minus 1.															
	31:14	_	reserved													

8.7.2.1 Receive FIFO Interrupt Enable (RIE)

The Receive FIFO Interrupt Enable (RIE) bit is used to mask or enable the receive FIFO service request interrupt. When RIE=0, the interrupt is masked and the interrupt controller ignores the state of the Receive FIFO Service Request (RFS) bit in the SSPC Status Register. When the RIE bit is set to a 1, the interrupt is enabled and, if the RFS bit is set to a 1, an interrupt request is made to the interrupt controller. If the RIE bit is set to 0 neither the RFS bit's current state nor the receive FIFO logic's ability to set and clear the RFS bit is affected. However, the interrupt request generation is blocked.

The RIE bit's state does not affect the receive FIFO DMA request generation that is asserted when the RFS bit is set to a 1.

8.7.2.2 Transmit FIFO Interrupt Enable (TIE)

The Transmit FIFO Interrupt Enable (TIE) bit is used to mask or enable the transmit FIFO service request interrupt. When the TIE bit is set to a 0, the interrupt is masked and the interrupt controller ignores the state of the Transmit FIFO Service Request (TFS) bit in the SSPC Status Register. When the TIE bit is set to a 1, the interrupt is enabled and, if the TFS bit is set to a 1, an interrupt request is made to the interrupt controller. If the TIE bit is set to a 0, neither the TFS bit's current state nor the transmit FIFO logic's ability to set and clear the TFS bit is affected. However, the interrupt request generation is blocked.

The TIE bit's state does not affect the transmit FIFO DMA request generation that is asserted when the TFS bit is set to a 1.

8.7.2.3 Loop Back Mode (LBM)

The loop back mode (LBM) bit is used to enable and disable the SSP transmit and receive logic. When LBM=0, the SSP operates normally. The transmit and receive data paths are separate and communicate via their respective pins. However, since the SSP is a master device, it must transmit in order to receive while in either SSP or SPI mode. When LBM=1, the transmit serial shifter's output is directly connected internally to the receive serial shifter's input.



Note: Loop back mode cannot be used with Microwire frame format.

8.7.2.4 Serial Clock Polarity (SPO)

The serial clock polarity bit (SPO) selects the SSPSCLK signal's inactive state in the Motorola SPI format (FRF=00). For SPO=0, the SSPSCLK is held low in the inactive or idle state when the SSP is not transmitting/receiving data. When the SPO bit is set to a 1, the SSPSCLK is held high during the inactive/idle state. The SPO bit's programmed setting alone does not determine which SSPSCLK edge is used to transmit or receive data. The SPO bit's setting combined with the SSPSCLK phase bit (SPH) determine which edge is used.

Note: The SPO bit is ignored for all data frame formats except for the Motorola SPI format (FRF=00).

8.7.2.5 Serial Clock Phase (SPH)

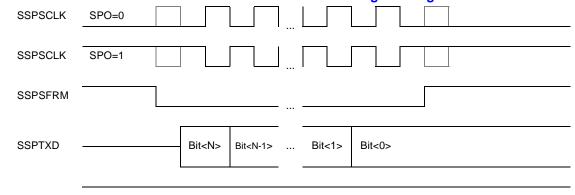
The serial clock (SSPSCLK) phase bit (SPH) determines the phase relationship between the SSPSCLK and the serial frame (SSPSFRM) pins for the Motorola SPI format (FRF=00). When SPH=0, SSPSCLK remains in its inactive/idle state (as determined by the SPO setting) for one full cycle after SSPSFRM is asserted low at the beginning of a frame. SSPSCLK continues to transition for the rest of the frame and is then held in its inactive state for one-half of an SSPSCLK period before SSPSFRM is deasserted high at the end of the frame. When SPH=1, SSPSCLK remains in its inactive/idle state (as determined by the SPO setting) for one-half cycle after SSPSFRM is asserted low at the beginning of a frame. SSPSCLK continues to transition for the rest of the frame and is then held in its inactive state for one full SSPSCLK period before SSPSFRM is deasserted high at the end of the frame.

The combination of the SPO and SPH settings determines when SSPSCLK is active during the assertion of SSPSFRM and which SSPSCLK edge is used to transmit and receive data on the SSPTXD and SSPRXD pins. When SPO and SPH are programmed to the same value, transmit data is driven on SSPSCLK's falling edge and receive data is latched on SSPSCLK's rising edge. When SPO and SPH are programmed to opposite values (one 0 and the other 1), transmit data is driven on SSPSCLK's rising edge and receive data is latched on SSPSCLK's falling edge.

The SPH is ignored for all data frame formats except the Motorola SPI format (FRF=00).

Figure 8-4 shows the pin timing for the four SPO and SPH programming combinations. SPO inverts the SSPSCLK signal's polarity and SPH determines the phase relationship between SSPSCLK and SSPSFRM, shifting the SSPSCLK signal one-half phase to the left or right during the SSPSFRM assertion.

Figure 8-4. Motorola SPI* Frame Formats for SPO and SPH Programming





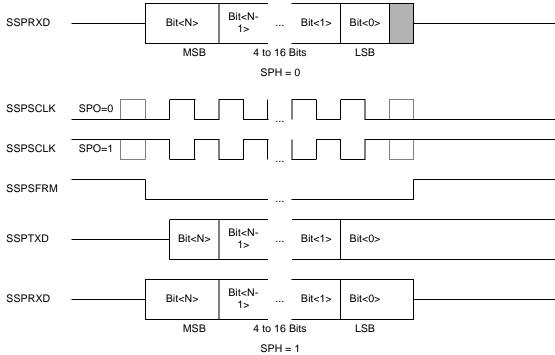


Figure 8-4. Motorola SPI* Frame Formats for SPO and SPH Programming

8.7.2.6 Microwire Transmit Data Size (MWDS)

The Microwire Transmit Data Size (MWDS) bit is used to select the 8- or 16-bit size for command word transmissions in the National Microwire frame format. When the MWDS is set to "0", 8-bit command words are transmitted. When the MWDS is set to "1", 16-bit command words are transmitted. The MWDS setting is ignored for all other frame formats.

8.7.2.7 Transmit FIFO Interrupt/DMA Threshold (TFT)

This 4-bit value sets the level at or below which the FIFO controller triggers a DMA service request and, if enabled, an interrupt request. Refer to Table 8-4 for suggested TFT values associated with DMA servicing.

8.7.2.8 Receive FIFO Interrupt/DMA Threshold (RFT)

This 4-bit value sets the level at or above which the FIFO controller triggers a DMA service interrupt and, if enabled, an interrupt request. Refer to Table 8-4 for suggested RFT values associated with DMA servicing.

Be careful not to set the RFT value too high for your system or the FIFO could overrun because of the bus latencies caused by other internal and external peripherals. This is especially the case for interrupt and polled modes that require a longer time to service.



Table 8-4. TFT and RFT Values for DMA Servicing

DMA Burst Size	TFT	Value	RFT Value						
	Min	Max	Min	Max					
8 Bytes	0	11	3	15					
16 Bytes	0	7	7	15					

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

8.7.3 SSP Data Register (SSDR)

SSDR, shown in Table 8-5, is a single address location that can be accessed by read/write data transfers. Transfers can be single transfers, 4 half-word bursts, or 8 half-word bursts.

As the system accesses the register, FIFO control logic transfers data automatically between register and FIFO as fast as the system moves data. The SSP Status Register has bits that indicate whether either FIFO is full, above/below a programmable threshold, or empty.

For transmit operations from SSPC to SSP peripheral, the CPU (using programmed I/O) may write to the SSDR when the transmit FIFO is below its threshold level.

When a data size less than 16-bits is selected, do not left-justify data before it is written to the SSDR. Transmit logic left-justifies the data and ignores any unused bits. Received data less than 16-bits is automatically right-justified in the receive FIFO.

When the SSPC is programmed for National Microwire frame format and the size for transmit data is 8-bits, as selected by the MWDS bit in the SSCR1, the most significant byte is ignored. SSCR0[DSS] controls receive data size.

Note: Both FIFOs are cleared when the SSPC is reset or a zero is written to the SSCR0[SSE] bit.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 8-5. SSDR Bit Definitions

		0x4100_0010 SSP Data Regis													gist	ister (SSDR) Synchronous Serial Port Controller															
Bit	31 3	0 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		reserved											Transmit/Receive Data																		
Reset		X												0x0000																	
	Bits Name													De	escr	ipti	on														
	15:0 Data (Low Word) Data word to be written to/re											/rea	d fro	m 1	「ran:	smit	/Re	ceiv	e FI	FO											
	31:16 — reserved																														



8.7.4 SSP Status Register (SSSR)

The SSP Status Register (SSSR) is shown in Table 8-6. The SSSR contains bits that signal overrun errors and transmit and receive FIFO service requests. These hardware-detected events signal an interrupt request to the interrupt controller. The status register also contains flags that indicate when the SSP is actively transmitting or receiving characters, when the transmit FIFO is not full, and when the receive FIFO is not empty (no interrupt generated).

Bits that cause an interrupt will signal the request as long as the bit is set. When the bit is cleared, the interrupt is cleared. Read/write bits are called status bits, read-only bits are called flags. Status bits are referred to as sticky (once set by hardware, they must be cleared by software). Writing a 1 to a sticky status bit clears it. Writing a 0 has no effect. Read-only flags are set and cleared by hardware. Writes have no effect. Some bits that cause interrupts have corresponding mask bits in the control registers.

All bits are read-only except ROR, which is read/write. ROR's reset state is zero. Writes to TNF, RNE, BSY, TFS, and RFS have no effect. Writes to reserved bits are ignored and reads from these bits are undetermined.



Table 8-6. SSSR Bit Definitions

		0x4100_0008	SSP Status Regi	ister (SSSR)	Synchrone	ous Seri	al P	ort C	ontro	ller								
Bit	31 30	29 28 27 26 25 24	23 22 21 20 19 18 17 1	6 15 14 13 12	11 10 9 8	7 6	5	4 3	2	1 0								
		rese	rved	RFL	TFL	ROR	TFS	BSY	TNF	reserved								
Reset)	(0xF	0x0	0 0	0	0 (1	0								
	Bits	Name		Descr	ription													
	1:0	1	reserved															
	2	TNF	Transmit FIFO Not Full (read 0 = Transmit FIFO is full 1 = Transmit FIFO is not ful															
	3	RNE	Receive FIFO Not Empty (re 0 = Receive FIFO is empty 1 = Receive FIFO is not em	•														
	4	BSY	SSP Busy (read only) 0 = SSP is idle or disabled 1 = SSP currently transmitti	ng or receiving a	frame													
	5	TFS	0 = Transmit FIFO level exc	ransmit FIFO Service Request (read only) = Transmit FIFO level exceeds TFT threshold, or SSP disabled = Transmit FIFO level is at or below TFT threshold, generate interrupt or DMA request														
	6	RFS	Receive FIFO Service Request (read only) 0 = Receive FIFO level exceeds RFT threshold, or SSP disabled 1 = Receive FIFO level is at or above RFT threshold, generate interrupt or DMA request															
	7	ROR	Receive FIFO Overrun (read 0 = Receive FIFO has not e 1 = Attempted data write to	experienced an o		rupt												
	11:8	TFL	Transmit FIFO Level (read o Number of entries in Transm empty or full and the softwar	it FIFO. Note: W)x0 is rea	ıd, th	ne FIF	O is	either								
	15:12	RFL	Receive FIFO Level (read or Number of entries minus one is either empty or full and the	in Receive FIFC			0xF	is rea	ad, the	e FIFO								
	31:16	_	reserved															

8.7.4.1 Transmit FIFO Not Full Flag (TNF)

This non-interruptible bit is set whenever the transmit FIFO is not full. TNF is cleared when the FIFO is completely full. This bit can be polled when using programmed I/O to fill the transmit FIFO over its threshold level. This bit does not request an interrupt.



8.7.4.2 Receive FIFO Not Empty Flag (RNE)

This non-interruptible bit is set when the receive FIFO contains one or more entries and is cleared when the FIFO is empty. Because CPU interrupt requests are only made when the Receive FIFO Threshold has been met or exceeded, the RNE bit can be polled when programmed I/O removes remaining bytes of data from the receive FIFO. This bit does not request an interrupt.

8.7.4.3 SSP Busy Flag (BSY)

This is a non-interruptible read-only bit that is set when the SSP is actively transmitting and/or receiving data and is cleared when the SSP is idle or disabled (SSE=0). This bit does not request an interrupt. Since the software can read this bit before the SSP starts to transmit data, software must read SSSP[TFL]=0x0 and SSSP[TNF]=0b1 and SSSP[BSY]=0b0 in order to insure that all data has transmitted completely.

8.7.4.4 Transmit FIFO Service Request Flag (TFS)

This bit contains a maskable interrupt and is set when the transmit FIFO is nearly empty and requires service to prevent an underrun. TFS is set any time the transmit FIFO has the same or fewer valid data entries than indicated by the Transmit FIFO Threshold. It is cleared when it has more valid data entries than the threshold value. When the TFS bit is set, an interrupt request is made unless the transmit FIFO interrupt request enable (TIE) bit is cleared. The TFS bit's setting indicates whether a DMA service has been requested from the DMA controller. The DMA request cannot be masked by the TIE bit. After the CPU or the DMA fills the FIFO such that it exceeds the threshold, the TFS flag (and the service request and/or interrupt) is automatically cleared.

8.7.4.5 Receive FIFO Service Request Flag (RFS)

This bit contains a maskable interrupt and is set when the receive FIFO is nearly filled and requires service to prevent an overrun. RFS is set any time the receive FIFO has the same or more valid data entries than indicated by the Receive FIFO Threshold. It is cleared when it has fewer entries than the threshold value. When the RFS bit is set, an interrupt request is made unless the receive FIFO interrupt request enable (RIE) bit is cleared. The RFS bit's setting indicates that DMA service has been requested from the DMA controller. This DMA request cannot be masked by the RIE bit. After the CPU or DMA reads the FIFO such that it has fewer entries than the RFT value, the RFS flag (and the service request and/or interrupt) is automatically cleared.

8.7.4.6 Receiver Overrun Status (ROR)

This is a non-interruptible bit that is set when the receive logic attempts to place data into the receive FIFO after it has been completely filled. Each time a new piece of data is received, the set signal to the ROR bit is asserted and the newly received data is discarded. This process is repeated for each new piece of data received until at least one empty FIFO entry exists. When the ROR bit is set, an interrupt request that cannot be locally masked by any SSPC register bit is made to the CPU. The ROR bit's setting does not generate any DMA service request. Writing 0b1 to this bit resets ROR status and its interrupt request. Writing a "0" does not affect ROR status.

8.7.4.7 Transmit FIFO Level (TFL)

This bit indicates the number of entries currently in the Transmit FIFO.



8.7.4.8 Receive FIFO Level (RFL)

This bit indicates the one less than number of entries in the Receive FIFO.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

8.8 SSP Controller Register Summary

Table 8-7 shows the SSP registers associated with the SSP controller and their physical addresses.

Table 8-7. SSP Controller Register Summary

Address	Abbreviation	Full Name
0x4100_0000	SSCR0	SSP Control Register 0
0x4100_0004	SSCR1	SSP Control Register 1
0x4100_0008	SSSR	SSP Status Register
0x4100_000C	_	reserved
0x4100_0010	SSDR (Write / Read)	SSP Data Write Register/SSP Data Read Register

Synchronous Serial Port Controller



intel® PC Bus Interface Unit

This chapter describes the Inter-Integrated Circuit (I2C) bus interface unit, including the operation modes and setup for the PXA255 processor.

Overview 9.1

The I²C bus was created by the Phillips Corporation and is a serial bus with a two-pin interface. The SDA data pin is used for input and output functions and the SCL clock pin is used to control and reference the I²C bus. The I²C unit allows the processor to serve as a master and slave device that resides on the I²C bus.

The I²C unit enables the processor to communicate with I²C peripherals and microcontrollers for system management functions. The I²C bus requires a minimum amount of hardware to relay status and reliability information concerning the processor subsystem to an external device.

The I²C unit is a peripheral device that resides on the processor internal bus. Data is transmitted to and received from the I²C bus via a buffered interface. Control and status information is relayed through a set of memory-mapped registers. Refer to The I²C-Bus Specification for complete details on I²C bus operation.

The I^2C unit does not support the hardware general call, 10-bit addressing, or CBUS compatibility.

9.2 **Signal Description**

The I²C unit signals are SDA and SCL. Table 9-1 describes each signal's function.

Table 9-1. I²C Signal Description

Signal Name	Input/Output	Description
SDA	Bidirectional	I ² C Serial Data/Address signal
SCL	Bidirectional	I ² C Serial Clock Line signal

Functional Description 9.3

The I²C bus defines a serial protocol for passing information between agents on the I²C bus using a two pin interface that consists of a Serial Data/Address (SDA) line and a Serial Clock Line (SCL). Each device on the I²C bus is recognized by a unique 7-bit address and can operate as a transmitter or as a receiver in master or slave mode. Table 9-2 lists the I²C operation modes.

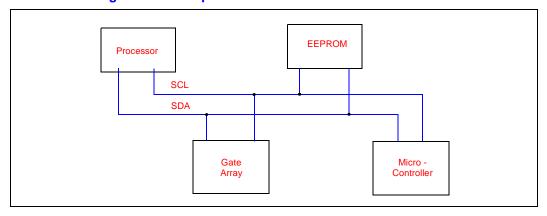


Table 9-2. I²C Bus Definitions

I ² C Device	Definition
Transmitter	Sends data to the I ² C bus.
Receiver	Receives data from the I ² C bus.
Master	Initiates a transfer, generates the clock signal, and terminates the transactions.
Slave	Device addressed by a master.
Multi-master	More than one master can attempt to control the bus at the same time without corrupting the message.
Arbitration	Ensures that only one master controls the bus when more than one master simultaneously tries to control the bus. This ensures that messages are not corrupted.

For example, when the processor I^2C unit acts as a master on the bus, it addresses an EEPROM as a slave to receive data (see Figure 9-1). When the I^2C unit is addressing the EEPROM, it is a master-transmitter and the EEPROM is a slave-receiver. When the I^2C reads data, it is a master-receiver and the EEPROM is a slave-transmitter. Whether it is a transmitter or receiver, the master generates the clock, initiates the transaction, and terminates the transaction.

Figure 9-1. I²C Bus Configuration Example



The I^2C bus allows for a multi-master system, which means more than one device can initiate data transfers at the same time. To support this feature, the I^2C bus arbitration relies on the wired-AND connection of all I^2C interfaces to the I^2C bus. Two masters can drive the bus simultaneously, provided they drive identical data. If a master tries to drive SDA high while another master drives SDA low, it loses the arbitration. The SCL line is a synchronized combination of clocks generated by the masters using the wired-AND connection to the SCL line.

The I²C bus serial operation uses an open-drain wired-AND bus structure, which allows multiple devices to drive the bus lines and to communicate status about events such as arbitration, wait states, error conditions, etc. For example, when a master drives the clock (SCL) line during a data transfer, it transfers a bit on every instance that the clock is high. When the slave is unable to accept or drive data at the rate that the master is requesting, the slave can hold the clock line low between the high states to insert a wait interval. The master's clock can only be altered by another master during arbitration or a slow slave peripheral that keeps the clock line low.

 I^2C transactions are either initiated by the processor as a master or received by the processor as a slave. Both conditions may result in reads, writes, or both to the I^2C bus.



9.3.1 Operational Blocks

The I²C unit is connected to the peripheral bus. The processor interrupt mechanism can be used to notify the CPU that there is activity on the I²C bus. Polling can be used instead of interrupts. The I²C unit consists of the two wire interface to the I²C bus, an 8-bit buffer for passing data to and from the processor, a set of control and status registers, and a shift register for parallel/serial conversions.

The I²C unit initiates an interrupt to the processor when a buffer is full, a buffer is empty, the I²C unit slave address is detected, arbitration is lost, or a bus error condition occurs. All interrupt conditions must be cleared explicitly by software. See Section 9.9.4 for details.

The 8-bit I^2C Data Buffer Register (IDBR) is loaded with a byte of data from the shift register interface to the I^2C bus when receiving data and from the processor internal bus when writing data. The serial shift register is not user accessible.

The I²C Control Register (ICR) and the I²C Status Register (ISR) are located in the I²C memory-mapped address space. The registers and their functions are defined in Section 9.9.

The I²C unit supports a fast mode operation of 400 Kbits/sec and a standard mode of 100 Kbits/sec. Refer to *The I²C-Bus Specification* for details.

9.3.2 I²C Bus Interface Modes

The I²C unit can accomplish a transfer in different operation modes. Table 9-3 summarizes the different modes.

Table 9-3. Modes of Operation

Mode	Description
Master - Transmit	 I²C unit acts as a master. Used for a write operation. I²C unit sends the data. I²C unit is responsible for clocking. Slave device in slave-receive mode
Master - Receive	 I²C unit acts as a master. Used for a read operation. I²C unit receives the data. I²C unit is responsible for clocking. Slave device in slave-transmit mode
Slave - Transmit	 I²C unit acts as a slave. Used for a master read operation. I²C unit sends the data. Master device in master-receive mode.
Slave - Receive (default)	 I²C unit acts as a slave. Used for a master write operation. I²C unit receives the data. Master device in master-transmit mode.

While the I²C unit is idle, it defaults to slave-receive mode. This allows the interface to monitor the bus and receive any slave addresses intended for the processor.



When the I^2C unit receives an address that matches the 7-bit address found in the I^2C Slave Address Register (ISAR) or the general call address (see Section 9.4.7), the interface either remains in slave-receive mode or transitions to slave-transmit mode. The Read/Write bit (R/nW) determines which mode the interface enters. The R/nW bit is the least significant bit of the byte containing the slave address. If the R/nW bit is low, the master that initiated the transaction intends write data and the I^2C unit remains in slave-receive mode. If the R/nW is high, the master that initiated the transaction intends to read data and the I^2C unit transitions to slave-transmit mode. Section 9.4.6 further defines slave operation.

When the I²C unit initiates a read or write on the I²C bus, it transitions from the default slave-receive mode to the master-transmit mode. If the transaction is a write, the I²C unit remains in master-transmit mode after the address transfer is completed. If the transaction is a read, the I²C unit transmits the start address, then transitions to master-receive mode. Section 9.4.5 further defines master operation.

9.3.3 Start and Stop Bus States

The I²C bus specification defines a transaction START, used at the beginning of a transfer, and a transaction STOP bus state, used at the end of a transfer. A START condition occurs if a high to low transition takes place on the SDA line when SCL is high. A STOP condition occurs if a low to high transition takes place on the SDA line when SCL is high.

The I²C unit uses the ICR[START] and ICR[STOP] bits to:

- Initiate an additional byte transfer
- Initiate a START condition on the I²C bus
- Enable Data Chaining (repeated START)
- Initiate a STOP condition on the I²C bus

Table 9-4 defines the START and STOP bits in the ICR.

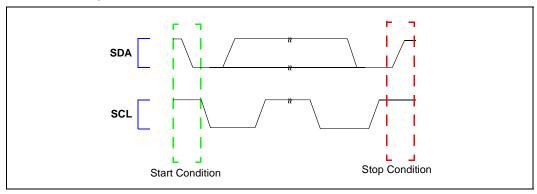
Table 9-4. START and STOP Bit Definitions

STOP bit	STAR T bit	Condition	Notes
0	0	No START or STOP	I ² C unit sends a no START or STOP condition. Used when multiple data bytes need to be transferred.
			I ² C unit sends a START condition and transmit the 8-bit IDBR's contents. The IDBR must contain the 7-bit address and the R/nW bit before a START is initiated.
0	1	START Condition and	For a Repeated Start, the IDBR contains the target slave address
	1	Repeated START	and the R/nW bit. This allows a master to make multiple transfers to different slaves without giving up the bus.
			The interface stays in master-transmit mode for writes and transitions to master-receive mode for reads.
			In master-transmit mode, the I ² C unit transmits the 8-bit IDBR and sends a STOP condition on the I ² C bus.
1	X	STOP Condition	In master-receive mode, the ICR[ACKNAK] must be changed to a negative ACK (see Section 9.4.3). The I ² C unit transmits the NAK bit, receives the data byte in the IDBR, and sends a STOP condition on the I ² C bus.



Figure 9-2 shows the relationship between the SDA and SCL lines for START and STOP conditions.

Figure 9-2. Start and Stop Conditions



9.3.3.1 START Condition

The START condition (ICR[START]=1, ICR[STOP]=0) initiates a master transaction or repeated START. Before it sets the START ICR bit, software must load the target slave address and the R/nW bit in the IDBR (see Section 9.9.2). The START and the IDBR contents are transmitted on the I²C bus after the ICR[TB] bit is set. The I²C bus stays in master-transmit mode for write requests and enters master-receive mode for read requests. For a repeated start, a change in read or write, or a change in the target slave address, the IDBR contains the updated target slave address and the R/nW bit. A repeated start enables a master to make multiple transfers to different slaves without surrendering the bus.

The START condition is not cleared by the I²C unit. If the I²C loses arbitration while initiating a START, it may re-attempt the START when the bus is freed. See Section 9.4.4 for details on how the I²C unit functions in those circumstances.

9.3.3.2 No START or STOP Condition

The no START or STOP condition (ICR[START]=0, ICR[STOP]=0) is used in master-transmit mode while the I²C unit is transmitting multiple data bytes (see Figure 9-2). Software writes the data byte and the I²C unit sets the ISR[ITE] bit and clears the ICR[TB] bit. The software then writes a new byte to the IDBR and sets the ICR[TB] bit, which initiates the new byte transmission. This process continues until the software sets the ICR[START] or ICR[STOP] bit. The ICR[START] and ICR[STOP] bits are not automatically cleared by the I²C unit after the transmission of a START, STOP, or repeated START.

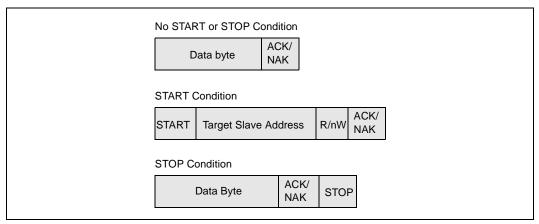
After each byte transfer, including the ICR[ACKNAK] bit, the I²C unit holds the SCL line low to insert wait states until the ICR[TB] bit is set. This action notifies the I²C unit to release the SCL line and allow the next information transfer to proceed.

9.3.3.3 STOP Condition

The STOP condition (ICR[START]=X, ICR[STOP]=1) terminates a data transfer. In master-transmit mode, the ICR[STOP] bit and the ICR[TB] bit must be set to initiate the last byte transfer (see Figure 9-2). In master-receive mode, the I²C unit must set the ICR[ACKNAK] bit, the ICR[STOP] bit, and the ICR[TB] bit to initiate the last transfer. Software must clear the ICR[STOP] condition after it is transmitted.



Figure 9-3. START and STOP Conditions





9.4 I²C Bus Operation

The I²C unit transfers data in 1-byte increments and always follows this sequence:

- 1) START
- 2) 7-bit Slave Address
- 3) R/nW Bit
- 4) Acknowledge Pulse
- 5) 8 Bits of Data
- 6) ACK/NAK Pulse
- 7) Repeat of Steps 5 and 6 for required number of bytes
- 8) Repeated START (Repeat Step 1) or STOP

9.4.1 Serial Clock Line (SCL) Generation

When the I^2C unit is in master-transmit or master-receive mode, it generates the I^2C clock output. The SCL clock is generated by setting the ICR[FM] bit for either 100KBit/sec or 400Kbit/sec operation.

9.4.2 Data and Addressing Management

The I^2C Data Buffer Register (IDBR) and the I^2C Slave Address Register (ISAR) manage data and slave addressing. The IDBR (see Section 9.9.2) contains one byte of data or a 7-bit slave address and the R/nW bit. The ISAR contains the processor programmable slave address. The I^2C unit puts received data in the IDBR after a full byte is received and acknowledged. To transmit data, the CPU writes to the IDBR, and the I^2C unit passes the information to the serial bus when the ICR[TB] bit is set. See Section 9.9.3.

When the I²C unit is in master- or slave-transmit mode:

- 1. Software writes data to the IDBR over the internal bus. This initiates a master transaction or sends the next data byte after the ISR[ITE] bit is set.
- 2. I²C unit transmits data from the IDBR when the ICR[TB] bit is set.
- 3. When enabled, an IDBR transmit empty interrupt is signalled when a byte is transferred on the I²C bus and the acknowledge cycle is complete.
- 4. When the I²C unit is ready to transfer the next byte before the CPU has written the IDBR and a STOP condition is not in place, the I²C unit inserts wait states until the CPU writes a new value into the IDBR and sets the ICR[TB] bit.

When the I²C unit is in master- or slave-receive mode:

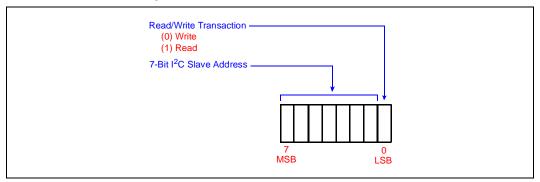
- 1. The processor reads IDBR data over the internal bus after the IDBR receive full interrupt is signalled.
- I²C unit transfers data from the shift register to the IDBR after the acknowledge cycle completes.
- 3. I²C unit inserts wait states until the IDBR is read. Refer to Section 9.4.3 for acknowledge pulse information in receiver mode.
- 4. After the CPU reads the IDBR, the I²C unit writes the ICR[ACKNAK] bit and the ICR[TB] bit, allowing the next byte transfer to proceed.



9.4.2.1 Addressing a Slave Device

As a master device, the I²C unit must compose and send the first byte of a transaction. This byte consists of the slave address for the intended device and a R/nW bit for transaction definition. The MSB is transmitted first. The slave address and the R/nW bit are written to the IDBR (see Figure 9-4).

Figure 9-4. Data Format of First Byte in Master Transaction



The first byte transmission must be followed by an ACK pulse from the addressed slave. When the transaction is a write, the I^2C unit remains in master-transmit mode and the addressed slave device stays in slave-receive mode. When the transaction is a read, the I^2C unit transitions to master-receive mode immediately following the ACK and the addressed slave device transitions to slave-transmit mode. When a NAK is returned, the I^2C unit aborts the transaction by automatically sending a STOP and setting the ISR[BED] bit.

When the I^2C unit is enabled and idle, it remains in slave-receive mode and monitors the I^2C bus for a START signal. When it detects a START pulse, the I^2C unit reads the first seven bits and compares them to those in the ISAR and the general call address (0x00). When the bits match those in the ISAR register, the I^2C unit reads the eighth bit (R/nW bit) and transmits an ACK pulse. The I^2C unit either remains in slave-receive mode (R/nW = 0) or transitions to slave-transmit mode (R/nW = 1). See Section 9.4.7 for actions when a general call address is detected.

9.4.3 I²C Acknowledge

Every I²C byte transfer must be accompanied by an acknowledge pulse that the master- or slave-receiver must generate. The transmitter must release the SDA line for the receiver to transmit the acknowledge pulse (see Figure 9-5).



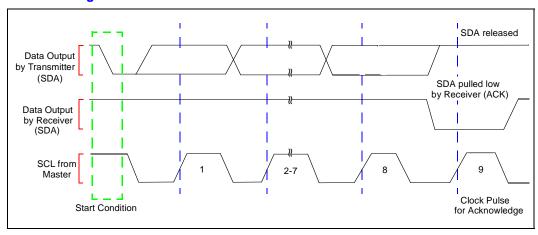


Figure 9-5. Acknowledge on the I²C Bus

In master-transmit mode, if the target slave-receiver device cannot generate the acknowledge pulse, the SDA line remains high. The lack of an acknowledge NAK causes the I²C unit to set the ISR[BED] bit and generate the associated interrupt when enabled. The I²C unit automatically generates a STOP condition and aborts the transaction.

In master-receive mode, the I²C unit sends a negative acknowledge (NAK) to signal the slave-transmitter to stop sending data. The ICR[ACKNAK] bit controls the ACK/NAK bit value that the I²C bus drives. As required by the I²C bus protocol, the ISR[BED] bit is not set for a master-receive mode NAK. The I²C unit automatically transmits the ACK pulse after it receives each byte from the serial bus. Before the unit receives the last byte, software must set the ICR[ACKNAK] bit to 1 (NAK). The NAK pulse is sent after the last byte to indicate that the last byte has been sent.

In slave mode, the I²C unit automatically acknowledges its own slave address, independent of the value in the ICR[ACKNAK] bit. In slave-receive mode, an ACK response automatically follows a data byte, independent of the value in the ICR[ACKNAK] bit. The I²C unit sends the ACK value after it receives the eighth data bit in a byte.

In slave-transmit mode, the I²C unit receives a NAK from the master to indicate the last byte has been transferred. The master then sends a STOP or repeated START. The ISR[UB] bit remains set until a STOP or repeated START is received.

9.4.4 Arbitration

The I²C bus' multi-master capabilities require I²C bus arbitration. Arbitration takes place when two or more masters generate a START condition in the minimum hold time.

Arbitration can take a long time. If the address bit and the R/nW are the same, the arbitration scheme considers the data. Because the I^2C bus has a wired-AND nature, a transfer does not lose data if multiple masters signal the same bus states. If the address and the R/nW bit or the data they contain are different, the master signals a high state loses arbitration and shuts off its data drivers. If the I^2C unit loses arbitration, it shuts off the SDA or SCL drivers for the rest of the byte transfer, sets the ISR[ALD] bit, and returns to slave-receive mode.



9.4.4.1 SCL Arbitration

Each master on the I²C bus generates its own clock on the SCL line for data transfers. As a result, clocks with different frequencies may be connected to the SCL line. Because data is valid when a clock is in the high period, bit-by-bit arbitration requires a defined clock synchronization procedure.

Clock synchronization is through the wired-AND connection of the I²C interfaces to the SCL line. When a master's clock changes from high to low, the master holds down the SCL line for its associated period (see Figure 9-6). A clock cannot switch from low to high if another master has not completed its period. The master with the longest low period holds down the SCL line. Masters with shorter periods are held in a high wait-state until the master with the longest period completes. After the master with the longest period completes, the SCL line changes to the high state and masters with the shorter periods continue the data cycle.

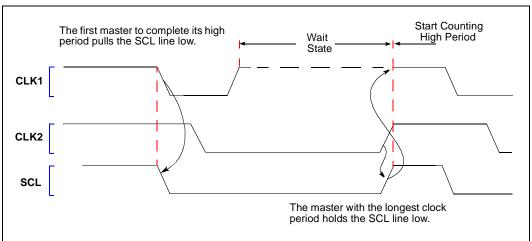


Figure 9-6. Clock Synchronization During the Arbitration Procedure

9.4.4.2 SDA Arbitration

Arbitration on the SDA line can continue for a long time because it starts with the address and R/nW bits and continues through the data bits. Figure 9-7 shows the arbitration procedure for two masters. More than two masters may be involved if more than two masters are connected to the bus. If the address bit and the R/nW are the same, the arbitration scheme considers the data. Because the I^2C bus has a wired-AND nature, a transfer does not lose data if multiple masters signal the same bus states. If the address and the R/nW bit or the data they contain are different, the master that sent the first low data bit loses arbitration and shuts off its data drivers. If the I^2C unit loses arbitration, it shuts off the SDA or SCL drivers for the rest of the byte transfer, sets the ISR[ALD] bit, and returns to slave-receive mode.



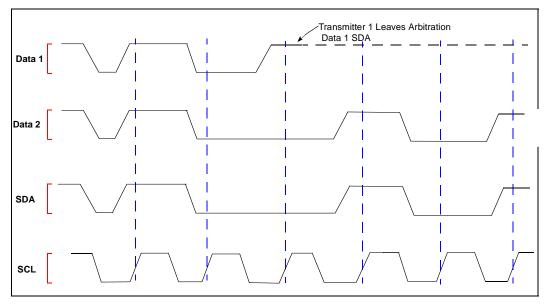


Figure 9-7. Arbitration Procedure of Two Masters

If the I^2C unit loses arbitration as the address bits are transferred and it is not addressed by the address bits, the I^2C unit resends the address when the I^2C bus becomes free. A resend is possible because the IDBR and ICR registers are not overwritten when arbitration is lost.

If the I²C unit loses arbitration because another bus master addresses the processor as a slave device, the I²C unit switches to slave-receive mode and overwrites the original data in the I²C data buffer register. Software can clear the start and re-initiate the master transaction.

Note: Software must prevent the I²C unit from starting a transaction to its own slave address because such a transaction puts the I²C unit in an indeterminate state.

Arbitration has boundary conditions in case an arbitration process is interrupted by a repeated START or STOP condition transmitted on the I^2C bus. To prevent errors, the I^2C unit acts as a master if no arbitration takes place in the following circumstances:

- Between a repeated START condition and a data bit
- Between a data bit and a STOP condition
- Between a repeated START condition and a STOP condition

These situations occur if different masters write identical data to the same target slave simultaneously and arbitration cannot be resolved after the first data byte transfer.

Note: Software ensures that arbitration is resolved quickly. For example, software can ensure that masters send unique data by requiring that each master transmit its I²C address as the first data byte of any transaction. When arbitration is resolved, the winning master sends a restart and begins a valid data transfer. The slave discards the master's address and use the other data.



9.4.5 Master Operations

When software initiates a read or write on the I^2C bus, the I^2C unit transitions from the default slave-receive mode to master-transmit mode. The 7-bit slave address and the R/nW bit follow the start pulse. After the master receives an acknowledge, the I^2C unit enters one of two master modes:

- Master-Transmit I²C unit writes data
- Master-Receive I²C unit reads data

The CPU writes to the ICR register to initiate a master transaction. Data is read and written from the I^2C unit through the memory-mapped registers. Table 9-5 describes the I^2C unit's responsibilities as a master device.

Table 9-5. Master Transactions (Sheet 1 of 2)

I ² C Master Action	Mode of Operation	Definition
Generate clock output	Master-transmit Master-receive	Master drives the SCL line. ICR[SCLE] bit must be set. ICR[IUE] bit must be set.
Write target slave address to IDBR	Master-transmit Master-receive	 CPU writes to IDBR bits 7-1 before a START condition enabled. First seven bits sent on bus after START. See Section 9.3.3.
Write R/nW Bit to IDBR	Master-transmit Master-receive	 CPU writes to least significant IDBR bit with target slave address. If low, master remains a master-transmitter. If high, master transitions to a master-receiver. See Section 9.4.2.
Signal START Condition	Master-transmit Master-receive	 See "Generate clock output" above. Performed after target slave address and R/nW bit are in IDBR. Software sets ICR[START] bit. Software sets ICR[TB] bit to initiate start condition. See Section 9.3.3.
Initiate first data byte transfer	Master-transmit Master-receive	 CPU writes byte to IDBR I²C unit transmits byte when ICR[TB] bit is set. I²C unit clears ICR[TB] bit and sets ISR[ITE] bit when transfer is complete.
Arbitrate for I ² C Bus	Master-transmit Master-receive	 If two or more masters signal a start within the same clock period, arbitration must occur. I²C unit arbitrates for as long as needed. Arbitration takes place during slave address and R/nW bit or data transmission and continues until all but one master loses the bus. No data lost. If I²C unit loses arbitration, it sets ISR[ALD] bit after byte transfer is completed and transitions to slave-receive mode. If I²C unit loses arbitration as it attempts to send target address byte, I²C unit attempts to resend it when the bus becomes free. System designer must ensure boundary conditions described in Section 9.4 do not occur.



Table 9-5. Master Transactions (Sheet 2 of 2)

I ² C Master Action	Mode of Operation	Definition
Write one data byte to the IDBR	Master-transmit only	 I²C master operation data transmit mode. Occurs when the ISR[ITE] bit is set and the ICR[TB] bit is clear. If the IDBR Transmit Empty Interrupt is enabled, it is signalled to the processor. CPU writes one data byte to the IDBR, sets the appropriate START/STOP bit combination, and sets the ICR[TB] bit to send the data. Eight bits are taken from the shift register and written to the serial bus. The eight bits are followed by a STOP, if requested.
Wait for Acknowledge from slave- receiver	Master-transmit only	 As a master-transmitter, the I²C unit generates the clock for the acknowledge pulse. The I²C unit releases the SDA line to allow slave-receiver ACK transmission. See Section 9.4.3.
Read one byte of I ² C Data from the IDBR	Master-receive only	 I²C master operation data receive mode. Eight bits are read from the serial bus, collected in the shift register then transferred to the IDBR after the ICR[ACKNAK] bit is read. The CPU reads the IDBR when the ISR[IRF] bit is set and the ICR[TB] bit is clear. If IDBR Receive Full Interrupt is enabled, it is signalled to the CPU. When the IDBR is read, if the ISR[ACKNAK] is clear (indicating ACK), the processor writes the ICR[ACKNAK] bit and set the ICR[TB] bit to initiate the next byte read. If the ISR[ACKNAK] bit is set (indicating NAK), ICR[TB] bit is clear, ICR[STOP] bit is set, and ISR[UB] bit is set, then the last data byte has been read into the IDBR and the I²C unit is sending the STOP. If the ISR[ACKNAK] bit is set (indicating NAK), ICR[TB] bit is clear, but the ICR[STOP] bit is clear, then the CPU has two options: 1. set the ICR[START] bit, write a new target address to the IDBR, and set the ICR[TB] bit which will send a repeated start condition or 2. set the ICR[MA] bit and leave the ICR[TB] bit clear which will send a STOP only.
Transmit Acknowledge to slave- transmitter	Master-receive only	 As a master-receiver, the I²C unit will generate the clock for the acknowledge pulse. The I²C unit is also responsible for driving the SDA line during the ACK cycle. If the next data byte is to be the last transaction, the CPU will set the ICR[ACKNAK] bit for NAK generation. See Section 9.4.3.
Generate a Repeated START to chain I ² C transactions	Master-transmit Master-receive	 If data chaining is desired, a repeated START condition is used instead of a STOP condition. This occurs after the last data byte of a transaction has been written to the bus. The CPU will write the next target slave address and the R/nW bit to the IDBR, set the ICR[START] bit, and set the ICR[TB] bit. See Section 9.3.3.
Generate a STOP	Master-transmit Master-receive	 Generated after the CPU writes the last data byte on the bus. CPU generates a STOP condition by setting the ICR[STOP] bit. See Section 9.3.3.

When the CPU needs to read data, the I^2C unit transitions from slave-receive mode to master-transmit mode to transmit the start address, R/nW bit, and the ACK pulse. After it sends the ACK pulse, the I^2C unit transitions to master-receive mode and waits to receive the read data from the slave device (see Figure 9-8).Multiple transactions can take place during an I^2C operation. For example, transitioning from master-receive to master-transmit through a repeated start.



Figure 9-8. Master-Receiver Read from Slave-Transmitter

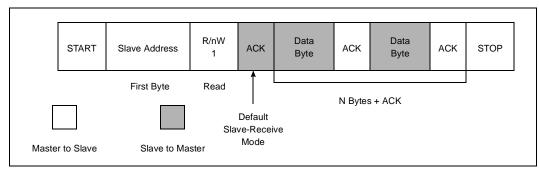


Figure 9-9. Master-Receiver Read from Slave-Transmitter / Repeated Start / Master-Transmitter Write to Slave-Receiver

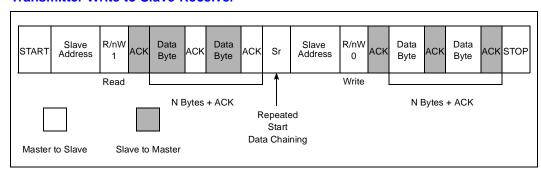
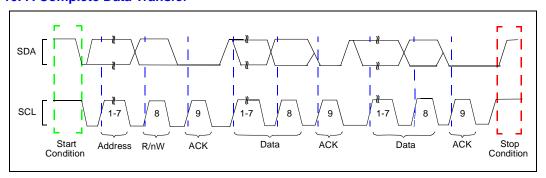


Figure 9-10. A Complete Data Transfer



9.4.6 Slave Operations

Table 9-6 describes how the I²C unit operates as a slave device.



Table 9-6. Slave Transactions

I ² C Slave Action	Mode of Operation	Definition
Slave-receive (default mode)	Slave-receive only	 I²C unit monitors all slave address transactions. ICR[IUE] bit must be set. I²C unit monitors bus for START conditions. When a START is detected, the interface reads the first 8 bits and compares the most significant seven bits with the 7-bit ISAR and the general call address (0x00). If there is a match, the I²C unit sends an ACK. If the first 8 bits are zero's, this is a general call address. If the ICR[GCD] bit is clear, both the ISR[GCAD] bit and the ISR[SAD] bit will be set. See Section 9.4.7. If the eighth bit of the first byte (R/nW bit) is low, the I²C unit stays in slave-receive mode and the ISR[SAD] bit is cleared. If R/nW bit is high, I²C unit switches to slave-transmit and ISR[SAD] bit is set.
Setting the Slave Address Detected bit	Slave-receive Slave-transmit	 Indicates the interface has detected an I²C operation that addresses the processor including the general call address. The processor can distinguish an ISAR match from a General Call by reading the ISR[GCAD] bit. An interrupt is signalled, if enabled, after the matching slave address is received and acknowledged.
Read one byte of I ² C Data from the IDBR	Slave-receive only	 Data receive mode of I²C slave operation. Eight bits are read from the serial bus into the shift register. When a full byte is received and the ACK/NAK bit is completed, the byte is transferred from the shift register to the IDBR. Occurs when the ISR[IRF] bit is set and the ICR[TB] bit is clear. If enabled, the IDBR Receive Full Interrupt is signalled to the CPU. Software reads one data byte from the IDBR. When the IDBR is read, the processor writes the desired ICR[ACKNAK] bit and sets the ICR[TB] bit. This causes the I²C unit to stop inserting wait states and let the master transmitter write the next piece of information.
Transmit Acknowledge to master- transmitter	Slave-receive only	 As a slave-receiver, the I²C unit pulls the SDA line low to generate the ACK pulse during the high SCL period. ICR[ACKNAK] bit controls the ACK data the I²C unit drives. See Section 9.4.3.
Write one byte of I ² C data to the IDBR	Slave-transmit only	 Data transmit mode of I²C slave operation. Occurs when ISR[ITE] bit is set and ICR[TB] bit is clear. If enabled, the IDBR Transmit Empty Interrupt is signalled to the processor. The processor writes a data byte to IDBR and sets ICR[TB] bit to start the transfer.
Wait for Acknowledge from master- receiver	Slave-transmit only	 As a slave-transmitter, the I²C unit releases the SDA line to allow the master-receiver to pull the line low for the ACK. See Section 9.4.3.

Figure 9-11 through Figure 9-13 are examples of I²C transactions and show the relationships between master and slave devices.



Figure 9-11. Master-Transmitter Write to Slave-Receiver

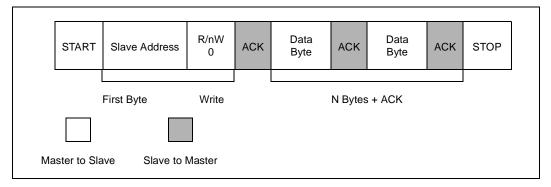


Figure 9-12. Master-Receiver Read to Slave-Transmitter

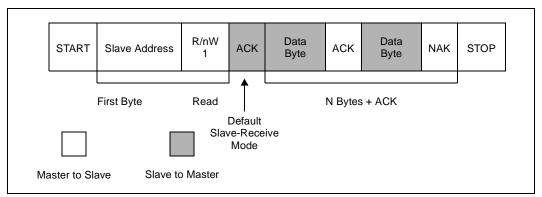
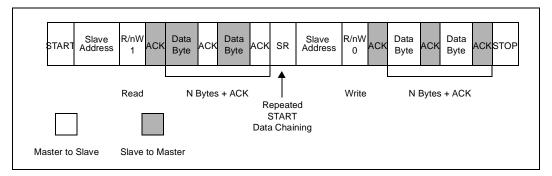


Figure 9-13. Master-Receiver Read to Slave-Transmitter, Repeated START, Master-Transmitter Write to Slave-Receiver



9.4.7 General Call Address

A general call address is a transaction with a slave address of 0x00. When a device requires the data from a general call address, it acknowledges the transaction and stays in slave-receiver mode. Otherwise, the device ignores the general call address. The other bytes in a general call transaction are acknowledged by every device that uses it on the bus. Devices that do not use these bytes must not send an ACK. The meaning of a general call address is defined in the second byte sent by the master-transmitter. Figure 9-14 shows a general call address transaction. The least significant bit of the second byte, called B, defines the transaction. Table 9-7 shows the valid values and definitions when B=0.



The I²C unit supports sending and receiving general call address transfers on the I²C bus. When software sends a general call message from the I²C unit, it must set the ICR[GCD] bit to prevent the I²C unit from responding as a slave. If the ICR[GCD] is not set, the I²C bus enters an indeterminate state.

If the I²C unit acts as a slave and receives a general call address while the ICR[GCD] bit is clear, it:

- Sets the ISR[GCAD] bit
- Sets the ISR[SAD] bit
- Interrupts the processor (if the interrupt is enabled)

If the I²C unit receives a general call address and the ICR[GCD] bit is set, it ignores the general call address.

Figure 9-14. General Call Address

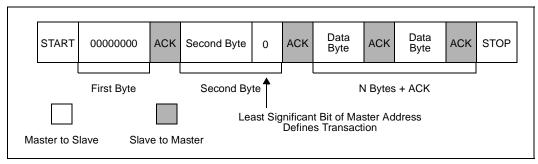


Table 9-7. General Call Address Second Byte Definitions

Least Significant Bit of Second Byte (B)	Second Byte Value	Definition
0	0x06	2-byte transaction in which the second byte tells the slave to reset and store this value in the programmable part of its address.
0	0x04	2-byte transaction in which the second byte tells the slave to store this value in the programmable part of its address. No reset.
0	0x00	Not allowed as a second byte
NOTE: Other v	alues are no	ot fixed and must be ignored.

Software must ensure that the I²C unit is not busy before it asserts a reset. Software must also ensure that the I²C bus is idle when the unit is enabled after reset. When directed to reset, the I²C unit, except for ISAR, returns to the default reset condition. ISAR is not affected by a reset.

When B=1, the sequence is a hardware general call and is not supported by the I^2C unit. Refer to the *The* I^2C -*Bus Specification* for information on hardware general calls.

I²C 10-bit addresses and CBUS compatibility are not supported.



9.5 Slave Mode Programming Examples

9.5.1 Initialize Unit

- 1. Set the slave address in the ISAR.
- 2. Enable desired interrupts in the ICR.
- 3. Set the ICR[IUE] bit to enable the I²C unit.

9.5.2 Write *n* Bytes as a Slave

- 1. When a Slave Address Detected interrupt occurs.
 Read ISR: Slave Address Detected (1), Unit Busy (1), R/nW bit (1), ACK/NAK (0)
- 2. Write a 1 to the ISR[SAD] bit to clear the interrupt.
- 3. Return from interrupt.
- 4. Load data byte to transfer in the IDBR.
- 5. Set ICR[TB] bit.
- 6. When a IDBR Transmit Empty interrupt occurs. Read ISR: IDBR Transmit Empty (1), ACK/NAK (0), R/nW bit (0)
- 7. Load data byte to transfer in the IDBR.
- 8. Set the ICR[TB] bit.
- 9. Write a 1 to the ISR[ITE] bit to clear interrupt.
- 10. Return from interrupt.
- 11. Repeat steps 6 to 10 for *n*-1 times. If, at any time, the slave does not have data, the I²C unit keeps SCL low until data is available.
- 12. When a IDBR Transmit Empty interrupt occurs.
 Read ISR: IDBR Transmit Empty (1), ACK/NAK (1), R/nW bit (0)
- 13. Write a 1 to the ISR[ITE] bit to clear interrupt.
- 14. Return from interrupt
- 15. When Slave Stop Detected interrupt occurs. Read ISR: Unit Busy (0), Slave STOP Detected (1)
- 16. Write a 1 to the ISR[SSD] bit to clear interrupt.

9.5.3 Read *n* Bytes as a Slave

- 1. When a Slave Address Detected interrupt occurs. Read ISR: Slave Address Detected (1), Unit busy (1), R/nW bit (0)
- 2. Write a 1 to the ISR[SAD] bit to clear the interrupt.
- 3. Return from interrupt.
- 4. Set ICR[TB] bit to initiate the transfer.



- 5. When an IDBR Receive Full interrupt occurs. Read ISR: IDBR Receive Full (1), ACK/NAK (0), R/nW bit (0)
- 6. Read IDBR to get the received byte.
- 7. Write a 1 to the ISR[IRF] bit to clear interrupt.
- 8. Return from interrupt.
- 9. Repeat steps 4 to 8 for *n*-1 times. Once the IDBR is full, the I²C unit will keep SCL low until the data is read.
- 10. Set ICR[TB] bit to release I²C bus and allow next transfer.
- 11. When a Slave Stop Detected interrupt occurs. Read ISR: Unit busy (0), Slave STOP Detected (1)
- 12. Write a 1 to the ISR[SSD] bit to clear interrupt.

9.6 Master Programming Examples

9.6.1 Initialize Unit

- 1. Set the slave address in the ISAR.
- 2. Enable desired interrupts in the ICR. Do not enable Arbitration Loss Detected interrupt
- 3. Set the ICR[IUE] and ICR[SCLE] bits to enable the I²C unit and SCL.

9.6.2 Write 1 Byte as a Master

- 1. Load target slave address and R/nW bit in the IDBR. R/nW must be 0 for a write.
- 2. Initiate the write.

 Set ICR[START], clear ICR[STOP], clear ICR[ALDIE], set ICR[TB]
- 3. When an IDBR Transmit Empty interrupt occurs. Read ISR: IDBR Transmit Empty (1), Unit Busy (1), R/nW bit (0)
- 4. Write a 1 to the ISR[ITE] bit to clear interrupt.
- 5. Write a 1 to the ISR[ALD] bit if set.

 If the master loses arbitration, it performs an address retry when the bus becomes free. The Arbitration Loss Detected interrupt is disabled to allow the address retry.
- 6. Load data byte to be transferred in the IDBR.
- 7. Initiate the write. Clear ICR[START], set ICR[STOP], set ICR[ALDIE], set ICR[TB]
- 8. When an IDBR Transmit Empty interrupt occurs (unit is sending STOP). Read ISR: IDBR Transmit Empty (1), Unit busy (x), R/nW bit (0)
- 9. Write a 1 to the ISR[ITE] bit to clear the interrupt.
- 10. Clear ICR[STOP] bit.



9.6.3 Read 1 Byte as a Master

- 1. Load target slave address and R/nW bit in the IDBR. R/nW must be 1 for a read.
- 2. Initiate the write.

 Set ICR[START], clear ICR[STOP], clear ICR[ALDIE], set ICR[TB]
- 3. When an IDBR Transmit Empty interrupt occurs. Read ISR: IDBR Transmit Empty (1), Unit busy (1), R/nW bit (1)
- 4. Write a 1 to the ISR[ITE] bit to clear the interrupt.
- 5. Initiate the read. Clear ICR[START], set ICR[STOP], set ICR[ALDIE], set ICR[ACKNAK], set ICR[TB]
- 6. When an IDBR Receive full interrupt occurs (unit is sending STOP). Read ISR: IDBR Receive Full (1), Unit Busy (x), R/nW bit (1), ACK/NAK bit (1)
- 7. Write a 1 to the ISR[IRF] bit to clear the interrupt.
- 8. Read IDBR data.
- 9. Clear ICR[STOP] and ICR[ACKNAK] bits

9.6.4 Write 2 Bytes and Repeated Start Read 1 Byte as a Master

- 1. Load target slave address and R/nW bit in the IDBR. R/nW must be 0 for a write.
- 2. Initiate the write.

 Set ICR[START], clear ICR[STOP], clear ICR[ALDIE], set ICR[TB]
- 3. When an IDBR Transmit Empty interrupt occurs. Read ISR: IDBR Transmit Empty (1), Unit Busy (1), R/nW bit (0)
- 4. Write a 1 to the ISR[ITE] bit to clear interrupt.
- 5. Load data byte to be transferred in the IDBR.
- 6. Initiate the write. Clear ICR[START], clear ICR[STOP], set ICR[ALDIE], set ICR[TB]
- 7. When an IDBR Transmit Empty interrupt occurs. Read ISR: IDBR Transmit Empty (1), Unit busy (1), R/nW bit (0)
- 8. Write a 1 to the ISR[ITE] bit to clear interrupt.
- 9. Repeat steps 5-8 one time.
- 10. Load target slave address and R/nW bit in the IDBR. R/nW must be 1 for a read.
- 11. Send repeated start as a master.
 Set ICR[START], clear ICR[STOP], clear ICR[ALDIE], set ICR[TB]
- 12. When an IDBR Transmit Empty interrupt occurs.

 Read ISR: IDBR Transmit Empty (1), Unit busy (1), R/nW bit (1)
- 13. Write a 1 to the ISR[ITE] bit to clear interrupt.
- 14. Initiate the read.

 Clear ICR[START], set ICR[STOP], set ICR[ALDIE], set ICR[ACKNAK], set ICR[TB]
- 15. When an IDBR Receive full interrupt occurs (unit is sending stop). Read ISR: IDBR Receive Full (1), Unit Busy (x), R/nW bit (1), ACK/NAK bit (1)



- 16. Write a 1 to the ISR[IRF] bit to clear the interrupt.
- 17. Read IDBR data.
- 18. Clear ICR[STOP] and ICR[ACKNAK] bits

9.6.5 Read 2 Bytes as a Master - Send STOP Using the Abort

- 1. Load target slave address and R/nW bit in the IDBR. R/nW must be 1 for a read.
- 2. Initiate the write.

 Set ICR[START], clear ICR[STOP], clear ICR[ALDIE], set ICR[TB]
- 3. When an IDBR Transmit Empty interrupt occurs. Read ISR: IDBR Transmit Empty (1), Unit busy (1), R/nW bit (1)
- 4. Write a 1 to the ISR[ITE] bit to clear interrupt.
- 5. Initiate the read Clear ICR[START], clear ICR[STOP], set ICR[ALDIE], clear ICR[ACKNAK], set ICR[TB]
- 6. When an IDBR Receive full interrupt occurs. Read ISR: IDBR Receive Full (1), Unit Busy (1), R/nW bit (1), ACK/NAK bit (0)
- 7. Write a 1 to the ISR[IRF] bit to clear the interrupt.
- 8. Read IDBR data.
- 9. Clear ICR[STOP] and ICR[ACKNAK] bits
- 10. Initiate the read.

 Clear ICR[START], clear ICR[STOP], set ICR[ALDIE], set ICR[ACKNAK], set ICR[TB] ICR[STOP] is not set because STOP or repeated start will be decided on the byte read.
- 11. When an IDBR Receive full interrupt occurs.

 Read ISR: IDBR Receive Full (1), Unit Busy (1), R/nW bit (1), ACK/NAK bit (1)
- 12. Write a 1 to the ISR[IRF] bit to clear the interrupt.
- 13. Read IDBR data.
- 14. Initiate STOP abort condition (STOP with no data transfer). Set ICR[MA]

Note: If a NAK is not sent in Step 11, the next transaction must involve another data byte read.

9.7 Glitch Suppression Logic

The I²C unit has built-in glitch suppression logic that suppresses glitches of 60ns or less. This is within the 50ns glitch suppression specification.

9.8 Reset Conditions

Software must ensure that the I^2C unit is not busy before it asserts a reset. Software must also ensure that the I^2C bus is idle when the unit is enabled after reset. When directed to reset, the I^2C unit, except for ISAR, returns to the default reset condition. ISAR is not affected by a reset.



When the ICR[UR] bit is set, the I^2C unit resets but the associated I^2C MMRs remain intact. When resetting the I^2C unit with the ICR's unit reset, use the following guidelines:

- 1. Set the reset bit in the ICR register and clear the remainder of the register.
- 2. Clear the ISR register.
- 3. Clear reset in the ICR.

9.9 Register Definitions

9.9.1 I²C Bus Monitor Register (IBMR)

The IBMR, shown in Table 9-8, tracks the status of the SCL and SDA pins. The values of these pins are recorded in this read-only IBMR so software can determine when the I²C bus is hung and the I²C unit must be reset.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 9-8. IBMR Bit Definitions

			PI			Ad _168	dres 80	SS		1 ² C Bus Monitor Register											I ² C Bus Interface Unit											
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	reserved																							SCLS	SDAS							
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1
1		31	:2			_			rese	erve	d																					
•		1				SC	LS		SCI	L Sta	atus	: Tr	is b	it co	ontin	uou	sly ı	refle	cts t	the	valu	e of	the	SC	L pi	n.						
		()			SD	AS		SD/	A St	atus	: Tł	nis b	it co	ontir	านดน	ısly	refle	ects	the	valu	ie of	the	SD	A pi	n.						

9.9.2 I²C Data Buffer Register (IDBR)

The processor uses the IDBR, shown in Table 9-9, to transmit and receive data from the I²C bus. The IDBR is accessed by the program I/O on one side and by the I²C shift register on the other. The IDBR receives data coming into the I²C unit after a full byte is received and acknowledged. The processor core writes data going out of the I²C unit to the IDBR and sends it to the serial bus.

When the I²C unit is in transmit mode (master or slave), the processor writes data to the IDBR over the internal bus. The processor writes data to the IDBR when a master transaction is initiated or when the IDBR Transmit Empty Interrupt is signalled. Data moves from the IDBR to the shift register when the Transfer Byte bit is set. The IDBR Transmit Empty Interrupt is signalled (if enabled) when a byte is transferred on the I²C bus and the acknowledge cycle is complete. If the IDBR is not written by the processor and a STOP condition is not in place before the I²C bus is ready to transfer the next byte packet, the I²C unit inserts wait states until the processor writes the IDBR and sets the Transfer Byte bit.

When the I²C unit is in receive mode (master or slave), the processor reads IDBR data over the internal bus. The processor reads data from the IDBR when the IDBR Receive Full Interrupt is signalled. The data moves from the shift register to the IDBR when the ACK cycle is complete. The I²C unit inserts wait states until the IDBR is read. Refer to Section 9.4.3 for more information



on the acknowledge pulse in receiver mode. After the processor reads the IDBR, the ACK/NAK Control bit is written and the Transfer Byte bit is written, allowing the next byte transfer to proceed to the $\rm I^2C$ bus. The IDBR register is 0x00 after reset.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 9-9. IDBR Bit Definitions

			P		ical)30_			SS		I ² C Data Buffer Register												I ² C Bus Interface Unit										
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	reserved																IDB															
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		31	:8			_	_		res	erve	ed																					
		7	:0			ID	В		I ² C Data Buffer: Buffer for I ² C bus send/receive data.																							

9.9.3 I²C Control Register (ICR)

The processor uses the ICR, shown in Table 9-10, to control the I²C unit.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 9-10. ICR Bit Definitions (Sheet 1 of 3)

			PI		ical)30_			S					I ² C	Co	ntro	l R	egis	ter						l ² C	Bus	Int	erfa	ice	Unit	:		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							r	ese	rved	ł							ЬM	UR	SADIE	ALDIE	SSDIE	BEIE	IRFIE	ITEIE	GCD	IUE	SCLE	MA	TB	ACKNAK	STOP	START
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
		31:	16			-	_			erve	-																					
			_			_				t M		-																				
		1	5			F	M		0 = 1 =	10 40	O KI	Bit/s Bit/s	ec. ec.	ope ope	ratic ratic	n n																
									Uni	t Re	set	:																				
		1.	4			U	R			No Re			I ² C	unit	onl	y.																
									Sla	ve A	١dd	res	s De	etec	ted	Inte	rruj	ot E	nab	le:												
		1	3			SAI	DIE				able	es th	ne l ²			o int	erru	pt th	ne p	roce	ssor	whe	en it	dete	ects	a sl	lave	ado	dres	s ma	atch	or
									Arb	itra	tior	Lo	ss l	Dete	ecte	d In	terr	upt	Ena	able	:											
		1:	2			ALI	DIE			Dis En					nit to	o int	erru	pt th	ne p	roce	esso	r wh	nen	it los	ses	arbit	tratio	on ir	n ma	ste	· mo	de.
									Sla	ve S	то	P D	ete	cted	Int	erru	ıpt E	nal	ble:													
		1	1			SSI	DIE				able		ne l ²		nit t	o int	terru	ıpt tl	he p	roc	esso	or w	hen	it de	etec	ts a	ST	OP (cond	ditio	n in	



Table 9-10. ICR Bit Definitions (Sheet 2 of 3)

Physical Address 4030_1690 I²C Control Register I²C Bus Interface Unit

Bit	31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	reserved	FM	NR .	SADIE	ALDIE	SSDIE	BEIE	IRFIE	ITEIE	GCD	IUE	SCLE	MA	TB	ACKNAK	STOP	START
_		_	_	_	_	_	_				/ _ /	_	_				

t 0	0 0 0	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
			Bus Error Interrupt Enable:
			0 = Disable interrupt. 1 = Enables the I ² C unit to interrupt the processor for the following I ² C bus errors:
	10	BEIE	As a master transmitter, no ACK was detected after a byte was sent.
			 As a slave receiver, the I²C unit generated a NAK pulse.
			NOTE: Software is responsible for guaranteeing that misplaced START and STOP conditions do not occur. See Section 9.7.
			IDBR Receive Full Interrupt Enable:
	9	IRFIE	 0 = Disable interrupt. 1 = Enables the I²C unit to interrupt the processor when the IDBR receives a data byte from the I²C bus.
			IDBR Transmit Empty Interrupt Enable:
	8	ITEIE	 0 = Disable interrupt. 1 = Enables the I²C unit to interrupt the processor after transmitting a byte onto the I²C bus.
			General Call Disable:
	7	GCD	$0 = $ Enables the I^2C unit to respond to general call messages.
			1 = Disables I ² C unit response to general call messages as a slave.
-			Must be set when the I ² C unit sends a master mode general call message.
			I ² C Unit Enable:
	6	IUE	 0 = Disables the unit and does not master any transactions or respond to any slave transactions. 1 = Enables the I²C unit (defaults to slave-receive mode).
			Software must ensure that the I ² C bus is idle before it sets this bit.
-			SCL Enable:
	5	SCLE	$0 = \text{Disables the } I^2\text{C}$ unit from driving the SCL line.
			1 = Enables the I ² C clock output for master mode operation.
			Master Abort : generates a STOP without transmitting another data byte when the I ² C unit is in master mode.
			0 = The I ² C unit transmits STOP using the STOP ICR bit only. 1 = The I ² C unit sends STOP without data transmission.
	4	MA	In master-transmit mode, after a data byte is sent, the ICR's Transfer Byte bit is cleared and IDBR Transmit Empty bit is set. When no more data bytes need to be sent, setting master abort bit sends the STOP. The Transfer Byte bit (03) must remain clear.
			In master-receive mode, when a NAK is sent without a STOP (STOP ICR bit was not set) and the processor does not send a repeated START, setting this bit sends the STOP. Once again, the Transfer Byte bit (03) must remain clear.
			Transfer Byte : used to send/receive a byte on the I ² C bus.
	3	ТВ	0 = Cleared by I ² C unit when the byte is sent/received. 1 = Send/receive a byte.
	-		The processor can monitor this bit to determine when the byte transfer is completed. In master or slave mode, after each byte transfer, including ACK/NAK bit, the I ² C unit holds the SCL line low (inserting wait states) until the Transfer Byte bit is set.



Table 9-10. ICR Bit Definitions (Sheet 3 of 3)

			P		cal Ad 30_16		SS					I ² C	Co	ntro	ol R	egis	ter						l ² C	Bus	Int	erfa	ce l	Unit	t		
Bit	31	30	29	28	27 26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						r	esei	rved	ł							FIM	N	SADIE	ALDIE	SSDIE	BEIE	IRFIE	ITEIE	GCD	IUE	SCLE	MA	TB	ACKNAK	STOP	START
Reset	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	•	
		;	2		ACH	(NAI	<	rece 0 = 1 = The whe	Th Th Th I ² C en it	mode I ² (e I ² (uni resp	de. Cui Cui tau oon	nit s nit s utom ds i	end end natic n sla	s ar s a ally ave-	n AC neg sen rece	CK p ative ids a eive	ulse e AC an A mod	afte CK (I CK de, i	er it NAk puls nde	rece () af se w pen	eives ter it hen dent	s a d t red it re	data ceive espo the /	bytes a ands	e. data to i	a by ts sl K co	te. ave	ado	mast dress it se	s or tting	J.
			1		ST	ОР		mas bit. 0 =	ster See Do	moc Sec	le. I ctio sei	In m n 9. nd a	aste 3.3.3 ST	er-re	ceiv	ve m	ode	, the	e AC	CK/N	١AK	con	trol						sferre		
			0		ST	ART		Sec 0 =	tion Do	9.3	.3.1 sei	for nd a	STA moi STA T.	re d	etail							ien i	n m	aste	er m	ode	. Se	е			

9.9.4 I²C Status Register (ISR)

The ISR, shown in Table 9-11, signals I^2C interrupts to the processor interrupt controller. Software can use the ISR bits to check the status of the I^2C unit and bus. ISR bits (bits 9-5) are updated after the ACK/NAK bit is completed on the I^2C bus.

The ISR also clears the following interrupts signalled from the I²C unit:

- IDBR Receive Full
- IDBR Transmit Empty
- Slave Address Detected
- Bus Error Detected
- STOP Condition Detect
- Arbitration Lost

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 9-11. ISR Bit Definitions (Sheet 1 of 2)

			P			al Ad 0_16	ddre 198	SS					I ² C	Sta	atus	Re	gis	er						I ² C	Bus	Int	erfa	ice	Unit			
Bit	31	30	29	28	27	7 26	25	2	4 2	3 22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										res	serv	ed										BED	GYS	GCAD	IRF	ЭШ	ALD	OSS	IBB	BՈ	ACKNAK	RWM
Reset	0	0	0	0	0	0	0	(0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		3	1:11				_		re	serve	ed																					
			10			E	BED		0 1 • • •	us Er = No = Th As As OTE:	err a m a sl Whe mi	or de C un aster ave r en an splac s bit,	tece it se r tra ece eri ed wri	ets tension of the service of the se	this I mitte r, the occu ART a 1 to	r, no e l ² (rs, l and	o A0 C ur ² C l	CK nit g	is de gene s trai	etec erate	ted s a tion	on th NAK s co	ne ir (pu ntin	nterfa Ise. ue. S	ace Soft	afte ware	er a l	byte ust e	e is s ensu	ent. ire tl		
			9			S	SAD		0	ave A = No = I ² O int clea	sla C un erru	ve ad it det pt is	ddre ect	ess ed a	dete a 7-l ed w	oit a /her	ddr	ess e S	tha ADII	t ma E int	tche erru	es th pt is	e ge	ener to a	al ca a 1.	all a	ıddr	ess	or IS	SAR	. Ar	1
			8			G	CAD		G	enera = No = I ² (al Ca	all Ao	d dr ca	ess	De	t ect	ecei			ess.												
			7			I	RF		0 1	BR F = Th = Th	Rece le ID le ID len l	BR IDBR IDBR IDBR IDBR	ras regi RFI	l: not ster E is	rec	eive eive to a	ed a	ne	w da	ata b	yte	or th	ne I ²	² C u ∋ I ² C	nit is	s idl	e. In in	terri	upt i	s siç	gnal	led
			6			ı	TE		0	BR T = Th = Th sig	ie da ie l ² gnall	ata by C un ed w	yte it ha her	is s as fi n the	till b inish e ITI	ed EIE	tran	sm	ittin	gad			e on	the	I ² C	bus	s. Ar	n int	erru	pt is	i	
			5			P	λLD		A 0 1	bitra = Cl = Se clea	tior eare	Los ed wh	s s C nen ne l	ete arb ² C	ecte itrati unit	d: u ion i lose	s w	on	or n	evei					ion.							
									SI	ave S	STO	P De	tec	ted	:																	

0 = No STOP detected.

I²C Bus Busy:

Unit Busy:

To clear this bit, write a 1 to it.

1 = Set when the I²C unit detects a STOP while in slave-receive or slave-transmit mode.

 $0 = I^2C$ unit not busy. 1 = Set when the I^2C unit is busy. Defined as the time between the first START and STOP.

 $0 = I^2C$ bus is idle or the I^2C unit is using the bus (i.e., unit busy). 1 = Set when the I^2C bus is busy but the I^2C unit is not involved in the transaction.

4

3

2

SSD

IBB

UB



Table 9-11. ISR Bit Definitions (Sheet 2 of 2)

			P	•		Add _169		SS					l ² (C St	atus	s Re	egis	ter						l ² C	Bus	Int	erfa	ce	Unit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										res	serv	ed										BED	SAD	GCAD	IRF	ITE	ALD	SSD	IBB	NB	ACKNAK	RWM
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		,	1		ļ	ACK	NAI	<	0 = 1 = Use	I ² (ed in	un un sla	it re it re ve-t	ceiv ceiv	ed smit	or se mo	ent a de t	a N <i>i</i> o de	AK. eterr	nine	wh	bus. en t	he ti				oyte	is th	ne la	ast c	one.		
		()			RV	VM		0 = 1 =	I ² (C un	it is it is	in n in n	nast nast	er-r	ecei	ive (or sl	ave-	-trar	ceive nsmi	t mc	ode.	ard	ware	e aft	er a	sto	p st	ate.		

9.9.5 I²C Slave Address Register (ISAR)

The ISAR, shown in Table 9-12, defines the I^2C unit's 7-bit slave address. In slave-receive mode, the processor responds when the 7-bit address matches the value in this register. The processor writes this register before it enables I^2C operations. The ISAR is fully programmable (no address is assigned to the I^2C unit) so it can be set to a value other than those of hard-wired I^2C slave peripherals in the system. If the processor is reset, the ISAR is not affected. The ISAR register default value is 0000000_2 .

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 9-12. ISAR Bit Definitions

			P		ical)30_			SS			l ²	C S	lave	Add	Ires	s R	egis	ster					I ² C	Bus	Int	erfa	ce	Unit	t		
Bit	31	30	29	28	27	26	25	24 2	3 2	2 2	21 20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											re	ser	ved															ISA	1		
Reset	0	0	0	0	0	0	0	0 () () (0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		31	1:7			_	_	re	eser	ved	i																				
		6	:0			IS	SA.		C S		e Ad	dre	ss: 7	7-bit	add	ress	tha	t the	e I ² (C un	it re	spo	nds	to v	vher	n in s	slav	e-re	cei	/e	





UARTs 10

This chapter describes the universal asynchronous receiver/transmitter (UART) serial ports. The serial ports are controlled via direct memory access (DMA) or programmed I/O. The PXA255 processor has four UARTs: Full Function UART (FFUART), Bluetooth UART (BTUART), Standard UART (STUART) and Hardware UART (HWUART). The HWUART is covered in Chapter 17. The UARTs use the same programming model.

10.1 Feature List

The UARTs share the following features:

- Functionally compatible with the 16550
- Ability to add or delete standard asynchronous communications bits (start, stop, and parity) in the serial data
- Independently controlled transmit, receive, line status, and data set interrupts
- Programmable baud rate generator that allows the internal clock to be divided by 1 to (2¹⁶–1) to generate an internal 16X clock
- Modem control pins that allow flow control through software. Each UART has different modem control capability.
- Fully programmable serial-interface:
 - 5-, 6-, 7-, or 8-bit characters
 - Even, odd, and no parity detection
 - 1, 1.5, or 2 stop bit generation
 - Baud rate generation up to 921 Kbps for the BTUART and HWUART. Up to 230 Kbps for other UARTs.
- 64-byte transmit FIFO
- 64-byte receive FIFO
- Complete status reporting capability
- Ability to generate and detect line breaks
- Internal diagnostic capabilities that include:
 - Loopback controls for communications link fault isolation
 - Break, parity, and framing error simulation
- Fully prioritized interrupt system controls
- Separate DMA requests for transmit and receive data services
- Slow infrared asynchronous interface that conforms to the Infrared Data Association (IRDA) standard



10.2 Overview

Each serial port contains a UART and a slow infrared transmit encoder and receive decoder that conforms to the IRDA Serial Infrared (SIR) Physical Layer Link Specification.

Each UART performs serial-to-parallel conversion on data characters received from a peripheral device or a modem and parallel-to-serial conversion on data characters received from the processor. The processor can read a UART's complete status during functional operation. Status information includes the type and condition of transfer operations and error conditions (parity, overrun, framing, or break interrupt) associated with the UART.

Each serial port operates in FIFO or non-FIFO mode. In FIFO mode, a 64-byte Transmit FIFO holds data from the processor until it is transmitted on the serial link and a 64-byte Receive FIFO buffers data from the serial link until it is read by the processor. In non-FIFO mode, the transmit and Receive FIFOs are bypassed.

Each UART includes a programmable baud rate generator that can divide the input clock by 1 to $(2^{16}-1)$. This produces a 16X clock that can be used to drive the internal transmitter and receiver logic. Software can program interrupts to meet its requirements. This minimizes the number of computations required to handle the communications link. Each UART operates in an environment that is controlled by software and can be polled or is interrupt driven.

10.2.1 Full Function UART

The FFUART supports modem control capability. The maximum tested baud rate on this UART is 230.4 kbps. The divisor programmed in the divisor latch registers must be equal to or greater than four.

10.2.2 Bluetooth UART

The BTUART is a high speed UART that supports baud rates up to 921.6 kbps and can be connected to a Bluetooth module. It supports the functions in the feature list, (see Section 10.1) but only supports two modem control pins (nCTS, nRTS).

10.2.3 Standard UART

The STUART supports all functions in the feature list (see Section 10.1), but does not support modem control capability. The UART's maximum tested baud rate is 230.4 kbps. The divisors programmed in divisor latch registers must be equal to or greater than four.

10.2.4 Compatibility with 16550

The processor UARTs are functionally compatible with the 16550 industry standard. Each UART supports all of the 16550's functions as well as the following features:

- DMA requests for transmit and receive data services
- Slow infrared asynchronous interface
- Non-Return-to-Zero (NRZ) encoding/decoding function



10.3 Signal Descriptions

Table 10-1 lists and describes each external signal that is connected to a UART module. The pins are connected through the System Integration Unit to GPIOs. Refer to Section 4.1, "General-Purpose I/O" on page 4-1 for details on the GPIOs.

Table 10-1. UART Signal Descriptions (Sheet 1 of 2)

Name	Туре	Description
RXD	Input	SERIAL INPUT: Serial data input to the receive shift register. In infrared mode, it is connected to the infrared receiver input. This signal is present on all three UARTs.
TXD	Output	SERIAL OUTPUT: Serial data output to the communications link- peripheral, modem, or data set. The TXD signal is set to the logic 1 state upon a Reset operation. It is connected to the output of the infrared transmitter in infrared mode. This signal is present all three UARTs.
nCTS	Input	CLEAR TO SEND: When low, indicates that the modem or data set is ready to exchange data. The nCTS signal is a modem status input and its condition can be tested by reading bit 4 (CTS) of the Modem Status Register. Bit 4 is the complement of the nCTS signal. Bit 0 (DCTS) of the Modem Status Register (MSR) indicates whether the nCTS input has changed state since the last time the Modem Status Register was read. nCTS has no effect on the transmitter. This signal is present on the FFUART and BTUART. When the CTS bit of the MSR changes state and the Modem Status intervent in capabled, an intervent in page total.
nDSR	Input	interrupt is enabled, an interrupt is generated. DATA SET READY: When low, indicates that the modem or data set is ready to establish a communications link with a UART. The nDSR signal is a Modem Status input and its condition can be tested by reading Bit 5 (DSR) of the MSR. Bit 5 is the complement of the nDSR signal. Bit 1 (DDSR) of the MSR indicates whether the nDSR input has changed state since the MSR was last read. This signal is only present on the FFUART. When the DSR bit of the MSR changes state, an interrupt is generated if the Modem Status interrupt is enabled.
nDCD	Input	DATA CARRIER DETECT: When low, indicates that the data carrier has been detected by the modem or data set. The nDCD signal is a modem status input and its condition can be tested by reading Bit 7 (DCD) of the MSR. Bit 7 is the complement of the nDCD signal. Bit 3 (DDCD) of the MSR indicates whether the nDCD input has changed state since the previous reading of the Modem Status Register. nDCD has no effect on the receiver. This signal is only present on the FFUART. When the DCD bit changes state and the Modem Status interrupt is enabled, an interrupt is generated.



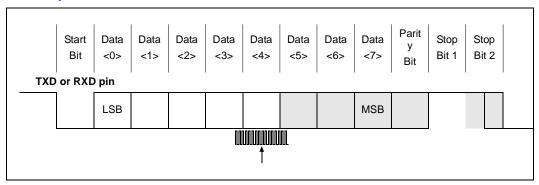
Table 10-1. UART Signal Descriptions (Sheet 2 of 2)

Name	Туре	Description
nRI	Input	RING INDICATOR: When low, indicates that the modem or data set has received a telephone ringing signal. The nRI signal is a Modem Status input whose condition can be tested by reading Bit 6 (RI) of the MSR. Bit 6 is the complement of the nRI signal. Bit 2, the trailing edge of ring indicator (TERI), of the MSR indicates whether the nRI input signal has changed from low to high since the MSR was last read. This signal is only present on the FFUART. When the RI bit of the MSR changes from a high to low state and the Modem Status interrupt is enabled, an interrupt is generated.
nDTR	Output	DATA TERMINAL READY: When low, signals the modem or the data set that the UART is ready to establish a communications link. The nDTR output signal can be set to an active low by programming Bit 0 (DTR) of the MSR to a 1. A Reset operation sets this signal to its inactive state. LOOP mode operation holds this signal in its inactive state. This signal is only present on the FFUART.
nRTS	Output	REQUEST TO SEND: When low, signals the modem or the data set that the UART is ready to exchange data. The nRTS output signal can be set to an active low by programming Bit 1 (RTS) of the Modem Control Register to a 1. A Reset operation sets this signal to its inactive (high) state. LOOP mode operation holds this signal in its inactive state. This signal is used by the FFUART and BTUART.

10.4 UART Operational Description

The format of a UART data frame is shown in Figure 10-1.

Figure 10-1. Example UART Data Frame



Receive data sample counter frequency is 16 times the value of the bit frequency. The 16X clock is created by the baud rate generator. Each bit is sampled three times in the middle. Shaded bits in Figure 10-1 are optional and can be programmed by software.

Each data frame is between seven and 12 bits long, depending on the size of the data programmed, whether parity is enabled, and the number of stop bits. A data frame begins by transmitting a start bit that is represented by a high to low transition. The start bit is followed by from five to eight bits of data that begin with the least significant bit (LSB). The data bits are followed by an optional parity bit. The parity bit is set if even parity is enabled and the data byte has an odd number of ones

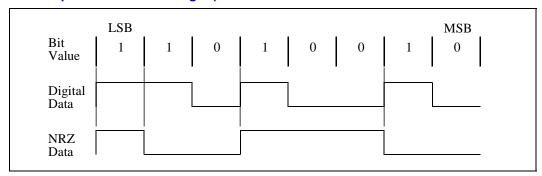


or if odd parity is enabled and the data byte has an even number of ones. The data frame ends with one, one and a half or two stop bits, as programmed by software. The stop bits are represented by one, one and a half, or two successive bit periods of a logic one.

Each UART has two FIFOs: one transmit and one receive. The transmit FIFO is 64 bytes deep and eight bits wide. The receive FIFO is 64 bytes deep and 11 bits wide. Three bits are used for tracking errors.

The UART can use NRZ coding to represent individual bit values. NRZ coding is enabled when the Interrupt Enable Register's (IER) bit 5, IER[5] is set to high. A one is represented by a line transition and a zero is represented by no line transition. Figure 10-2 shows the data byte 0b 0100 1011 in NRZ coding. The byte's LSB is transmitted first.

Figure 10-2. Example NRZ Bit Encoding - (0b0100 1011



10.4.1 Reset

The UARTs are disabled on reset. To enable a UART, Software must program the GPIO registers (see Section 4.1, "General-Purpose I/O" on page 4-1) then set IER[UUE]. When the UART is enabled, the receiver waits for a frame start bit and the transmitter sends data if it is available in the Transmit Holding Register. Transmit data can be written to the Transmit Holding Register before the UART unit is enabled. In FIFO mode, data is transmitted from the FIFO to the Transmit Holding Register before it goes to the pin.

When the UART unit is disabled, the transmitter or receiver finishes the current byte and stops transmitting or receiving more data. Data in the FIFO is not cleared and transmission resumes when the UART is enabled.

10.4.2 Internal Register Descriptions

Each UART has 13 registers: 12 registers for UART operation and one register for slow infrared configuration. The registers are all 32-bit registers but only the lower eight bits have valid data. The 12 UART operation registers share nine address locations in the I/O address space. Table 10-2 shows the registers and their addresses as offsets of a base address. The base address for each UART is 32 bits. The state of the SLCR[DLAB] bit affects the selection of some UART registers. To access the Baud Rate Generator Divisor Latch registers, software must set the SLCR[DLAB] bit high.



Table 10-2. UART Register Addresses as Offsets of a Base

UART Register Addresses (Base + offset)	DLAB Bit Value	Register Accessed
Base	0	Receive Buffer (read only)
Base	0	Transmit Buffer (write only)
Base + 0x04	0	Interrupt Enable (read/write)
Base + 0x08	Х	Interrupt Identification (read only)
Base + 0x08	Х	FIFO Control (write only)
Base + 0x0C	Х	Line Control (read/write)
Base + 0x10	Х	Modem Control (read/write)
Base + 0x14	Х	Line Status (read only)
Base + 0x18	Х	Modem Status (read only)
Base + 0x1C	Х	Scratch Pad (read/write)
Base + 0x20	Х	Infrared Selection (read/write)
Base	1	Divisor Latch Low (read/write)
Base + 0x04	1	Divisor Latch High (read/write)

10.4.2.1 Receive Buffer Register (RBR)

In non-FIFO mode, the RBR, shown in Table 10-3, holds the character received by the UART's Receive Shift Register. If the RBR is configured to use fewer than eight bits, the bits are right-justified and the most significant bits (MSB) are zeroed. Reading the register empties the register and clears the Data Ready (DR) bit in the Line Status Register (LSR) to a 0.

In FIFO mode, the RBR latches the value of the data byte at the front of the FIFO.

This is a read-only register. Ignore reads from reserved bits.

Table 10-3. RBR Bit Definitions

				Bas	e (D	LAE	3=0))					Rec	eive	But	ffer	Reg	iste	r							UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											ı	rese	rve	d											RBR7	RBR6	RBR5	RBR4	RBR3	RBR2	RBR1	RBR0
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
•		В	its			Na	me												D	esci	riptio	on										
•		31	1:8			_	_		res	erve	ed																					
•		7	:0		F	RBR	[7:0]	Dat	ta b	yte	rece	ivec	llea	st s	ignii	ficar	nt bit	firs	t.												



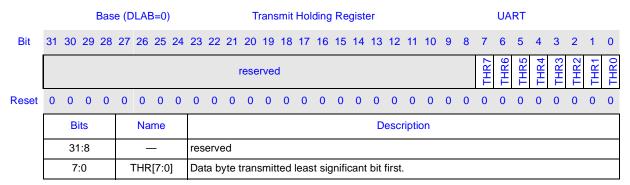
10.4.2.2 Transmit Holding Register (THR)

In non-FIFO mode, the THR, shown in Table 10-4, holds the data byte that is to be transmitted next. When the TSR is emptied, the contents of the THR are loaded in the TSR and the LSR[TDRQ] is set to a 1.

In FIFO mode, a write to the THR puts data into the top of the FIFO. The data at the front of the FIFO is loaded to the TSR when that register is empty.

This is a write-only register. Write zeros to reserved bits.

Table 10-4. THR Bit Definitions



10.4.2.3 Divisor Latch Registers (DLL and DLH)

Each UART contains a programmable baud rate generator that can take the 14.7456 MHz fixed input clock and divide it by 1 to $(2^{16}-1)$. For the FFUART and the STUART, the divisor is from 4 to $(2^{16}-1)$. The baud rate generator output frequency is 16 times the baud rate. Two 8-bit latch registers, shown in Table 10-5 and Table 10-6, store the divisor in a 16-bit binary format. Load these divisor latches during initialization to ensure that the baud rate generator operates properly. If each Divisor Latch is loaded with a 0, the 16X clock stops. The Divisor Latches are accessed with a word write.

The baud rate of the data shifted in to or out of a UART is given by the formula:

$$BaudRate = \frac{14.7456 \text{ MHz}}{(16xDivisor)}$$

For example: if the divisor is 24, the baud rate is 38400 bps.

The divisor's reset value is 0x0002. For the FFUART and the STUART, the divisor must be set to at least 0x0004 before the UART unit is enabled.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 10-5. DLL Bit Definitions

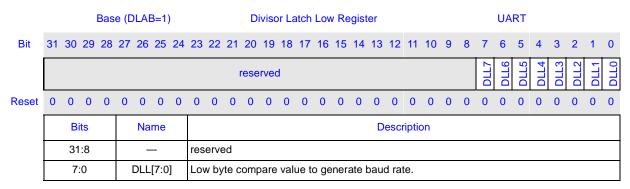


Table 10-6. DLH Bit Definitions

			Bas	se+0)x04	(DI	LAB	=1)				Di	viso	r La	tch	Higl	n Re	gist	er							UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											ı	rese	rve	d											DLH15	DLH14	DLH13	DLH12	DLH11	DLH10	БГН9	DLH8
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
·		В	its			Na	me												D	esci	iptio	on										
		31	:8			_			res	erve	ed																					
•		7	:0		[DLH	[7:0]	Hig	h by	yte o	com	pare	e val	lue t	o ge	ener	ate	bau	d ra	te.											

10.4.2.4 Interrupt Enable Register (IER)

The IER, shown in Table 10-7, enables the five types of interrupts that set a value in the Interrupt Identification Register (IIR). To disable an interrupt, software must clear the appropriate bit in the IER. Software can enable some interrupts by setting the appropriate bit.

The Character Timeout Indication interrupt is separated from the Received Data Available interrupt to ensure that the processor and the DMA controller do not service the receive FIFO at the same time. When a Character Timeout Indication interrupt occurs, the processor must handle the data in the Receive FIFO through programmed I/O.

An error interrupt is used when DMA requests are enabled. The interrupt is generated when LSR bit 7 is set to a 1, because a receive DMA request is not generated when the receive FIFO has an error. The error interrupt tells the processor to handle the data in the receive FIFO through programmed I/O. The error interrupt is enabled when DMA requests are enabled and it can not be masked. Receiver Line Status interrupts occur when the error is at the front of the FIFO.

Note: When DMA requests are enabled and an interrupt occurs, software must first read the LSR to see if an error interrupt exists, then check the IIR for the source of the interrupt. When the last error byte is read from the FIFO, DMA requests are automatically enabled. Software is not required to check for the error interrupt if DMA requests are disabled because an error interrupt only occurs when DMA requests are enabled.



Bit 7 of the IER is used to enable DMA requests. The IER also contains the unit enable and NRZ coding enable control bits. Bits 7 through 4 are used differently from the standard 16550 register definition.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 10-7. IER Bit Definitions

			Bas	se+C)x04	(DL	.AB	=0)				lr	nterr	upt	Ena	able	Reg	jiste	er							UA	RT				
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 0
											r	ese	rvec	l											DMAE	UUE	NRZE	RTOIE	MIE	RLSE	TIE
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
		Bi	ts			Nar	ne												D	esc	ripti	on									
		31	:8						rese	erve	d																				
		7	7			DM.	ΑE		0 –	D١		equ	ests	are	e dis	sable															
		6	6			UU	JE		0 –	Th	Jnit e ur e ur	it is	dis	able																	
		Ę	5			NR	ZE		mod 0 –	de. I NF	•	slo odir	ow ii ng d	nfrai isab	red oled	rece		_		_		•			UAF NRZ						ared
		2	1			RTC	DIE		0 –	Ch	ara	ter	Tim	eou	ıt In	ion l idica	tion	inte	errup	ot di	sab										
		3	3			MI	E		0 –	Mo		n St	tatus	int	erru	ıpt d															
		2	2			RLS	SE		0 –	Re	ceiv	er I	_ine	Sta	tus	rupt inte inte	rrup	t dis	sable												
		,	1			TI	E		0 –	Tra	ansn	nit F	-IFC) Da	ıta I	errup Requ Requ	uest	inte	errup												
		()			RA\	/IE		0 –	Re	ceiv	er l	Data	Av	aila		(Trig	ger	leve	el re					t dis						

Note: To ensure that the DMA controller and programmed I/O do not access the same FIFO, software must not set the DMAE while the TIE or RAVIE bits are set to a 1.

10.4.2.5 Interrupt Identification Register (IIR)

The UART prioritizes interrupts in four levels (see Table 10-8) and records them in the IIR. The IIR, shown in Table 10-9, stores information that indicates that a prioritized interrupt is pending and identifies the source of the interrupt.



In FIFO mode, the "Received Data is available" interrupt (Priority Level 2) takes priority over the "Character Timeout Indication" interrupt (Priority Level 2). For example, if the UART is in FIFO mode and FIFO Control Register[ITL] = 0b00, this will cause the UART to generate an interrupt when there is one byte in the FIFO. In this scenario, if there is one byte in the FIFO, an interrupt is generated, and IIR[3:0] = 0b0100, which indicates that Received Data is available. If data remains in the FIFO and if a Character Timeout occurs (no data has been sent for 4 character times), then the interrupt status does not change to IIR[3:0] = 0b1100 (Character Timeout Indication).

The error interrupt is reported separately in the LSR. In DMA mode, software must check for the error interrupt before it checks the IIR.

If additional data is received before a Character Timeout Indication interrupt is serviced, the interrupt is deasserted.

Table 10-8. Interrupt Conditions

Priority Level	Interrupt origin
1 (highest)	Receiver Line Status: one or more error bits were set.
2	Received Data is available. In FIFO mode, trigger level was reached. In non-FIFO mode, RBR has data.
2	Character Timeout Indication occurred. Occurs only in FIFO mode, when data is in the receive FIFO but no data has been sent for a set time period.
3	Transmitter requests data. In FIFO mode, the transmit FIFO is at least half empty. In non-FIFO mode, the THR has been transmitted.
4 (lowest)	Modem Status: one or more modem input signal has changed state.

This is a read-only register. Ignore reads from reserved bits.

Table 10-9. IIR Bit Definitions (Sheet 1 of 2)

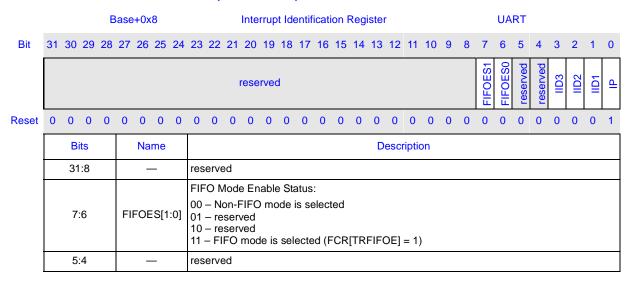




Table 10-9. IIR Bit Definitions (Sheet 2 of 2)

				В	ase	+0x	8					Inte	rrup	t Ide	entifi	icati	on F	Regi	ster							UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											ı	ese	rvec	i											FIFOES1	FIFOES0	reserved	reserved	IID3	IID2	IID1	Б
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
		Bi	its			Nai	me												D	esci	riptio	n										
		3	3			TC (IIE	_			No	Ch		cter	Tim	eou	t Ind	dica	tion	inte		ot is p				mod	de o	nly)					
		2:	:1			IID[2	2:1]		00 - 01 - 10 -	- Mo - Tra - Re	ode ans ecei	mit F ved	tatus FIFC Data	s (C) re a A	TS, que: vaila	DS st da able	ata				dem oreal					d st	tate)					
		()			IF	5		Inte 0 – 1 –	Int	erru	endi ıpt is erru	s pe				/e lc	w)														

Table 10-10. Interrupt Identification Register Decode (Sheet 1 of 2)

Int	errup	t ID B	its		lr	nterrupt SET/RESET Function	n
3	2	1	0	Priority	Туре	Source	Cleared By
0	0	0	1	_	None	No interrupt is pending	_
0	1	1	0	Highest	Receiver Line Status	Overrun error, parity error, framing error, break interrupt	Reading the LSR
0	1	0	0	Second Highest	Received Data Available	Non-FIFO mode: Receive buffer is full FIFO mode: Trigger level was reached	Non-FIFO mode: Reading the Receiver Buffer Register FIFO mode: Reading bytes until Receiver FIFO drops below trigger level or setting FCR[RESETRF].



Table 10-10. Interrupt Identification Register Decode (Sheet 2 of 2)

Int	errup	t ID B	its		Ir	nterrupt SET/RESET Function	n
1	1	0	0	Second Highest	Character Timeout Indication	FIFO mode only: At least one character is in the Receive FIFO and no data has been sent for four character times.	Reading the Receiver FIFO, setting FCR[RESETRF] or a new start bit is received
0	0	1	0	Third Highest	Transmit FIFO Data Request	Non-FIFO mode: Transmit Holding register empty FIFO mode: Transmit has half, or less than half, data	Non-FIFO mode: Reading the IIR Register (if the source of the interrupt) or writing into the Transmit Holding Register FIFO mode: Reading the IIR Register (if the source of the interrupt) or writing to the Transmitter FIFO
0	0	0	0	Fourth Highest	Modem Status	Clear to Send, Data Set Ready, Ring Indicator, Data Carrier Detect	Reading the Modem Status Register

10.4.2.6 FIFO Control Register (FCR)

The FCR, shown in Table 10-11, is a write-only register that is located at the same address as the IIR, which is a read-only register. The FCR enables/disables the transmitter/receiver FIFOs, clears the transmitter/receiver FIFOs, and sets the receiver FIFO trigger level.

This is a write-only register. Write zeros to reserved bits.

Table 10-11. FCR Bit Definitions (Sheet 1 of 2)

				Ва	ase+	+0x0	80						FIF	00	ont	rol F	Regi	ster								UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rved	d											E		reserved	reserved	reserved	RESETTF	RESETRF	TRFIFOE
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Na	me												D	esci	riptio	on										
		31	:8			_	_		res	erve	d																					
		7	:6			ΙΤ	L		ena rec 0b0 0b0	ger abled eive 00 – 01 – 00 –	leve d via DM 1 by 8 by	el pro the 1A re yte o ytes bytes	ogra e IEI eque or m or r	mm R, a est i ore more	n in s al in l e in	into terro so g FIFC FIF n FI	this upt in ene car Car Car	field s ge rate uses ause caus	d an nera d wl s inte es in	d thated hen erru terr	e R I and the pt (I upt	s in teceid ap trigg Not vand tand	ved prop ger /alid DM d DN	Date oriate leve of the level o	a Ave e bit of the bit	vaila ts ar eac A mo est rest	able re se hed	Inte et in	rrup	t is		
		5	:3			_	_		res	erve	d																					



Table 10-11. FCR Bit Definitions (Sheet 2 of 2)

				Ba	se+	-0x0	8					F	FO	Co	ontro	Re	gis	ter								UA	RT					
Bit	31	30	29	28	27	26	25	24	23 2	2 2′	1 2	20 1	9 18	8	17 1	6 1	5	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											re	eserv	ed												Ē	1	reserved	reserved	reserved	RESETTF	RESETRF	TRFIFOE
Reset	0	0	0	0	0	0	0	0	0 0	0		0 0	0)	0) ()	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bit	S			Naı	me												D	esc	ripti	ion										
		2			R	ESE	ΞTΤ	F	Reserrance claim and it $0 - \sqrt{1 - 1}$	eare ipt, con Vriti	ed. if t npl ng	The he Tiletes 0 ha	TDF E bi the s no	RC it ii cu	the bit in the surrent of the series of the	n the IEF trai	e L R re nsn	SR gis niss	is s ter i	et a s se	and	the	IIR s	shov	vs a	trai	nsmi	itter	requ	uest	s da	ta
		1			RI	ESE	ETR	F	Reserved in the set. The Availation of the set. The availation of the set.	ed. T LSF he r ble, Vriti	The R a rec , it ng	e DR are cl eiver is cle 0 ha	bit in eare shif areas s no	n ted. it r id.	the LS Any egist	SR i erro er is	s re or b	ese oits,	t to OE	0. A i, Pl	dl tł ≣, F	ne er E or	ror BI,	bits that	in th	e F d be	IFO en s	and set in	the n LS	FIF R a	OE I re st	bit
		0			Т	RFI	FOE	_	Trans receiv TRFII all by versa this re 0 - F 1 - F	er F OE es i da gist	FIF n b ta i er	Os. \ 0, th both I is aut are \ are d	Vhe e FII FIFC oma vritte isab	n FC Os. atio en le	TRFI Ds are . Whe cally or th d	FOE bo en c clea	E = th har	1, disangir ngir d fra	both able ng fr om t	n FII d (r om the	FIF	s are -FIF(FO m	ena D M lode This	able ode to i bit	d (F). W non- mus	IFO ritin FIF	Moo g a (O m	de). 0 to ode	Wh this and	en bit d d vic	е	

10.4.2.7 Line Control Register (LCR)

The LCR, shown in Table 10-12, specifies the format for the asynchronous data communications exchange. The serial data format consists of a start bit, five to eight data bits, an optional parity bit, and one, one and a half, or two stop bits. The LCR has bits that allow access to the Divisor Latch and bits that can cause a break condition.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 10-12. LCR Bit Definitions

				Ва	ase-	+0x(С						Lin	e Co	ontr	ol R	Regi	ster								UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	erve	d											DLAB	SB	STKYP	EPS	PEN	STB	WLS1	WLS0
Docot	^	•	^	^	^	^	^	^	•	^	^	^	_	^	^	^	^	•	^	^	^	^	^	^	^	^	•	^	^	^	^	•

Bits	Name	Description
31:8	_	reserved
7	DLAB	Divisor Latch Access Bit: Must be set high (logic 1) to access the Divisor Latches of the Baud Rate Generator during a READ or WRITE operation. Must be set low (logic 0) to access the Receiver Buffer, the Transmit Holding Register, or the IER. 0 – access Transmit Holding register (THR), Receive Buffer register (RBR) and IER.
		1 - access Divisor Latch registers (DLL and DLH)
6	SB	Set Break: Causes a break condition to be transmitted to the receiving UART. Acts only on the TXD pin and has no effect on the transmitter logic. In FIFO mode, wait until the transmitter is idle, LSR[TEMT]=1, to set and clear SB.
		0 - no effect on TXD output 1 - forces TXD output to 0 (space)
5	STKYP	Sticky Parity: Forces the bit value at the parity bit location to be the opposite of the EPS bit, rather than the parity value. This stops parity generation. If PEN = 0, STKYP is ignored. 0 - no effect on parity bit 1 - forces parity bit to be opposite of EPS bit value
4	EPS	Even Parity Select: Even parity select bit. If PEN = 0, EPS is ignored. 0 - sends or checks for odd parity 1 - sends or checks for even parity
3	PEN	Parity Enable: Enables a parity bit to be generated on transmission or checked on reception. 0 - no parity 1 - parity
2	STB	Stop Bits: Specifies the number of stop bits transmitted and received in each character. When receiving, the receiver only checks the first stop bit. 0 - 1 stop bit 1 - 2 stop bits, except for 5-bit character then 1-1/2 bits
1:0	WLS[1:0]	Word Length Select: Specifies the number of data bits in each transmitted or received character. 00 – 5-bit character 01 – 6-bit character 10 – 7-bit character 11 – 8-bit character



10.4.2.8 Line Status Register (LSR)

The LSR, shown in Table 10-13, provides data transfer status information to the processor.

In non-FIFO mode, LSR[4:2]: parity error, framing error, and break interrupt, show the error status of the character that has just been received.

In FIFO mode, LSR[4:2] show the status bits of the character that is currently at the front of the FIFO.

LSR[4:1] produce a receiver line status interrupt when the corresponding conditions are detected and the interrupt is enabled. In FIFO mode, the receiver line status interrupt only occurs when the erroneous character reaches the front of the FIFO. If the erroneous character is not at the front of the FIFO, a line status interrupt is generated after the other characters are read and the erroneous character becomes the character at the front of the FIFO.

The LSR must be read before the erroneous character is read. LSR[4:1] bits are set until software reads the LSR.

See Section 10.4.5 for details on using the DMA to receive data.

This is a read-only register. Ignore reads from reserved bits.

Table 10-13. LSR Bit Definitions (Sheet 1 of 3)

			Ва	ase+	+0x1	14						Lir	ne S	tatu	ıs R	egis	ter								UA	RT					
Bit	31	30 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										r	ese	erve	d											FIFOE	TEMT	TDRQ	B	FE	PE	OE	DR
Reset	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0	0
		Bits			Nai	me												D	esc	riptio	on										
		31:8			_	-		res	erve	ed																					
		7			FIF	OE		the in the error 7 is required have a process of the error of the erro	re is he F oned s set uest ve bo re bor in No	at left out of the control of the co	leas cha l) a gene clea ests upt.	or no	e pa cess ers I IFO ed e I by not	arity sor i hav DE is ven rea ena	erroread e be s set who ding able	or, from the took took took took took took took too	amir he L ead , the ie re FIF	ng e SR from e er ceiv O, E it7 i	doe m th ror i ve F DM/ s se	r, or es no ee FI inter IFO A rec	brea FO. rupt rea ques 0), F	ak ir set If D is g che ts a	ndicathis MA gene s the	ation bit. requerate e trig e-en	for FIFO Lest d and ger able	any OE is s are nd n leve	is s of the s res e en lo re el. O utom s not	he o set o able ceiv nce natio	char whe ed (I re D the cally	acte n all ER I MA erro	ers I bit ors
		6			TEI	MT		Reg Tra trar	giste nsm nsm Th	er ar nitter itter nere	e bor Sh FIF is d	oth on the other of the other o	emp Regis nd t in th	oty. I ster he ne T	t is con Tran rans	clea tain: smit	red v s a d Shi	whe lata ft R t Re	en ei cha egist egist	ither arac ster : ter, t	the ter. I are I he T	Tra n F both ran	nsm IFO em	nit Ho mod pty.	oldii de,	ng F TEM	nsmi Regis IT is egiste	ster set	or tl	he en th	



Table 10-13. LSR Bit Definitions (Sheet 2 of 3)

				Ва	ase-	+0x′	14						Lir	ne S	tatu	s Re	egis	ter								UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											I	rese	rve	t											FIFOE	TEMT	TDRQ	BI	FE	ЭE	OE	DR
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0	0
•		В	its			Na	me												D	esci	iptio	on										

Bits	Name	Description
5	TDRQ	Transmit Data Request: Indicates that the UART is ready to accept a new character for transmission. In addition, this bit causes the UART to issue an interrupt to the processor when the transmit data request interrupt enable is set high and generates the DMA request to the DMA controller if DMA requests and FIFO mode are enabled. The TDRQ bit is set when a character is transferred from the Transmit Holding register into the Transmit Shift register. The bit is cleared with the loading of the Transmit Holding Register. In FIFO mode, TDRQ is set to 1 when half of the characters in the FIFO have been loaded into the Shift register or the RESETTF bit in FCR has been set. It is cleared when the FIFO has more than half data. If more than 64 characters are loaded into the FIFO, the excess characters are lost. 0 — There is data in Holding register or FIFO waiting to be shifted out
		1 - Transmit FIFO has half or less than half data
4	ВІ	Break Interrupt: BI is set when the received data input is held low for longer than a full word transmission time (that is, the total time of Start bit + data bits + parity bit + stop bits). The Break indicator is reset when the processor reads the LSR. In FIFO mode, only one character equal to 0x00, is loaded into the FIFO regardless of the length of the break condition. BI shows the break condition for the character at the front of the FIFO, not the most recently received character.
		0 - No break signal has been received1 - Break signal received
3	FE	Framing Error: FE indicates that the received character did not have a valid stop bit. FE is set when the bit following the last data bit or parity bit is detected to be 0. If the LCR had been set for two stop bit mode, the receiver does not check for a valid second stop bit. The FE indicator is reset when the processor reads the LSR. The UART will resynchronize after a framing error. To do this it assumes that the framing error was due to the next start bit, so it samples this "start" bit twice and then reads in the "data". In FIFO mode, FE shows a framing error for the character at the front of the FIFO, not for the most recently received character. 0 — No Framing error 1 — Invalid stop bit has been detected



Table 10-13. LSR Bit Definitions (Sheet 3 of 3)

									•			,																	
				Ba	se-	+0x14				I	Line S	Statu	s Re	gist	ter								UA	RT					
Bit	31	30	29	28	27	26 25	24	23 22	21	20 1	9 18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
									ı	reserv	ed											FIFOE	TEMT	TDRQ	ВІ	FE	PE	OE	DR
Reset	0	0	0	0	0	0 0	0	0 0	0	0 (0 0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0	0
		Bi	ts			Name										De	esci	riptio	on										
		2	!			PE		Parity E odd pa error a error fo 0 – No 1 – Pa	rity, nd is or th	as se s clea e cha arity ei	lected red w racter ror	d by then the at the	the e the p ne fro	ever oroc	n pa cess	rity or re	sele ead:	ect b	it. P e LS	E is	s set In F	t upo IFO	on d mod	etec de, F	tion E s	of a	a pa /s a	rity pari	ity
		1				OE		Overru not rea lost. In receive condition 0 – No 1 – Re	d by FIF ed by on a	y the post of the	oroce: de, O s bee eared s bee	ssor E ind en dis whe n los	befo licate scare n the	re t es t ded e pr	he r hat a	next all 6 e O	cha 4 by E in	arac /tes dica	ter v of thator	vas ne F is se	rece FIFO et up	eive are	d. The full	he n and	ew the	cha mo	racte st re	er is ecen	ntly
		0	١			DR		Data R into the receive been re 0 - No 1 - Data	Re buread	ceive ffer is from lata	Bufferead. RBR) bee	er Re In Fl or th n rec	giste FO e FI eive	er o mod FO d	r the de, [is re	FIF OR i	FÖ. is cl	In r eare	on-led if	FIF(the	O mo	ode, O is	DR	is c	lear	ed v	whe	n the	е



10.4.2.9 Modem Control Register (MCR)

The MCR, shown in Table 10-14, uses the modem control pins nRTS and nDTR to control the interface with a modem or data set. The MCR also controls the Loopback mode. Loopback mode must be enabled before the UART is enabled. The differences between UARTs specific to this register are described in Section 10.5.1.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 10-14. MCR Bit Definitions (Sheet 1 of 2)

				В	ase-	+0x	10					N	lode	m (Con	trol	Reg	gist	er							UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	serv	ed													LOOP	OUT2	OUT1	RTS	DTR
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	Re	ad/\	Vrite	9																												
		Ві	its			Na	me												De	scr	ipti	on										
		31	:5			_	_		rese	erve	d																					
		4	1			LO	OP		UAI is set the moor the Corr (bits the Loo The bits In Ico proof and actification of the Loo O' The Door actification of the Loo O' The Door and actification of the	RT. et to Tra demonstrate de la composition della composition dell	Who a long out of the control of the	en L ogici ntro n co it of n the its ir moc four 1 fc = 1 f k m c y M ank l al c al c	cOO : 1 sir r Sh I inpontrol the e Moon thee bits orce force force code, rerify on the mitted ART moon through the moon throu	P is tate ift results and of the control of the con	set set in the second s	to a e recter is TS, pins ck matus confi MCF to a a 1 to a at is nsmith the to the to on the the to on the the to on the the the to on the	logiceives "logiceives "logiceives "logiceives "logiceives (nF) ode Registration (nF) od	ic 1. er s ope SR, RTS ma giste ed t eco ope den rec ion	, the erial documents of the control	following following following for the ctechnical following for the ctechnical following followin	owiinto	ng was disso the nRI of the nRI o	rill oconicately attely as of a the attely as of a	ccui nective eivee dis ced dis ced the enal er fo	r: The ted of the ted	ne tra from hift re- neer heir stivat MSF I. Mode	ensr the egis ted inaction R is The ontro an a	mitter i pin ter i i from tive of the treact of the transfer i interior of	trupt	rial cout. The out. The pinse. e. elta lee to egis allov t, re ts a ansi	outp put e fo s an oits cle ter ter	of our nd ear he ve ed
		(3			OL	JT2		unit 0 – 1 – Whe	. W U U en L M	hen ART ART -OC SR[I	LO inte inte P= DCI	OP= erruperrup 1, in 0] fo 0] fo	=0: ot is ot is terro rceo	disa ena upts d to	able abled alw a 0	d. d.								,				- r.			



Table 10-14. MCR Bit Definitions (Sheet 2 of 2)

				Ва	ase-	+0x1	10					M	ode	m (Con	trol	Reg	jiste	er							UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	serv	ed													LOOP	OUT2	OUT1	RTS	DTR
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	Re	ad/\	Vrite	е																												
		В	its			Na	me												De	escr	iptio	on										
		2	2			OU	JT1			Fo	. Thi rce rce	MS	R[R	I] to	0	nly ir	n Lo	opb	ack	mod	de. I	t is i	gno	red	oth	erwi	se.					
		,	1			RT	ГS		0 –	nR	t to	pin	is 1																			
		()			DI	ΓR			nΣ	ermii DTR DTR	pin	is 1																			

10.4.2.10 Modem Status Register (MSR)

The MSR, shown in Table 10-15, provides the current state of the control lines from the modem or data set (or a peripheral device emulating a modem) to the processor. In addition to this current state information, four bits of the MSR provide change information. MSR[3:0] are set when a control input from the Modem changes state. They are cleared when the processor reads the MSR. Differences between UARTs specific to this register are described in Section 10.5.1.

The status of the modem control lines do not affect the FIFOs. To use these lines for flow control, IER[MIE] must be set. When an interrupt on one of the flow control pins occurs, the interrupt service routine must disable the UART. The UART will continue transmission/reception of the current character and then stop. The contents of the FIFOs will be preserved. If the UART is reenabled, transmission will continue where it stopped. Interrupts from the flow control pins will not come through the UART unit if the unit is disabled. When disabling the unit because of flow control, interrupts must be enabled in the processor Interrupt Controller for the flow control pins. The Interrupt Controller will still trigger interrupts if the pins are in Alternate Function Mode.

Note: When bit 0, 1, 2, or 3 is set, a Modern Status interrupt is generated if IER[MIE] is set.

This is a read-only. Ignore reads from reserved bits.

4

3

2

1

0

CTS

DDCD

TERI

DDSR

DCTS

0 - nCTS pin is 1 1 - nCTS pin is 0

Delta Data Carrier Detect:

Trailing Edge Ring Indicator:

1 - nDSR pin has changed state

Delta Data Set Ready:

Delta Clear To Send:



Та	ble	10-	-15	. N	ISR	Bit	De	efini	ion	S																					
				В	ase+0	x18						Mod	lem	Sta	itus	Reg	iste	r							UA	RT					
Bit	31	30	29	28	27 2	6 2	5 2	24 23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										I	rese	rvec	t											DCD	R	DSR	CTS	DDCD	TERI	DDSR	DCTS
Reset	0	0	0	0	0 () ()	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bit	S		N	lam	е		Description																						
		31:	8			_		res	serv	ed																					
		7			1	OCE)	M	CR[C - n[DUT.	2] if pin	MC	R[L		plem P] is			he [Data	Cai	rier	Det	tect	(nD(CD)	inpi	ut. E	qui	/aler	nt to)
		6				RI		M(ng Ir CR[L - nf - nf	.00 RI pi	P] is n is	s set 1		eme	ent o	f the	e Ri	ng I	ndic	ator	(nR	tl) ir	put.	. Equ	uiva	lent	to N	//CR	[OU	T1]	if
		5				DSF	R	M(0 -		OTR: OSR] if N pin	ICR is 1	[LO		ment		he	Data	s Se	t Re	ady	(nE)SR)) inp	ut. I	Equi	vale	ent to	O		
		4				CTS	3		ear T CR[L					eme	ent c	of the	e CI	ear	to S	end	(nC	TS)	inp	ut. E	qui	vale	nt to) MC	CR[F	RTS] if

0 – No change in nDCD pin since last read of MSR
1 – nDCD pin has changed state

0-No change in nDSR pin since last read of MSR

 $0-\,$ nRI pin has not changed from 0 to 1 since last read of MSR $1-\,$ nRI pin has changed from 0 to 1



10.4.2.11 Scratchpad Register (SPR)

The SPR, shown in Table 10-16, has no effect on the UART. It is intended as a scratchpad register for use by the programmer. It is included for 16550 compatibility.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 10-16. SPR Bit Definitions

				Ba	ase	+0x1	С						Sci	ratch	n Pa	ad R	egis	ster								UA	RT					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											ı	rese	rve	d														ď	5			
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Na	me												D	esci	riptic	n										
		31	:8			_	_		res	erve	ed																					
		7	:0			S	Р		Scr No			ad on U	ART	- fun	ctio	nali	ty															

10.4.3 FIFO Interrupt Mode Operation

This section describes how to service interrupts in FIFO mode.

10.4.3.1 Receive Interrupt

For a receive interrupt to occur, the receive FIFO and receive interrupts must be enabled. IIR[IID] changes to show that receive data is available when the FIFO reaches its trigger level. IIR[IID] changes to show the next waiting interrupt when the FIFO drops below the trigger level. A change in IIR[IID] triggers an interrupt to the core. Software reads IIR[IID] to determine the cause of the interrupt.

The receiver line status interrupt (IIR = 0xC6) has the highest priority and the received data available interrupt (IIR = 0xC4) is lower. The line status interrupt occurs only when the character at the front of the FIFO has errors.

The data ready bit (DR in the LSR) is set when a character is transferred from the shift register to the Receive FIFO. The DR bit is cleared when the FIFO is empty.

10.4.3.2 Character Timeout Indication Interrupt

A character timeout indication interrupt occurs when the receive FIFO and character timeout indication interrupt are enabled and all of the following conditions exist:

- At least one character is in the FIFO.
- The most recently received character was received more than four continuous character times ago. If two stop bits are programmed, the second is included in this interval.
- The most recent FIFO read was performed more than four continuous character times ago.



After the processor reads one character from the receive FIFO or a new start bit is received, the character timeout indication interrupt is cleared and the timeout is reset. If a character timeout indication interrupt has not occurred, the timeout is reset when a new character is received or the processor reads the receive FIFO.

10.4.3.3 Transmit Interrupt

Transmit interrupts can only occur when the transmit FIFO and transmit interrupt are enabled. The transmit data request interrupt occurs when the transmit FIFO is at least half empty. The interrupt is cleared when the THR is written or the IIR is read.

10.4.4 FIFO Polled Mode Operation

When the FIFOs are enabled, setting IER[7] and IER[4:0] to all zeroes puts the serial port in the FIFO polled mode of operation. The receiver and the transmitter are controlled separately. Either one or both can be in the polled mode. In polled mode, software checks receiver and transmitter status via the LSR.

10.4.5 DMA Requests

The FIFO data is one byte wide. DMA requests are either transmit data service requests or receive data service requests. DMA requests are only generated in FIFO mode.

The transmit DMA request is generated when the transmit FIFO is at least half empty and IER[DMAE] is set. After the transmit DMA request is generated, the DMA Controller (DMAC) writes data to the FIFO. For each DMA request, the DMAC sends 8, 16, or 32 bytes of data to the FIFO. The number of bytes to be transmitted is programmed in the DMA channel.

The receive DMA request is generated when the receive FIFO reaches its trigger level with no errors in its entries and the IER[DMAE] is set. A receive DMA request is not generated if the trigger level is set to 1.

The DMAC then reads data from the FIFO. For each DMA request, the DMA controller can read 8, 16 or 32 bytes of data from the FIFO. The number of bytes to be read is programmed in the DMA channel.

Note: Do not program the channel to read more data than the FIFO trigger level.

If DMA requests are enabled and an erroneous character is received, the receive DMA requests are automatically disabled and an error interrupt is generated. The erroneous character is placed in the receive FIFO. If the UART was requesting a receive DMA transaction, the request is immediately cancelled. This prevents the DMAC from attempting to access the FIFOs while software clears the error.

When all the errors in the receive FIFO are cleared, receive DMA requests are automatically enabled and can be generated when the trigger level is reached.



Note: Ensure that the DMAC has finished previous receive DMA requests before the error interrupt handler begins to clear the errors from the FIFO.

10.4.5.1 Trailing Bytes in the Receive FIFO

Trailing bytes occur when the number of entries in the receive FIFO is less than its trigger level and no more data is being received. In such a case, a receive DMA request is not generated. To read the trailing bytes follow these steps:

- 1. Wait for a character timeout indication interrupt. The character timeout indication interrupt must be enabled.
- 2. Disable the receive DMA channel and wait for it to stop.
- 3. Read one byte at a time. The FIFO is empty when LSR[DR] is cleared.
- 4. Re-enable the receive DMA channel.

10.4.6 Slow Infrared Asynchronous Interface

The Slow Infrared (SIR) interface is used with the STUART to support two-way wireless communication that uses infrared transmission. The SIR provides a transmit encoder and receive decoder to support a physical link that conforms to the IRDA Serial Infrared Specification Version 1.1.

The SIR interface does not contain the actual IR LED driver or the receiver amplifier. The I/O pins attached to the SIR only have digital CMOS level signals. The SIR supports two-way communication, but full duplex communication is not possible because reflections from the transmit LED enter the receiver. The SIR interface supports frequencies up to 115.2 kbps. Because the input clock is 14.7456 MHz, the baud divisor must be eight or more.

10.4.6.1 Infrared Selection Register (ISR)

The IRDA module is managed through the UART to which it is attached. The ISR, shown in Table 10-17, controls IRDA functions.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 10-17. ISR Bit Definitions

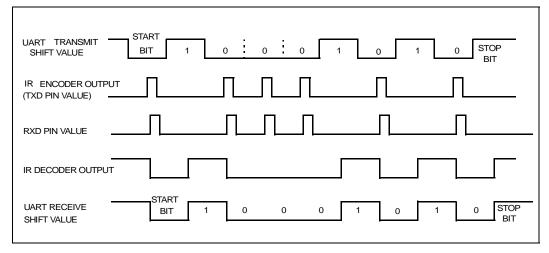
Ia	Die	10-1		J	וע או	· D	5111		OHS																							
			E	За	se+0	(20					In	frar	ed S	Se	election	n Re	gis	ter								UA	RT					
Bit	31	30 29	28	3 2	27 26	3 25	24	23	22	21	20	19	18	1	17 16	15	14	13	12		11 1	0	9	8	7	6	5	4	3	2	1	0
												res	serv	/e	d													RXPL	TXPL	XMODE	RCVEIR	XMITIR
Reset	0	0 0	0		0 0	0	0	0	0	0	0	0	0		0 0	0	0	0	0		0)	0	0	0	0	0	0	0	0	0	0
		Bits			N	ame												D	esc	rij	ption											
		31:5						re	serv	ed																						
		4			R	XPL		0 -		R d	eco	der i	inter	rp	rets po																	
		3			T	XPL		0 -		Rе	nco	der	gene	er	ates a																	
		2			ΧM	1ODE	=	clo ge a l	ock the nerabit tir	he II ites ne v	RDA 1.6 vide mit p	tra μs p , an	nsm oulse d th	nit es e id	elect: and res (that receive th is 3	ecei are e de /16 (ve 1 3/1 eco of a	logio 6 of der	a bi expe	he t t	en X time cts p	MC at	DDE 115	is: 5.2 k	set, bps	the) ins	tran stead	smi d of	eno puls	code	r	of
								Re	eceiv e IRI	er S	SIR E leco	Enal der	ole: befo	W or	th is 1 Vhen F e it is f	CV ed t	EIR o th	ne U	ART	Γ.	If RC	VE	ΞIR	is c	lear	ed,	then	is p	roce	esse king	d by	y he
		1			RC	VEIF	₹	0 -	- R	ecei	ver i	is in	UA	R	d and T mod red mo	le	RX	D pi	n is	fe	ed di	rec	tly	to th	ie U	AR	Г.					
								is clo	proc	esse g to	ed b	y the	e IR	D	: Whei A enc coder	odei	r be	fore	it is	f	ed to	th	ie d	evic	e pi	n. If	ΧM	ITIR	isc	lear	ed,	all
		0			ΧN	/IITIR	2	sta	ate, v	vill s	switc	h to	an	10	Enable rmally d befor	low	de	faul	t sta	te	outpu e. Th	it p	oin, can	whicau	ch is ise a	in a a fal	a no se s	rma tart	ly h bit ι	gh c inles	lefa s th	ult ne
								_						_	ART m frared																	

10.4.6.2 Operation

The SIR modulation technique works with 5-, 6-, 7-, or 8-bit characters with an optional parity bit. The data is preceded by a zero value start bit and ends with one or more stop bits. The encoding scheme is to set a pulse 3/16 of a bit wide in the middle of every zero bit and send no pulses for bits that are ones. The pulse for each zero bit must occur, even for consecutive bits with no edge between them.





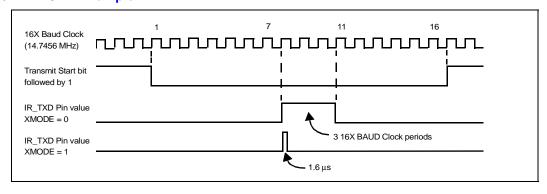


The top line in Figure 10-3 shows an asynchronous transmission as it is sent from the UART. The second line shows the pulses generated by the IR encoder at the TXD pin. A pulse is generated in the middle of the START bit and any data bit that is a zero. The third line shows the values received at the RXD input pin. The fourth line shows the receive decoder's output. The receive decoder drives the receiver data line low when it detects a pulse. The bottom line shows how the UART's receiver interprets the decoder's action. This last line is the same as the first, but it is shifted half a bit period.

When XMODE is cleared, each zero bit has a pulse width of 3/16 of a bit time. When XMODE is set, a pulse of $1.6~\mu s$ is generated in the middle of each zero bit. The shorter infrared pulse generated when XMODE is set reduces the LEDs' power consumption. At 2400 bps, the LED is normally on for $78~\mu s$ for each zero bit that is transmitted. When XMODE is set, the LED is on only $1.6~\mu s$. XMode changes the behavior of the receiver. The receiver expects pulses of the correct pulse width. If the transceiver crops the incoming pulse, then Xmode must be set.

Note:

Figure 10-4. XMODE Example



Note: Note: The SIR TXD output pin is automatically held deasserted when the RCVEIR bit is set. Before setting the RCVEIR bit, check that the TEMT bit is 1. While receiving, any data placed in



the Transmit FIFO will not be held. Only add data to the Transmit FIFO while not receiving. To start transmission, the RCVEIR bit must be cleared.

To disable SIR, disable the IrDA LED first, if possible. Second, set the TXD GPIO pin to the infrared LED's default state using the GPCR/GPSR registers. Next, change the TXD pin from alternate function to GPIO mode. Now the SIR can be disabled without causing spurious transmit pulses.

10.5 UART Register Summary

Table 10-18, Table 10-19, and Table 10-20 contain the register addresses for the FFUART, BTUART, and STUART.

Table 10-18. FFUART Register Summary

Register Addresses	DLAB Bit Value	Name	Description
0x4010_0000	0	FFRBR	Receive Buffer register (read only)
0x4010_0000	0	FFTHR	Transmit Holding register (write only)
0x4010_0004	0	FFIER	IER (read/write)
0x4010_0008	Х	FFIIR	Interrupt ID register (read only)
0x4010_0008	Х	FFFCR	FCR (write only)
0x4010_000C	Х	FFLCR	LCR (read/write)
0x4010_0010	Х	FFMCR	MCR (read/write)
0x4010_0014	Х	FFLSR	LSR (read only)
0x4010_0018	Х	FFMSR	MSR (read only)
0x4010_001C	Х	FFSPR	Scratch Pad Register
0x4010_0020	Х	FFISR	Infrared Selection register (read/write)
0x4010_0000	1	FFDLL	Divisor Latch Low register (read/write)
0x4010_0004	1	FFDLH	Divisor Latch High register (read/write)

Table 10-19. BTUART Register Summary (Sheet 1 of 2)

Register Addresses	DLAB Bit Value	Name	Description
0x4020_0000	0	BTRBR	Receive Buffer register (read only)
0x4020_0000	0	BTTHR	Transmit Holding register (write only)
0x4020_0004	0	BTIER	IER (read/write)
0x4020_0008	Х	BTIIR	Interrupt ID register (read only)
0x4020_0008	Х	BTFCR	FCR (write only)
0x4020_000C	Х	BTLCR	LCR (read/write)
0x4020_0010	Х	BTMCR	MCR (read/write)
0x4020_0014	Х	BTLSR	LSR (read only)
0x4020_0018	Х	BTMSR	MSR (read only)



Table 10-19. BTUART Register Summary (Sheet 2 of 2)

Register Addresses	DLAB Bit Value	Name	Description
0x4020_001C	Х	BTSPR	Scratch Pad Register
0x4020_0020	Х	BTISR	Infrared Selection register (read/write)
0x4020_0000	1	BTDLL	Divisor Latch Low register (read/write)
0x4020_0004	1	BTDLH	Divisor Latch High register (read/write)

Table 10-20. STUART Register Summary

Register Addresses	DLAB Bit Value	Name	Description
0x4070_0000	0	STRBR	Receive Buffer register (read only)
0x4070_0000	0	STTHR	Transmit Holding register (write only)
0x4070_0004	0	STIER	IER (read/write)
0x4070_0008	Х	STIIR	Interrupt ID register (read only)
0x4070_0008	Х	STFCR	FCR (write only)
0x4070_000C	Х	STLCR	LCR (read/write)
0x4070_0010	Х	STMCR	MCR (read/write)
0x4070_0014	Х	STLSR	LSR (read only)
0x4070_0018	Х	STMSR	reserved (no modem control pins)
0x4070_001C	Х	STSPR	Scratch Pad Register
0x4070_0020	Х	STISR	Infrared Selection register (read/write)
0x4070_0000	1	STDLL	Divisor Latch Low register (read/write)
0x4070_0004	1	STDLH	Divisor Latch High register (read/write)



10.5.1 UART Register Differences

The default descriptions for BTMCR, BTMSR and STMCR are modified as shown in Table 10-21.

Table 10-21. Flow Control Registers in BTUART and STUART

	Bit7-5	Bit4	Bit3	Bit2	Bit1	Bit0
BTMCR	reserved	LOOP	OUT2	reserved	RTS	reserved
BTMSR	reserved	CTS	reserved	reserved	reserved	DCTS
STMCR	reserved	LOOP	OUT2	reserved	reserved	reserved

Fast Infrared Communication Port 11

The Fast Infrared Communications Port (FICP) for the PXA255 processor operates at half-duplex and provides direct connection to commercially available Infrared Data Association (IrDA) compliant LED transceivers. The FICP is based on the 4-Mbps IrDA standard and uses four-position pulse modulation (4PPM) and a specialized serial packet protocol developed for IrDA transmission. To support the standard, the FICP has:

- A bit encoder/decoder,
- A serial-to-parallel data engine
- A transmit FIFO 128 entries deep and 8 bits wide
- A receive FIFO 128 entries deep and 11 bits wide

The FICP shares GPIO pins for transmit and receive data with the Standard UART. Only one of the ports can be used at a time. To support a variety of IrDA transceivers, both the transmit and receive data pins can be individually configured to communicate using normal or active low data.

11.1 Signal Description

The FICP signals are IRRXD and IRTXD. Table 11-1 describes each signal's function. Most IrDA transceivers also have enable and speed pins. Use GPIOs to enable the transceiver and select the speed. See Section 4, "System Integration Unit" on page 4-1 for more information.

Table 11-1. FICP Signal Description

Signal Name	Input/Output	Description
IRRXD	Input	Receive pin for FICP
IRTXD	Output	Transmit pin for FICP

11.2 FICP Operation

The FICP is disabled and does not have control of the port's pins after a reset. Before software enables the FICP for high-speed operation, it must set the control registers to reflect the desired operating mode. After the control registers are set, software can either preload the FICP's transmit FIFO with up to 128 bytes, or leave the FIFO empty and use the DMA to service it after the FICP is enabled. Once the FICP is enabled, transmit/receive data can be sent on the transmit and receive pins.

The transmit/receive data is modulated according to the 4PPM IrDA standard and converted to serial or parallel data. The modulation technique and the frame format are discussed in the sections that follow.



11.2.1 4PPM Modulation

Four-position pulse modulation (4PPM) is used to transmit data at the high-speed rate, 4.0 Mbps. Data bits are encoded two at a time by placing a single 125 ns light pulse in one of four timeslots. The four timeslots are collectively termed a chip. Bytes are encoded one at a time. They are divided into four individual 2-bit pairings called nibbles. The least significant nibble is transmitted first. Figure 11-1 shows the 4PPM encoding for the possible 2-bit combinations and Figure 11-2 shows an example of 4PPM modulation for the byte 0b10110001, which is constructed with four chips. Bits within each nibble are not reordered, but nibble 0 (least significant) is transmitted first and nibble 3 (most significant) is transmitted last.

Figure 11-1. 4PPM Modulation Encodings

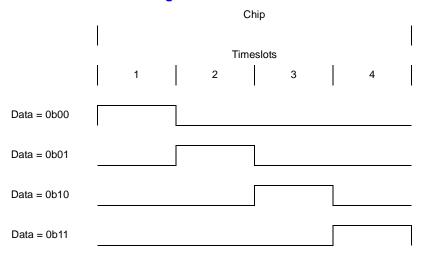
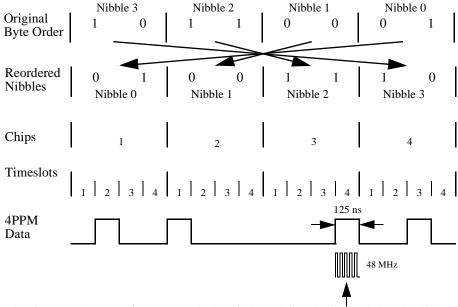


Figure 11-2. 4PPM Modulation Example



Receive data sample counter frequency = 6/pulse width. Each timeslot is sampled on the third clock.



11.2.2 Frame Format

The frame format used with 4-Mbps transmission is shown in Figure 11-3.

Figure 11-3. Frame Format for IrDA Transmission (4.0 Mbps)

64 chips	8 chips	4 chips (8 bits)	4 chips (8 bits)	8180 chips max (2045 bytes)	16 chips (32 bits)	8 chips								
Preamble	Start Flag	Address (optional)	Control (optional)	Data	CRC-32	Stop Flag								
Preamble -	1000 0000	00 1010 1000 repeated at least 16 times												
Start flag -	0000 1100 0000 1100 0110 0000 0110 0000													
Stop Flag -	0000 1100 0	0000 1100 0000 1100 0000 0110 0000 0110												

The preamble, start, and stop flags are a mixture of chips that contain 0, 1, or 2 pulses in their timeslots. Chips with 0 and 2 pulses are used to construct flags because the chips represent invalid data bit pairings. The preamble contains 16 repeated transmissions of the chips: 1000 0000 1010 1000. The start flag contains one transmission of eight chips: 0000 1100 0000 1100 0110 0000 0110 0000. The stop flag contains one transmission of eight chips: 0000 1100 0000 1100 0000 1100 0000 0110. The address, control, data, and CRC-32 use the standard 4PPM chip encoding to represent two bits per chip.

11.2.3 Address Field

A transmitter uses the 8-bit address field to target a receiver when multiple stations are connected to the same set of serial lines. The address allows up to 255 stations to be uniquely addressed (0x00 to 0xFE). The broadcast address 0xFF is used to send messages to all of the connected stations.

For reception, FICP control register 1 (ICCR1) is used to program a unique receive address. The AME bit in the FICP control register 0 (ICCR0) determines the address match function. The received frames' addresses are stored in the receive FIFO with normal data.

11.2.4 Control Field

The control field is an optional 8-bit field that is defined by software. The FICP does not provide hardware decode support for the control byte. It treats all bytes between the address and the CRC as data.

11.2.5 Data Field

The data field can have a length from 0 to 2045 bytes. Application requirements and target system's transmission characteristics affect the data field's length. Software must determine the length of the data to maximize the amount that can be transmitted in each frame while allowing the CRC to detect all errors during transmission. The serial port does not contain hardware that restricts the maximum amount of data that can be transmitted or received. If a data field that is not a multiple of eight bits is received, an abort is signalled.



11.2.6 CRC Field

The FICP uses a 32-bit Cyclic Redundancy Check (CRC) to detect bit errors that occur during transmission. The CRC is generated from the address, control, and data fields, and is included in each frame. Transmit and receive logic have separate CRC generators. The CRC computation logic is set to all ones before each frame is transmitted or received and the result is inverted before it is used for comparison or transmission. The transmitter calculates a CRC as data is transmitted and places the inverse of the resulting 32-bit value at the end of each frame until the stop flag is transmitted. The receiver also calculates a CRC and inverts it for each data frame that it receives. The receiver compares the calculated CRC to the expected CRC value at the end of each frame.

If the calculated value does not match the expected value, the CRC error bit that corresponds to the last data byte received is set. When this byte reaches the trigger level range, an interrupt is generated.

Note: Unlike the address, control, and data fields, the 32-bit inverted CRC value is transmitted and received most significant nibble first.

The cyclic redundancy checker uses the 32-term polynomial:

$$CRC(x) = (x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^{8} + x^{7} + x^{5} + x^{4} + x^{2} + x + 1)$$

11.2.7 Baud Rate Generation

The baud rate is derived by dividing a fixed 48-MHz clock by six. Using a digital PLL, the 8-MHz baud (or timeslot) clock for the receive logic is synchronized with the 4PPM data stream each time a transition is detected on the receive data line. To encode a 4-Mbps data stream, the required chip frequency is 2.0 MHz, with four timeslots per chip at a frequency of 8.0 MHz. Receive data is sampled halfway through each timeslot period by counting three of the six 48-MHz clock periods that make up each timeslot (see Figure 11-2). The chips are synchronized during the reception preamble. The pattern of four chips repeated 16 times is used to identify the first timeslot (or the beginning of a chip) and resets the 2-bit timeslot counter logic.

11.2.8 Receive Operation

The IrDA standard specifies that all transmission occurs at half-duplex. This restriction forces software to enable one direction at a given time. Either the transmit or receive logic can be enabled, but not both. The FICP's hardware does not impose such a restriction. The software can enable both the transmitter and receiver. This feature is used with the FICP's loopback mode, which internally connects the output of the transmit serial shifter to the input of the receive serial shifter.

After the FICP is enabled, the receiver logic begins and selects an arbitrary chip boundary, uses a serial shifter to receive four incoming 4PPM chips from the receive data pin, and latches and decodes the chips one at a time. If the chips do not have the correct preamble, the timeslot counter's clock skips one 8-MHz period and effectively delays the timeslot count by one. This process is called hunt mode and is repeated until the chips have the correct preamble, which indicates that the timeslot counter is synchronized. The preamble can be repeated as few as 16 times or can be continuously repeated to indicate an idle transmit line.



After 16 preambles are transmitted, the start flag is received. The start flag is eight chips long. If any portion of the start flag does not match the encoding, the receive logic signals a framing error and the receive logic returns to hunt mode.

When the correct start flag is recognized, each following group of four chips is decoded into a data byte and placed in a 5-byte temporary FIFO that is used to prevent the CRC from being placed in the receive FIFO. When the temporary FIFO is full, data values are transferred to the receive FIFO one at a time. A frame's first data byte is the address. If receiver address matching is enabled, the received address is compared to the address in the address match value field in ICCR1. If the values match or the incoming address contains all ones, all following data bytes, including the address byte, are stored in the receive FIFO. If the values do not match, the receiver logic does not store any data in the receive FIFO, ignores the remainder of the frame, and searches for the next preamble. If receiver address matching is not enabled, the frame's first data byte is stored in the FIFO as normal data. The frame's second data byte can contain an optional control field and must be decoded in software.

The IrDA standard limits frames to any amount of data up to a 2047 bytes (including the address and control bytes). The FICP does not limit frame size. Software must ensure that each incoming frame does not exceed 2047 bytes.

When the receive FIFO reaches its trigger level, an interrupt (if enabled) and DMA transfer request (if no errors are detected in the data) are signalled. If the data is not removed quickly enough to prevent the FIFO from completely filling, the receive logic attempts to place additional data into the full FIFO and an overrun error is signalled. When the FIFO is full, all subsequent data bytes received are lost and all FIFO contents remain intact.

If the data field contains any invalid chips (such as 0011, 1010, 1110) the frame aborts and the oldest byte in the temporary FIFO is moved to the receive FIFO, the remaining temporary FIFO entries are discarded, the end-of-frame (EOF) tag is set in the FIFO entry that holds the last valid byte of data, and the receiver logic searches for the preamble.

The receive logic continuously searches for the 8-chip stop flag. When the stop flag is recognized, the last byte that was placed within the receive FIFO is flagged as the frame's last byte and the data in the temporary FIFO is removed and used as the CRC value for the frame. The receive logic compares the frame's CRC value to the CRC-32 value, which is continuously calculated from the incoming data stream. If CRC and CRC-32 values do not match, the last byte that was placed in the receive FIFO is also tagged with a CRC error. The frame's CRC value is not placed in the receive FIFO. If the stop flag is not properly detected, an abort is signalled.

If software disables the FICP's receiver while it is operating, the data byte being received stops immediately, the serial shifter and receive FIFO are cleared, the System Integration Unit (SIU) takes control of the receive data pin, and the clocks used by the receive logic are shut off to conserve power. The receive data input polarity must be reprogrammed if the receive data pin is used as a GPIO input.

11.2.9 Transmit Operation

Before it enables the FICP for transmission, the software can either preload the transmit FIFO by filling it with data or allow service requests to cause the CPU or DMA to fill the FIFO after the FICP is enabled. When the FICP is enabled, the transmit logic issues a service request if its FIFO requires more data.



A minimum of 16 preambles are transmitted for each frame. If data is not available after the sixteenth preamble, additional preambles are transmitted until a byte of valid data resides in the bottom of the transmit FIFO. The preambles are followed by the start flag and the data from the transmit FIFO. Groups of four chips (eight bits) are encoded and loaded in a serial shift register. The contents of the serial shift register are sent out on the transmit data pin, which is clocked by the 8-MHz baud clock. The preamble, start and stop flags, and CRC value are transmitted automatically.

When the transmit FIFO has 32 or more empty entries, an interrupt (if enabled) and DMA service request are sent. If new data does not arrive quickly enough to prevent the FIFO from becoming empty, the transmit logic attempts to transfer additional data from the empty FIFO. Software determines whether to interpret the data underrun (a lack of data) as a signal of normal frame completion or as an unexpected frame termination.

When software selects normal frame completion and an underrun occurs, the transmit logic transmits the CRC value that was calculated during data transmission, including the address and control bytes, followed by the stop flag that marks the end of the frame. The transmitter then continuously transmits preambles until data is available in the FIFO. When data is available, the transmitter starts to transmit the next frame.

When software selects unexpected frame termination and an underrun occurs, the transmit logic transmits an abort and interrupts the CPU. The transmitter continues to send the abort until data is available in the transmit FIFO. When data is available, the FICP transmits 16 preambles and a start flag and starts the new frame. The off-chip receiver can choose to ignore the abort and continue to receive data or signal the FICP to attempt to transmit the aborted frame again.

At the end of each transmitted frame, the FICP sends a pulse called the serial infrared interaction pulse (SIP). A SIP must be sent at least every 500 ms to ensure that low-speed devices (115.2 Kbps and slower) do not interfere with devices that transmit at higher speeds. The SIP simulates a start bit that causes low-speed devices to stay off the air for at least another 500 ms. The SIP pulse forces the transmit data pin high for $1.625~\mu s$ and low for $7.375~\mu s$ (the total SIP period is $9.0~\mu s$). After the SIP period, the preamble is transmitted continuously to indicate to the off-chip receiver that the FICP's transmitter is in the idle state. The preamble is transmitted until new data is available in the transmit FIFO or the FICP's transmitter is disabled. At least one frame must be completed every 500 ms to ensure that an SIP pulse can keep low-speed devices from interrupting the transmission. Because most IrDA compatible devices produce an SIP after each frame transmitted, software only needs to ensure that a frame is either transmitted or received by the FICP every 500 ms. Frame length does not represent a significant portion of the 500 ms timeframe in which an SIP must be produced. At 4.0 Mbps, the longest frame allowed is 16,568 bits, which takes just over 4 ms to transmit. The FICP also issues an SIP when the transmitter is first enabled. This ensures that low-speed devices do not interfere as the FICP transmits its data.

If software disables the FICP's transmitter during operation, data transmission stops immediately, the serial shifter and transmit FIFO are cleared, and the SIU takes control of the transmit data pin. The transmit data output's polarity must be properly reprogrammed if the pin is used as a GPIO output.

11.2.10 Transmit and Receive FIFOs

The transmit FIFO is 128 entries deep and 8 bits wide. The receive FIFO is 128 entries deep, 11 bits wide. The receive FIFO uses 3 bits of its entries as status bits. The transmit FIFO and the receive FIFO use two separate, dedicated DMA requests.



When the transmit FIFO has 32 or more empty bytes, the transmit DMA request and an interrupt (if enabled) are generated and tell the processor to send more data to the FIFO. When the transmit FIFO is full, any more data from the processor is lost. When the receive FIFO reaches its trigger level (programmed in ICCR2), the receive DMA request (if no errors are found within the entries) and an interrupt (if enabled) are generated and tell the processor to remove the data from the FIFO. If an error is found in the FIFO's trigger level range, DMA requests are disabled and an interrupt is generated to ensure that the DMAC does not read the error bytes.

The number of bytes being transferred for each DMA request is programmed in the DMAC and can be 8, 16, or 32 bytes. The receive FIFO's trigger level must be set so the FIFO has enough data for the DMAC to read. The transmit FIFO does not have programmable trigger levels. Its DMA request is generated when the FIFO has 32 or more empty bytes, regardless of the DMA transfer size.

The DMA controller must not service the receive FIFO when the processor tries to respond to a receive error interrupt. The error interrupt may be set high before the DMA controller finishes the previous request. The processor can not remove the error bytes until the DMAC has completed its transaction.

11.2.11 Trailing or Error Bytes in the Receive FIFO

When the number of bytes in the receive FIFO is less than the trigger level and no more data is being received, the bytes in the FIFO are called trailing bytes. Trailing bytes do not trigger a receive DMA request. Instead they trigger the end/error in FIFO, ICSR0[EIF] interrupt, which is nonmaskable. When ICSR0[EIF] is set, DMA requests are disabled. The core must read bytes from the FIFO until ICSR0[EIF] is cleared.

The core must also read bytes from the FIFO until ICSR0[EIF] is cleared if there are errors in FIFO entries below the DMA trigger level. When the entries below the DMA trigger level no longer contain status flags, DMA requests are enabled.

11.3 FICP Register Definitions

The FICP has six registers: three control registers, one data register, and two status registers. The FICP registers are 32 bits wide, but only the lower 8 bits have valid data. The FICP does not support byte or half-word operations. CPU reads and writes to the FICP registers must be word wide.

The control registers determine: IrDA transmission rate, address match value, how transmit FIFO underruns are handled, normal or active low transmit and receive data, whether transmit and receive operations are enabled, the FIFO interrupt service requests, receive address matching, and loopback mode.

The data register addresses the top of the transmit FIFO and the bottom of the receive FIFO. Reads to the data register access the receive FIFO. Writes to the data register access the transmit FIFO.

The status registers contain: CRC, overrun, underrun, framing, and receiver abort errors; the transmit FIFO service request; the receive FIFO service request; and end-of-frame conditions. Each of these hardware-detected events signals an interrupt request to the interrupt controller. The status registers also contain flags for transmitter busy, receiver synchronized, receive FIFO not empty, and transmit FIFO not full (no interrupt generated).



FICP Control Register 0 (ICCR0) 11.3.1

The ICCR0, shown in Table 11-2, contains eight valid bit fields that control various functions for 4 Mbps IrDA transmission. The FICP must be disabled (RXE=TXE=0) when ICCR0[ITR] and ICCR0[LBM] are changed. To allow various modes to be changed during active operation, ICCR0[7:2] may be written when the FICP is enabled.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Ta	ble	11	-2.	IC	CR	0 B	sit C	Defi	nit	ior	าร (Sh	eet	110	of 2	2)																
				0x4	1080	0_00)00				F			ared ol R						ort		F	ast I	Infra	red	Cor	nmu	ınica	atior	n Po	rt	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												rese	erve	b											AME	TE	RIE	RXE	TXE	TUS	LBM	ITR
Reset	0	0	0	0	0	0	0	0												0	0	0	0	0	0	0	0	0	0	0	0	0
															De	esci	ipti	on														
	[31:8] — reserved																															
	7 AME Address match enable. Receive logic will compatible in ICCR1. 0 = Disable receiver address is recogniz													pare addr iddre	ess ess	mat	tch f	unc unct	tion	. Sto	re o	data put	in r	ecei a in t	ve F	FIFC).					
			6			T	ΊE		0 = 1 = Set	Tr Tr	ansı ansı TIE	mit I mit I	FIFC FIFC es n) se	rvic rvic lear	e re e re	que que S or	st ge pre	ener vent	rate: t TF	s an	inte	errup bein	ot.	enei				·	t. tran	smit	t

1 = Receive FIFO service request generates an interrupt

Receive FIFO interrupt enable.

0 = Receive FIFO service request, ICSR0[RFS], does not generate an interrupt.

5

4

RIE

RXE



Table 11-2. ICCR0 Bit Definitions (Sheet 2 of 2)

0x4080_0000 Fast Infrared Communication Port Control Register 0 (ICCR0) Fast Infrared Communication Port

Bits	Name	Description
3	TXE	Transmit enable. 0 = FICP transmit logic disabled. 1 = FICP transmit logic enabled if ICCR0[ITR] is set. All other control bits must be configured before TXE is set. An SIP is transmitted immediately after the transmitter is enabled. If the transmit FIFO is empty, preambles are sent until data is placed in the FIFO. If TXE is cleared while it transmits data, transmission stops immediately, all data in the transmit FIFO and serial output shifter is cleared, and the SIU takes control of the transmit data pin. While communication is normally half-duplex, it is possible to transmit and receive data at the same time. Duplex communication is used for testing in Loopback Mode. If TXE is used to clear the transmitter, check ICSR1[TBY] to ensure the transmitter is not busy before the transmitter is re-enabled.
2	TUS	Transmit FIFO underrun select. A transmit FIFO underrun can either end the current frame normally, or transmit an abort. 0 = Transmit FIFO underrun causes CRC, stop flag, and SIP to be transmitted, and masks transmit underrun interrupt generation. 1 = Transmit FIFO underrun causes abort to be transmitted, and generates an interrupt. Clearing ICCR0[TUS] does not affect the current state of ICSR0[TUR] or prevent TUR from being set or cleared by the transmit FIFO. After an abort, a SIP is transmitted followed by 16 preambles. Preambles continue until data is in the FIFO.
1	LBM	Loopback mode. Used for testing FICP. 0 = Normal FICP operation enabled. 1 = Output of transmit serial shifter is connected to input of receive serial shifter.
0	ITR	IrDA transmission. 0 = ICP unit is not enabled. 1 = ICP unit is enabled.



11.3.2 FICP Control Register 1 (ICCR1)

The ICCR1, shown in Table 11-3, contains the 8-bit address match value field that the FICP uses to selectively receive frames. To allow the address match value to be changed during active receive operation, ICCR1 may be written while the FICP is enabled.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 11-3. ICCR1 Bit Definitions

				0x4	1080	0_00)04				F			ared ol R						ort		F	ast I	Infra	ired	Cor	mmu	ınica	atio	n Po	ort	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d														A۱	ΛV			
Reset	Bits Name								0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Na	me												De	escr	ipti	on										
	[31:8] —								res	erve	ed																					
		[7	:0]			ΑN	ΜV		The and and the	e 8-l d AN d da nex	bit va //V nata in kt pr	nato rec eam	use ches ceive nble	ed by the e FIF	ado FO.	res If th	s of e ad	the ddre	inco ss d	omir loes	ng fr s not	ame ma	, sto	ore ign	the i	fram the t	ne ad fram	ddre ne ai	ess, nd s	con	trol,	



11.3.3 FICP Control Register 2 (ICCR2)

The ICCR2, shown in Table 11-4, contains two bit fields that control the polarity of the transmit and receive data pins and two bits that determine the trigger level for the receive FIFO. The FICP must be disabled (RXE=TXE=0) when these bits are changed.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 11-4. ICCR2 Bit Definitions

				0x4	080	0_00	800				F					mm ster				ort		F	ast l	nfra	red	Cor	nmı	unica	atior	n Po	rt	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													r	ese	rve	ed													RXP	TXP	TP IC	2
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	0
		Bit	s			Naı	me												De	escr	ipti	on										
		[31:	4]			_	_		res	erve	d																					
		3				RΣ	〈Ρ		0 =	Da Da	ta f ata f	rom	the	rec	eiv	t. e da e da											the	FICI	P ur	nit.		
		2				ΤX	(P		0 =	Da Da	ta f ata f	rom	the	FIC	Ρi	ct. s inv s no														ո.		
		[1:0	0]			TR	IG		The and size be 0	e red I has e as disa 00- r	eives no the bled ece	erro Reo d. ive f ive f	FIFC	gene n its e FI O se	era da FO rvic	tes sata. T trigg ce re ce re	The I ger I que	OM/ evel st is st is	co . To ger	ntro cha nera nera	ller ange ated ated	data the whe whe	trar trig en th	nsfe Iger ne F	r siz leve IFO IFO	e m el, th has has	ust l le R 8 b 16	be so ecei ytes byte	et to ive F s or i	the FIFC mor	e sar O mi e ore	me

0b11- reserved



11.3.4 FICP Data Register (ICDR)

The ICDR, shown in Table 11-5, is a 32-bit register and its lower 8 bits are the top entry of the transmit FIFO when the register is written and the bottom entry of the receive FIFO when the register is read.

Reads to ICDR access the lower 8 bits of the receive FIFO's bottom entry. As data enters the top of the receive FIFO, bits 8 – 10 are used as tags to indicate conditions that occur as each piece of data is received. The tag bits are transferred down the FIFO with the data byte that encountered the condition. When data reaches the bottom of the FIFO, bit 8 of the FIFO entry is transferred to the end-of-frame (EOF) flag, bit 9 to the CRC error (CRE) flag, and bit 10 to the receiver overrun (ROR) flag. All these flags are in FICP status register 1. These flags can be read to determine whether the value at the bottom of the FIFO represents the frame's last byte or an error that was encountered during reception. After the flags are checked, the FIFO value can be read. This causes the data in the next location of the receive FIFO to be transferred to the bottom entry and its EOF, CRE, and ROR bits to be transferred to the status register.

The end/error in FIFO (EIF) flag is set in status register 0 when a tag bit is set in any of the receive FIFO's bottom eight, 16, or 32 entries, as determined by the trigger level. The EIF flag is cleared when no error bits are set in the FIFO's bottom entries. When EIF is set, an interrupt is generated and the receive FIFO DMA request is disabled. Software must empty the FIFO and check for the EOF, CRE, and ROR error flags in ICSR1 before it removes each data value from the FIFO. After each entry is removed, the EIF bit must be checked to determine if any set end or error tag remains and the procedure is repeated until all set tags are flushed from the FIFO's bottom entries. When EIF is cleared, DMA service for the receive FIFO is re-enabled.

Both FIFOs are cleared when the processor is reset. The transmit FIFO is cleared when TXE is 0. The receive FIFO is cleared when RXE is 0.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 11-5. ICRD Bit Definitions

				0x4	1080	0_00	00C				Fa	ast I				nmı ster			n Po	ort		Fa	ast I	Infra	red	Cor	nmı	unica	atior	n Po	ort	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rvec	ł														DA	TA			
Reset													0	0	0	0	0	0	0	0	0	0	0	0	0	0	0					
	Bits Name																		De	escr	ipti	on										
•	[31:8] —								res	erve	d																					
	[7:0] DATA										Rea	ad d	ata	fron	n fro	nt o	f red	eiv	e FII													



11.3.5 FICP Status Register 0 (ICSR0)

The ICSR0, shown in Table 11-6, contains bits that signal the transmit FIFO service request, receive FIFO service request, receiver abort, transmit FIFO underrun, framing error, and the end/error in receive FIFO conditions. Each of these hardware-detected events signal an interrupt request to the interrupt controller.

If a bit signals an interrupt request, it signals the interrupt request as long as the bit is set. When the bit is cleared, the interrupt is cleared. Read/write bits are called status bits. Read-only bits are called flags. Status bits that must be cleared by software after they are set by hardware are called sticky status bits. Writing a 1 to a sticky status bit clears it. Writing a 0 to a sticky status bit has no effect. Read-only flags are set and cleared by hardware. Writes to read-only flags have no effect. Some bits that cause interrupts have corresponding mask bits in the control registers.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 11-6. ICSR0 Bit Definitions (Sheet 1 of 2)

	0x4080_0014						4	Fast Infrared Communication Port Status Register 0 (ICSR0) Fast Infrared Communication Port											t												
Bit	31	30	29	28	27 2	6 2	5 2	4 23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rvec	ł												FRE	RFS	TFS	RAB	TUR	EIF
Reset	0	0	0	0	0 0) (0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts		N	am	е											De	escr	iptio	on										
		[31	:6]			_		res	erve	ed																					
	[31:6] — reserved Framing error. 5 FRE 0 = No framing errors encountered in the receipt of this data. 1 = Framing error occurred. A preamble was followed by something other than anoth preamble or start flag, request interrupt.										her																				
		4	ļ		F	RFS	;	0 =	= Re	e FIF eceiv eceiv ques	/e F	IFC	has has	no s rea	t rea	che ed it	ed it s tri	trig gger	ger r lev	el a	nd re	ecei	ver i	s er	nabl	led.			ervice	Э	
		3	3		-	ΓFS	;	Tra 0 = 1 =	Tr Tr	nit FI ansr ansr ques	nit I nit I	FIFC) ha) ha	s m s 96	ore or	thar less	n 96 ent	ent ries	ries of o	data	and	trai	nsm	itter	is e	enab	led.		IA se	ervio	е



Table 11-6. ICSR0 Bit Definitions (Sheet 2 of 2)

				0x4	4080	0_00)14				F		Infra Status							ort		F	ast I	nfra	red	Cor	nmu	ınica	atior	n Po	rt	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													rese	rve	ł												FRE	RFS	TFS	RAB	TUR	FIF
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Bits	Name	Description
		Receiver abort.
2	RAB	 0 = No abort has been detected for the incoming frame. 1 = Abort detected during receipt of incoming frame. Two or more chips containing no pulses or any invalid chips were detected on the receive pin. EOF bit set on last piece of "good" data received before the abort, interrupt requested.
		Transmit FIFO underrun.
		0 – Transmit FIFO has not experienced an underrun.
1	TUR	1 – Transmit logic attempted to fetch data from transmit FIFO while it was empty. Interrupt request signalled if not masked by ICCR0[TUS].
		Underruns are not generated when the FICP transmitter is first enabled and is idle.
		End/error in FIFO (read-only).
		0 – Bits 8–10 are not set within any of the entries at or below the trigger level of the receive FIFO. Receive FIFO DMA service requests are enabled.
0	EIF	1 – One or more tag bits (8 – 10) are set within the entries at or below the trigger level of the receive FIFO. Request interrupt, disable receive FIFO DMA service requests.
		This interrupt is not maskable in the FICP. Once the bad bytes have been removed from the FIFO and EIF is cleared, DMA requests are automatically enabled.



11.3.6 FICP Status Register 1 (ICSR1)

ICSR1, shown in Table 11-7, contains flags that indicate that the receiver is synchronized, the transmitter is active, the transmit FIFO is not full, the receive FIFO is not empty, and that an EOF, CRE, or underrun error has occurred.

This is a read-only register. Ignore reads from reserved bits.

Table 11-7, ICSR1 Bit Definitions

	0x4080_0018					118					t Infra Statu							ort		F	ast l	Infra	red	Cor	nmı	ınica	atior	ı Po	rt		
Bit	31	30	29	28	27	26	25	24	23 2	2 2	20	0 19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	eserv	red												ROR	CRE	EOF	TNF	RNE	ТВУ	RSY
Reset	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0
		В	ts			Nar	ne											De	esci	ipti	on										
		[31	:7]			_			reser	/ed																					
	Receive FIFO overrun (re							•																							
		(6			RC)R			Rece	ive the	logic FIFC	atte) is f	mpt ull a	ed t ire l	o pla ost.	ace	data	a int	o re	ceiv										
									Each entry									וטט	lOIII	OI II	ie ie	ecer	ver	IFU	, DIL	. 101	1011	ı trie	lasi	LFII	-0
	entry is transferred to the ROR bit. CRC error (read-only). 0 = CRC not encountered yet or no CRC check errors encountered in the receipt of dat 1 = CRC calculated on the incoming data. Does not match CRC value contained within the received frame. Each time an 11-bit value reaches the bottom of the receive FIFO, bit 9 from the last FIF										1																				
									entry									DOL	lon	OI I	iie i	ece	ive		J, DI	11 9 1	10111	uie	iasi		O
		4	1			EC)F		Each	Current he nclu time	ent valu ding e an	frame ue at t g abo	e has the b rted it va	s no ootto fran lue i	om c nes. reac	of the ches	e red the	ceiv													
									entry																						
		;	3			TN	IF		Trans 0 = 1 1 = 1	ran	smi	t FIFC	O is	tull.		• ,		rupt	t ge	nera	ited)).									
	1 = Transmit FIFO is not full (no interrupt generated). Receive FIFO not empty (read-only). 2 RNE 0 = Receive FIFO is empty. 1 = Receive FIFO is not empty (no interrupt generated).																														
			1			ТВ	Υ		Trans 0 = 1 1 = 1	ran ran	smit smit	tter is	idle c is c	(co curre	ntin ently	uou / tra	nsm	ittin						, coi	ntrol	l, da	ta, C	CRC	, or	sta	rt/
		()			RS	SY.		Rece		•			•	•		•	,	ed.				. ,								

1 = Receiver logic is synchronized with the incoming data (no interrupt generated).



11.4 FICP Register Summary

Table 11-8 shows the registers associated with the FICP block and the physical addresses used to access them.

Table 11-8. FICP Register Summary

Address	Name	Description
0x4080_0000	ICCR0	FICP control register 0
0x4080_0004	ICCR1	FICP control register 1
0x4080_0008	ICCR2	FICP control register 2
0x4080_000C	ICDR	FICP data register
0x4080_0010	_	reserved
0x4080_0014	ICSR0	FICP status register 0
0x4080_0018	ICSR1	FICP status register 1



USB Device Controller

This section describes the Universal Serial Bus (USB) protocol and its implementation-specific options for device controllers for the PXA255 processor. These options include endpoint number, type, and function; interrupts to the core; and a transmit/receive FIFO interface. A working knowledge of the USB standard is vital to using this section effectively. The Universal Serial Bus Device Controller (UDC) is USB-compliant and supports all standard device requests issued by the host. UDC operation summaries and quick reference tables are provided. Refer to the *Universal Serial Bus Specification*, revision 1.1, for a full description of the USB protocol. The *Universal Serial Bus Specification* is available at http://www.usb.org.

12.1 USB Overview

The UDC supports 16 endpoints and can operate half-duplex at a rate of 12 Mbps (as a slave only, not as a host or hub controller). The UDC supports four device configurations. Configurations 1, 2, and 3 each support two interfaces. Alternate interface settings are not supported. This allows the host to accommodate dynamic changes in the physical bus topology. A configuration is a specific combination of USB resources available on the device. An interface is a related set of endpoints that present a device feature or function to the host.

The UDC transmits serial information that contains layers of communication protocols. Fields are the most basic protocol. UDC fields include: sync, packet identifier (PID), address, endpoint, frame number, data, and Cyclic Redundancy Check (CRC). Fields are combined to produce packets. A packet's function determines the combination and number of fields that make up the packet. Packet types include: token, start of frame, data, and handshake. Packets are assembled into groups to produce transactions. Transactions fall into four groups: bulk, control, interrupt, and isochronous. Endpoint 0 is used only to communicate the control transactions that configure the UDC. Endpoint 0's responsibilities include: connection, address assignment, endpoint configuration, bus enumeration, and disconnection.

The UDC uses a dual-port memory to support FIFO operations. Each Bulk and Isochronous Endpoint FIFO structure is double buffered to enable the endpoint to process one packet as it assembles another. The DMA and the Megacell can fill and empty the FIFOs. An interrupt or DMA service request is generated when a packet has been received. The DMA engine services the UDC FIFOs in 32-byte increments. Interrupts are also generated when the FIFO encounters a short packet or zero-length packet. Endpoint 0 has a 16-entry long, 8-bit wide FIFO that can only be read or written by the processor.

For endpoints 1-15, the UDC uses its dual-ported memory to hold data for a Bulk OUT transaction while the transaction is checked for errors. If the Bulk OUT transaction data is invalid, the UDC sends a NAK handshake to request the host to resend the data. The software is not notified that the OUT data is invalid until the Bulk OUT data is received and verified. If the host sends a NAK handshake in response to a Bulk IN data transmission, the UDC resends the data. Because the FIFO maintains a copy of the data, the software does not have to reload the data.

The external pins dedicated to the UDC interface are UDC+ and UDC-. The USB protocol uses differential signalling between the two pins for half-duplex data transmission. A 1.5 k Ω pull-up resistor must be connected to the USB cable's D+ signal to pull the UDC+ pin high when it is not driven. Pulling the UDC+ pin high when it is not driven allows the UDC to be a high-speed,



12-Mbps device and provides the correct polarity for data transmission. The serial bus uses differential signalling to transmit multiple states simultaneously. These states are combined to produce transmit data and various bus conditions, including: Idle, Resume, Start of Packet, End of Packet, Disconnect, Connect, and Reset.

12.2 Device Configuration

Table 12-1 shows the device's configuration.

Table 12-1. Endpoint Configuration

Endpoint Number	Туре	Function	FIFO Size (bytes) X number of FIFOs
0	Control	IN/OUT	16
1	Bulk	IN	64x2
2	Bulk	OUT	64x2
3	Isochronous	IN	256x2
4	Isochronous	OUT	256x2
5	Interrupt	IN	8
6	Bulk	IN	64x2
7	Bulk	OUT	64x2
8	Isochronous	IN	256x2
9	Isochronous	OUT	256x2
10	Interrupt	IN	8
11	Bulk	IN	64x2
12	Bulk	OUT	64x2
13	Isochronous	IN	256x2
14	Isochronous	OUT	256x2
15	Interrupt	IN	8

Data flow is relative to the USB host. IN packets represent data flow from the UDC to the host. OUT packets represent data flow from the host to the UDC.

The FIFOs for the Bulk and Isochronous endpoints are double-buffered so one packet can be processed as the next is assembled. While the UDC transmits an IN packet from a particular endpoint, the Megacell can load the same endpoint for the next frame transmission. While the Megacell unloads an OUT endpoint, the UDC can continue to process the next incoming packet to that endpoint.

12.3 USB Protocol

After a core reset or when the USB host issues a USB reset, the UDC configures all endpoints and is forced to use the USB default address, zero. After the UDC configures the endpoints, the host assigns the UDC a unique address. At this point, the UDC is under the host's control and responds to commands that use control transactions to transmit to endpoint 0.



12.3.1 Signalling Levels

USB uses differential signalling to encode data and to indicate various bus conditions. The USB specification refers to the J and K data states to differentiate between high- and low-speed transmissions. Because the UDC supports only 12 Mbps transmissions, references are only made to actual data state 0 and actual data state 1.

By decoding the polarity of the UDC+ and UDC- pins and using differential data, four distinct states are represented. Two of the four states are used to represent data. A 1 indicates that UDC+ is high and UDC- is low. A 0 indicates that UDC+ is low and UDC- is high. The two remaining states and pairings of the four encodings are further decoded to represent the current state of the USB. Table 12-2 shows how differential signalling represents eight different bus states.

Table 12-2. USB States

Bus State	UDC+/UDC- Pin Levels
Idle	UDC+ high, UDC- low (same as a 1).
Suspend	Idle state for more than 3 ms.
Resume	UDC+ low, UDC- high (same as a 0).
Start of Packet	Transition from idle to resume.
End of Packet	UDC+ AND UDC- low for 2 bit times followed by an idle for 1 bit time.
Disconnect	UDC+ AND UDC- below single-ended low threshold for more than 2.5 µs. (Disconnect is the static bus condition that results when no device is plugged into a hub port.)
Connect	UDC+ OR UDC- high for more than 2.5 μs.
Reset	UDC+ AND UDC- low for more than 2.5 μs . (Reset is driven by the host controller and sensed by a device controller.)

Hosts and hubs have pull-down resistors on both the D+ and D- lines. When a device is not attached to the cable, the pull-down resistors cause D+ and D- to be pulled down below the single-ended low threshold of the host or hub. This creates a state called single-ended zero (SE0). The host detects a disconnect when an SE0 persists for more than 2.5 μ s (30 bit times). When the UDC is connected to the USB cable, the pull-up resistor on the UDC+ pin causes D+ to be pulled above the single-ended high threshold level. After 2.5 μ s, the host detects a connect.

After the host detects a Connect, the bus is in the Idle state because UDC+ is high and UDC- is low. The bus transitions from the Idle state to the Resume state (a 1 to 0 transition) to signal the Start of Packet (SOP). Each USB packet begins with a Sync field that starts with the 1-to-0 transition (see Section 12.3.1). After the packet data is transferred, the bus signals the End of Packet (EOP) state by pulling both UDC+ and UDC- low for 2 bit times followed by an Idle state for 1 bit time. If the idle persists for more than 3 ms, the UDC enters Suspend state and is placed in low-power mode. The host can awaken the UDC from the Suspend state by signalling a reset or by switching the bus to the resume state via normal bus activity. Under normal operating conditions, the host periodically signals an Start of Frame (SOF) to ensure that devices do not enter the suspend state.

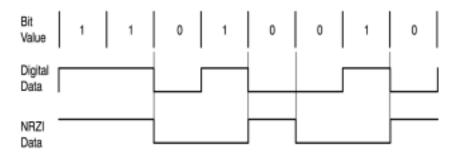
12.3.2 Bit Encoding

USB uses nonreturn to zero inverted (NRZI) to encode individual bits. Both the clock and the data are encoded and transmitted in the same signal. Data is represented by transitions rather than by the signal's state. A zero is represented by a transition, and a one is represented by no transition, which produces the data. Each time a zero occurs, the receiver logic synchronizes the baud clock to the



incoming data, which produces the clock. To ensure the receiver is periodically synchronized, six consecutive ones in the serial bit stream trigger the transmitter to insert a zero. This procedure is known as bit stuffing. The receiver logic detects stuffed bits and removes them from incoming data. Bit stuffing causes a transition on the incoming signal at least once every 7 bit times to ensure the baud clock is locked. Bit stuffing is enabled for an entire packet from the time the SOP is detected until the EOP is detected (enabled during the Sync field through the CRC field). Figure 12-1 shows the NRZI encoding of the data byte 0b1101 0010.

Figure 12-1. NRZI Bit Encoding Example



12.3.3 Field Formats

Individual bits are assembled into groups called fields. Fields are used to construct packets and packets are used to construct frames or transactions. There are seven USB field types: Sync, PID, Address, Endpoint, Frame Number, Data, and CRC.

A Sync is preceded by the Idle state and is the first field of every packet. The first bit of a Sync field signals the SOP to the UDC or host. A Sync is 8 bits wide and consists of seven zeros followed by a one (0x80). Bits are transmitted to the bus least significant bit first in every field, except the CRC field.

The PID is 1 byte wide and always follows the sync field. The first four bits contain an encoded value that represents packet type (Token, Data, Handshake, and Special), packet format, and type of error detection. The last four bits contain a check field that ensures the PID is transmitted without errors. The check field is generated by performing a ones complement of the PID. The UDC XORs the PID and CRC fields and takes the action prescribed in the USB standard if the result does not contain all ones, which indicates an error has occurred in transmission.

The Address and Endpoint fields are used to access the UDC's 16 endpoints. The Address field contains seven bits and permits 127 unique devices to be placed on the USB. After the USB host signals a reset, the UDC and all other devices are assigned the default address, zero. The host is then responsible for assigning a unique address to each device on the bus. Addresses are assigned in the enumeration process, one device at a time. After the host assigns the an address to the UDC, the UDC only responds to transactions directed to that address. The Address field follows the PID in every packet transmitted.

When the UDC detects a packet that is addressed to it, it uses the Endpoint field to determine which of the UDC's endpoints is being addressed. The Endpoint field contains four bits. Encodings for endpoints 0 (0000b) through 15 (1111b) are allowed. The Endpoint field follows the Address field.



The Frame Number is an 11-bit field incremented by the host each time a frame is transmitted. When it reaches its maximum value, 2047 (0x7FF), its value rolls over. Frame Number is transmitted in the SOF packet, which the host outputs in 1 ms intervals. Device controllers use the Frame Number field to control isochronous transfers. Data fields are used to transmit the packet data between the host and the UDC. A data field consists of 0 to 1023 bytes. Each byte is transmitted least significant bit first. The UDC generates an interrupt to indicate that a Start of Frame event has occurred.

CRC fields are used to detect errors introduced during token and data packet transmission, and are applied to all the fields in the packet except the PID field. The PID contains its own 4-bit ones complement check field for error detection. Token packets use a 5-bit CRC (x^5+x^2+1) called CRC5 and Data packets use a 16-bit CRC ($x^{16}+x^{15}+x^2+1$) called CRC16. For both CRCs, the checker resets to all ones at the start of each packet.

12.3.4 Packet Formats

USB supports four packet types: Token, Data, Handshake, and Special. A PRE (Preamble) PID precedes a low-speed (1.5 Mbps) USB transmission. The UDC supports high-speed (12 Mbps) USB transfers only. PRE packets that signify low-speed devices and the low-speed data transfer that follows such PRE packets are ignored.

12.3.4.1 Token Packet Type

A Token packet is placed at the beginning of a frame and is used to identify OUT, IN, SOF, and SETUP transactions. OUT and IN frames are used to transfer data, SOF packets are used to time isochronous transactions, and SETUP packets are used for control transfers to configure endpoints. A Token packet consists of a sync, a PID, an address, an endpoint, and a CRC5 field (see Table 12-3). For OUT and SETUP transactions, the address and endpoint fields are used to select the UDC endpoint that receives the data. For an IN transaction, the address and endpoint fields are used to select the UDC endpoint that transmits data.

Table 12-3. IN, OUT, and SETUP Token Packet Format

8 bits	8 bits	7 bits	4 bits	5 bits
Sync	PID	Address	Endpoint	CRC5

12.3.4.2 Start of Frame Packet Type

An SOF is a special type of Token packet that the host issues at a nominal interval of once every 1 ms +/- 0.0005 ms. SOF packets consist of a Sync, a PID, a Frame Number (incremented after each frame is transmitted), and a CRC5 field (see Table 12-4). The presence of SOF packets every 1 ms prevents the UDC from entering Suspend mode.

Table 12-4. SOF Token Packet Format

8 bits	8 bits	11 bits	5 bits
Sync	PID	Frame Number	CRC5



12.3.4.3 Data Packet Type

Data packets follow Token packets and are used to transmit data between the host and UDC. The PID specifies two types of Data packets: DATA0 and DATA1. These Data packets are used to provide a mechanism to ensure that the data sequence between the transmitter and receiver is synchronized across multiple transactions. During the handshake phase, the transmitter and receiver determine which data token type to transmit first. For each subsequent packet transmitted, the data packet type is toggled (DATA0, DATA1, DATA0, and so on). A Data packet consists of a Sync, a PID, from 0 to 1023 bytes of data, and a CRC16 field (see Table 12-5). The UDC supports a maximum of eight bytes of data for an Interrupt IN data payload, a maximum of 64 bytes of data for a Bulk data payload and a maximum of 256 bytes of data for an Isochronous data payload.

Table 12-5. Data Packet Format

8 bits	8 bits	0-1023 bytes	16 bits
Sync	PID	Data	CRC16

12.3.4.4 Handshake Packet Type

Handshake packets consist of a Sync and a PID. Handshake packets do not contain a CRC because the PID contains its own check field. Handshake packets are used to report data transaction status, including confirmation that data was successfully received, flow control, and stall conditions. Only transactions that support flow control can return handshakes. The three types of handshake packets are: ACK, NAK, and STALL. ACK indicates that a data packet was received without bit stuffing, CRC, or PID check errors. NAK indicates that the UDC was unable to accept data from the host or has no data to transmit. STALL indicates that the UDC was unable to transmit or receive data, and requires host intervention to clear the stall condition. The receiving unit signals Bit stuffing, CRC, and PID errors by omitting a handshake packet. Table 12-6 shows the format of a handshake packet.

Table 12-6. Handshake Packet Format

8 bits	8 bits
Sync	PID

12.3.5 Transaction Formats

Packets are assembled into groups to form transactions. The USB protocol uses four different transaction formats. Each transaction format is specific to a particular type of endpoint: Bulk, Control, Interrupt, or Isochronous. Endpoint 0, by default, is a control endpoint and receives only control transactions. All USB transactions are initiated by the host controller and transmitted in one direction at a time (known as half-duplex) between the host and UDC.

12.3.5.1 Bulk Transaction Type

Bulk transactions guarantee error-free data transmission between the host and UDC by using packet error detection and retry. The host schedules bulk packets when the bus has available time. Bulk transactions are made up of three packet types: Token, Data, and Handshake. The eight types of bulk transactions are based on data direction, error, and stall conditions. The types of bulk transactions are shown in Table 12-7. Packets sent from the UDC to the host are highlighted in boldface type and packets sent from the host to the UDC are not.



Table 12-7. Bulk Transaction Formats

Action	Token Packet	Data Packet	Handshake Packet
Host successfully received data from UDC	IN	DATA0/DATA1	ACK
UDC temporarily unable to transmit data	IN	None	NAK
UDC endpoint needs host intervention	IN	None	STALL
Host detected PID, CRC, or bit stuff error	IN	DATA0/DATA1	None
UDC successfully received data from host	OUT	DATA0/DATA1	ACK
UDC temporarily unable to receive data	OUT	DATA0/DATA1	NAK
UDC endpoint needs host intervention	OUT	DATA0/DATA1	STALL
UDC detected PID, CRC, or bit stuff error	OUT	DATA0/DATA1	None
	Packets	from UDC to host are	boldface

12.3.5.2 Isochronous Transaction Type

Isochronous transactions ensure constant rate, error-tolerant transmission of data between the host and UDC. The host schedules isochronous packets during every frame. USB protocol allows isochronous transfers to take up to 90% of the USB bandwidth. Unlike bulk transactions, if corrupted data is received, the UDC will continue to process the corrupted data that corresponds to the current start of frame indicator. Isochronous transactions do not support a handshake phase or retry capability. Two packet types are used to construct isochronous transactions: token and data. The types of isochronous transactions based on data direction are shown in Table 12-8.

Table 12-8. Isochronous Transaction Formats

Action	Token Packet	Data Packet
Host received data from UDC	IN	DATA0
UDC received data from host	OUT	DATA0
	Packets from UDC t	o host are boldface

12.3.5.3 Control Transaction Type

The host uses control transactions to configure endpoints and query their status. Like bulk transactions, control transactions begin with a setup packet, followed by an optional data packet, then a handshake packet. Control transactions, by default, use DATA0 type transfers. Table 12-9 shows the four types of control transactions.

Table 12-9. Control Transaction Formats

Action	Token Packet	Data Packet	Handshake Packet
UDC successfully received control from host	SETUP	DATA0	ACK
UDC temporarily unable to receive data	SETUP	DATA0	NAK
UDC endpoint needs host intervention	SETUP	DATA0	STALL
UDC detected PID, CRC, or bit stuff error	SETUP	DATA0	None
	Packets	from UDC to host are	boldface



To assemble control transfers, the host sends a control transaction to tell the UDC what type of control transfer is taking place (control read or control write), followed by one or more data transactions. The setup is the first stage of the control transfer. The device must respond with an ACK or no handshake (if the data is corrupted). The control transaction, by default, uses a DATA0 transfer and each subsequent data transaction toggles between DATA1 and DATA0 transfers. A control write to an endpoint uses OUT transactions. Control reads use IN transactions. The transfer direction is the opposite of the last data transaction. The transfer direction is used to report status and functions as a handshake. For a control write, the last transaction is an IN from the UDC to the host. For a control read, the last transaction is an OUT from the host to the UDC. The last data transaction always uses a DATA1 transfer, even if the previous transaction used DATA1.

12.3.5.4 Interrupt Transaction Type

The host uses interrupt transactions to query the status of the device. Like bulk transactions, interrupt transactions begin with a setup packet, followed by an optional data packet, then a handshake packet. Interrupt transactions, by default, use DATA0 type transfers. Figure 12-10 shows the four types of interrupt transactions.

Table 12-10. Interrupt Transaction Formats

Action	Token Packet	Data Packet	Handshake Packet
Host successfully received data from UDC	IN	DATA0	ACK
UDC temporarily unable to transmit data	IN	None	NAK
UDC endpoint needs host intervention	IN	None	STALL
Host detected PID, CRC, or bit stuff error	IN	DATA0	None
	Packets	from UDC to host are	boldface

12.3.6 UDC Device Requests

The UDC uses its control, status, and data registers to control and monitor the transmit and receive FIFOs for endpoints 1 - 15. The host controls all other UDC configuration and status reporting using device requests that are sent as control transactions to endpoint 0 via the USB. Each setup packet to Endpoint 0 is 8 bytes long and specifies:

- Data transfer direction: host to device, device to host
- Data transfer type: standard, class, vendor
- Data recipient: device, interface, endpoint, other
- Number of bytes to transfer
- · Index or offset
- Value: used to pass a variable-sized data parameter
- Device request

Table 12-11 shows a summary of all device requests. Refer to the *Universal Serial Bus Specification Revision 1.1* for a full description of host device requests.



Request	Name
SET_FEATURE	Enables a specific feature such as device remote wake-up or endpoint stalls.
CLEAR_FEATURE	Clears or disables a specific feature.
SET_CONFIGURATION	Configures the UDC for operation. Used after a reset of the Megacell or after a reset has been signalled via the USB.
GET_CONFIGURATION	Returns the current UDC configuration to the host.
SET_DESCRIPTOR	Sets existing descriptors or add new descriptors. Existing descriptors include: device, configuration, string, interface, and endpoint.
GET_DESCRIPTOR	Returns the specified descriptor, if it exists.
SET_INTERFACE	Selects an alternate setting for the UDC's interface.
GET_INTERFACE	Returns the selected alternate setting for the specified interface.
GET_STATUS	Returns the UDC's status including: remote wake-up, self-powered, data direction, endpoint number, and stall status.
SET_ADDRESS	Sets the UDC's 7-bit address value for all future device accesses.
SYNCH_FRAME	Sets then reports an endpoint's synchronization frame.

The UDC decodes most standard device commands with no intervention required by the user. The following commands are not passed to the user are: Set Address, Set Feature, Clear Feature, Get Configuration, Get Status, Get Interface, and Sync Frame. The Set Configuration and Set Interface commands are passed to the user to indicate that the host set the specified configuration or interface and the software must take any necessary actions. Alternate interfaces settings are not supported; the host must set the alternate settings field in SET_INTERFACE requests to zero. If the UDC receives a SET_INTERFACE request with the alternate settings field non-zero, the UDC responds with a STALL. The Get Descriptor and Set Descriptor commands are passed to the user to be decoded.

Because the Set Feature and Clear Feature commands are not passed on, the user is not able decode the device remote wakeup feature commands. To solve this problem, the status bit UDCCS0:DRWF indicates whether or not the device remote wakeup feature is enabled. UDCCS0: DRWF is a read-only bit. When the bit is set to 1, the device remote wakeup feature is enabled. When the bit is set to 0, the feature is not enabled.

12.3.7 Configuration

In response to the GET_DESCRIPTOR command, the user device sends back a description of the UDC configuration. The UDC can physically support more data channel bandwidth than the USB specification allows. When the device responds to the host, it must specify a legal USB configuration. For example, if the device specifies a configuration of six isochronous endpoints of 256 bytes each, the host is not be able to schedule the proper bandwidth and does not take the UDC out of Configuration 0. The user device determines which endpoints to report to the host. If an endpoint is not reported, it is not used. Another option, attractive for use with isochronous endpoints, is to describe a configuration of a packet with a maximum size less than 256 bytes to the host. For example, if software responds to the GET_DESCRIPTOR command that Endpoint 3 only supports 64 bytes maximum packet Isochronous IN data, the user device must set the UDCCS3[TSP] bit after it loads 64 bytes for transmission. Similarly, if Endpoint 4 is described as supporting 128 bytes maximum packet Isochronous OUT data, the UDC recognizes the end of the packet, sets UDCCS4[RPC], and an interrupt is generated.



The direction of the endpoints is fixed. Physically, the UDC only supports interrupt endpoints with a maximum packet size of 8 bytes or less, bulk endpoints with a maximum packet size of 64 bytes or less, and isochronous endpoints with a maximum packet size of 256 bytes or less.

To make the processor more adaptable, the UDC supports a total of four configurations. Each of these configurations are identical in the UDC, software can make three distinct configurations, each with two interfaces. Configuration 0 is a default configuration of Endpoint 0 only and cannot be defined as any other arrangement.

After the host completes a SET_CONFIGURATION or SET_INTERFACE command, the software must decode the command to empty the OUT endpoint FIFOs and allow the Megacell to set up the proper power/peripheral configurations.

12.4 UDC Hardware Connection

This section explains how to connect the USB interface for a variety of devices.

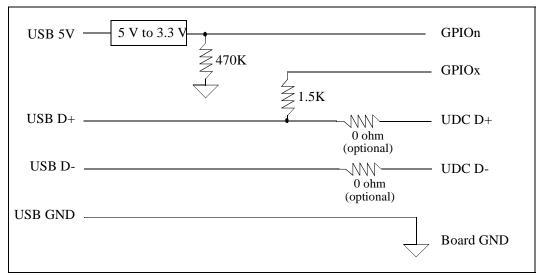
12.4.1 Self-Powered Device

Figure 12-2 shows how to connect the USB interface for a self-powered device. The $0\,\Omega$ resistors are optional and if they are not used, USB D+ must connect directly to the device UDC D+ and connect USB D- must connect directly to the device UDC D-. The UDC D+ and UDC D- pins are designed to match the impedance of a USB cable, $90\,\Omega$, without external series resistors. To allow minor impedance corrections to compensate for the impedance that results from the board trace, $0\,\Omega$ resistors are recommended on the board.

The "5 V to 3.3 V" device is required because the input pins of the processor can only tolerate 3.3 V. The device can be implemented in a number of ways. The most robust and expensive solution is a Power-On-Reset device such as a MAX6348. This solution produces a clean signal edge and minimizes signal bounce. A more inexpensive solution is a 3.3 V line buffer with inputs that can tolerate 5 V. This solution does not reduce signal bounce, so software must compensate by reading the GPIO repeatedly until it proves to be stable. A third solution is a signal bounce minimization circuit that can tolerate 5 V but produces a 3.3 V signal to the GPIO pin.







12.4.1.1 When GPIOn and GPIOx are Different Pins

The GPIOn and GPIOx pins can be any GPIO pins. GPIOn must be a GPIO that can wake the device from sleep mode. After a reset, GPIOx is configured as an input. This causes the UDC+ line to float. GPIOn is configured to act as an input and to cause an interrupt on a rising or falling edge. When an interrupt occurs, software must read the GPIOn pin to determine if the cable is connected. The GPIOn pin is set to a 1 when the cable is connected and a 0 when the cable is disconnected. When a USB connect is detected, software must enable the UDC peripheral and drive a 1 to the GPIOx pin to indicate to the host PC that a high-speed USB device is connected. When a USB disconnect is detected, software must configure the GPIOx pin as an input, configure the GPIOn pin to detect a wakeup event, and put the disconnected peripheral in sleep mode, if desired.

If software must put a peripheral in sleep mode, it configures the GPIOx pin as an input. This causes the UDC+ line to float, which appears to be a disconnect to the host PC. The peripheral is put in sleep mode. When the peripheral comes out of sleep mode, software must drive a 1 to the GPIOx pin to indicate to the host PC that a high-speed USB peripheral is connected.

12.4.1.2 When GPIOn and GPIOx are the Same Pin

After a reset, GPIOn is configured to act as an input and to cause an interrupt on a rising or falling edge. When an interrupt occurs, software must read the GPIOn pin to determine if the cable is connected. GPIOn is set to a 1 when the cable is connected and a 0 when the cable is disconnected. If a USB connect is detected, software must enable the UDC peripheral before the host PC sends the first USB command. If a USB disconnect is detected, software must configure the GPIOn pin to detect a wakeup event and put the peripheral in sleep mode, if desired.

When GPIOn and GPIOx are the same pin, do not put a peripheral in sleep mode if the USB cable is connected to the device. During sleep, the USB controller is in reset and does not respond to the host PC. When it returns from sleep mode, the peripheral does not respond to its host-assigned address.



12.4.2 Bus-Powered Devices

The processor does not support bus-powered devices because it is required to consume less that 500 μA when the host issues a suspend (see Section 7.2.3 of the *USB Specification*, version 1.1). The processor cannot limit the amount of current it consumes to 500 μA unless it enters sleep mode. When processor enters sleep mode it resets the USB registers and does not respond to its host-assigned address.

12.5 UDC Operation

When a USB interrupt is received, software is directed to the USB ISR. USIR is level sensitive. Be sure to clear USIR as the last step before exiting the ISR. Upon exiting the ISR, the user should always clear the interrupt source bit, then clear the USIR. On power up or after a reset, software initially enables only EP0 interrupt. The other interrupts are enabled as required by the SET CONFIG command.

12.5.1 Case 1: EP0 Control Read

- 1. When software starts, it initializes a software state machine to EP0_IDLE. The software state machine is used to track endpoints stages when software communicates with the host PC.
- 2. The host PC sends a SETUP command.
- 3. UDC generates an EP0 Interrupt.
- 4. Software determines that the UDCCS0[SA] and UDCCS0[OPR] bits are set. This indicates that a new OUT packet is in the EP0 Buffer and identifies a SETUP transaction.
- 5. Software reads the data into a buffer while UDCCS0[RNE] bit (receiver not empty) is set.
- 6. Software parses the command in the buffer and determines that it is a Control Read.
- 7. Software starts to load the UDDR0 register FIFO with the first data packet and sets the internal state machine to EP0_IN_DATA_PHASE.
- 8. After it reads and parses the data, software clears the UDCCS0[SA] and the UDCCS0[OPR] bits and sets the UDCCS0[IPR] bit, if transmitting less than MAX_PACKET bytes, which prompts the UDC to transmit the data on the next IN. The UDC sends NAKs to all requests on this EP until the UDCCS0[IPR] bit is set.
- 9. Software clears the UDC interrupt bit and returns from the interrupt service routine.
- 10. The host PC issues an IN packet, which the UDC sends data back to the host. After the host PC sends an ACK to the UDC, the UDC clears the UDDCS0[IPR] bit and generates an interrupt.
- 11. Software enters the ISR routine and examines its internal state machine. It determines that it is in the EPO_IN_DATA_PHASE state and must transmit more data. Software loads the next amount of data, sets the UDCCS0[IPR] bit if necessary, and returns from the interrupt. The internal state machine is not affected.
- 12. Repeat Steps 10 and 11 until all the data is transmitted or the last data packet is a short packet.
- 13. If the last packet software sends is a short packet, it sets its internal state machine to EP0_END_XFER. If the last data packet ends on a 16-byte boundary, software sets UDCCS0[IPR] to send a zero-length packet without loading data in the FIFO. After it sends the zero-length packet, software sets the internal state machine to EP0_END_XFER.



- 14. When the host executes the STATUS OUT stage (zero-length OUT), the UDC sets the UDDCS0[OPR] bit, which causes an interrupt.
- 15. Software enters the ISR routine and determines that the UDCCS0[OPR] bit is set, the UDCCS0[SA] bit is clear, and its internal state machine is EP0_END_XFER. Software clears the UDCCS0[OPR] bit and transfers its internal state machine to EP0_IDLE.
- 16. Software clears the UDC interrupt bit and returns from the interrupt service routine.

If the host sends another SETUP command during these steps, the software must terminate the first SETUP command and start the new command.

12.5.2 Case 2: EP0 Control Read with a Premature Status Stage

Case 2 occurs during an enumeration cycle when the host PC sends a premature status stage during a Get Device Descriptor command.

- 1. When software starts, it initializes a software state machine to EP0_IDLE. The software state machine is used to track endpoints stages when software communicates with the host PC.
- 2. The host PC sends a SETUP command.
- 3. UDC generates an EP0 Interrupt.
- 4. Software determines that the UDCCS0[SA] and UDCCS0[OPR] bits are set. This indicates that a new OUT packet is in the EP0 Buffer and identifies a SETUP transaction.
- 5. Software reads the data into a buffer while UDCCS0[RNE] bit (receiver not empty) is set.
- 6. Software parses the command in the buffer and determines that it is a Control Read.
- Software starts to load the UDDR0 register FIFO with the first data packet and sets the internal state machine to EP0 IN DATA PHASE.
- 8. After it reads and parses the data, software clears the UDCCS0[SA] and the UDCCS0[OPR] bits and sets the UDCCS0[IPR] bit, if transmitting less than MAX_PACKET bytes, which prompts the UDC to transmit the data on the next IN. The UDC sends NAKs to all requests on this EP until the UDCCS0[IPR] bit is set.
- 9. Software clears the UDC interrupt bit and returns from the interrupt service routine.
- 10. The host PC issues an IN packet, which the UDC sends back to the host. After the host PC sends an ACK to the UDC, the UDC clears the UDDCS0[IPR] bit and generates an interrupt.
- 11. Software enters the ISR routine and examines its internal state machine. It determines that it is in the EPO_IN_DATA_PHASE state and must transmit more data. Software loads the next amount of data, sets the UDCCS0[IPR] bit, if transmitting less than MAX_PACKET bytes, and returns from the interrupt. The internal state machine is not affected.
- 12. Repeat Steps 10 and 11 until all the data is transmitted or the last data packet is a short packet.
- 13. As Steps 10 and 11 are repeated, the host sends a premature STATUS OUT stage, which indicates that the host PC can not accept more data, instead of an IN packet.
- 14. When the EP0 interrupt occurs, software determines that the UDCCS0[OPR] bit is set, the UDCCS0[SA] bit is cleared, and its machine state is EP0_IN_DATA_PHASE. This indicates that a premature STATUS OUT occurred.
- 15. Software clears the UDCCS0[OPR] bit and changes the pin's state to EP0_IDLE. The software writes to the UDCCS0[FTF] bit to clean up any buffer pointers and empty the transmit FIFO.



16. Software clears the UDC interrupt bit and returns from the interrupt service routine.

If the host sends another SETUP command during these steps, the software must terminate the first SETUP command and start the new command.

12.5.3 Case 3: EP0 Control Write With or Without a Premature Status Stage

- 1. When software starts, it initializes a software state machine to EP0_IDLE. The software state machine is used to track stages when software communicates with the host PC.
- 2. The host PC sends a SETUP command.
- 3. UDC generates an EP0 Interrupt.
- 4. Software determines that the UDCCS0[SA] and UDCCS0[OPR] bits are set. This indicates that a new OUT packet is in the EP0 Buffer and identifies a SETUP transaction.
- 5. Software reads the data into a buffer while UDCCS0[RNE] bit (receiver not empty) is set.
- 6. Software parses the command in the buffer and determines that it is a Control Write (such as Set Descriptor).
- 7. Software sets the internal to EP0_OUT_DATA_PHASE and clears the UDCCS0[OPR] and UDCCS0[SA] bits.
- 8. To allow a premature STATUS IN stage, software sets the UDCCS0[IPR] bit and loads a zero-length packet in the transmit FIFO.
- 9. Software clears the UDC interrupt bit and returns from the interrupt service routine.
- 10. The host PC issues an OUT packet and the UDC issues an EP0 interrupt.
- 11. Software enters the ISR routine and determines that it is in the EPO_OUT_DATA_PHASE state, the UDCCS0[OPR] bit is set, and the UDCCS0[SA] bit is clear. This indicates that there is more data to receive.
- 12. Software reads the data into a buffer while UDCCS0[RNE] bit is set and clears the UDDCCS0[OPR] bit.
- 13. Software sets the UDCCS0[IPR] bit to allow a premature STATUS IN stage.
- 14. Software clears the UDC interrupt bit and returns from the interrupt service routine.
- 15. Steps 11 through 14 are repeated until all of the data is received.
- 16. As Steps 11 through 14 are repeated, the host sends a STATUS IN stage, which indicates that the host PC can not send more data, instead of an OUT packet. The STATUS IN stage may be premature or not.
- 17. Because software loaded a zero-length packet (see Step 8), the UDC responds to the STATUS IN by sending a a zero-length packet back to the host PC. This causes an interrupt.
- 18. Software enters the ISR routine and determines that it is in the EP0_OUT_DATA_PHASE state and the UDCCS0[OPR] and UDCCS0[IPR] bits are clear. This indicates that a STATUS IN stage occurred.
- 19. Software determines how many bytes were received before the interrupt and compares the number of received bytes to the wLength field in the original SETUP packet. If the correct amount of data was sent, software parses the data and performs the action the data indicates. If



the wrong amount of data was sent, software cleans up any buffer pointers and disregards the received data.

- 20. Software changes its internal state machine to EP0_IDLE.
- 21. Software clears the UDC interrupt bit and returns from the interrupt service routine.

If the host sends another SETUP command during these steps, the software must terminate the first SETUP command and start the new command.

12.5.4 Case 4: EP0 No Data Command

- 1. When software starts, it initializes a software state machine to EP0_IDLE. The software state machine is used to track stages when software communicates with the host PC.
- 2. The host PC sends a SETUP command.
- 3. UDC generates an EP0 Interrupt.
- 4. Software determines that the UDCCS0[SA] and UDCCS0[OPR] bits are set. This indicates that a new OUT packet is in the EP0 Buffer and identifies a SETUP transaction.
- 5. Software reads the data into a buffer while UDCCS0[RNE] bit (receiver not empty) is set.
- 6. Software parses the data in the buffer and determines that it is a No Data command.
- 7. Software executes the command and sets its internal state machine to EP0_IDLE. Software clears the UDCCS0[IPR] and UDCCS0[SA] bits. If the command is a No-Data-Phase Standard command, then do not set the UDCCS0[IPR] bit. If the command is not a No-Data-Phase Standard command, e.g a No-Data-Phase Vendor command or No-Data-Phase Class command, then software must set the UDCCS0[IPR] bit.
- 8. When the host PC executes the STATUS IN stage, the UDC sends back a zero-length packet, which indicates a successful handshake. This does not cause an interrupt.

If the host sends another SETUP command during these steps, the software must terminate the first SETUP command and start the new command.

12.5.5 Case 5: EP1 Data Transmit (BULK-IN)

The procedure in Case 5 can also be used to operate Endpoints 6 and 11.

In Case 5, the Transmit Short Packet is only set if a packet size of less than 64 bytes is sent. If the packet size is 64 bytes, the system arms when the 64th byte is loaded. Loading the 64th byte and setting the UDCCS1[TSP] bit produces one 64-byte packet and one zero-length packet.

When software receives a SETUP VENDOR command to set up an EP1 BULK IN transaction, it may take one of two courses of action, as appropriate for the chosen operating model:

- Configure the DMA engine and disable the EP1 interrupt to allow the DMA engine to handle the transaction.
- Enable the EP1 interrupt to allow the Megacell to directly handle the transaction.

12.5.5.1 Software Enables the DMA

If software enables the DMA engine, use the following steps:



- 1. During the SETUP VENDOR command, software enables the DMA engine and masks the EP1 interrupt. The DMA start address must be aligned on a 16-byte boundary.
 - a. If the packet size is 64 bytes, software transfers the all the data in one DMA descriptor and sets the UDCCS1[TSP] bit in the second DMA descriptor.
 - b. If the packet size is less than 64 bytes, software sets up a string of descriptors in which the odd numbered descriptors point to the data and the even numbered descriptors are writes to the UDCCS1[TSP] bit.
- 2. The host PC sends a BULK-IN and the UDC sends a data packet back to the host PC.
- 3. The UDC generates an interrupt that is masked from the Megacell.
- 4. The DMA engine fills the EP1 data FIFO (UDDR1) with data and sets the UDCCS1[TSP] bit if the data packet is a short packet.
- 5. Steps 2 through 4 repeat until all the bulk data is sent to the host PC.

12.5.5.2 Software Enables the EP1 Interrupt

If software enables the EP1 interrupt to allow the Megacell to directly handle the transaction:

- 1. During the SETUP VENDOR command, software fills the EP1 data FIFO (UDDR1) with data and clears the UDCCS1[TPC] bit. If the data packet is a short packet, software also sets the UDCCS1[TSP] bit.
- 2. The host PC sends a BULK-IN and the UDC sends a data packet back to the host PC and generates an EP1 Interrupt.
- 3. Software fills the EP1 data FIFO (UDDR1) with data and clears the UDCCS1[TPC] bit. If the data packet is a short packet, software also sets the UDCCS1[TSP] bit.
- 4. Return from interrupt.
- 5. Steps 2 through 4 repeat until all of the data is sent to the host PC.

12.5.6 Case 6: EP2 Data Receive (BULK-OUT)

This procedure can also be used to operate Endpoints 7 and 12.

When software receives a SETUP VENDOR command to set up an EP2 BULK OUT transaction, it may take one of two courses of action, as appropriate for the chosen operating model:

- Enable the DMA engine to handle the transaction.
- Allow the Megacell to directly handle the transaction.

12.5.6.1 Software Enables the DMA

If software enables the DMA engine to handle the transaction:

- 1. During the SETUP VENDOR command, software sets up the DMA engine and sets the UDCCS2[DME] bit.
 - a. If the packet size is 32 or 64 bytes, software sets up a string of descriptors, each with a length of modulo 32 or 64. Software sets the interrupt bit for the appropriate descriptor.
 - b. If the packet size is less than 32 bytes, software uses interrupt mode.



- 2. The host PC sends a BULK-OUT.
- 3. The DMA engine reads data from the EP2 data FIFO (UDDR2).
- 4. Steps 2 and 3 repeat until all the data has been read from the host.
- 5. If the software receives an EP2 interrupt it completes the following process:
 - a. If UDCCS2[RNE] is clear and UDCCS2[RSP] is set, the data packet was a zero-length packet.
 - b. If UDCCS2[RNE] is set, the data packet was a short packet and software must use the UDCWC2 count register to read the proper amount of data from the EP2 data FIFO (UDDR2).
 - c. Software clears the UDCCS2[RPC] bit.
- 6. Return from interrupt.

12.5.6.2 Software Allows the Megacell to Handle the Transaction

If software allows the Megacell to handle the transaction:

- 1. During the SETUP VENDOR command, software clears the UDCCS2[DME] bit.
- 2. The host PC sends a BULK-OUT and the UDC generates an EP2 Interrupt.
- 3. If UDCCS2[RNE] is clear and UDCCS2[RSP] is set, the data packet was a zero-length packet.
- 4. If UDCCS2[RNE] is set, software uses the UDCWC2 count register to read the proper amount of data from the EP2 data FIFO (UDDR2).
- 5. Software clears the UDCCS2[RPC] bit.
- 6. Return from interrupt.
- 7. Steps 2 through 6 repeat until all the data has been read from the host.

12.5.7 Case 7: EP3 Data Transmit (ISOCHRONOUS-IN)

The procedure in Case 7 can also be used to operate Endpoints 8 and 13.

In Case 7, the Transmit Short Packet is only set if a packet size of less than 256 bytes is sent. If the packet size is 256 bytes, the system arms when the 256th byte is loaded. Loading the 256th byte and setting the UDCCS3[TSP] bit produces one 256-byte packet and one zero-length packet.

When software receives a SETUP VENDOR command to set up an EP3 ISOCHRONOUS IN transaction, it may take one of three courses of action, as appropriate for the chosen operating model:

- Configure the DMA engine and disable the EP3 interrupt to allow the DMA engine to handle the transaction.
- Enable the EP3 interrupt to allow the Megacell to directly handle the transaction.
- Enable the SOF interrupt to handle the transaction on a frame count basis.

12.5.7.1 Software Enables DMA

If software enables the DMA engine to handle the transaction:



- 1. During the SETUP VENDOR command, software enables the DMA engine and masks the EP3 interrupt. The DMA start address must be aligned on a 16-byte boundary.
 - a. If the packet size is 256 bytes, software transfers the all the data in one DMA descriptor.
 - b. If the packet size is less than 256 bytes, software sets up a string of descriptors in which the odd numbered descriptors point to the data and the even numbered descriptors are writes to the UDCCS1[TSP] bit.
- 2. The host PC sends an ISOC-IN and the UDC sends a data packet back to the host PC.
- 3. The UDC generates an interrupt that is masked from the Megacell.
- 4. The DMA engine fills the EP3 data FIFO (UDDR3) with data and sets the UDCCS3[TSP] bit if the data packet is a short packet.
- 5. Steps 2 through 4 repeat until all the data has been sent to the host.

12.5.7.2 Software Enables the EP3 Interrupt

If software enables the EP3 interrupt to allow the Megacell to directly handle the transaction:

- 1. During the SETUP VENDOR command, software fills the EP3 data FIFO (UDDR3) with data and clears the UDCCS3[TPC] bit. If the data packet is a short packet, software also sets the UDCCS3[TSP] bit.
- 2. The host PC sends a ISOC-IN command and the UDC sends a data packet back to the host PC and generates an EP3 Interrupt.
- 3. Software fills the EP3 data FIFO (UDDR3) with data and clears the UDCCS3[TPC] bit. If the data packet is a short packet, software also sets the UDCCS3[TSP] bit.
- 4. Return from interrupt.
- 5. Steps 2 through 4 repeat until all of the data is sent to the host PC.

12.5.7.3 Software Enables the SOF Interrupt

If software enables the SOF interrupt to handle the transaction on a frame count basis:

- 1. Software disables the UDCCS3 Interrupt by setting UICR0[IM3] to a 1 and enables the SOF interrupt in the UFNHR register by setting UFNHR[SIM] to a 0.
- 2. When the host PC sends an SOF, the UDC sets the UFNHR[SIR] bit, which causes an SOF interrupt.
- 3. Software checks the UDCCS3[TFS] bit to determine if there is room for a data packet. If there is room, software fills the EP3 data FIFO (UDDR3) with data and clears the UDCCS3[TPC] bit. If the data packet is a short packet, software sets the UDCCS3[TSP] bit.
- 4. Software clears the UFNHR[SIR] bit.
- 5. Return from interrupt.
- 6. Steps 2 through 5 repeat until all the data is sent to the host PC.

12.5.8 Case 8: EP4 Data Receive (ISOCHRONOUS-OUT)

The procedure in Case 8 can also be used to operate Endpoints 9 and 14.



When software receives a SETUP VENDOR command to set up an EP4 ISOCHRONOUS OUT transaction, it may take one of three courses of action, as appropriate for the chosen operating model:

- Configure the DMA engine and disable the EP4 interrupt to allow the DMA engine to handle the transaction.
- Enable the EP4 interrupt to allow the Megacell to directly handle the transaction.
- Enable the SOF interrupt to handle the transaction on a frame count basis.

12.5.8.1 Software Enables the DMA

If software enables the DMA engine, use the following steps:

- 1. During the SETUP VENDOR command, software enables the DMA engine and sets the UDCCS4[DME] bit. ISO packet sizes are not restricted, but a packet size of modulo 32 is highly recommended efficiency.
 - a. If the packet size is between 32 and 256 bytes modulo 32, software determines the number of descriptors needed and sets up a string of descriptors. Software sets the interrupt bit for the appropriate descriptor.
 - b. If the packet size is between 32 bytes and 256 bytes, but not modulo 32 bytes, software sets up a descriptor to receive each data packet, then reads the remaining data on each UDCCS2[RSP] bit interrupt and sets up another descriptor.
 - c. If the packet size is less than 32 bytes, software must use interrupt mode.
- 2. The host PC sends a ISOC-OUT.
- 3. The DMA engine reads the data from the EP4 data FIFO (UDDR4).
- 4. Steps 2 and 3 repeat until all the data has been read from the host.
- 5. If the software receives an EP4 interrupt it completes the following process:
 - a. If UDCCS4[RNE] is clear and UDCCS4[RSP] is set, the data packet was a zero-length packet.
 - b. If UDCCS4[RNE] is set, the data packet was a short packet and software uses the UDCWC4 count register to read the proper amount of data from the EP4 data FIFO (UDDR4).
 - c. Software clears the UDCCS4[RPC] bit.
- 6. Return from interrupt.

12.5.8.2 Software Allows the Megacell to Handle the Transaction

If software allows the Megacell to handle the transaction:

- 1. During the SETUP VENDOR command, software clears the UDCCS4[DME] bit.
- 2. The host PC sends a ISOC-OUT and the UDC generates an EP4 Interrupt.
- 3. If UDCCS4[RNE] is clear and UDCCS4[RSP] is set, the data packet was a zero-length packet.
- 4. If UDCCS4[RNE] is set, software uses the UDCWC4 count register to read the proper amount of data from the EP4 data FIFO (UDDR4).
- 5. Software clears the UDCCS4[RPC] bit.



- 6. Return from interrupt.
- 7. Steps 2 through 6 repeat until all the data has been read from the host.

12.5.8.3 Software Enables the SOF Interrupt

If software enables the SOF interrupt to handle the transaction on a frame count basis:

- 1. Software disables the UDCCS4 Interrupt by setting UICR0[IM4] to a 1 and enables the SOF interrupt in the UFNHR register by setting UFNHR[SIM] to a 0.
- 2. When the host PC sends an SOF, the UDC sets the UFNHR[SIR] bit, which causes an SOF interrupt.
- 3. If UDCCS4[RNE] is clear and UDCCS4[RSP] is clear, no data packet was received.
- 4. If UDCCS4[RNE] is clear and UDCCS4[RSP] is set, the data packet was a zero-length packet.
- 5. If UDCCS4[RNE] is set, the data packet was a short packet and software uses the UDCWC4 count register to read the proper amount of data from the EP4 data FIFO (UDDR4).
- 6. Software clears the UDCCS4[RPC] and UFNHR[SIR] bits.
- 7. Return from interrupt.
- 8. Steps 2 through 7 repeat until all the data is sent to the host PC.

12.5.9 Case 9: EP5 Data Transmit (INTERRUPT-IN)

The procedure in Case 9 can also be used to operate Endpoints 10 and 15.

In Case 9, the Transmit Short Packet is only set if a packet size of less than 8 bytes is sent. If the packet size is 8 bytes, the system arms when the 8th byte is loaded. Loading the 8th byte and setting the UDCCS5[TSP] bit produces one 8-byte packet and one zero-length packet.

When software receives a SETUP VENDOR command to set up an EP5 INTERRUPT-IN transaction, it can only allow the Megacell to handle the transaction:

- 1. During the SETUP VENDOR command, software fills the EP5 data FIFO (UDDR5) with data and clears the UDCCS5[TPC] bit.
- 2. The host PC sends an INTERRUPT-IN and the UDC generates an EP5 Interrupt.
- 3. Software fills the EP5 data FIFO (UDDR5) with data and clears the UDCCS5[TPC] bit. If the data packet is a short packet, software also sets the UDCCS5[TSP] bit.
- 4. Return from interrupt.
- 5. Steps 2 through 4 repeat until all the data is sent to the host PC.

12.5.10 Case 10: RESET Interrupt

- 1. After a system reset, software loads the registers with the required values.
- 2. Software enables the UDC by setting the UDCCR[UDE] bit and immediately reads the UDCCR[UDA] bit to determine if a USB reset is currently on the USB bus.
 - a. If UDCCR[UDA] is a 0, there is currently a USB reset on the bus and software clears the interrupt by writing a 1 to the UDCCR[RSTIR] bit. Software enables future reset interrupts by clearing the UDCCR[REM] bit.



- b. If UDCCR[UDA] is a 1, there is currently no USB reset on the bus and software enables future reset interrupts by clearing the UDCCR[REM] bit.
- 3. Return from interrupt.
- 4. The host either asserts a USB reset or negates a USB reset.
- 5. The UDC generates a Reset Interrupt.
- 6. Software determines that the UDCCR[RSTIR] bit is set and clears the interrupt by writing a 1 to the UDCCR[RSTIR] bit. Software then examines the UDCCR[UDA] bit to determine the type of reset that took place:
 - a. If UDCCR[UDA] is a 0, a Reset Assertion took place. Software returns from the interrupt and waits for the Reset Negation interrupt.
 - b. If UDDCR[UDA] is a 1, a Reset Negation took place. Software sets any initialization that is necessary.
- 7. Return from interrupt.

12.5.11 Case 11: SUSPEND Interrupt

- 1. As software starts, it clears the UDCCR[SRM] bit to allow a USB suspend interrupt.
- 2. The host PC asserts a USB suspend by stopping activity on the UDC+ and UDC- signals.
- 3. The UDC generates a Suspend Interrupt.
- 4. Software determines that the UDCCR[SUSIR] bit is set. This indicates that a USB suspend has occurred and software takes any necessary actions to turn off other peripherals, clean up internal buffers, perform power management, and perform similar functions. Software must not disable the UDC and must not allow the processor to go into sleep mode while the USB cable is attached.

12.5.12 Case 12: RESUME Interrupt

- 1. As software starts, it clears the UDCCR[SRM] bit to allow a USB resume.
- 2. The host PC asserts a USB resume by resuming activity after a suspend state on the UDC+ and UDC- signals.
- 3. The UDC generates a Resume Interrupt.
- 4. Software determines that the UDCCR[RESIR] bit is set. This indicates that a USB resume has occurred and the OS may take any necessary actions to turn on other peripherals, initialize internal buffers, perform power management, and perform similar functions.

12.6 UDC Register Definitions

All configuration, request/service, and status reporting is controlled by the USB host controller and is communicated to the UDC via the USB. The UDC has registers that control the interface between the UDC and the software. A control register enables the UDC and masks the interrupt sources in the UDC. A status register indicates the state of the interrupt sources. Each of the sixteen endpoints (control, OUT, and IN) have a control or status register. Endpoint 0 (control) has an



address for the 16 x 8 data FIFO that can be used to transmit and receive data. Endpoint 0 also has a write count register that is used to determine the number of bytes the USB host controller has sent to Endpoint 0.

12.6.1 UDC Control Register (UDCCR)

UDCCR, shown in Table 12-12, contains seven control bits: one to enable the UDC, one to show activity, and five to show status and associated control functions.

Table 12-12. UDCCR Bit Definitions

	0x 4060_00								0x 4060_0000										UDCCR										Read/Write and Read-Only									
Bit	31	30	29	28	2	7	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0					
												r	esei	ve	d											REM	RSTIR	SRM	SUSIR	RESIR	RSM	NDA	UDE					
	X	x	x	X		X	X	x	X	X	X	x	X	X	X	x	x	X	X	X	X	x	x	x	X	1	0	1	0	0	0	0	0					
		В	its				Naı	me												De	escr	ipti	on															
		31	1:8				_	_		res	erve	ed																										
			7				RE	М		0 =	Re	eset	rupt inte inte	rrup	ot er	nabl	ed.	ite)																				
		(6				RS	ΓIR					rupt vas						9 1 t	o cl	ear)																	
		ļ	5				SR	RM		0 =	Su	ispe	esum end/r end/r	esı	ıme	inte	rrup	t en	abl	ed.	te).																	
			4			;	SUS	SIR					terri																									
		;	3			ı	RES	SIR					terru ecei	•			`					,	ost.															
		:	2				RS	SM		0 =	Ma	ainta	sumo ain L UD(DC	su	spei	nd s	tate	et)																			
			1				UE	DΑ		0 =	UE	OC o	e (rea	ntly	/ red	ceivi						i.																
		(0				UE	ÞΕ		0 =	UE	OC d	e.(re disab enab	le.		te)																						

12.6.1.1 UDC Enable (UDE)

Enables the UDC. When UDE is set to a 1, the UDC is enabled for USB serial transmission or reception. When UDE is set to a 0, the UDC is disabled and the UDC+ and UDC- pins are tristated. This means that the UDC ignores all activity on the USB bus.

If UDE is set to a 0 the entire UDC design is reset. If the reset occurs while the UDC is actively transmitting or receiving data, it stops immediately and the remaining bits in the transmit or receive serial shifter are reset. All entries in the transmit and receive FIFO are also reset.



12.6.1.2 **UDC** Active (**UDA**)

This read-only bit can be read to determine if the UDC is currently active or in a USB reset. This bit is only valid when the UDC is enabled. A zero indicates that the UDC is currently receiving a USB reset from the host. A one indicates that the UDC is currently involved in a transaction.

12.6.1.3 UDC Resume (RSM)

When the UDC is in a suspend state, this bit can be written to force the UDC into a non-idle state (K state) for 3 ms to perform a remote wakeup operation. If the host PC does not start a wakeup sequence in 3 ms, the UDC returns to the suspend mode. This bit is a trigger bit for the UDC and is automatically cleared.

12.6.1.4 Resume Interrupt Request (RESIR)

The resume interrupt request bit is set if the SRM bit in the UDC control register is cleared, the UDC is currently in the suspended state, and the USB is driven with resume signalling.

12.6.1.5 Suspend Interrupt Request (SUSIR)

The suspend interrupt request register is set when the USB remains idle for more than 6 ms. The SUSIR bit retains state so software can determine that the USB is idle. If SRM is zero, SUSIR being set will not generate an interrupt but status continues to be updated.

12.6.1.6 Suspend/Resume Interrupt Mask (SRM)

This bit masks or enables the suspend interrupt request to the interrupt controller. When SRM is 1, the interrupt is masked and the setting of SUSIR will not generate an interrupt. When SRM is 0, the setting of SUSIR generates an interrupt when the USB is idle for more than 6ms. Programming SRM does not affect the state of SUSIR.

12.6.1.7 Reset Interrupt Request (RSTIR)

The reset interrupt request register is set when the host issues a reset. When the host issues a reset, the entire UDC is reset. The RSTIR bit retains its state so software can determine that the design was reset. If REM is zero, RSTIR being set does not generate an interrupt but status continues to be updated.

12.6.1.8 Reset Interrupt Mask (REM)

This bit masks or enables the reset interrupt request to the interrupt controller. When REM is 1, the interrupt is masked and the setting of RSTIR does not generate an interrupt. When REM is 0, the RSTIR setting generates an interrupt when the USB host controller issues an UDC reset. Programming REM does not affect the state of RSTIR.

The UDE bit is cleared to zero, which disables the UDC following a Megacell reset. Writes to reserved bits are ignored and reads return zeros.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



12.6.2 UDC Control Function Register (UDCCFR)

The UDC Control Function register (UDCCFR) contains 1 mode bit and 1 enable bit that lets software delay sending back an ACK response to SET_CONFIG or SET_INTERFACE commands from the host. The remaining bits are reserved and should not be written.

Table 12-13. UDC Control Function Register

		0h 4	4060 0008		UDCCF	R	Read/V	/rite and Rea	d-Only				
Bit	31:8	7	6	5	4	3	2	1	0				
	Reserved	AREN	Reserved	Reserved	Reserved	Reserved	ACM	UDA	UDE				
Reset	X	1	1	0 1 1 1 1									
		Bits	Name			Desci	ription						
		31:8	_	Reserved -	Read as unl	known and m	ust be writte	n as zero.					
		7	AREN	0 - Send SET_ 1 - Send	ONSE ENAE NAK respon INTERFACE ACK respon INTERFACE	se to SET_C commands se to SET_C	ONFIGURAT						
		6	_	Reserved -	Read as unl	known and m	ust be writte	n as zero.					
		5	_	Reserved -	Read as unl	known and m	ust be writte	n as zero.					
		4	_	Reserved -	Read as unl	known and m	ust be writte	n as zero.					
		3	_	Reserved -	Read as unl	known and m	ust be writte	n as zero.					
		2	ACM	0 - Send SET_ 1 - Send	ROL MODE ACK respon- INTERFACE NAK respon- INTERFACE	se to SET_C commands v se to SET_C	ONFIGURAT with no user on ONFIGURAT	intervention TON and					
		1	_	1 – Rese	rved – Read	as unknown	and must be	written as zer	O.				
		0	_	1 – Rese	rved – Read	as unknown	and must be	written as zer	O.				

12.6.2.1 ACK Control Mode

The ACK control mode (ACM) bit enables user control of the ACK response to the status IN requests of SET_CONFIGURATION and SET_INTERFACE commands. When ACM is set to 0, the UDC automatically responds to the STATUS IN request following a SET_CONFIGURATION and SET_INTERFACE with an ACK. When ACM is set to 1, the UDC responds to the STATUS IN request following a SET_CONFIGURATION and SET_INTERFACE command with a NAK until AREN is set to 1. When the user sets AREN to 1, the UDC responds with an ACK to the next STATUS IN request.

12.6.2.2 ACK Response Enable

When ACM = 1, the ACK response enable (AREN) bit enables user control of the ACK response to the status IN requests of SET_CONFIGURATION and SET_INTERFACE commands. When ACM is set to 1, the UDC responds to the STATUS IN request following a SET_CONFIGURAION and SET_INTERFACE command with a NAK until AREN is set to 1. When the user sets AREN to 1, the UDC responds with an ACK to the next STATUS IN request. AREN is cleared by the UDC when another SETUP command is received.



12.6.3 UDC Endpoint 0 Control/Status Register (UDCCS0)

UDCCS0, shown in Table 12-14, contains seven bits that are used to operate endpoint zero, the control endpoint.

Table 12-14. UDCCS0 Bit Definitions

		0x 4060_0010									UDCCS0												US	SBI	Devi	ice (Con	trol	ler		USB Device Controller									
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	7 16	15	14	4 13	12	11	10	9	8	7	6	5	4	3	2	1										
										re	se	rve	d											SA	RNE	FST	SST	DRWF	FTF	IPR										
x	x	X	X	X	x	X	x	x	x	x	X	X	X	х	x	x	Х	x	X	x	x	X	x	0	0	0	0	0	0	0										
	В	its			Na	me												D	esc	ripti	on																			
		0			OF	PR			•	acke JT p		•	`		write	1 to	o c	lear)																						
		1			ΙP	PR				cet re		, ,		ıys	read	0/w	rite	e 1 to	set).																				
		2			F٦	ΓF									ead (of Tx			1 to s	et)																					
		3			DR'	WF		0 =	: De	evice	Re	emo	te V	Nal	keup	Fea	tur	d-only re is o	, disa																					
		4			SS	ST									clea Ishak																_									
		5			FS	ST				stall (o se ke	t).																								
		6			RI	NE		0 =	Re	eceiv	e F	FIFC	en (npt	(rea y. mpty		nly)																							
		7			S	Α									1 to			e US	В																					
	-							1																							_									

12.6.3.1 OUT Packet Ready (OPR)

reserved

31:8

The OUT packet ready bit is set by the UDC when it receives a valid OUT packet to endpoint zero. When this bit is set, the USIR0[IR0] bit will be set in the UDC status/interrupt register if endpoint zero interrupts are enabled. This bit is cleared by writing a one. The UDC is not allowed to enter the data phase of a transaction until this bit is cleared.

12.6.3.2 IN Packet Ready (IPR)

The IN packet ready bit is set by the core if less than max_packet bytes (16) have been written to the endpoint 0 FIFO to be transmitted. The core must not set this bit if a max_packet is to be transmitted. The UDC clears this bit when the packet has been successfully transmitted, the UDCCS0[FTF] bit has been set, or a control OUT is received. When this bit is cleared due to a successful IN transmission or the reception of a control OUT, the USIR0[IR0] bit in the UDC interrupt register is set if the endpoint 0 interrupt is enabled via UICR0[IM0]. The Intel® XScaleTM microarchitecture is not able to clear UDCCS0[IPR] and always reads back a zero



When software enables the status stage for Vendor/Class commands and control data commands such as GET_DESCRIPTOR, GET_CONFIGURATION, GET_INTERFACE, GET_STATUS, and SET_DECSCRIPTOR, software must also set IPR. The data in the Transmit FIFO must be transmitted and the interrupt must be processed before the IPR is set for the status stage.

The status stage for all other USB Standard Commands that do not have a data stage, such as SET_ADDRESS, SET_CONFIGURATION, SET_INTERFACE, SET_FEATURE, and CLEAR FEATURE, is handled by the UDC and the software must not set IPR.

12.6.3.3 Flush Tx FIFO (FTF)

The Flush Tx FIFO bit triggers the reset of the endpoint 0 transmit FIFO. It is set when software writing a one or when the UDC receives an OUT packet from the host on endpoint 0. This bit always reads back a zero value.

12.6.3.4 Device Remote Wakeup Feature (DRWF)

The host indicates the state of the device remote wakeup feature by sending a Set Feature command or a Clear Function command. The UDC decodes the command sent by the host and sets this bit to a 1 if the feature is enabled and a 0 if the feature is disabled. This bit is read-only.

12.6.3.5 **Sent Stall (SST)**

The sent stall bit is set by the UDC when FST successfully forces a software-induced STALL on the USB bus. This bit is not set if the UDC detects a protocol violation from the host when a STALL handshake is returned automatically. In this event, there is no intervention by the core and the UDC clears the STALL status before the host sends the next SETUP command. When the UDC sets this bit, the transmit FIFO is flushed. The core writes a one to this bit to clear it.

12.6.3.6 Force Stall (FST)

The force stall bit can be set by the core to force the UDC to issue a STALL handshake. The UDC issues a STALL handshake for the current setup control transfer and the bit is cleared by the UDC because endpoint zero can not remain in a stalled condition.

12.6.3.7 Receive FIFO Not Empty (RNE)

The receive FIFO not empty bit indicates that the receive FIFO contains unread data. To determine if the FIFO has data in it, this bit must be read when the UDCCS0[OPR] bit is set. The receive FIFO must continue to be read until this bit clears or the data will be lost.

If UDCCS0[RNE] is not set when an interrupt generated by UDCCS0[OPR] is initially serviced, it indicates that a zero-length OUT packet was received.

12.6.3.8 Setup Active (SA)

The Setup Active bit indicates that the current packet in the FIFO is part of a USB setup command. This bit generates an interrupt and becomes active at the same time as UDCCS0[OPR]. Software must clear this bit by writing a 1 to it. Both UDCS0[OPR] and UDCCS0[SA] must be cleared.



Bit

12.6.4 UDC Endpoint x Control/Status Register (UDCCS1/6/11)

UDCCS1/6/11, shown in Table 12-15, contains 6 bits that are used to operate endpoint(x), a Bulk IN endpoint).

Table 12-15. UDCCS1/6/11 Bit Definitions

0x 4060_0014	UDCCS1	
0x 4060_0028	UDCCS6	USB Device Controller
0x 4060_003C	UDCCS11	

t	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d											dST	reserved	FST	SST	TUR	FTF	TPC	TFS
	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	Х	0	0	0	0	0	0	0	1

Bit	Name	Description
31:8	_	reserved
7	TSP	Transmit short packet 1 = Short packet ready for transmission.
6	_	reserved
5	FST	Force STALL 1 = Issue STALL handshakes to IN tokens.
4	SST	Sent STALL 1 = STALL handshake was sent.
3	TUR	Transmit FIFO underrun 1 = Transmit FIFO experienced an underrun.
2	FTF	Flush Tx FIFO 1 = Flush Contents of TX FIFO
1	TPC	Transmit packet complete 0 = Error/status bits invalid. 1 = Transmit packet has been sent and error/status bits are valid.
0	TFS	Transmit FIFO service 0 = Transmit FIFO has no room for new data 1 = Transmit FIFO has room for at least 1 complete data packet

12.6.4.1 Transmit FIFO Service (TFS)

The transmit FIFO service bit is active if one or fewer data packets remain in the transmit FIFO. TFS is cleared when two complete packets of data remain in the FIFO. A complete packet of data is signified by loading 64 bytes of data or by setting UDCCSx[TSP].

12.6.4.2 Transmit Packet Complete (TPC)

The transmit packet complete bit is set by the UDC when an entire packet is sent to the host. When this bit is set, the IRx bit in the appropriate UDC status/interrupt register is set if transmit interrupts are enabled. This bit can be used to validate the other status/error bits in the endpoint(x) control/status register. The UDCCSx[TPC] bit is cleared by writing a 1 to it. This clears the interrupt source for the IRx bit in the appropriate UDC status/interrupt register, but the IRx bit must also be cleared.



Setting this bit does not prevent the UDC from transmitting the next buffer. The UDC issues NAK handshakes to all IN tokens if this bit is set and neither buffer has been triggered by writing 64-bytes or setting UDCCSx[TSP].

When DMA loads the transmit buffers, the interrupt generated by UDCCSx[TPC] can be masked to allow data to be transmitted without core intervention.

12.6.4.3 Flush Tx FIFO (FTF)

The Flush Tx FIFO bit triggers a reset for the endpoint's transmit FIFO. The Flush Tx FIFO bit is set when software writes a 1 to it or when the host performs a SET_CONFIGURATION or SET INTERFACE. The bit's read value is zero.

12.6.4.4 Transmit Underrun (TUR)

The transmit underrun bit is set if the transmit FIFO experiences an underrun. When the UDC experiences an underrun, NAK handshakes are sent to the host. UDCCSx[TUR] does not generate an interrupt and is for status only. UDCCSx[TUR] is cleared by writing a 1 to it.

12.6.4.5 **Sent STALL (SST)**

The sent stall bit is set by the UDC in response to FST successfully forcing a user induced STALL on the USB bus. This bit is not set if the UDC detects a protocol violation from the host PC when a STALL handshake is returned automatically. In either event, the core does not intervene and the UDC clears the STALL status when the host sends a CLEAR_FEATURE command. The endpoint operation continues normally and does not send another STALL condition, even if the UDCCSx[SST] bit is set. To allow the software to continue to send the STALL condition on the USB bus, the UDCCSx[FST] bit must be set again. The core writes a 1 to the sent stall bit to clear it.

12.6.4.6 Force STALL (FST)

The core can set the force stall bit to force the UDC to issue a STALL handshake to all IN tokens. STALL handshakes continue to be sent until the core clears this bit by sending a Clear Feature command. The UDCCSx[SST] bit is set when the STALL state is actually entered, but this may be delayed if the UDC is active when the UDCCSx[FST] bit is set. The UDCCSx[FST] bit is automatically cleared when the UDCCSx[SST] bit is set. To ensure that no data is transmitted after the Clear Feature command is sent and the host resumes IN requests, software must clear the transmit FIFO by setting the UDCCSx[FTF] bit.

12.6.4.7 Bit 6 Reserved

Bit 6 is reserved for future use.

12.6.4.8 Transmit Short Packet (TSP)

The software uses the transmit short packet bit to indicate that the last byte of a data transfer to the FIFO has occurred. This indicates to the UDC that a short packet or zero-sized packet is ready to transmit. Software must not set this bit if a 64-byte packet is to be transmitted. When the data packet is successful transmitted, the UDC clears this bit.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



12.6.5 UDC Endpoint x Control/Status Register (UDCCS2/7/12)

UDCCS2/7/12, shown in Table 12-16, contains 7 bits that are used to operate endpoint x, a Bulk OUT endpoint.

Table 12-16. UDCCS2/7/12 Bit Definitions

0x 4060_0018	UDCCS2	
0x 4060_002C	UDCCS7	USB Device Controller
0x 4060 0040	UDCCS12	

Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

										re	sei	rved												RSP	RNE	FST	SST	DME	reserved	RPC	RFS
x)	(x	x	x	x	x	x	x	x	x	x	X	x	x	X	X	x	x	x	x	x	x	X	0	0	0	0	0	0	0	0
	•	Bit	•		Na	me		Description																							
	31:8				_				reserved																						
	7				R	SP		Receive short packet (read-only). 1 = Short packet received and ready for reading.																							
	6				RI	NE		Receive FIFO not empty (read-only). 0 = Receive FIFO empty. 1 = Receive FIFO not empty.																							
	5				F	ST		Force stall (read/write). 1 = Issue STALL handshakes to OUT tokens.																							
		4			S	ST						l/writ ndsh																			
	3				DI	ME		0 =	DMA Enable(read/write) 0 = Send data received interrupt after EOP received 1 = Send data received interrupt after EOP received and Receive FIFO has < 32 bytes of data																						
		2			-	_		rese	erve	t																					
		1			RI	PC		Receive packet complete (read/write 1 to clear). 0 = Error/status bits invalid. 1 = Receive packet has been received and error/status bits are valid.																							
		0			R	FS		Receive FIFO service (read-only). 0 = Receive FIFO has less than 1 data packet. 1 = Receive FIFO has 1 or more data packets.																							

12.6.5.1 Receive FIFO Service (RFS)

The receive FIFO service bit is set if the receive FIFO has one complete data packet in it and the packet has been error checked by the UDC. A complete packet may be 64 bytes, a short packet, or a zero packet. This bit is not cleared until all data has been read from both buffers.

12.6.5.2 Receive Packet Complete (RPC)

The receive packet complete bit is set by the UDC when an OUT packet is received. When this bit is set, the IRx bit in the appropriate UDC status/interrupt register is set, if receive interrupts are enabled. This bit must be used to validate the other status/error bits in the endpoint(x) control/status register. Status bits are not updated until RPC is set. Status bits stay set until RPC is cleared. The



exception is RNE which will get set with RPC but will clear itself once the active FIFO is empty. After clearing RPC, the next buffer will become active and the status bits will be updated accordingly, including RPC. The UDCCSx[RPC] bit is cleared by writing a 1 to it. The UDC issues NAK handshakes to all OUT tokens while this bit is set and both buffers have unread data.

12.6.5.3 Bit 2 Reserved

Bit 2 is reserved for future use.

12.6.5.4 DMA Enable (DME)

The dma enable is used by the UDC to control the timing of the data received interrupt. If the bit is set, the interrupt is asserted if the end of packet has been received and the receive FIFO has less than 32 bytes of data remaining in it. If the bit is not set, the interrupt is asserted when the end of packet is received and all of the received data is still in the receive FIFO.

12.6.5.5 **Sent Stall (SST)**

The sent stall bit is set by the UDC in response to FST successfully forcing a user induced STALL on the USB bus. This bit is not set if the UDC detects a protocol violation from the host PC when a STALL handshake is returned automatically. In either event, the core does not intervene and the UDC clears the STALL status when the host sends a CLEAR_FEATURE command. Any valid data in the FIFO remains valid and the software must unload it. The endpoint operation continues normally and does not send another STALL condition, even if the UDCCSx[SST] bit is set. To allow the software to continue to send the STALL condition on the USB bus, the UDCCSx[FST] bit must be set again. The core writes a 1 to the sent stall bit to clear it.

12.6.5.6 Force Stall (FST)

The core can set the force stall bit to force the UDC to issue a STALL handshake to all OUT tokens. STALL handshakes continue to be sent until the core clears this bit by sending a Clear Feature command. The UDCCSx[SST] bit is set when the STALL state is actually entered, but this may be delayed if the UDC is active when the UDCCSx[FST] bit is set. The UDCCSx[FST] bit is automatically cleared when the UDCCSx[SST] bit is set. To ensure that no data is transmitted after the Clear Feature command is sent and the host resumes IN requests, software must clear the transmit FIFO by setting the UDCCSx[FTF] bit.

12.6.5.7 Receive FIFO Not Empty (RNE)

The receive FIFO not empty bit indicates that unread data remains in the receive FIFO. This bit must be polled when the UDCCSx[RPC] bit is set to determine if there is any data in the FIFO that the DMA did not read. The receive FIFO must continue to be read until this bit clears or data will be lost.

12.6.5.8 Receive Short Packet (RSP)

The UDC uses the receive short packet bit to indicate that the received OUT packet in the active buffer currently being read is a short packet or zero-sized packet. This bit is updated by the UDC after the last byte is read from the active buffer and reflects the status of the new active buffer. If UDCCSx[RSP] is a one and UDCCSx[RNE] is a 0, it indicates a zero-length packet. If a zero-length packet is present, the core must not read the data register. UDCCSx[RSP] is cleared when the next OUT packet is received.



These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

12.6.6 UDC Endpoint x Control/Status Register (UDCCS3/8/13)

USCCS3/8/13, shown in Table 12-17, contains 4 bits that are used to operate endpoint(x), an Isochronous IN endpoint.

Table 12-17, UDCCS3/8/13 Bit Definitions

					0x4	4060 4060 4060	0_00	030							Ü	UDC	CS CS CC13	8						US	SB C)evi	ce (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												rese	rve	d											TSP	res	serv	/ed	TUR	Ħ	TPC	TFS
Reset	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	0	0	0	0	0	0	0	1
		В	it			Nai	me												De	escr	ipti	on										

Bit	Name	Description
31:8	_	reserved
7	TSP	Transmit short packet (read/write 1 to set). 1 = Short packet ready for transmission.
6:4	_	reserved
3	TUR	Transmit FIFO underrun (read/write 1 to clear) 1 = Transmit FIFO experienced an underrun.
2	FTF	Flush Tx FIFO (always read 0/ write a 1 to set) 1 = 1 - Flush Contents of TX FIFO
1	TPC	Transmit packet complete (read/write 1 to clear). 0 = Error/status bits invalid. 1 = Transmit packet has been sent and error/status bits are valid.
0	TFS	Transmit FIFO service (read-only). 0 = Transmit FIFO has no room for new data 1 = Transmit FIFO has room for at least 1 complete data packet

12.6.6.1 Transmit FIFO Service (TFS)

The transmit FIFO service bit is be set if one or fewer data packets remain in the transmit FIFO. UDCCSx[TFS] is cleared when two complete data packets are in the FIFO. A complete packet of data is signified by loading 256 bytes or by setting UDCCSx[TSP].

12.6.6.2 Transmit Packet Complete (TPC)

The the UDC sets transmit packet complete bit when an entire packet is sent to the host. When this bit is set, the IRx bit in the appropriate UDC status/interrupt register is set if transmit interrupts are enabled. This bit can be used to validate the other status/error bits in the endpoint(x) control/status register. The UDCCSx[TPC] bit gets cleared by writing a one to it. This clears the interrupt source for the IRx bit in the appropriate UDC status/interrupt register, but the IRx bit must also be cleared.

Setting this bit does not prevent the UDC from transmitting the next buffer. The UDC issues NAK handshakes to all IN tokens if this bit is set and neither buffer has been triggered by writing 64 bytes or setting UDCCSx[TSP].



When DMA is used to load the transmit buffers, the interrupt generated by UDCCSx[TPC] can be masked to allow data to be transmitted without core intervention.

12.6.6.3 Flush Tx FIFO (FTF)

The Flush Tx FIFO bit triggers a reset for the endpoint's transmit FIFO. The Flush Tx FIFO bit is set when software writes a 1 to it or when the host performs a SET_CONFIGURATION or SET_INTERFACE. The bit's read value is zero.

12.6.6.4 Transmit Underrun (TUR)

The transmit underrun bit is be set if the transmit FIFO experiences an underrun. When the UDC experiences an underrun, UDCCSx[TUR] generates an interrupt. UDCCSx[TUR] is cleared by writing a 1 to it.

12.6.6.5 Bits 6:4 Reserved

Bits 6:4 are reserved for future use.

12.6.6.6 Transmit Short Packet (TSP)

Software uses the transmit short packet to indicate that the last byte of a data transfer has been sent to the FIFO. This indicates to the UDC that a short packet or zero-sized packet is ready to transmit. Software must not set this bit if a packet of 256 bytes is to be transmitted. When the data packet is successfully transmitted, this bit is cleared by the UDC.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

12.6.7 UDC Endpoint x Control/Status Register (UDCCS4/9/14)

UDCCS4/9/14, shown in Table 12-18, contains six bits that are used to operate endpoint(x), an Isochronous OUT endpoint.



Table 12-18. UDCCS4/9/14 Bit Definitions

0x 4060_0020 **UDCCS4** 0x 4060_0034 **UDCCS9 USB Device Controller** 0x4060_0048 UDCCS14

Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

											r	ese	rve	d											RSP	RNE	reserved		DME	ROF	RPC	RFS
Reset	x	X	x	x	x	x	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	0	0	0	0	0	0	0	0

Bits	Name	Description
31:8	_	reserved
7	RSP	Receive short packet (read-only) 0 = Short packet received and ready for reading.
6	RNE	Receive FIFO not empty (read-only). 0 = Receive FIFO empty. 1 = Receive FIFO not empty.
5:4	_	reserved
3	DME	DMA Enable(read/write) 0 = Send receive interrupt after EOP receive 1 = Send data received interrupt after EOP received and Receive FIFO has < 32 bytes of data
2	ROF	Receive overflow (read/write 1 to clear) 1 = Isochronous data packets are being dropped from the host because the receiver is full.
1	RPC	Receive packet complete (read/write 1 to clear). 0 = Error/status bits invalid. 1 = Receive packet has been received and error/status bits are valid.
0	RFS	Receive FIFO service (read-only). 0 = Receive FIFO has less than 1 data packet. 1 = Receive FIFO has 1 or more data packets.

12.6.7.1 Receive FIFO Service (RFS)

The receive FIFO service bit is set if the receive FIFO has one complete data packet in it and the packet has been error checked by the UDC. A complete packet may be 256 bytes, a short packet, or a zero packet. UDCCSx[RFS] is not cleared until all data is read from both buffers.

12.6.7.2 **Receive Packet Complete (RPC)**

The receive packet complete bit gets set by the UDC when an OUT packet is received. When this bit is set, the IRx bit in the appropriate UDC status/interrupt register is set if receive interrupts are enabled. This bit can be used to validate the other status/error bits in the endpoint(x) control/status register. The UDCCSx[RPC] bit is cleared by writing a 1 to it.

12.6.7.3 Receive Overflow (ROF)

The receive overflow bit generates an interrupt on IRx in the appropriate UDC status/interrupt register to alert the software that Isochronous data packets are being dropped because neither FIFO buffer has room for them. This bit is cleared by writing a 1 to it.



12.6.7.4 DMA Enable (DME)

The DMA enable is used by the UDC to control the timing of the data received interrupt. If the bit is set, the interrupt is asserted when the end of packet is received and the receive FIFO has less than 32 bytes of data in it. If the bit is not set, the interrupt is asserted when the end of packet is received and all of the received data is still in the receive FIFO.

12.6.7.5 Bits 5:4 Reserved

Bits 5:4 are reserved for future use.

12.6.7.6 Receive FIFO Not Empty (RNE)

The receive FIFO not empty bit indicates that the receive FIFO has unread data in it. When the UDCCSx[RPC] bit is set, this bit must be read to determine if there is any data in the FIFO that DMA did not read. The receive FIFO must continue to be read until this bit clears or data will be lost.

12.6.7.7 Receive Short Packet (RSP)

The receive short packet bit is used by the UDC to indicate that the received OUT packet in the active buffer currently being read is a short packet or zero-sized packet. This bit is updated by the UDC after the last byte is read from the active buffer and reflects the status of the new active buffer. If UDCCSx[RSP] is a one and UDCCSx[RNE] is a zero, it indicates a zero-length packet. If a zero-length packet is present, the core must not read the data register. UDCCSx[RSP] clears when the next OUT packet is received.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

12.6.8 UDC Endpoint x Control/Status Register (UDCCS5/10/15)

UDCCS5/10/15, shown in Table 12-19 contains 6 bits that are used to operate endpoint(x), an Interrupt IN endpoint.

Table 12-19. UDCCS5/10/15 Bit Definitions (Sheet 1 of 2)

reserved

Та	ble	12	2-19). L	JD(CCS	35/	10/	15	Bit	De	fin	itio	ns	(S	hee	et 1	of	2)													
					0 x	406 406 406	0_0	038							ı	JDC JDC	cs	5						US	8B C	evi	ce (Con	troll	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	erve	d											TSP	reserved	FST	SST	TUR	FTF	TPC	TFS
Reset	x	X	X	X	X	X	X	X	X	X	x	x	X	X	x	x	x	X	X	X	x	X	x	X	0	0	0	0	0	0	0	1
•		В	its			Na	me												De	escr	ipti	on										
•		3′	1:8			_	_		res	erve	ed																					
•			7			TS	SP							ket reac).												

6



Table 12-19. UDCCS5/10/15 Bit Definitions (Sheet 2 of 2)

 0x 4060_0024
 UDCCS5

 0x 4060_0038
 UDCCS5
 USB Device Controller

 0x 4060_004C
 UDCCS15

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

reserved

reserved

Bits	Name	Description
5	FST	Force STALL (read/write). 1 = Issue STALL handshakes to IN tokens.
4	SST	Sent STALL (read/write 1 to clear). 1 = STALL handshake was sent.
3	TUR	Transmit FIFO underrun (read/write 1 to clear) 1 = Transmit FIFO experienced an underrun.
2	FTF	Flush Tx FIFO (always read 0/ write a 1 to set) 1 = Flush Contents of TX FIFO
1	TPC	Transmit packet complete (read/write 1 to clear). 0 = Error/status bits invalid. 1 = Transmit packet has been sent and error/status bits are valid.
0	TFS	Transmit FIFO service (read-only). 0 = Transmit FIFO has no room for new data 1 = Transmit FIFO has room for 1 complete data packet

12.6.8.1 Transmit FIFO Service (TFS)

The transmit FIFO service bit is set if the FIFO does not contain any data bytes and UDCCSx[TSP] is not set.

12.6.8.2 Transmit Packet Complete (TPC)

The transmit packet complete bit is be set by the UDC when an entire packet is sent to the host. When this bit is set, the IRx bit in the appropriate UDC status/interrupt register is set if transmit interrupts are enabled. This bit can be used to validate the other status/error bits in the endpoint(x) control/status register. The UDCCSx[TPC] bit is cleared by writing a 1 to it. This clears the interrupt source for the IRx bit in the appropriate UDC status/interrupt register, but the IRx bit must also be cleared.

The UDC issues NAK handshakes to all IN tokens if this bit is set and the buffer is not triggered by writing 8 bytes or setting UDCCSx[TSP].

12.6.8.3 Flush Tx FIFO (FTF)

The Flush Tx FIFO bit triggers a reset for the endpoint's transmit FIFO. The Flush Tx FIFO bit is set when software writes a 1 to it or when the host performs a SET_CONFIGURATION or SET_INTERFACE. The bit's read value is zero.



12.6.8.4 Transmit Underrun (TUR)

The transmit underrun bit is be set if the transmit FIFO experiences an underrun. When the UDC experiences an underrun, NAK handshakes are sent to the host. UDCCSx[TUR] does not generate an interrupt and is for status only. UDCCSx[TUR] is cleared by writing a 1 to it.

12.6.8.5 Sent STALL (SST)

The sent stall bit is set by the UDC in response to FST successfully forcing a user induced STALL on the USB bus. This bit is not set if the UDC detects a protocol violation from the host PC when a STALL handshake is returned automatically. In either event, the core does not intervene and the UDC clears the STALL status when the host sends a CLEAR_FEATURE command. The endpoint operation continues normally and does not send another STALL condition, even if the UDCCSx[SST] bit is set. To allow the software to continue to send the STALL condition on the USB bus, the UDCCSx[FST] bit must be set again. The core writes a 1 to the sent stall bit to clear it.

12.6.8.6 Force STALL (FST)

The core can set the force stall bit to force the UDC to issue a STALL handshake to all IN tokens. STALL handshakes continue to be sent until the core clears this bit by sending a Clear Feature command. The UDCCSx[SST] bit is set when the STALL state is actually entered, but this may be delayed if the UDC is active when the UDCCSx[FST] bit is set. The UDCCSx[FST] bit is automatically cleared when the UDCCSx[SST] bit is set. To ensure that no data is transmitted after the Clear Feature command is sent and the host resumes IN requests, software must clear the transmit FIFO by setting the UDCCSx[FTF] bit.

12.6.8.7 Bit 6 Reserved

Bit 6 is reserved for future use.

12.6.8.8 Transmit Short Packet (TSP)

Software uses the transmit short to indicate that the last byte of a data transfer has been sent to the FIFO. This indicates to the UDC that a short packet or zero-sized packet is ready to transmit. Software must not set this bit if a packet of 8 bytes is to be transmitted. When the data packet is successfully transmitted, the UDC clears this bit.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

12.6.9 UDC Interrupt Control Register 0 (UICR0)

UICR0, shown in Table 12-20, contains 8 control bits to enable/disable interrupt service requests from data endpoints 0 - 7. All of the UICR0 bits are reset to a 1 so interrupts are not generated on initial system reset.



Table 12-20. UICR0 Bit Definitions

					0x	406	0_0	050								UIC	R0							US	SB C	Devi	ce (Con	trol	er		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	1	1 10	9	8	7	6	5	4	3	2	1	0
											re	se	rve	d											M7	9WI	IM5	™	IM3	IM2	M	IM0
Reset	x	x	x	x	x	x	x	x	x	x	x	x	x	x	x	x	x	x	x	X	X	x	x	x	1	1	1	1	1	1	1	1
		Bi	ts			Na	me												De	SCI	ip	tion										
		31	:8			_	_		res	erve	ed																					
	Interrupt Mask for Endpoint 7 0 = Receive interrupt enabled 1 = Receive interrupt disabled Interrupt Mask for Endpoint 6 0 = Transmit interrupt enabled																															
	1 = Receive interrupt disabled Interrupt Mask for Endpoint 6 0 = Transmit interrupt enabled 1 = Transmit interrupt disabled Interrupt mask for Endpoint 5																															
	6 IM6 0 = Transmit interrupt enabled 1 = Transmit interrupt disabled Interrupt mask for Endpoint 5 5 IM5 0 = Transmit interrupt enabled 1 = Transmit interrupt disabled																															
															ena	bled																
		3	3			IN	13		0 =	Tra	ot ma ansn ansn	nit i	nte	rrupt	ena	able																
	1 = Transmit in Interrupt Mask f 0 = Receive int 1 = Receive int Interrupt Mask f												nter	rupt	ena	bled																
		1	l			IN	11		0 =	Tra	ot Ma ansn ansn	nit i	inter	rrupt	ena	able																
		()			IN	10		0 =	Er	ot ma ndpo ndpo	nt :	zero	inte	erru _l	ot er																

12.6.9.1 Interrupt Mask Endpoint x (IMx), Where x is 0 through 7

The UICR0[IMx] bit is used to mask or enable the corresponding endpoint interrupt request, USIR0[IRx]. When the mask bit is set, the interrupt is masked and the corresponding bit in the USIR0 register is not allowed to be set. When the mask bit is cleared and an interruptible condition occurs in the endpoint, the appropriate interrupt bit is set. Programming the mask bit to a 1 does not affect the current state of the interrupt bit. It only blocks future zero to one transitions of the interrupt bit.

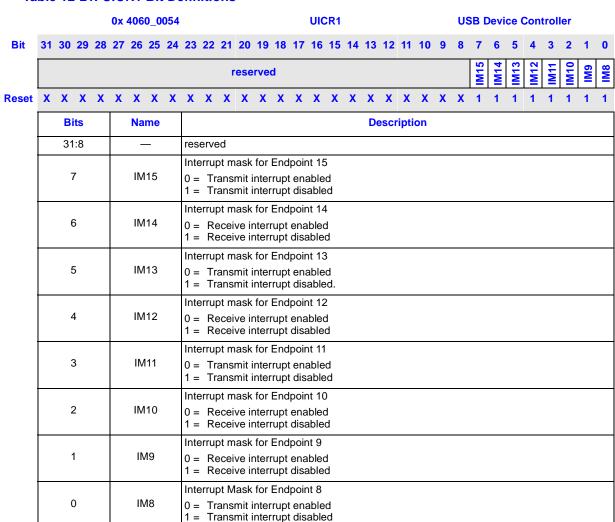
These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



12.6.10 UDC Interrupt Control Register 1 (UICR1)

UICR1, shown in Table 12-21, contains 8 control bits to enable/disable interrupt service requests from endpoints 8 - 15. The UICR1 bits are reset to 1 so interrupts are not generated on initial system reset.

Table 12-21. UICR1 Bit Definitions



12.6.10.1 Interrupt Mask Endpoint x (IMx), where x is 8 through 15.

The UICR1[IMx] bit is used to mask or enable the corresponding endpoint interrupt request, USIR1[IRx]. When the mask bit is set, the interrupt is masked and the corresponding bit in the USIR1 register is not allowed to be set. When the mask bit is cleared and an interruptible condition occurs in the endpoint, the appropriate interrupt bit is set. Programming the mask bit to a 1 does not affect the current state of the interrupt bit. It only blocks future zero to one transitions of the interrupt bit.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



12.6.11 UDC Status/Interrupt Register 0 (USIR0)

USIR0, shown in Table 12-22, and USIR1, shown in Table 12-23, contain bits that are used to generate the UDC's interrupt request. Each bit in the UDC status/interrupt registers is logically ORed together to produce one interrupt request. When the ISR for the UDC is executed, it must read the UDC status/interrupt register to determine why the interrupt occurred. USIRx is level sensitive. Be sure to clear USIRx as the last step before exiting the ISR.

The bits in USIR0 and USIR1 are controlled by a mask bit in the UDC Interrupt Control Register (UICR0/1). The mask bits, when set, prevent a status bit in the USIRx from being set. If the mask bit for a particular status bit is cleared and an interruptible condition occurs, the status bit is set. To clear status bits, the core must write a 1 to the position to be cleared. The interrupt request for the UDC remains active as long as the value of the USIRx is non-zero.

Table 12-22, USIR0 Bit Definitions

				0 x	406	0_0	058								US	IR0							US	SB C)evi	ce (Con	troll	er		
Bit	31	30 29	28	27	26	25	24	23	22	21	20	19	18	1	7 16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										r	ese	rve	d											IR7	IR6	IR5	IR4	IR3	IR2	IR1	<u>₹</u>
Reset	X	х х	X	X	X	x	X	X	X	X	X	X	X	X	СХ	X	x	x	X	x	X	x	x	0	0	0	0	0	0	0	0
		Bits			Na	me												De	esci	ipti	on										
	31:8 — reserved 7 IR7 Interrupt Request Endpoint 7 (read/write 1 to clear)																														
	7 IR7 Interrupt Request Endpoint 7 (read/write 1 to clear) 1 = Endpoint 7 needs service.																														
		5			IF	₹5									int 5 ervice		d/wı	rite '	1 to	clea	ar)										
		4			IF	R4									int 4 ervice		d/wı	rite '	1 to	clea	ar)										
		3			IF	3									int 3 ervice		d/wı	rite '	1 to	clea	ar)										
		2			IF	R2									int 2 ervice		d/wı	rite '	1 to	clea	ar)										
		1			IF	R1									int 1 ervice		d/wı	rite '	1 to	clea	ar)										
		0			IF	RO									int 0 ervice		d/wı	rite '	1 to	clea	ar)										

12.6.11.1 Endpoint 0 Interrupt Request (IR0)

The endpoint 0 interrupt request is set if the IM0 bit in the UDC control register is cleared and, in the UDC endpoint 0 control/status register, the OUT packet ready bit is set, the IN packet ready bit is cleared, or the sent STALL bit is set. The IR0 bit is cleared by writing a 1 to it.

12.6.11.2 Endpoint 1 Interrupt Request (IR1)

The interrupt request bit is set if the IM1 bit in the UDC interrupt control register is cleared and the IN packet complete (TPC) in UDC endpoint 1 control/status register is set. The IR1 bit is cleared by writing a 1 to it.



12.6.11.3 Endpoint 2 Interrupt Request (IR2)

The interrupt request bit is set if the IM2 bit in the UDC interrupt control register is cleared and the OUT packet ready bit (RPC) in the UDC endpoint 2 control/status register is set. The IR2 bit is cleared by writing a 1 to it.

12.6.11.4 Endpoint 3 Interrupt Request (IR3)

The interrupt request bit is set if the IM3 bit in the UDC interrupt control register is cleared and the IN packet complete (TPC) or Transmit Underrun (TUR) in UDC endpoint 3 control/status register is set. The IR3 bit is cleared by writing a 1 to it

12.6.11.5 Endpoint 4 Interrupt Request (IR4)

The interrupt request bit is set if the IM4 bit in the UDC interrupt control register is cleared and the OUT packet ready (RPC) or receiver overflow (ROF) in the UDC endpoint 4 control/status register or the Isochronous Error Endpoint 4 (IPE4) in the UFNHR are set. The IR4 bit is cleared by writing a 1 to it.

12.6.11.6 Endpoint 5 Interrupt Request (IR5)

The interrupt request bit is set if the IM5 bit in the UDC interrupt control register is cleared and the IN packet complete (TPC) in UDC endpoint 5 control/status register is set. The IR5 bit is cleared by writing a 1 to it.

12.6.11.7 Endpoint 6 Interrupt Request (IR6)

The interrupt request bit gets set if the IM6 bit in the UDC interrupt control register is cleared and the IN packet complete (TPC) in UDC endpoint 6 control/status register gets set. The IR6 bit is cleared by writing a one to it.

12.6.11.8 Endpoint 7 Interrupt Request (IR7)

The interrupt request bit is set if the IM7 bit in the UDC interrupt control register is cleared and the OUT packet ready bit (RPC) in the UDC endpoint 7 control/status register is set. The IR7 bit is cleared by writing a 1 to it.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



12.6.12 UDC Status/Interrupt Register 1 (USIR1)

Table 12-23. USIR1 Bit Definitions

					0x 4	4060	0_0	05C								US	IR1							US	SB C)evi	ce (Con	trol	er		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d											IR15	IR14	IR13	IR12	IR11	IR10	IR9	IR8
	X	X	x	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	x	X	X	x	x	X	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	scr	iptio	on										
		31	:8			-	-		res	erve	ed																					
		7	7			IR	15							End need				ad/w	rite	1 to	cle	ar)										
		6	6			IR	14							End need				ad/w	rite	1 to	cle	ar)										
		5	5			IR	13							End need				ad/w	rite	1 to	cle	ar)										
		2	1			IR	12							End need				ad/w	rite	1 to	cle	ar)										
		3	3			IR	11							End need				ad/w	rite	1 to	cle	ar)										
		2	2			IR	10							End				ad/w	rite	1 to	cle	ar)										
					_																											-

Interrupt Request Endpoint 9 (read/write 1 to clear)

Interrupt Request Endpoint 8 (read/write 1 to clear)

1 = Endpoint 9needs service.

1 = Endpoint 8 needs service.

12.6.12.1 Endpoint 8 Interrupt Request (IR8)

IR9

IR8

1

0

The interrupt request bit is set if the IM8 bit in the UDC interrupt control register is cleared and the IN packet complete (TPC) or Transmit Underrun (TUR) in UDC endpoint 8 control/status register is set. The IR8 bit is cleared by writing a 1 to it.

12.6.12.2 Endpoint 9 Interrupt Request (IR9)

The interrupt request bit is set if the IM9 bit in the UDC interrupt control register is cleared and the OUT packet ready (RPC) or receiver overflow (ROF) in the UDC endpoint 9 control/status register or the Isochronous Error Endpoint 9 (IPE9) in the UFNHR are set. The IR9 bit is cleared by writing a 1 to it.

12.6.12.3 Endpoint 10 Interrupt Request (IR10)

The interrupt request bit is set if the IM10 bit in the UDC interrupt control register is cleared and the IN packet complete (TPC) or in UDC endpoint 10 control/status register is set. The IR10 bit is cleared by writing a 1 to it.



12.6.12.4 Endpoint 11 Interrupt Request (IR11)

The interrupt request bit is set if the IM11 bit in the UDC interrupt control register is cleared and the IN packet complete (TPC) in UDC endpoint 11 control/status register is set. The IR11 bit is cleared by writing a 1 to it.

12.6.12.5 Endpoint 12 Interrupt Request (IR12)

The interrupt request bit is set if the IM12 bit in the UDC interrupt control register is cleared and the OUT packet ready bit (RPC) in the UDC endpoint 12 control/status register is set. The IR12 bit is cleared by writing a 1 to it.

12.6.12.6 Endpoint 13 Interrupt Request (IR13)

The interrupt request bit is set if the IM13 bit in the UDC interrupt control register is cleared and the IN packet complete (TPC) or Transmit Underrun (TUR) in UDC endpoint 13 control/status register is set. The IR13 bit is cleared by writing a 1 to it.

12.6.12.7 Endpoint 14 Interrupt Request (IR14)

The interrupt request bit is set if the IM14 bit in the UDC interrupt control register is cleared and the OUT packet ready (RPC) or receiver overflow (ROF) in the UDC endpoint 14 control/status register or the Isochronous Error Endpoint 14 (IPE14) in the UFNHR are set. The IR14 bit is cleared by writing a 1 to it.

12.6.12.8 Endpoint 15 Interrupt Request (IR15)

The interrupt request bit is set if the IM15 bit in the UDC interrupt control is set. The IR15 bit is cleared by writing a 1 to it.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

12.6.13 UDC Frame Number High Register (UFNHR)

UFNHR, shown in Table 12-24, holds the three most significant bits of the frame number contained in the last received SOF packet, the isochronous OUT endpoint error status, and the SOF interrupt status/interrupt mask bit.



Table 12-24. UFNHR Bit Definitions

				Leserved SiR SIR PE 4 IPE 9															er													
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d											SIR	SIM	IPE14	IPE9	IPE4	F No	3-Bi ram umb NSE	ie er
Reset	x	X	x	X	x	X	X	x	x	x	X	X	x	X	X	X	X	x	X	X	X	X	x	X	0	1	0	0	0	0	0	0
		Bi	ts	Name Description — reserved																												
		31	:8			_	_		res	erve	d																					
		7	·																													
		6	6			SI	M		0 =	SC)F i	upt i nter nter	rupt	ena																		
		5	5			IPE	14									Endata	•		•					,								
		4	ļ			IPI	E9									Endata									ted							
		3	3			IPI	E4									Endata	•		•					,	ted							
		2:	0		ı	FNN	/ISB					nbe ican			of '	11-b	it fra	me	nun	nber	ass	socia	ated	with	n las	t re	ceiv	e S0	OF.			

12.6.13.1 UDC Frame Number MSB (FNMSB)

The UFNHR[FNMSB] is the three most significant bits of the 11-bit frame number contained in the last received SOF packet. The remaining bits are located in the UFNLR. This information is used for isochronous transfers. These bits are updated every SOF.

12.6.13.2 Isochronous Packet Error Endpoint 4 (IPE4)

The isochronous packet error for Endpoint 4 is set if Endpoint 4 is loaded with a data packet that is corrupted. This status bit is used in the interrupt generation of endpoint 4. To maintain synchronization, the software must monitor this bit when it services an SOF interrupt and reads the frame number. This bit is not set if the token packet is corrupted or if the sync or PID fields of the data packet are corrupted.

12.6.13.3 Isochronous Packet Error Endpoint 9 (IPE9)

The isochronous packet error for Endpoint 9 is set if Endpoint 9 is loaded with a data packet that is corrupted. This status bit is used in the interrupt generation of endpoint 9. To maintain synchronization, software must monitor this bit when it services the SOF interrupt and reads the frame number. This bit is not set if the token packet is corrupted or if the sync or PID fields of the data packet are corrupted.



12.6.13.4 Isochronous Packet Error Endpoint 14 (IPE14)

The isochronous packet error for Endpoint 14 is set if Endpoint 14 is loaded with a data packet that is corrupted. This status bit is used in the interrupt generation of endpoint 14. To maintain synchronization, software must monitor this bit when it services the SOF interrupt and reads the frame number. This bit is not set if the token packet is corrupted or if the sync or PID fields of the data packet are corrupted.

12.6.13.5 Start of Frame Interrupt Mask (SIM)

The UFNHR[SIM] bit is used to mask or enable the SOF interrupt request. When UFNHR[SIM]=1, the interrupt is masked and the SIR bit is not allowed to be set. When UFNHR[SIM]=0, the interrupt is enabled and when an interruptible condition occurs in the receiver, the UFNHR[SIR] bit is set. Setting UFNHR[SIM] to a 1 does not affect the current state of UFNHR[SIR]. It only blocks future zero to one transitions of UFNHR[SIR].

12.6.13.6 Start of Frame Interrupt Request (SIR)

The interrupt request bit is set if the UFNHR[SIM] bit is cleared and an SOF packet is received. The UFNHR[SIR] bit is cleared by writing a 1 to it.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

12.6.14 UDC Frame Number Low Register (UFNLR)

UFNLR, shown in Table 12-25, is the eight least significant bits of the 11-bit frame number contained in the last received SOF packet. The three remaining bits are located in the UFNHR. This information is used for isochronous transfers. These bits are updated every SOF.

This is a read-only register. Ignore reads from reserved bits.

Table 12-25, UFNLR Bit Definitions

					0x	406	0_0	064								UFI	NLR							US	SBI	Devi	ice (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•		reserved 8-Bit Frame Numb															nbe	r LS	SB													
Reset	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	0	0	0	0	0	0	0	0
•		Ві	its			Na	me												De	escr	ipti	on										
,		31	:8			-	_		res	erve	ed																					
		7:	:0			FNI	LSB		Fra Lea					_	of f	ram	ie ni	ımb	er a	SSO	ciate	ed w	ith l	ast	rec	eive	d S0	OF.				

12.6.15 UDC Byte Count Register x (UBCR2/4/7/9/12/14)

UBCR2/4/7/9/12/14, shown in Table 12-26, maintains the remaining byte count in the active buffer of OUT endpoint(x).



Table 12-26. UBCR2/4/7/9/12/14 Bit Definitions

					0x 0x 0x 0x	406 406 406 406 406 406	0_0 0_0 0_0 0_0	06C 070 074 078							ı	UB UB UB JBC	CR2 CR4 CR7 CR9 CR1	2						US	SB [Devi	ice (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											ı	ese	rve	d														В	С			
Reset	x	X	x	x	X	X	x	X	x	x	x	x	x	x	x	x	x	x	x	x	x	x	x	x	0	0	0	0	0	0	0	0
		В	its			Na	me												De	escr	ipti	on										
		31	1:8			_	_		res	erve	ed																					
		7	:0			В	С		Bvt	e C	oun	t (re	ad-	only). N	umb	er c	f bv	tes	in th	ne F	IFO	is E	Svte	Cou	ınt r	olus	1 (E	3C+	1).		

12.6.15.1 Endpoint x Byte Count (BC)

The byte count is updated after each byte is read. When software receives an interrupt that indicates the endpoint has data, it can read the byte count register to determine the number of bytes that remain to be read. The number of bytes that remain in the input buffer is equal to the byte count +1.

This is a read-only register. Ignore reads from reserved bits.

12.6.16 UDC Endpoint 0 Data Register (UDDR0)

UDDR0, shown in, Table 12-27, is a 16-entry by 8-bit bidirectional FIFO. When the host transmits data to the UDC Endpoint 0, the core reads the UDC endpoint 0 register to access the data. When the UDC sends data to the host, the core writes the data to be sent in the UDC endpoint 0 register. The core can only read and write the FIFO at specific points in a control sequence. The direction that the FIFO flows is controlled by the UDC. Normally, the UDC is in an idle state, waiting for the host to send commands. When the host sends a command, the UDC fills the FIFO with the command from the host and the core reads the command from the FIFO when it arrives. The only time the core may write the endpoint 0 FIFO is after a valid command from the host is received and it requires a transmission in response.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 12-27. UDDR0 Bit Definitions

					0:	¢ 40	060_	800	0								UDI	OR0							USE	De	vice	e Co	ontr	olle	r	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d											Во	(fe	or R dpo	ead	ls) 1 0 Fl	int Γορ FO	of	
Reset	X	X	X	X	X	X	X	X	x	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	0	0	0	0	0	0	0	0
·		Ві	its			Na	me												De	scr	ipti	on										
•		31	:8			_	-		res	erve	d																					
·		7:	:0			DA	TΑ		Rea	ad –	Bot	ttom	of o	end	ooin	t 0 I	O da FIFO O da) da	ta.													

12.6.17 UDC Endpoint x Data Register (UDDR1/6/11)

UDDR1/6/11, shown in Table 12-28, is a double-buffered bulk IN endpoint that is 64 bytes deep. Data can be loaded via DMA or direct core writes. Because it is double buffered, up to two packets of data may be loaded for transmission.

These are write-only registers. Write zeros to reserved bits.

Table 12-28, UDDR1/6/11 Bit Definitions

					0x 4 0x 4	406	0_0	600								UDI	DR1 DR6 DR11							US	SB [Devi	ice (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•		reserved 8-Bit Data															ta															
Reset	X	x	X	X	X	X	X	x	X	x	x	X	X	X	X	x	X	X	X	X	x	X	X	x	X	x	0	0	0	0	0	0
•		В	its			Na	me												De	esci	ipti	on										
•		31	1:8			-			res	erve	ed																					
		7	:0			DA	λTΑ		Top	of	end	poir	nt da	ıta c	urre	ntly	beir	ng lo	oade	ed												

12.6.18 UDC Endpoint x Data Register (UDDR2/7/12)

UDDR2/7/12, shown in Table 12-29, is a double-buffered bulk OUT endpoint that is 64 bytes deep. The UDC will generate either an interrupt or DMA request as soon as the EOP is received. Since it is double buffered, up to two packets of data may be ready. Via DMA or by direct read from the core, the data can be removed from the UDC. If one packet is being removed and the packet behind it has already been received, the UDC will issue a NAK to the host the next time it sends an OUT packet to endpoint(x). This NAK condition will remain in place until a full packet space is available in the UDC at Endpoint(x).



These are read-only registers. Ignore reads from reserved bits.

Table 12-29, UDDR2/7/12 Bit Definitions

					0x	4060 4060 4060	0_0	680								UD	DR2 DR7 DR12							US	SB [Devi	ce	Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d													8	-bit	Dat	a		
Reset	X	X	X	X	X	x	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	0	0	0	0	0	0	0	0
		Ві	its			Na	me												De	escr	ipti	on										
		31	:8			-	-		res	erve	ed																					
		7	:0			DA	TA		Тор	of o	end	ooin	t da	ta c	urre	ently	bei	ng re	ead													

12.6.19 UDC Endpoint x Data Register (UDDR3/8/13)

UDDR3/8/13, shown in Table 12-30, is a double-buffered isochronous IN endpoint that is 256 bytes deep. Data can be loaded via DMA or direct core writes. Because it-is double buffered, up to two packets of data may be loaded for transmission.

These are write-only registers. Write zeros to reserved bits.

Table 12-30. UDDR3/8/13 Bit Definitions

					0x 4 0x 4	406	0_0	700								UDI UDI UDE	DR8							US	SB [Devi	ice (Con	itrol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d													8	-bit	Dat	а		
Reset	x	x	X	x	x	x	x	x	x	x	x	X	x	X	x	x	x	x	x	x	x	x	x	x	0	0	0	0	0	0	0	0
		Ві	its			Na	me												De	escr	ipti	on										
		31	:8			-	-		res	erve	ed																					
		7	:0			DA	TΑ		Top	of o	end	poin	t da	ta c	urre	ntly	bei	ng lo	oade	ed												

12.6.20 UDC Endpoint x Data Register (UDDR4/9/14)

UDDR4/9/14, shown in Table 12-31, is a double-buffered isochronous OUT endpoint that is 256 bytes deep. The UDC generates an interrupt or DMA request when the EOP is received. Because it is double-buffered, up to two packets of data may be ready. The data can be removed from the UDC via DMA or by a direct read from the Megacell. If one packet is being removed and the packet behind it has already been received, the UDC issues a NAK to the host the next time it sends an OUT packet to Endpoint(x). This NAK condition remains in place until a full packet space is available in the UDC at Endpoint(x).

These are read-only registers. Ignore reads from reserved bits.



Table 12-31. UDDR4/9/14 Bit Definitions

					0 x	406 406 406	0_0	900								UDI UDI	DR9)						US	B [Devi	ce (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•											r	ese	rve	d													8	-bit	Dat	a		
Reset	x	x	X	X	x	X	X	x	X	X	X	x	x	X	X	x	X	X	X	X	x	x	X	x	0	0	0	0	0	0	0	0
·		В	its			Na	me												De	escr	ipti	on										
,		31	1:8			-	_		res	erve	ed																					
•		7	:0			DA	TΑ		Top	of of	end	poir	nt da	ıta c	urre	ently	bei	ng lo	oade	ed												

12.6.21 UDC Endpoint x Data Register (UDDR5/10/15)

UDDR5/10/15, shown in Table 12-32, is an interrupt IN endpoint that is 8 bytes deep. Data must be loaded via direct Megacell writes. Because the USB system is a host initiator model, the host must poll Endpoint 5 to determine interrupt conditions. The UDC can not initiate the transaction.

Table 12-32. UDDR5/10/15 Bit Definitions

					0 x	4060 4060 4060	0_0	OC0							ı	UDI UDE	R1	0						US	SB I	Devi	ce (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•											r	ese	erve	d													8	-bit	Dat	a		
Reset	x	X	x	x	X	x	x	X	x	x	x	x	x	x	x	x	x	X	X	x	X	x	X	x	0	0	0	0	0	0	0	0
•		В	its			Na	me												De	escr	ipti	on										
•		31	1:8			_	-		res	erve	ed																					
•		7	:0			DA	TΑ		Top	of	end	poir	nt da	ıta c	urre	ently	bei	ng lo	oade	ed												

12.7 USB Device Controller Register Summary

Table 12-33 shows the registers associated with the UDC and the physical addresses used to access them.

Table 12-33. USB Device Controller Register Summary (Sheet 1 of 3)

Address	Name	Description
0x4060_0000	UDCCR	UDC Control Register
0x4060_0004	_	reserved for future use
0x4060_0008	UDCCFR	UDC Control Function Register
0x4060_000C	_	reserved for future use



Table 12-33. USB Device Controller Register Summary (Sheet 2 of 3)

Address	Name	Description
0x4060_0010	UDCCS0	UDC Endpoint 0 Control/Status Register
0x4060_0014	UDCCS1	UDC Endpoint 1 (IN) Control/Status Register
0x4060_0018	UDCCS2	UDC Endpoint 2 (OUT) Control/Status Register
0x4060_001C	UDCCS3	UDC Endpoint 3 (IN) Control/Status Register
0x4060_0020	UDCCS4	UDC Endpoint 4 (OUT) Control/Status Register
0x4060_0024	UDCCS5	UDC Endpoint 5 (Interrupt) Control/Status Register
0x4060_0028	UDCCS6	UDC Endpoint 6 (IN) Control/Status Register
0x4060_002C	UDCCS7	UDC Endpoint 7 (OUT) Control/Status Register
0x4060_0030	UDCCS8	UDC Endpoint 8 (IN) Control/Status Register
0x4060_0034	UDCCS9	UDC Endpoint 9 (OUT) Control/Status Register
0x4060_0038	UDCCS10	UDC Endpoint 10 (Interrupt) Control/Status Register
0x4060_003C	UDCCS11	UDC Endpoint 11 (IN) Control/Status Register
0x4060_0040	UDCCS12	UDC Endpoint 12 (OUT) Control/Status Register
0x4060_0044	UDCCS13	UDC Endpoint 13 (IN) Control/Status Register
0x4060_0048	UDCCS14	UDC Endpoint 14 (OUT) Control/Status Register
0x4060_004C	UDCCS15	UDC Endpoint 15 (Interrupt) Control/Status Register
0x4060_0050	UICR0	UDC Interrupt Control Register 0
0x4060_0054	UICR1	UDC Interrupt Control Register 1
0x4060_0058	USIR0	UDC Status Interrupt Register 0
0x4060_005C	USIR1	UDC Status Interrupt Register 1
0x4060_0060	UFNHR	UDC Frame Number Register High
0x4060_0064	UFNLR	UDC Frame Number Register Low
0x4060_0068	UBCR2	UDC Byte Count Register 2
0x4060_006C	UBCR4	UDC Byte Count Register 4
0x4060_0070	UBCR7	UDC Byte Count Register 7
0x4060_0074	UBCR9	UDC Byte Count Register 9
0x4060_0078	UBCR12	UDC Byte Count Register 12
0x4060_007C	UBCR14	UDC Byte Count Register 14
0x4060_0080	UDDR0	UDC Endpoint 0 Data Register
0x4060_0100	UDDR1	UDC Endpoint 1 Data Register
0x4060_0180	UDDR2	UDC Endpoint 2 Data Register
0x4060_0200	UDDR3	UDC Endpoint 3 Data Register
0x4060_0400	UDDR4	UDC Endpoint 4 Data Register
0x4060_00A0	UDDR5	UDC Endpoint 5 Data Register
0x4060_0600	UDDR6	UDC Endpoint 6 Data Register
0x4060_0680	UDDR7	UDC Endpoint 7 Data Register
0x4060_0700	UDDR8	UDC Endpoint 8 Data Register
0x4060_0900	UDDR9	UDC Endpoint 9 Data Register



Table 12-33. USB Device Controller Register Summary (Sheet 3 of 3)

Address	Name	Description
0x4060_00C0	UDDR10	UDC Endpoint 10 Data Register
0x4060_0B00	UDDR11	UDC Endpoint 11 Data Register
0x4060_0B80	UDDR12	UDC Endpoint 12 Data Register
0x4060_0C00	UDDR13	UDC Endpoint 13 Data Register
0x4060_0E00	UDDR14	UDC Endpoint 14 Data Register
0x4060_00E0	UDDR15	UDC Endpoint 15 Data Register



AC'97 Controller Unit

13.1 Overview

The AC'97 Controller Unit (ACUNIT) of the PXA255 processor supports the AC'97 revision 2.0 features listed in Section 13.2. The ACUNIT also supports audio controller link (AC-link). AC-link is a serial interface for transferring digital audio, modem, mic-in, CODEC register control, and status information.

The AC'97 CODEC sends the digitized audio samples that the ACUNIT stores in memory. For playback or synthesized audio production, the processor retrieves stored audio samples and sends them to the CODEC through the AC-link. The external digital-to-analog converter (DAC) in the CODEC then converts the audio samples to an analog audio waveform.

This chapter describes the programming model for the ACUNIT. The information in this chapter requires an understanding of the AC'97 revision 2.0 specification.

Note: The ACUNIT and the I²S Controller cannot be used at the same time.

13.2 Feature List

The processor ACUNIT supports the following AC'97 features:

- Independent channels for stereo Pulse Code Modulated (PCM) In, Stereo PCM Out, modem-out, modem-in and mono mic-in
 - All of the above channels support only 16-bit samples in hardware. Samples less than 16 bits are supported through software.
- Multiple sample rate AC'97 2.0 CODECs (48 kHz and below). The ACUNIT depends on the CODEC to control the varying rate.
- Read/write access to AC'97 registers
- · Secondary CODEC support
- Three Receive FIFOs (32-bit, 16 entries)
- Two Transmit FIFOs (32-bit, 16 entries)

The processor ACUNIT does not support these optional AC'97 features:

- Double-rate sampling (n+1 sample for PCM L, R & C)
- 18- and 20-bit sample lengths



13.3 Signal Description

The AC'97 signals form the AC-link, which is a point-to-point synchronous serial interconnect that supports full-duplex data transfers. All digital audio streams, Modem line CODEC streams, and command/status information are communicated over the AC-link. The AC-link uses General Purpose I/Os (GPIOs). Software must reconfigure the GPIOs to use them as the AC-link. The AC-link pins are listed and described in Table 13-1.

Table 13-1. External Interface to CODECs

Name	Direction	Description summary
nACRESET	0	Active-low CODEC reset. The CODEC's registers reset when nACRESET is asserted.
GP28/BITCLK	1	12.288 MHz bit-rate clock.
GP31/SYNC	0	48 kHz frame indicator and synchronizer.
GP30/SDATA_OUT	0	Serial audio output data to CODEC for digital-to-analog conversion.
GP29/SDATA_IN_0	I	Serial audio input data from Primary CODEC.
GP32/SDATA_IN_1	I	Serial audio input data from Secondary CODEC.

13.3.1 Signal Configuration Steps

- 1. Configure SYNC and SDATA_OUT as outputs.
- 2. Configure BITCLK, SDATA IN 0, and SDATA IN 1 as inputs.
- nACRESET is a dedicated output. It remains asserted on power-up. Complete these steps to deassert nACRESET:
 - a. Configure the other AC'97 signals as previously described.
 - b. In the Global Control Register (GCR), Set the GCR[COLD_RST] bit. Refer to Table 13-7 for more details.

Note: Refer to Section 4.1.3, "GPIO Register Definitions" on page 4-6 for details on programing the GPDR and GAFR for use with the ACUNIT.

13.3.2 Example AC-link

Figure 13-1 shows an example interconnect for an AC-link. The ACUNIT supports one or two CODECs on the AC-link. SDATA_IN_1 is not needed if only a Primary CODEC is connected.



AC'97 Controller Unit (ACUNIT)

AC-link

AC'97 Primary CODEC

Figure 13-1. Data Transfer Through the AC-link

13.4 AC-link Digital Serial Interface Protocol

Each AC'97 CODEC incorporates a five-pin digital serial interface that links it to the ACUNIT. AC-link is a full-duplex, fixed-clock, PCM digital stream. It employs a time division multiplexed (TDM) scheme to handle control register accesses and multiple input and output audio streams. The AC-link architecture divides each audio frame into 12 outgoing and 12 incoming data streams. Each stream has 20-bit sample resolution, but only 16-bit samples are supported in hardware. Each stream requires a DAC or an analog-to-digital converter (ADC), both having a minimum 16-bit resolution. The ACUNIT supports the data streams shown in Table 13-2.

Table 13-2. Supported Data Stream Formats (Sheet 1 of 2)

Channel	Slots	Comments
PCM Playback	Two output slots	Two-channel composite PCM output stream
PCM Record data	Two input slots	Two-channel composite PCM input stream
CODEC control	Two output slots	Control register write port
CODEC status	Two input slots	Control register read port
Modem Line CODEC Output	One output slot	Modem line CODEC DAC input stream
Modem Line CODEC Input	One input slot	Modem line CODEC ADC output stream



Table 13-2. Supported Data Stream Formats (Sheet 2 of 2)

Channel	Slots	Comments
Dedicated Microphone Input	One input slot	Dedicated microphone input stream in support of stereo AEC and other voice applications.
I/O Control	One output slot	One slot dedicated to GPOs on the modem CODEC.
I/O Status	One input slot	One slot dedicated to status from GPIs on the modem CODEC. Data is returned on every frame.

The ACUNIT provides synchronization for all data transaction on the AC-link. A data transaction is made up of 256 bits of information broken up into groups of 13 time slots and is called a frame. Time slot 0 is called the Tag Phase and is 16 bits long. The other 12 time slots are called the Data Phase. The Tag Phase contains one bit that identifies a valid frame and 12 bits that identify the time slots in the Data Phase that contain valid data. Each time slot in the Data Phase is 20 bits long.

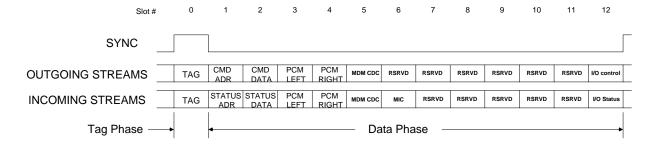
A frame begins when SYNC goes high. The amount of time that SYNC is high corresponds to the Tag Phase. AC'97 frames occur at fixed 48 kHz intervals and are synchronous to the 12.288 MHz bit rate clock, BITCLK.

The ACUNIT and the CODEC use the SYNC and BITCLK to determine when to send transmit data and when to sample receive data. A transmitter transitions the serial data stream on each rising edge of BITCLK and a receiver samples the serial data stream on each falling edge of BITCLK. The transmitter must tag the valid slots in its serial data stream. The valid slots are tagged in slot 0.

Serial data on the AC-link is ordered most significant bit (MSB) to least significant bit (LSB). The Tag Phase's first bit is bit 15 and the first bit of each slot in Data Phase is bit 19. The last bit in any slot is bit 0.

Figure 13-2 shows Tag and Data Phase organization for the ACUNIT and the CODEC. The figure also lists the slot definitions that the ACUNIT supports.

Figure 13-2. AC'97 Standard Bidirectional Audio Frame

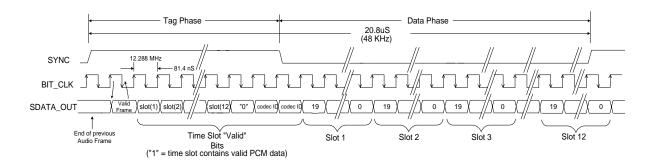


13.4.1 AC-link Audio Output Frame (SDATA_OUT)

The audio output frame data stream corresponds to the multiplexed bundles that make up the digital output data that targets the AC'97 DAC inputs and control registers. Each audio output frame supports up to twelve 20-bit outgoing data time slots. The ACUNIT does not generate samples larger than 16 bits. The four least significant bits are padded with zeroes.

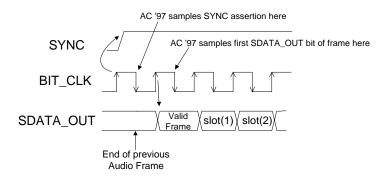


Figure 13-3. AC-link Audio Output Frame



A new audio output frame begins with a low-to-high SYNC transition synchronous to BITCLK's rising edge. BITCLK's falling edge immediately follows and AC'97 samples SYNC's assertion. BITCLK's falling edge marks the instance that AC-link's sides are each aware that a new audio frame has started. On BITCLK's next rising edge, the ACUNIT transitions SDATA_OUT into the slot 0's first bit position (valid frame bit). Each new bit position is presented to AC-link on a BITCLK rising edge and then sampled by AC'97 on the following BITCLK falling edge. This sequence ensures that data transitions and subsequent sample points for both incoming and outgoing data streams are time aligned.

Figure 13-4. Start of Audio Output Frame



The SDATA_OUT composite stream is MSB-justified (MSB first). The ACUNIT fills all non-valid slot bit positions with zeroes. If fewer than 20 valid bits exist in an assigned valid time slot, the ACUNIT stuffs all trailing non-valid bit positions of the 20-bit slot with zeroes.

For example, if a 16-bit sample stream is being played to an AC'97 DAC, the first 16 bit positions are presented to the DAC MSB-justified. They are followed by the next four bit positions that the ACUNIT stuffs with zeroes. This process ensures that the least significant bits do not introduce any DC biasing, regardless of the implemented DAC's resolution (16-, 18-, or 20-bit).

Note: When the ACUNIT transmits mono audio sample streams, software must ensure that the left and right sample stream time slots are filled with identical data.



13.4.1.1 Slot 0: Tag Phase

In slot 0, the first bit is a global bit (SDATA_OUT slot 0, bit 15) that flags the validity for the entire audio frame. If the valid frame bit is a 1, the current audio frame contains at least one slot time of valid data. The next 12 bit positions sampled by AC'97 indicate which of the corresponding 12 time slots contain valid data. Bits 0 and 1 of slot 0 are used as CODEC ID bits for I/O reads and writes to the CODEC registers as described in the next section. This way, data streams of differing sample rates can be transmitted across AC-link at its fixed 48 kHz audio frame rate. The CODEC can control the output sample rate of the ACUNIT using the SLOTREQ bits as described later (in the Input frame description).

13.4.1.2 Slot 1: Command Address Port

Slot 1 is the Command Address Port. Slot 1 (in conjunction with the Command Data Port of Slot 2) controls features and monitors status for AC'97 functions including, but not limited to, mixer settings and power management (refer to AC'97 Specification revision 2.0 for more details).

The control-interface architecture supports up to sixty-four16-bit read/write registers, addressable on even byte boundaries. Only accesses to even registers (0x00, 0x02, etc.) are valid. Accesses to odd registers (0x01, 0x03, etc.) are not valid.

Audio output frame slot 1 communicates control register address and write/read command information to the ACUNIT.

Two CODECs are connected to the single SDATA_OUT. To address the primary and secondary CODECs individually, follow these steps:

To access the primary CODEC:

- 1. Set the Valid Frame bit (slot 0, bit 15)
- 2. Set the valid bits for slots 1 and 2 (slot 0, bits 14 and 13)
- 3. Write 0b00 to the CODEC ID field (slot 0, bits 1 and 0)
- 4. Specify the read/write direction of the access (slot 1, bit 19).
- 5. Specify the index to the CODEC register (slot 1, bits 18-12)
- 6. If the access is a write, write the data to the command data port (slot 2, bits 19-4)

To access the secondary CODEC:

- 1. Set the Valid Frame bit (slot 0, bit 15)
- 2. Clear the valid bits for slots 1 and 2 (slot 0, bits 14 and 13)
- 3. Write a non-zero value (0b01, 0b10, 0b11) to the CODEC ID field (slot 0, bits 1 and 0)
- 4. Specify the read/write direction of the access (slot 1, bit 19).
- 5. Specify the index to the CODEC register (slot 1, bits 18-12)
- 6. If the access is a write, write the data to the command data port (slot 2, bits 19-4).



Table 13-3. Slot 1 Bit Definitions

Bit	Name	Description
Bit(19)	RW	1 = read, 0 = write
Bit(18:12)	IDX	Code register index
Bit(11:0)	reserved	Stuff with 0s

Only one I/O cycle can be pending across the AC-link at any time. The ACUNIT uses write and read posting on I/O accesses across the link. For example, read data from a CODEC register is not sent over the AC-link (Slot 2 of incoming stream) within the same frame that the read request is sent.

For CODEC reads, the ACUNIT gives the CODEC a maximum of four subsequent frames to respond -- if no response is received, the ACUNIT returns a dummy read completion (0xFFFF_FFFF) to the CPU and sets the Read Completion Status (RDCS) bit of the Global Status Register (GSR).

The CAIP bit of the CODEC Access Register (CAR) is used to assure that only one I/O cycle occurs across the AC-link at any time. Software must read the CAIP bit before initiating an I/O cycle. If the CAIP bit reads as a one, another driver is performing an I/O cycle; if the CAIP bit reads as a zero, a new I/O cycle can be initiated.

The exception to posted accesses is reads to the CODEC GPIO Pin Status register (address 0x54). CODEC GPIO Pin Status read data is sent by the CODEC over the AC-link in the same frame that the read request was sent to the CODEC. The CODEC GPIO Pin Status read data is sent in Slot 12 of the incoming stream. A CODEC with a GPIO Pin Status register must constantly send the status of the register in slot 12.

13.4.1.3 Slot 2: Command Data Port

Slot 2 is the Command Data Port. Slot 2 (in conjunction with the Command Address Port of Slot 1) delivers 16-bit control register write data in the event that the current command port operation is a write cycle (as indicated by slot 1, bit 19).

Table 13-4. Slot 2 Bit Definitions

Bit	Name	Description
Bit(19:4)	Control register write data	Stuffed with 0s if current operation is a read
Bit(3:0)	reserved	Stuffed with 0s

If the current command port operation is a read, the ACUNIT fills Slot 2 with zeroes.

13.4.1.4 Slot 3: PCM Playback Left Channel

Slot 3 contains the composite digital audio left playback stream. If the playback stream contains an audio sample with a resolution that is less than 20 bits, the ACUNIT fills all trailing non-valid bit positions with zeroes.



13.4.1.5 Slot 4: PCM Playback Right Channel

Slot 4 is the composite digital audio right playback stream. If the playback stream contains an audio sample with a resolution that is less than 20 bits, the ACUNIT fills all trailing non-valid bit positions with zeroes.

13.4.1.6 Slot 5: Modem Line CODEC

Slot 5 contains the MSB justified modem DAC input data if the modem line CODEC is supported. The optional modem DAC input resolution can be implemented as 16, 18, or 20 bits. If the modem line CODEC is supported, the ACUNIT driver determines the DAC resolution at boot time. During normal runtime operation, the ACUNIT fills all trailing non-valid bit positions in the Slot 5 with zeroes. The modem line CODEC may be a separate CODEC on the secondary line or it may be integrated with the audio CODEC.

13.4.1.7 Slots 6-11: Reserved

These slots are reserved for future use. The ACUNIT fills them with zeroes.

13.4.1.8 Slot 12: I/O Control

Slot 12 contains 16 MSB bits for GPO Status (output). The following rules govern the use of Slot 12:

- 1. Slot 12 is initially marked invalid by default.
- 2. A write to address 0x54 in CODEC IO space (using Slot 1 and Slot 2 in the outgoing stream of the present frame) results in the same write data (sent in Slot 2 of the present outgoing frame) being sent in Slot 12 of the next outgoing frame, where Slot 12 is then marked as valid.
- 3. After the first write to address 0x54, Slot 12 remains valid for all subsequent frames. The data transmitted on Slot 12 is the data last written to address 0x54. Any subsequent write to the register sends the new data out on the next frame.
- 4. Following a system reset or AC'97 cold reset, Slot 12 is invalidated. Slot 12 remains invalid until the next write to address 0x54.

13.4.2 AC-link Audio Input Frame (SDATA_IN)

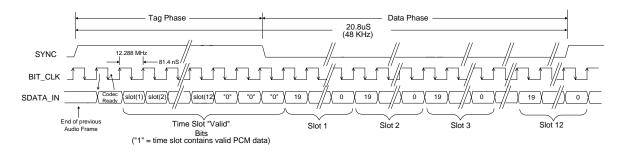
The ACUNIT has two SDATA_IN lines, one primary and one secondary. Each line can have CODECs attached. The type of CODEC attached determines which slots are valid or invalid. The data slots on the two inputs are completely orthogonal, i.e., no two data slots at the same location will be valid on both lines.

Multiple input data streams are received and multiplexed on slot boundaries as dictated by the slot valid bits in each stream. Each AC-link audio input frame consists of twelve 20-bit time slots. Slot 0 is reserved and contains 16 bits that are used for AC-link protocol infrastructure.

Software must poll the first bit in the audio input frame (SDATA_IN slot 0, bit 15) for an indication that the CODEC is in the CODEC ready state before it places the ACUNIT into data transfer operation. When the "CODEC is ready" state is sampled, the next 12 sampled bits indicate which of the 12 time slots are assigned to input data streams and whether they contain valid data. Figure 13-5, "AC'97 Input Frame" illustrates the time slot-based AC-link protocol.

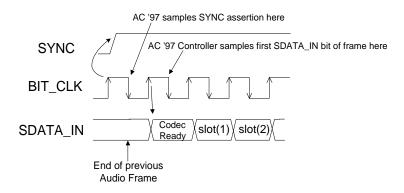


Figure 13-5. AC'97 Input Frame



A new audio input frame begins when SYNC transitions from low to high. The low to high transition is synchronous to BITCLK's rising edge. On BITCLK's next falling edge, AC'97 samples SYNC's assertion. This falling edge marks the moment that AC-link's sides are each aware that a new audio frame has started. The next time BITCLK rises, the ACUNIT transitions SDATA_IN to the first bit position in slot 0 (CODEC ready bit). Each new bit position is presented to AC-link on a BITCLK's rising edge and then sampled by ACUNIT on the following BITCLK's falling edge. This sequence ensures that data transitions and subsequent sample points are time aligned for both incoming and outgoing data streams.

Figure 13-6. Start of an Audio Input Frame



The SDATA_IN composite stream is MSB justified (MSB first) and the AC'97 CODEC fills non-valid bit positions with zeroes. SDATA IN data is sampled on BITCLK's falling edge.

13.4.2.1 Slot 0: Tag Phase

In Slot 0, the first bit is a global bit (SDATA_IN slot 0, bit 15) that indicates whether or not the CODEC is in the CODEC ready state. If the CODEC Ready bit is a 0, the CODEC is not ready for operation. This condition is normal after power is asserted on reset, i.e., while the CODEC voltage references are settling. When the AC-link CODEC Ready indicator bit is a one, the AC-link and AC'97 control and status registers are fully operational. The ACUNIT must probe the CODEC Powerdown Control/Status register to determine which subsections are ready.



CODEC Ready, sent by the CODEC on its data out stream in bit 15 of Slot 0, is not expected to change during normal operation. The AC'97 Specification revision 2.0 requires that a CODEC only change its CODEC Ready status in response to a power down (PR) state change issued by the ACUNIT. The ACUNIT's hardware by itself does not monitor the CODEC Ready for the purpose of sending or receiving data. The ACUNIT stores CODEC Ready in the PCR bit of the GSR for a primary CODEC and SCR bit of the GSR for a secondary CODEC. Software should monitor PCR or SCR to trigger a DMA or a programmed I/O operation. The ACUNIT only samples CODEC Ready valid once and then ignores it for subsequent frames. CODEC Ready is only resampled after a PR state change.

13.4.2.2 Slot 1: Status Address Port/SLOTREQ bits

Slot 1 is the Status Address Port. Slot 1 monitors the status of ACUNIT functions including, but not limited to, mixer settings and power management.

Slot 1 echoes the control register index for the data to be returned in Slot 2, if the ACUNIT tags Slot 1 and Slot 2 as valid during Slot 0.

The ACUNIT only accepts status data (Slot 2 of incoming stream) if the accompanying control register index (Slot 1 of incoming stream) matches the last valid control register index that was sent in Slot 1 of the outgoing stream of the most recent previous frame.

For multiple sample rate output, the CODEC examines its sample-rate control registers, its FIFOs' states, and the incoming SDATA_OUT tag bits at the beginning of each audio output frame to determine which SLOTREQ bits to set active (low). SLOTREQ bits asserted during the current audio input frame indicate which output slots require data from the ACUNIT in the next audio output frame. For fixed 48 kHz operation, the SLOTREQ bits are set active (low), and a sample is transferred during each frame.

For multiple sample-rate input, the tag bit for each input slot indicates whether valid data is present.

Again, Slot 1 delivers a CODEC control register index and multiple "sample-rate slot request flags" for all output slots. AC'97 defines the ten least significant bits of Slot 1 as CODEC ondemand data request flags for outgoing stream Slots 3-12. For two-channel audio, only data-request flags corresponding to slots 3 and 4 are meaningful.

Table 13-5. Input Slot 1 Bit Definitions (Sheet 1 of 2)

Bit	Description
19	reserved (Filled with zero)
18-12	Control register Index (Filled with zeroes if AC'97 tags it invalid)
11	Slot 3 request: PCM Left channel
10	Slot 4 request: PCM Right channel
9	Slot 5 request: Modem Line 1
8	Slot 6 request: NA
7	Slot 7 request: NA
6	Slot 8 request: NA
5	Slot 9 request: NA



Table 13-5. Input Slot 1 Bit Definitions (Sheet 2 of 2)

4	Slot 10 request: NA
3	Slot 11 request: NA
2	Slot 12 request: NA
1,0	reserved (Filled with zero)

SLOTREQ bits are independent of the Control Register Index bits.

Note: Slot requests for Slot 3 and Slot 4 are always set or cleared in tandem (both are set or both are cleared).

13.4.2.3 Slot 2: Status Data Port

Slot 2 delivers 16-bit control register read data.

Table 13-6. Input Slot 2 Bit Definitions

Bit	Name	Description
Bit(19:4)	Control register read data	Filled with data
Bit(3:0)	reserved	Filled with zeroes

Note: If Slot 2 is tagged invalid, the CODEC fills the entire slot with zeroes.

13.4.2.4 Slot 3: PCM Record Left Channel

Slot 3 contains the CODEC left channel output.

The CODEC transmits its ADC output data (MSB first) and fills any trailing non-valid bit positions with zeroes.

13.4.2.5 Slot 4: PCM Record Right Channel

Slot 4 contains the CODEC right-channel output.

The CODEC transmits its ADC output data (MSB first), and fills any trailing non-valid bit positions with zeroes.

13.4.2.6 Slot 5: Optional Modem Line CODEC

Slot 5 contains MSB justified line modem ADC output data (if the line modem CODEC is supported).

The ACUNIT only supports a 16-bit ADC output resolution from the optional line modem.

13.4.2.7 Slot 6: Optional Dedicated Microphone Record Data

Slot 6 contains an optional third PCM system-input channel available for dedicated use by a microphone. This input channel supplements a true stereo output to enable a more precise echocancellation algorithm for speakerphone applications.



The ACUNIT only supports a 16-bit resolution from the microphone.

13.4.2.8 Slots 7-11: Reserved

Slots 7-11 are reserved for future use. The ACUINT ignores them.

13.4.2.9 Slot 12: I/O Status

GPIOs which are configured as inputs return their status in Slot 12 of every frame. Only the 16 MSBs are used to return GPIO status. Bit 0 in the LSBs indicates a GPI Input Interrupt event. See the AC'97 revision 2.0 spec for more information.

The data returned on the latest frame is also accessible to software through the CODEC register at address 0x54 in the modem CODEC I/O space. Data received in Slot 12 is stored internally in the ACUNIT. So when software initiates a read of the CODEC register at address 0x54 in the modem CODEC I/O space, the read data is already inside the ACUNIT.

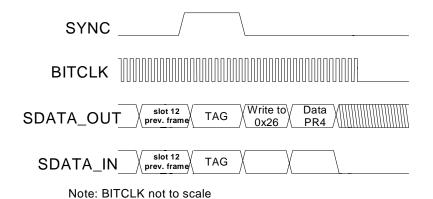
13.5 AC-link Low Power Mode

Software must set the ACLINK_OFF bit of the GCR to one before the processor enters Sleep Mode. The ACUNIT then drives SYNC and SDATA_OUT to a logic low level . The ACUNIT maintains nACRESET high when ACLINK_OFF is set to one.

13.5.1 Powering Down the AC-link

The AC-link signals enter a low power mode after the PR4 bit of the AC'97 CODEC Powerdown Register (0x26) is set to a 1 (by writing 0x1000). Then, the Primary CODEC drives both BITCLK and SDATA_IN to a logic low voltage level. The sequence follows the timing diagram shown in Figure 13-7.

Figure 13-7. AC-link Powerdown Timing





The ACUNIT transmits the write to the Powerdown Register (0x26) over the AC-link. Set up the ACUNIT so that it does not transmit data in Slots 3-12 when it writes 0x1000 to the PR4 bit of the Powerdown Register. AC'97 revision 2.0 does not require the CODEC to process other data when it receives a power down request. When the CODEC receives the power down request, it immediately transitions BITCLK and SDATA_IN to a logic low level.

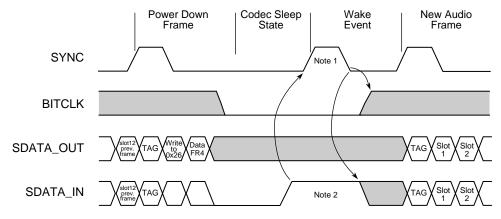
13.5.2 Waking up the AC-link

13.5.2.1 Wake up triggered by the CODEC

To wake up the AC-link, a CODEC drives SDATA_IN to a logic high level. The rising edge triggers the Resume Interrupt if that CODEC's resume enable bit is set to a one. The CPU then wakes up the CODEC using the cold or warm reset sequence. The ACUNIT uses a warm reset to wake up the primary CODEC. The CODEC detects a warm reset when SYNC is driven high for a minimum of one microsecond and the BITCLK is absent. The CODEC must wait until it samples SYNC low before it can start BITCLK. The CODEC that signaled the wake event must keep its SDATA_IN high until it detects that a warm reset has been completed. The CODEC can then transition its SDATA_IN low.

Figure 13-8 shows the AC-link timing for a wake up triggered by a CODEC. Because the processor may need to be awakened, the Power Management unit detects the AC'97 wake-up event (SDATA_IN high for more than one microsecond). When the ACUNIT is ready, it responds to the wake-up event by asserting a warm or cold reset (see Figure 13-8). A Modem CODEC may require the capacity to wake up the AC-link to report events such as Caller-ID and wake-up-on-ring.

Figure 13-8. SDATA IN Wake Up Signaling



Notes:

- After SDATA_IN goes high, SYNC must be held for a minimum of 1 μSec.
- The minimum SDATA_IN wake up pulse width is 1 μSec.
- BITCLK not to scale.

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NOTES:

- 1. After SDATA_IN goes high, SYNC must be held for a minimum of 1 μ Sec.
- 2. The minimum SDATA_IN wake up pulse width is 1 μ Sec.
- 3. BITCLK not to scale



13.5.2.2 Wake Up Triggered by the ACUNIT

AC-link protocol provides for a cold AC'97 reset and a warm AC'97 reset. The current power-down state ultimately dictates which AC'97 reset is used. Registers must stay in the same state during all power-down modes unless a cold AC'97 reset is performed. In a cold AC'97 reset, the AC'97 registers are initialized to their default values.

After a power down, the AC-link must wait for a minimum of four audio frame times after the frame in which the power down occurred before it can be reactivated by reasserting the SYNC signal. When AC-link powers up, it indicates readiness through the CODEC ready bit (input slot 0, bit 15).

13.5.2.2.1 Cold AC'97 Reset

A cold reset is generated when the nACRESET pin is asserted by software setting the COLD_RST bit of the GCR to zero. Asserting and deasserting nACRESET activates SDATA_OUT and causes the CODEC to activate BITCLK. All AC'97 control registers are initialized to their default power-on reset values. nACRESET is an asynchronous input to the AC'97 CODEC.

13.5.2.2.2 Warm AC'97 Reset

A warm AC'97 reset reactivates the AC-link without altering the current AC'97 register values. A warm reset is generated by software setting the WARM_RST bit of the GCR to one. When BITCLK is absent, generation of a warm reset results in SYNC being driven high for a minimum of one microsecond. The CODEC must not activate BITCLK until it samples SYNC low again. This prevents a new audio frame from being falsely detected.

13.6 **ACUNIT Operation**

The ACUNIT can be accessed through the processor or the DMA controller. The processor uses programmed I/O instructions to access the ACUNIT, and it can access four register types:

- ACUNIT registers: Accessible at 32-bit boundaries. They are listed in Section 13.8.3.
- CODEC registers: An audio or modem CODEC can contain up to sixty-four 16-bit registers. A
 CODEC uses a 16-bit address boundary for registers. The ACUNIT supplies access to the
 CODEC registers by mapping them to its 32-bit address domain boundary. Section 13.8.3.17
 describes the mapping from the 32-bit to 16-bit boundary. A write or read operation that
 targets these registers is sent across the AC-link.
- Modem CODEC GPIO register: If the ACUNIT is connected to a modem CODEC, the CODEC GPIO register can also be accessed. The CODEC GPIO register uses access address 0x0054 in the CODEC domain. The GPIO write operation goes across the AC-link, but a read does not. The GPIO register contents are continuously updated into a shadow register in the ACUNIT when a frame is received from the CODEC. When the processor tries to read the CODEC GPIO register, this shadow register is read instead.
- ACUNIT FIFO data: The ACUNIT has two Transmit FIFOs for audio-out and modem-out and three receive FIFOs for audio-in, modem-in, and mic-in. Data enters the transmit FIFOs by writing to either the PCM Data Register (PCDR) or the Modem Data Register (MODR).



Receive FIFO entries are read through the PCDR, the MODR, or the Mic-in Data Register (MCDR).

Note: After it is enabled, the ACUNIT requests the DMA to fill the transmit FIFO.

Note: The ACUNIT registers do not store the status of DMA requests or information regarding the number of data samples in each FIFO. As a result, programmed I/O must not be used in place of DMA requests for data transfers.

Only the DMA can access the FIFOs. The DMA controller accesses FIFO data in 8-, 16-, or 32-byte blocks. The DMA request thresholds are not programmable. The ACUNIT makes a transmit DMA request when the transmit FIFO has less than 32 bytes. The ACUNIT makes a receive DMA request when the receive FIFO has 32 bytes or more. Regardless of burst size, the DMA descriptor length must be a multiple of 32 bytes to prevent audio artifacts from being introduced onto the AC-link.

The DMA controller responds to the following ACUNIT DMA requests:

- PCM FIFO transmit and receive DMA requests made when the PCM transmit and receive FIFOs are half full.
- Modem FIFO transmit and receive DMA requests made when the modem transmit and receive FIFOs are half full.
- Mic-in receive DMA requests made when the Mic-in receive FIFO is half full.

13.6.1 Initialization

The AC'97 CODEC and ACUNIT are reset on power up. After power up, the nACRESET signal remains asserted until the audio or modem driver sets the COLD_RST bit of the GCR to one. During operation, clearing the COLD_RST bit to zero resets the ACUNIT and CODEC. To initialize the ACUNIT follow theses steps:

- 1. Program the GPIO Direction register and GPIO Alternate Function Select register to assign proper pin directions for the ACUNIT ports. Refer to Section 13.3 for details.
- 2. Set the COLD_RST bit of the GCR to one to deassert nACRESET. Until this is done, all other registers remain in a reset state. Deasserting nACRESET has the following effects:
 - a. Frames filled with zeroes are transmitted because the transmit FIFO is still empty. This situation does not cause an error condition.
 - b. The ACUNIT records zeroes until the CODEC sends valid data.
 - c. DMA requests are enabled.
- 3. Enable the Primary Ready Interrupt Enable and/or the Secondary Ready Interrupt Enable by setting the PRIRDY_IEN bit and/or the SECRDY_IEN bit of the GCR to one.
- 4. Software enables DMA operation in response to primary and secondary ready interrupts.
- 5. The ACUNIT triggers transmit DMA requests. The DMA fills the transmit FIFO in response.
- 6. The ACUNIT continues to transmit zeroes until the transmit FIFO is half full. When it is half full, valid transmit FIFO data is sent across the AC-link.

Note: When nACRESET is deasserted, a read to the CODEC Mixer register returns the type of hardware that resides in the CODEC. If the CODEC is not present or if the AC'97 is not supported, the



ACUNIT does not set the CODEC-ready bit, GCR[PCRDY] for the Primary CODEC or GCR[SCRDY] for the Secondary CODEC.



13.6.2 Trailing bytes

Trailing bytes in the transmit and receive FIFOs are handled as follows:

If the transmit buffers do not have 32-byte resolution, the trailing bytes in the Transmit FIFO are not transmitted. A transmit buffer must be padded with zeroes if it is smaller than a multiple of 32 bytes. Regardless of burst size, the DMA descriptor length must be a multiple of 32 bytes to prevent audio artifacts from being introduced onto the AC-link.

If the CODEC transmitted data has a total buffer size smaller than a multiple of 32 bytes, zeroes are recorded. A receive DMA request is made when the receive FIFO is half-full.

13.6.3 Operational Flow for Accessing CODEC Registers

Software accesses the CODEC registers by translating a 32-bit processor physical address to a 7-bit CODEC address. For details regarding the address translation, refer to Section 13.8.3.17.

Software must read the CODEC Access Register (CAR) to lock the AC-link. The AC-link is free if the CAIP bit of the CAR is zero. For details about the CAR, refer to Table 13-13.

A read access to the CAR sets the CAIP bit. The ACUNIT clears the CAIP bit when the CODEC-write or CODEC-read operation completes. Software can also clear the CAIP bit by writing a zero to it.

After it locks the AC-link, software can write or read a CODEC register using the appropriate processor physical address.

The ACUNIT sets the CDONE bit of the GSR to one after the completion of a CODEC write operation. For details, refer to Table 13-8. Software clears this bit by writing a 1 to it.

To read a CODEC, the software must complete the following steps:

- 1. Software issues a dummy read to the CODEC register. The ACUNIT responds to this read operation with invalid data. The ACUNIT then initiates the read access across the AC-link.
- 2. When the CODEC read operation completes, the ACUNIT sets the SDONE bit of the GSR to one. For details, refer to Table 13-8. Software clears this bit by writing a 1 it.
- 3. Software repeats the read operation as detailed in Step 1. The ACUNIT now returns the data sent by the CODEC. The second read operation also initiates a read access across the AC-link.
- 4. The ACUNIT times-out the read operation if the CODEC fails to respond in four SYNC frames. In this case, the second read operation returns a timed-out data value of 0x0000_FFFF.

13.7 Clocks and Sampling Frequencies

By default, the ACUNIT transmits and receives data at a sampling frequency of 48 kHz. It can, however, sample data at frequencies less than 48 kHz if the CODEC supports on-demand slot requests. The CODEC in this case executes a certain algorithm and informs the ACUNIT not to transmit valid data in certain frames. For example, if the ACUNIT sends out 480 frames, and the CODEC instructs the ACUNIT not to send valid data in 39 of those 480 frames, the CODEC would have in effect sampled data at 44.1 kHz. When the CODEC transmits data (ACUNIT-receive mode), it can use the same algorithm to transmit valid frames with some empty ones mixed in between.



All data transfers across the AC-link are synchronized to SYNC's rising edge. The ACUNIT divides the BITCLK by 256 to generate the SYNC signal. This calculation yields a 48 kHz SYNC signal, and its period defines a frame. Data is transitioned on AC-link on every BITCLK rising edge and subsequently sampled on AC-link's receiving side on each following BITCLK falling edge. For a timing diagram see Figure 13-3.

The ACUNIT synchronizes data between two different clock domains: the BITCLK and an internal system clock. This internal system clock is always half the run mode frequency. The run mode frequency is equal to or greater than eight times the BITCLK frequency.

13.8 Functional Description

The functional description section applies to all channels.

13.8.1 FIFOs

The ACUNIT has five FIFOs:

- PCM Transmit FIFO, with sixteen 32-bit entries.
- PCM Receive FIFO, with sixteen 32-bit entries.
- Modem Transmit FIFO, with sixteen 32-bit entries (upper 16 bits must always be zero).
- Modem Receive FIFO, with sixteen 32-bit entries (upper 16 bits are always zero).
- Mic-in Receive FIFO, with sixteen 32-bit entries (upper 16 bits are always zero).

A receive FIFO triggers a DMA request when the FIFO has eight or more entries. A transmit FIFO triggers a DMA request when it holds less than eight entries. A transmit FIFO must be half full (filled with eight entries) before any data is transmitted across the AC-link.

13.8.1.1 Transmit FIFO Errors

Channel-specific status bits are updated during transmit under-run conditions and will trigger interrupts if enabled. Refer to Table 13-11 and Table 13-20 for details regarding the status bits. During transmit under-run conditions, the last valid sample is continuously sent out across the AC-link. A transmit under-run can occur under the following conditions:

- Valid transmit data is still available in memory but the DMA controller starves the transmit FIFO because it is servicing other higher priority peripherals.
- The DMA controller has transferred all valid data from memory to the transmit FIFO. This prompts the last valid sample to be echoed across the AC-link until nACRESET is asserted to turn off the ACUNIT.

13.8.1.2 Receive FIFO Errors

Channel-specific status bits are updated during receive overrun conditions and trigger interrupts when enabled. Refer to Table 13-12, Table 13-16, and Table 13-21 for details regarding the status bits. During receive over-run conditions, data that the CODEC sends is not recorded.



13.8.2 Interrupts

The following status bits interrupt the processor when the interrupts are enabled:

- Mic-in FIFO error: Mic-in Receive FIFO's over-run or under-run error.
- Modem-in FIFO error: Modem Receive FIFO's over-run or under-run error.
- PCM-in FIFO error: Audio Receive FIFO's over-run or under-run error.
- Modem-out FIFO error: Modem Transmit FIFO's over-run or under-run error.
- PCM-out FIFO error: Audio Transmit FIFO's over-run or under-run error.
- Modem CODEC GPI status change interrupt: Interrupts the CPU if bit 0 of Slot 12 is set. This
 indicates a change in one of the bits in the modem CODEC's GPIO register.
- Primary CODEC resume interrupt: Sets a status register bit when the Primary CODEC resumes from a lower power mode. Software writes a one to this bit to clear it.
- Secondary CODEC resume interrupt: Sets a status register bit when the Secondary CODEC resumes from a lower power mode. Software writes a one to this bit to clear it.
- CODEC command done interrupt: Interrupts the CPU when a CODEC register's command is completed. Software writes a one to this bit to clear it.
- CODEC status done interrupt: Interrupts the CPU when a CODEC register's status address and data reception are completed. Software writes a one to this bit to clear it.
- Primary CODEC ready interrupt: Sets a status register bit when the Primary CODEC is ready. The CODEC sets bit 0 of Slot 0 on the input frame to signal that it is ready. Software clears the PRIRDY_IEN bit of the GCR to clear this interrupt.
- Secondary CODEC ready interrupt: Sets a status register bit when the Secondary CODEC is ready. The CODEC sets bit 0 of Slot 0 on the input frame to signal that it is ready. Software clears the SECRDY IEN bit of the GCR to clear this interrupt.

13.8.3 Registers

The ACUNIT and CODEC registers are mapped in addresses ranging from 0x4050_0000 through 0x405F_FFFF. All ACUNIT registers are 32-bit addressable. Even though a CODEC has up to sixty-four 16-bit registers that are 16-bit addressable, they are accessed via a 32-bit address map and translated to 16-bits for the CODEC.

Programmed I/O and DMA bursts can access the following registers:

- Global registers: The ACUNIT has three global registers: Status, Control, and CODEC access registers that are common to the audio and modem domains.
- Channel-specific audio ACUNIT registers refer to PCM-out, PCM-in, and mic-in channels.
- Channel-specific Modem ACUNIT registers refer to modem-out and modem-in channels.
- Audio CODEC registers
- Modem CODEC registers



Channel specific data registers are for FIFO accesses and the PCM, Modem, and Mic-in FIFOs each have a register. A write access to one of these registers updates the written data in the corresponding Transmit FIFO. A read access to one of these registers flushes out an entry from the corresponding Receive FIFO.

Note: Register tables show organization and individual bit definitions. All reserved bits are read as unknown values and must be written with a 0. A question mark indicates the value is unknown at reset.

Note: Some register bits receive status from CODECs. The CODEC status sets the bit and software clears the bit (write a one to clear). The status can come in at any time, even when the bit is set or during a software clear. If software clears the bit as the CODEC status updates the bit, the CODEC status event takes higher priority. The term "interruptible" denotes bits that can be affected by this condition.

13.8.3.1 Global Control Register (GCR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 13-7. GCR Bit Definitions (Sheet 1 of 2)

Physical Address

			PI		ical A 050_00		SS						GC	R R	egis	ster							AC'	97 (Con	trol	ler l	Jnit			
Bit	31	30	29	28	27 20	3 25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
					res	erve	∤d					CDONE_IE	SDONE_IE			r	ese	rved	d			SECRDY_IEN	PRIRDY_IEN	reserved		SECRES_IEN	PRIRES_IEN	ACLINK_OFF	WARM_RST	COLD_RST	GIE
Reset	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
•		Bi	ts		N	ame												De	escr	ipti	on										
•		31:	20							_																					
		31:20 — reserved Command Done Interrupt Enable (CDONE_IE): 0 = The ACUNIT does not trigger an interrupt to the CPU after sending the comaddress and data to the CODEC. 1 = The ACUNIT triggers an interrupt to the CPU after sending the command a data to the CODEC. Status Done Interrupt Enable (SDONE_IE):																		ınd											
		18	8		SDC	NE_	_IE	0 =	Int En	erru	ipt is	s dis	sable	ed				_	,	ring	the	statı	us ac	ddre	ess	and	data	a fro	m tl	ne	
•		17:	10			_		rese	erve	:d																					
		9)		SECF	:DY_	<u>I</u> EN	0 =	Int En	erru able	ipt is	s dis ın int	able terru	ed opt t	t En o oc _1 pi	cur	`			_	,	у С	ODE	:C s	end	ls th	e C0	ODE	C F	REA	DY
•								Prin	•		•				nabl	e (F	PRIF	DY.	_IEI	N):											
		8	3		PRIR	DY_	IEN	0 = 1 =	En	able	es a	s dis in int ATA	terru	ipt t		cur	whe	n th	e P	rima	ary C	OD	EC s	sen	ds th	he C	OD	EC	REA	NDY	bit
•		7:	6			_		rese	erve	d																					



Table 13-7. GCR Bit Definitions (Sheet 2 of 2)

	Bits Name Description Secondary Resume Interrupt Enable: 5 SECRES_IEN 0 = Interrupt is disabled 1 = Enables an interrupt to occur when the Secondary CODEC causes a resume ever the AC-link Primary Resume Interrupt Enable: 4 PRIRES_IEN 0 = Interrupt is disabled																											
Bit	AC 97 Controller Onit AC 97 Controller Onit															1	0											
Dit.	reserved																											
				re	serve	d					1 1			res	serv	red			SECRDY_IE	PRIRDY_IEI	pennesar		SECRES_IE	PRIRES_IEI	ACLINK_OF	WARM_RS	COLD_RS1	GIE
Reset	0	0 0	0	0	0 0	0	0	0 0	0	0	0	0	0	0 () (0 0	0	0	0	0	0	0	0	0	0	0	0	0
		Bits			Name											Desc	ripti	ion										
	5 SECRES_IEN 0 = Interrupt is disabled 1 = Enables an interrupt to occur when the Secondary CODEC causes a resume even the AC-link Primary Resume Interrupt Enable: 0 = Interrupt is disabled 1 = Enables an interrupt to occur when the Primary CODEC causes a resume event of the secondary CODE															ent	on											
	4 PRIRES_IEN PRIRES_IEN PRIRES_IEN AC-link Primary Resume Interrupt Enable: 0 = Interrupt is disabled 1 = Enables an interrupt to occur when the Primary CODEC causes a resume event of AC-link AC-link Shut Off:															on t	the											
	Primary Resume Interrupt Enable: 0 = Interrupt is disabled 1 = Enables an interrupt to occur when the Primary CODEC causes a resume event of AC-link AC-link Shut Off: 0 = If the AC-link was off, turns it back on, otherwise this bit has no effect. 1 = Causes the ACUNIT to drive SDATA_OUT and SYNC outputs low and turn off in buffer enables. The reset output is however maintained high. The AC-link will no allowed to access any of the FIFOs. Setting this bit does not ensure a clean shut down. Software must make sure that all															ot be												
		2		WA	RM_R		0 = 1 = If so igno	A wa Caus susp ftwar	rm r ses a ende e atte	eset war ed Co empt ne bit	is no m re ODE s to t will	eset t EC wi perfo not	o oc thou orm a chan	cur out	arin rm i This	g it's reset bit is	inte whil s sel	rnal e Bli f cle	regi TCL arin	ster: _K is g i.e	s. run ., it	nin(rem	g, th ains	e w	rite v	will b		set
		1		со	LD_R		0 = 1 = Defa after	Caus and A co aults t susp erated	es a he (d re o 0. end	colo CODI set is After s, he	EC v s not resence,	vill be gen et, th if th	e los erate e dri	t. ed. iver r	nus	t to s	et th	is bi	t to	a 1.1	Γhe	valu	ıe of	f this	s bit	is re		
		0			GIE		This inter 0 =	bit co rupt. If this gene If this inter	ontro s bit rate s bit	is no d. is se	heth t set	er the t, bit e cha	e chains of nge	ange the (Glob alue	oal St	atus GPI	Reg (as i	giste	er is a	set,	but	an i	nter	rupt	is r	ot	an

13.8.3.2 Global Status Register (GSR)



Table 13-8. GSR Bit Definitions (Sheet 1 of 2)

	Pł	nysical 4050_								GS	R Re	gist	er							AC'	97 (Con	trol	ler l	Jnit	:		
Bit	31 30 29	28 27	26 2	25 24	23	22	21	20	19	18	17	16 1	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		re	eserv	/ed					CDONE	SDONE	reserved		RDCS	BIT3SLT12	BIT2SLT12	BIT1SLT12	SECRES	PRIRES	SCR	PCR	MINT	POINT	PIINT	reserved		MOINT	MINT	GSCI

0 0 0 0	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
Bits	Name	Description
31:20	_	reserved
		Command Done (CDONE):
19	CDONE	 0 = ACUNIT has not sent command address and data to the CODEC. 1 = ACUNIT has sent command address and data to the CODEC.
		This bit is cleared by software writing a '1' to this location (interruptible)
		Status Done (SDONE):
18	SDONE	 0 = ACUNIT has not received status address and data from the CODEC. 1 = ACUNIT has received status address and data from the CODEC.
		This bit is cleared by software writing a '1' to this location (interruptible)
17:16	_	reserved
		Read Completion Status:
		This bit indicates the status of CODEC read completions.
15	RDCS	0 = The CODEC read completed normally1 = The CODEC read resulted in a timeout.
		The bit remains set until cleared by software. This bit is cleared by software writing a '1' to this location.
14	BIT3SLT12	Bit 3 of slot 12:
14	BIT33LI 12	Display Bit 3 of the most recent valid slot 12
13	BIT2SLT12	Bit 2 of slot 12:
15	DITZOLITZ	Display Bit 2 of the most recent valid slot 12
12	BIT1SLT12	Bit 1 of slot 12:
12	DITTOLITZ	Display Bit 1 of the most recent valid slot 12
		Secondary Resume Interrupt:
11	SECRES	0 = A resume event has not occurred on the SDATA_IN_1. 1 = A resume event occurred on the SDATA_IN_1.
		This bit is cleared by software writing a '1' to this location (interruptible).
		Primary Resume Interrupt:
10	PRIRES	0 = A resume event has not occurred on the SDATA_IN_0. 1 = A resume event occurred on the SDATA_IN_0.
		This bit is cleared by software writing a '1' to this location (interruptible).
9	SCR	Secondary CODEC Ready (SCR)
9	SCR	Reflects the state of the CODEC ready bit in SDATA_IN_1 (interruptible)
8	PCR	Primary CODEC Ready (PCR)
0	FUR	Reflects the state of the CODEC ready bit in SDATA_IN_0 (interruptible)
		Mic In Interrupt (MINT)
7	MINT	0 = None of the mic-in channel interrupts occurred.
		1 = One of the mic-in channel interrupts occurred.
		When the specific interrupt is cleared, this bit will be cleared (interruptible).



Table 13-8. GSR Bit Definitions (Sheet 2 of 2)

			P			al <i>A</i> 60_0			SS						GS	R R	egis	ster							AC'	97 (Con	trol	ler l	Unit			
Bit	31	30	29	28	2	27 2	6	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						re	sei	rve	d					CDONE	SDONE	portugada		RDCS	BIT3SLT12	BIT2SLT12	BIT1SLT12	SECRES	PRIRES	SCR	PCR	MINT	POINT	PIINT	recerved	00000	MOINT	L	GSCI
Reset	0	0	0	0	(0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			ı	lar	me												De	esci	iptio	on										
	PCM Out Interrupt (POINT) 0 = None of the PCM out channel interrupts occurred. 1 = One of the PCM out channel interrupts occurred.																																
	0 = None of the PCM out channel interrupts occurred.																																
	POINT 0 = None of the PCM out channel interrupts occurred. 1 = One of the PCM out channel interrupts occurred. When the specific interrupt is cleared, this bit will be cleared (interruptible).																																
	1 = One of the PCM out channel interrupts occurred. When the specific interrupt is cleared, this bit will be cleared (interruptible). PCM In Interrupt (PIINT)																																
	1 = One of the PCM out channel interrupts occurred. When the specific interrupt is cleared, this bit will be cleared (interruptible).																																
										Wh	en t	he s	pec	ific	inter	rup	t is o	clea	red,	this	bit	will	be c	lea	red ((inte	errup	tible	€).				
		4	:3				-	-		rese		_																					
			2			M	101	INT		1 =	No Or	ne o	of th	e Mo	lode oden	m c	ut c	han ann	el in	terr	upts	ts oc	curre	ed.		(inte	errup	otible	e).				
										Мо	dem	-In I	nte	rrup	t (M	IINT)												-				
			1			N	ИIII	NT														000 000											
													•									will	be c	lea	red ((inte	errup	tible	€).				
				_			_	_	_						s Ch	`	,	nterr	upt	(GS	CI)		_	_		_			_			_	
			0			(GS	CI		0 = 1 =	Bit	0 о	fslo	t 12		et.	This				hat	one	of th	ne G	Pl's	s ch	ang	ed s	tate	and	d tha	t th	е
										The	bit	is c	ear	ed b	y so	oftw	are	writi	ng a	a "1'	' to '	his I	oit lo	ocat	ion ((inte	errup	otible	э).				

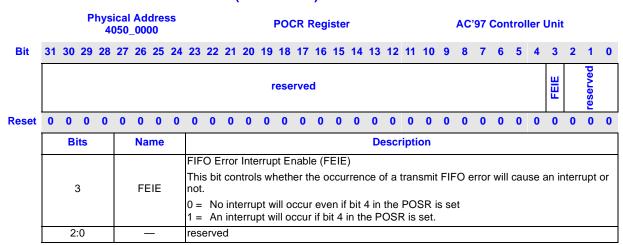
13.8.3.3 PCM-Out Control Register (POCR)

Table 13-9. POCR Bit Definitions (Sheet 1 of 2)

			P			Ad _000		SS						200	RF	legi	stei							AC'	97 (Con	trol	ler I	Unit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													r	ese	rved	ł													FEIE		reserved	
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	escr	ipti	on										
		31	:4			_			rese	erve	ed																					



Table 13-9. POCR Bit Definitions (Sheet 2 of 2)



13.8.3.4 PCM-In Control Register (PICR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 13-10. PICR Bit Definitions

			PI	hysi 40	ical)50_			S						PIC	RR	egi	ster							AC'	97 (Con	trol	ler I	Unit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
·													r	ese	rve	d													FEIE		reserved	
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
•		Bi	ts			Na	me												De	scr	iptic	n										
,		31	:4			_			rese	erve	d																					
•									FIF	O E	rror	Inte	rrup	ot Eı	nabl	e (F	EIE)														
		3	3			FE	ΙE		This		con	trol	s wł	neth	er th	ne o	ccur	ren	ce o	faı	recei	ive I	FIF	O er	ror	will	caus	se a	n int	erru	pt o	r
																					PISF		set									
		2:	:0			-	-		rese	erve	d																					



13.8.3.5 PCM-Out Status Register (POSR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 13-11. POSR Bit Definitions

		Physical Address 4050_0010 POSR 1 30 29 28 27 26 25 24 23 22 21 20 19 18 1 reserved															ste	•						AC'	'97(Con	trol	ler l	Jnit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	serv	ed													FIFOE		reserved		
Reset	0	D O O O O O O O O O															0	0	0	0	0	0	0	0	0	0	0	0				
		Bi	ts			Na	me												De	escr	ipti	on										
		31	:5			_	_		rese	erve	d																					
		4	1			FIF	OE		0 = 1 =	No A t this inc	s cas rem	nsm smit se, ento lo m	it FIF The the ed.7 nore	FO e last This val	rror vali cou id b avai	occ d sa ld h uffei able	urre impl app dat bu	d. T e is en d a av	his rep ue t aila	bit i etiti to: able conti	rolle	sen	t ou	it an	d th	е ро	und ointe ndw	ers a	are n	not		
		3:	0			_	_		rese	erve	d																					

13.8.3.6 PCM_In Status Register (PISR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 13-12. PISR Bit Definitions

			P			Ad _00′	dres 14	S						PIS	RR	legi	ster							AC	97 (Con	trol	ler l	Jnit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	serv	red													FIFOE		reserved		
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	its			Na	me												De	esci	riptio	on										
		31	:5			_			rese	erve	d																					
		4	4			FIF	OE		0 = 1 =	No A r cas wri exc	ece se, t itten cess	ceiv ive he f into sive n tin	e F FIF(FIF(the bar ne.	, IFO O er O po e FI ndwi	ror ointe FO dth	occi ers d and requ	urred lon't will uirer	d. Ti inci be I nen	remo	oit is ent, Thi nd h	s set the s co nenc	inco uld l	mir hap	ng d pen	ata f due	rom to I	the	AC A co	-link	is r ller l	ot navi	ng
		3	:0			-	_		rese	erve	d																					



13.8.3.7 CODEC Access Register (CAR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 13-13. CAR Bit Definitions

	Physical Address 4050_0020 CAR Register AC'97 Controller Unit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 reserved 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																															
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																CAIP															
Reset																0	0	0														
		Bi	ts			Na	me												De	escr	ipti	on										
		31	:1			-	_		res	erve	d																					
		()			CA	ΝP			No res co thi Ind	is r cyc serv mplo s bit dica	ead cle is es t ete, loc tes t	by s s in he r this atio that tly a	progright bit n if and	ware gres for is a it de	e to ss ar the utor cide r dri	che nd th soft nation es no ver i	ck where a ware cally of to	ct of e dri clea per erfor	f realiver ared forn	iding to p I. So n a 0 g a	g the erfo oftwa COE	e re rm f are DEC	giste the can O IO	er se IO c clea cycl	is contents the second in the second is contents to the second in the se	his b . Or is bi ter h	oit to ice t t by navii s the	o '1'. the d writing re e A0	Thi cycle ting ead C-lin	s e is a '0 this k a)' to s bit. nd

13.8.3.8 PCM Data Register (PCDR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 13-14. PCDR Bit Definitions

			P			Ad _004		SS					I	PCE	R F	Regi	istei							AC'	97 (Con	trol	ler I	Unit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•	PCM_RDATA 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																	PC	M_I	LDA	TA											
Reset	PCM_RDATA															0	0	0														
		В	its			Na	me												De	scr	ipti	on										
•		31	:16		PC	CM_	RDA	AΤΑ	PCI	M ri	ght-	chai	nnel	dat	а																	
		15	5:0		PC	CM_	LDA	ΛTΑ	PCI	M le	ft-cl	hanr	nel d	lata																		

Writing a 32-bit sample to this register updates the data into the PCM Transmit FIFO. Reading this register gets a 32-bit sample from the PCM Receive FIFO.



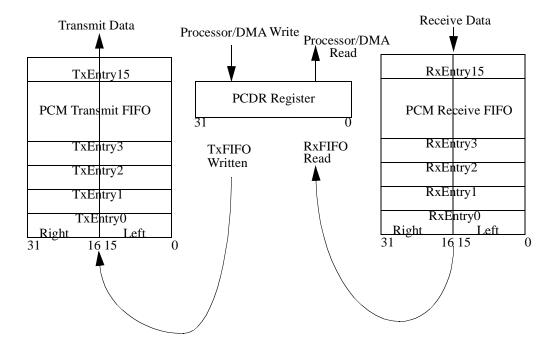
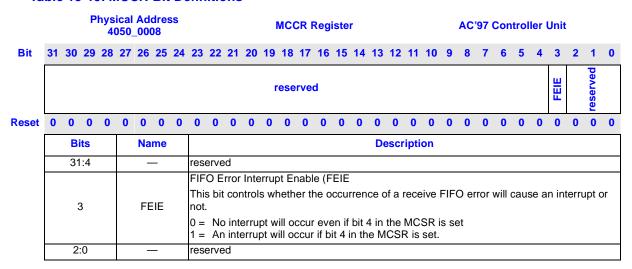


Figure 13-9. PCM Transmit and Receive Operation

13.8.3.9 Mic-In Control Register (MCCR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

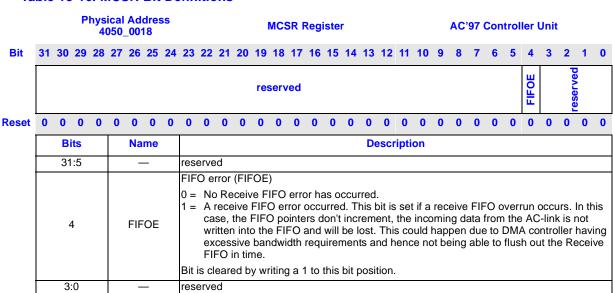
Table 13-15. MCCR Bit Definitions



13.8.3.10 Mic-In Status Register (MCSR)



Table 13-16, MCSR Bit Definitions



13.8.3.11 Mic-In Data Register (MCDR)

The Mic-In Data Register is a read-only register. A write to this register has no effect. A read to this register gets a 32-bit sample from the Mic-in Receive FIFO.

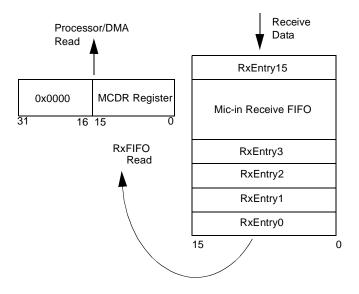
This is a read-only register. Ignore reads from reserved bits.

Table 13-17, MCDR Bit Definitions

			P	hysi 4(Ad _006		SS						MCI	OR F	Regi	iste	r						AC'	97 (Con	trol	ler l	Jnit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							r	ese	rve	d													MI	C_II	N_D	ΑT						
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	its			Na	me												De	scr	ipti	on										
		31:	:16			-	_		rese	erve	d																					
		15	5:0		MI	C_II	N_D	AT	mic	-in c	data																					



Figure 13-10. Mic-in Receive-Only Operation



13.8.3.12 Modem-Out Control Register (MOCR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

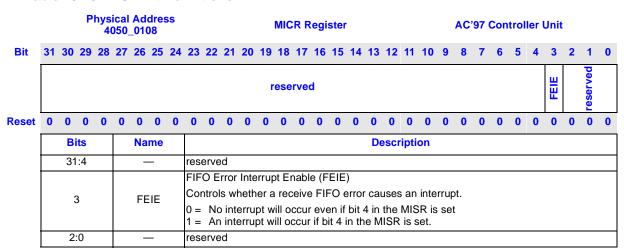
Table 13-18. MOCR Bit Definitions

			Pl	•	ical 050 ₋		dres 00	SS						MO	CR I	Reg	iste	r								AC	'97					
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													r	ese	erve	d													FEIE		reserved	
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
	Bits Name Description													esci	ripti	on																
		31	:4			_	_		res	erve	d																					
									FIF	O E	rror	Inte	rrup	ot E	nabl	e (F	EIE)														
		3	3			FE	ΙE				con	trol	s wł	neth	er th	ne o	ccui	ren	ce c	of a	tran	smit	FIF	O e	rror	will	cau	se a	an ir	terr	upt	or
	FEIE not. 0 = No interrupt will occur even if bit 4 in the M 1 = An interrupt will occur if bit 4 in the MOSR																et															
		2:	0			_	_		rese	erve	:d																					

13.8.3.13 Modem-In Control Register (MICR)



Table 13-19. MICR Bit Definitions



13.8.3.14 Modem-Out Status Register (MOSR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 13-20. MOSR Bit Definitions

			Ph		ical / 050_			ss					ı	MOS	SR F	Reg	iste	r						AC	97 (Con	trol	ler l	Jnit			
Bit	31	30 2	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	serv	ed													FIFOE		reserved		
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bits	5			Na	me												De	esci	ipti	on										
		31:	5			-	-		rese	erve	d																					
		4				FIF	OE		1 =	No A t thi inc	trans rans s ca rem c. N	nsm smit se, ento No m	it FIF the ed.T nore	FO e last his val	rror vali cou id b	occ d sa ld h uffe able	urre amp app da bu	e occed. The decision of the d	his rep lue /aila /A c	bit i etiti to: able cont	rolle	sen tran:	it ou	ut ar	nd th	е ро	ointe	ers a	are r	not		
		3:0	1			_	_		rese		_	u D	y wi	ıuııç	ja	ιο	11115	bit p	0051	uon	•											_

13.8.3.15 Modem-In Status Register (MISR)



Table 13-21. MISR Bit Definitions

			P	hysi 4(Add _011		S						MIS	R R	egi	ster							AC'	97 (Con	trol	ler l	Jnit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	serv	ed													FIFOE		reserved		
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		В	its			Na	me												De	escr	iptio	on										
		31	:5			_	_		rese	rve	d																					
		4	4			FIF	OE		FIF(0 = 1 =	No A r cas wri exc	Rece se, t itten cess	ceiv ive l he f into sive n tin	e F FIF(FIF(the bar ne.	IFO O er O po e FI ndwi	ror o ointe FO a dth	rs d and requ	urred on't will uirer	d. Ti incr be l nen	nis b eme ost. ts ai	oit is ent, This nd h	the s cor enc	inco uld l	mir nap	ig da pen	FIF ata f due able	rom to I	the DM/	AC A co	-link ntro	is r ller l	not havi	ng
		3	:0			-			rese	rve	d																					

13.8.3.16 Modem Data Register (MODR)

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 13-22. MODR Bit Definitions

			P			Add _014		SS						MOI	OR F	Regi	iste	r						AC	97	Con	trol	ler I	Unit			
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							r	ese	rvec	t													MC	DE	M_E	DAT						
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi												De	escr	ipti	on															
	31:16 —								rese	erve	d																					
	15:0 MODE					DEI	M_C	DAT	Mod	dem	da	ta																				

A 32-bit sample write to this register updates the data into the Modem Transmit FIFO. A read to this register gets a 32-bit sample from the Modem Receive FIFO.



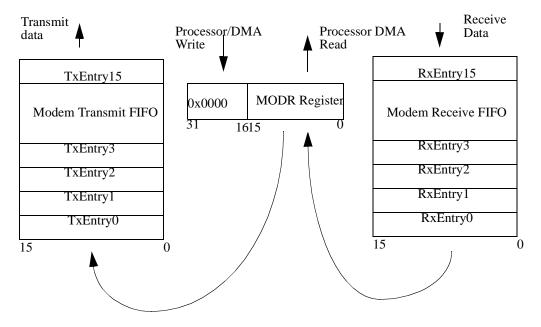


Figure 13-11. Modem Transmit and Receive Operation

13.8.3.17 Accessing CODEC Registers

Each CODEC has up to sixty-four 16-bit registers that are addressable internal to the CODEC at half-word boundaries (16-bit boundaries). Because the processor only supports internal register accesses at word boundaries (32-bit boundaries), software must select the one of the following formulas to translate a 7-bit CODEC address into a 32-bit processor address:

- Processor physical address for a Primary Audio CODEC
 = 0x4050-0200 + Shift Left Once(Internal 7-bit CODEC Register Address)
- Processor physical address for a Secondary Audio CODEC
 = 0x4050-0300 + Shift Left Once(Internal 7-bit CODEC Register Address)
- Processor physical address for a Primary Modem CODEC = 0x4050-0400 + Shift_Left_Once(Internal 7-bit CODEC Register Address)
- Processor physical address for a Secondary Modem CODEC
 = 0x4050-0500 + Shift_Left_Once(Internal 7-bit CODEC Register Address)

In the equations, Shift_Left_Once() shifts the 7-bit CODEC address left by one bit and shifts a 0 to the LSB. The address translations are shown in Table 13-23.



Table 13-23. Address Mapping for CODEC Registers (Sheet 1 of 2)

Audio CODEC Audio CODEC	
0x00 0x4050_0200 0x4050_0300	0x4050_0400
0x02 0x4050_0204 0x4050_0304	0x4050_0404 0x4050_0504
0x04 0x4050_0208 0x4050_0308	0x4050_0408
0x06	0x4050_040C
0x08 0x4050_0210 0x4050_0310	0x4050_0410 0x4050_0510
0x0A 0x4050_0214 0x4050_0314	0x4050_0414 0x4050_0514
0x0C 0x4050_0218 0x4050_0318	0x4050_0418
0x0E 0x4050_021C 0x4050_031C	0x4050_041C
0x10 0x4050_0220 0x4050_0320	0x4050_0420 0x4050_0520
0x12 0x4050_0224 0x4050_0324	0x4050_0424
0x14 0x4050_0228 0x4050_0328	0x4050_0428
0x16	0x4050_042C
0x18 0x4050_0230 0x4050_0330	0x4050_0430
0x1A 0x4050_0234 0x4050_0334	0x4050_0434
0x1C 0x4050_0238 0x4050_0338	0x4050_0438
0x1E 0x4050_023C 0x4050_033C	0x4050_043C
0x20 0x4050_0240 0x4050_0340	0x4050_0440
0x22 0x4050_0244 0x4050_0344	0x4050_0444
0x24 0x4050_0248 0x4050_0348	0x4050_0448
0x26 0x4050_024C 0x4050_034C	0x4050_044C
0x28 0x4050_0250 0x4050_0350	0x4050_0450
0x2A 0x4050_0254 0x4050_0354	0x4050_0454
0x2C 0x4050_0258 0x4050_0358	0x4050_0458
0x2E 0x4050_025C 0x4050_035C	0x4050_045C
0x30 0x4050_0260 0x4050_0360	0x4050_0460
0x32 0x4050_0264 0x4050_0364	0x4050_0464
0x34 0x4050_0268 0x4050_0368	0x4050_0468
0x36 0x4050_026C 0x4050_036C	0x4050_046C
0x38	0x4050_0470
0x3A 0x4050_0274 0x4050_0374	0x4050_0474
0x3C 0x4050_0278 0x4050_0378	0x4050_0478
0x3E 0x4050_027C 0x4050_037C	0x4050_047C
0x40 0x4050_0280 0x4050_0380	0x4050_0480
0x42 0x4050_0284 0x4050_0384	0x4050_0484



Table 13-23. Address Mapping for CODEC Registers (Sheet 2 of 2)

7-bit CODEC Address	Processor Physical Address for a Primary Audio CODEC	Processor Physical Address for a Secondary Audio CODEC	Processor Physical Address for a Primary Modem CODEC	Processor Physical Address for a Secondary Modem CODEC
0x44	0x4050_0288	0x4050_0388	0x4050_0488	0x4050_0588
0x46	0x4050_028C	0x4050_038C	0x4050_048C	0x4050_058C
0x48	0x4050_0290	0x4050_0390	0x4050_0490	0x4050_0590
0x4A	0x4050_0294	0x4050_0394	0x4050_0494	0x4050_0594
0x4C	0x4050_0298	0x4050_0398	0x4050_0498	0x4050_0598
0x4E	0x4050_029C	0x4050_039C	0x4050_049C	0x4050_059C
0x50	0x4050_02A0	0x4050_03A0	0x4050_04A0	0x4050_05A0
0x52	0x4050_02A4	0x4050_03A4	0x4050_04A4	0x4050_05A4
0x54	0x4050_02A8	0x4050_03A8	0x4050_04A8	0x4050_05A8
0x56	0x4050_02AC	0x4050_03AC	0x4050_04AC	0x4050_05AC
0x58	0x4050_02B0	0x4050_03B0	0x4050_04B0	0x4050_05B0
0x5A	0x4050_02B4	0x4050_03B4	0x4050_04B4	0x4050_05B4
0x5C	0x4050_02B8	0x4050_03B8	0x4050_04B8	0x4050_05B8
0x5E	0x4050_02BC	0x4050_03BC	0x4050_04BC	0x4050_05BC
0x60	0x4050_02C0	0x4050_03C0	0x4050_04C0	0x4050_05C0
0x62	0x4050_02C4	0x4050_03C4	0x4050_04C4	0x4050_05C4
0x64	0x4050_02C8	0x4050_03C8	0x4050_04C8	0x4050_05C8
0x66	0x4050_02CC	0x4050_03CC	0x4050_04CC	0x4050_05CC
0x68	0x4050_02D0	0x4050_03D0	0x4050_04D0	0x4050_05D0
0x6A	0x4050_02D4	0x4050_03D4	0x4050_04D4	0x4050_05D4
0x6C	0x4050_02D8	0x4050_03D8	0x4050_04D8	0x4050_05D8
0x6E	0x4050_02DC	0x4050_03DC	0x4050_04DC	0x4050_05DC
0x70	0x4050_02E0	0x4050_03E0	0x4050_04E0	0x4050_05E0
0x72	0x4050_02E4	0x4050_03E4	0x4050_04E4	0x4050_05E4
0x74	0x4050_02E8	0x4050_03E8	0x4050_04E8	0x4050_05E8
0x76	0x4050_02EC	0x4050_03EC	0x4050_04EC	0x4050_05EC
0x78	0x4050_02F0	0x4050_03F0	0x4050_04F0	0x4050_05F0
0x7A	0x4050_02F4	0x4050_03F4	0x4050_04F4	0x4050_05F4
0x7C	0x4050_02F8	0x4050_03F8	0x4050_04F8	0x4050_05F8
0x7E	0x4050_02FC	0x4050_03FC	0x4050_04FC	0x4050_05FC



13.9 AC'97 Register Summary

All AC'97 registers are word-addressable (32 bits wide) and increment in units of 0x00004. The registers in the CODEC are half-word addressable (16 bits wide), and increment in units of 0x00002. These register sets are mapped in the address range of 0x4050_000 through 0x405F_FFFF.

Table 13-24. Register Mapping Summary

Address	Name	Description
0x4050_0000	POCR	PCM Out Control Register
0x4050_0004	PICR	PCM In Control Register
0x4050_0008	MCCR	Mic In Control Register
0x4050_000C	GCR	Global Control Register
0x4050_0010	POSR	PCM Out Status Register
0x4050_0014	PISR	PCM In Status Register
0x4050_0018	MCSR	Mic-In Status Register
0x4050_001C	GSR	Global Status Register
0x4050_0020	CAR	CODEC Access Register
0x4050_0024 - 0x4050_003C	_	reserved
0x4050_0040	PCDR	PCM FIFO Data Register
0x4050_0044 - 0x4050_005C	_	reserved
0x4050_0060	MCDR	Mic-in FIFO Data Register
0x4050_0064 - 0x4050_00FC	_	reserved
0x4050_0100	MOCR	Modem-Out Control Register
0x4050_0104	_	reserved
0x4050_0108	MICR	Modem-In Control Register
0x4050_010C	_	reserved
0x4050_0110	MOSR	Modem-Out Status Register
0x4050_0114	_	reserved
0x4050_0118	MISR	Modem-In Status Register
0x4050_011C - 0x4050_013C	_	reserved
0x4050_0140	MODR	Modem FIFO Data Register
0x4050_0144 - 0x4050_01FC	_	reserved
(0x4050_0200 - 0x4050_02FC)	_	Primary Audio CODEC registers
with all in increments of 0x00004		1 mary rudio GGB LO registers
(0x4050_0300 - 0x4050_03FC) with all in increments of 0x00004	_	Secondary Audio CODEC registers
(0x4050_0400 - 0x4050_04FC) with all in increments of 0x0000_0004	_	Primary Modem CODEC registers
(0x4050_0500 - 0x4050_05FC) with all in increments of 0x00004	_	Secondary Modem CODEC registers





Inter-Integrated-Circuit Sound (PS) Controller

14

 I^2S is a protocol for digital stereo audio. The I^2S Controller (I2SC) functional block for the PXA255 processor controls the I^2S link (I2SLINK), which is a low-power four-pin serial interface for stereo audio. The I^2S interface and the Audio CODEC '97 (AC'97) interface may not be used at the same time.

14.1 Overview

The I2SC consists of buffers, status and control registers, serializers, and counters for transferring digitized audio between the processor system memory and an external I²S CODEC.

For playback of digitized audio or production of synthesized audio, the I2SC retrieves digitized audio samples from processor system memory and sends them to a CODEC through the I2SLINK. The external digital-to-analog converter in the CODEC then converts the audio samples into an analog audio waveform.

For recording of digitized audio, the I2SC receives digitized audio samples from a CODEC (through the I2SLINK) and stores them in processor system memory.

The I²S controller supports the normal-I²S and the MSB-Justified-I²S formats. Four, or optionally five, pins connect the controller to an external CODEC:

- A bit-rate clock, which can use either an internal or an external source.
- A formatting or "Left/Right" control signal.
- Two serial audio pins, one input and one output.
- The bit-rate clock, an optional system clock also sent to the CODEC by the I2SC.

The I^2S data can be stored to and retrieved from system memory either by the DMA controller or by programmed I/O.

For I²S systems, additional pins are required to control the external CODEC. Some CODECs use an L3 control bus, which requires 3 signals — L3_CLK, L3_DATA, and L3_MODE — for writing bytes into the L3-bus register. The I2SC supports the L3 bus protocol via software control of the general-purpose I/O (GPIO) pins. The I2SC does not provide hardware control for the L3 bus protocol.

Two similar protocols exist for transmitting digitized stereo audio over a serial path: Normal-I²S and MSB-Justified-I²S. Both work with a variety of clock rates, which can be obtained by dividing the PLL clock by a programmable divider, or from an external clock source. For further details regarding clock rates, see Table 14-2.



14.2 Signal Descriptions

SYSCLK is the clock on which all other clocks in the I²S unit are based. SYSCLK generates a frequency between approximately 2 MHz and 12.2 MHz by dividing down the PLL clock with a programmable divisor. This frequency is always 256 times the audio sampling frequency. SYSCLK is driven out of the processor system only if BITCLK is configured as an output.

BITCLK supplies the serial audio bit rate, which is the basis for the external CODEC bit-sampling logic. BITCLK is one-quarter the frequency of SYSCLK and is 64 times the audio sampling frequency. One bit of the serial audio data sample is transmitted or received each BITCLK period. A single serial audio sample comprises a "left" and "right" signal, each containing either 8, 16 or 32 bits.

SYNC is BITCLK divided by 64, resulting in an 8 kHz to 48 kHz signal. The state of SYNC is used to denote whether the current serial data samples are "Left" or "Right" channel data.

The SDATA_IN and SDATA_OUT data pins are used to send/receive the serial audio data to/from the CODEC.

Table 14-1 lists the signals between the I²S and an external CODEC device.

Table 14-1. External Interface to CODEC

Name	Direction	Description
GP32/SYSCLK	0	System Clock = BITCLK * 4 used by the CODEC only.
GP28/BITCLK	I or O	bit-rate clock = SYNC * 64
GP31/SYNC	0	Left/Right identifier
GP30/SDATA_OUT	0	Serial audio output data to CODEC
GP29/SDATA_IN	I	Serial audio input data from CODEC

BITCLK can be configured either as an input or as an output. To program the direction, follow these steps:

- 1. Program SYSUNIT's GPIO Direction Register (GPDR). See Section 4.1.3.2, "GPIO Pin Direction Registers (GPDR0, GPDR1, GPDR2)" on page 4-8 for details regarding the GPDR.
- 2. Program SYSUNIT's GPIO Alternate Function Select Register (GAFR). See Section 4.1.3.6, "GPIO Alternate Function Register (GAFR0_L, GAFR0_U, GAFR1_L, GAFR1_U, GAFR2_L, GAFR2_U)" on page 4-16 for details regarding the GAFR.
- 3. Program the BCKD bit in the I2SC's Serial Audio Control Register. See Section 14.6.1, "Serial Audio Controller Global Control Register (SACR0)" for more details.

Note: Modifying the status of the SACR0[BCKD] bit during normal operation can cause jitter on the BITCLK and can affect serial activity.

If BITCLK is an output, SYSCLK must be configured as an output. If BITCLK is supplied by the CODEC, the GPIO pin GP32 can be used for an alternate function. To configure the pin as an output, follow these steps:

1. Program SYSUNIT's GPIO Direction Register (GPDR). See Section 4.1.3.2, "GPIO Pin Direction Registers (GPDR0, GPDR1, GPDR2)" on page 4-8 for details regarding the GPDR.



 Program SYSUNIT's GPIO Alternate Function Select Register (GAFR). See Section 4.1.3.6, "GPIO Alternate Function Register (GAFR0_L, GAFR0_U, GAFR1_L, GAFR1_U, GAFR2_L, GAFR2_U)" on page 4-16 for details regarding the GAFR.

To configure SYNC and SDATA_OUT as outputs, follow these steps:

- 1. Program SYSUNIT's GPIO Direction Register (GPDR). See Section 4.1.3.2, "GPIO Pin Direction Registers (GPDR0, GPDR1, GPDR2)" on page 4-8 for details regarding the GPDR.
- 2. Program SYSUNIT's GPIO Alternate Function Select Register (GAFR). See Section 4.1.3.6, "GPIO Alternate Function Register (GAFR0_L, GAFR0_U, GAFR1_L, GAFR1_U, GAFR2_L, GAFR2_U)" on page 4-16 for details regarding the GAFR.

To configure SDATA_IN as an input, follow these steps:

- 1. Program SYSUNIT's GPIO Direction Register (GPDR). See Section 4.1.3.2, "GPIO Pin Direction Registers (GPDR0, GPDR1, GPDR2)" on page 4-8 for details regarding the GPDR.
- Program SYSUNIT's GPIO Alternate Function Select Register (GAFR). See Section 4.1.3.6, "GPIO Alternate Function Register (GAFR0_L, GAFR0_U, GAFR1_L, GAFR1_U, GAFR2_L, GAFR2_U)" on page 4-16 for details regarding the GAFR.

14.3 Controller Operation

The I²S Controller (I2SC) can be accessed either by the processor or by the DMA controller.

The processor uses programmed I/O instructions to access the I2SC and can access the following types of data:

- I2SC registers data All registers are 32 bits wide and are aligned to word boundaries. See Section 14.6 for further details.
- I2SC FIFO data An entry is placed into the Transmit FIFO by writing to the I2SC's Serial Audio Data register (SADR). Writing to SADR updates a Transmit FIFO entry. Reading SADR flushes out a Receive FIFO entry.
- I²S CODEC data The CODEC registers can be accessed through the L3 bus. The L3 bus operation is emulated by software controlling 3 GPIO pins.

The DMA controller can only access the FIFOs. Accesses are made through the SADR, as explained in the previous paragraph. The DMA controller accesses FIFO data in blocks of 8, 16, or 32 bytes. The DMA controller responds to the following DMA requests made by the I2SC:

- The Transmit FIFO request is based on the transmit threshold (TFTH) setting and is asserted if
 the Transmit FIFO has less than TFTH+1 entries. See Table 14-3 for further details regarding
 TFTH.
- The Receive FIFO request is based on the receive threshold (RFTH) setting and is asserted if the Receive FIFO has RFTH+1 or more entries. See Table 14-3 for further details regarding RFTH.

14.3.1 Initialization

1. Set the BITCLK direction by programming the SYSUNIT's GPIO Direction register, the SYSUNIT's GPIO Alternate Function Select register, and bit 2 of the I2SC's Serial Audio Controller Global Control Register (SACR0).



- 2. Choose between Normal I²S or MSB-Justified modes of operation. This can be done by programming bit 0 of Serial Audio Controller I²S/MSB-Justified Control Register (SACR1). For further details, see Section 14.6.2.
- 3. Optional: Programmed I/O may be used for priming the Transmit FIFO with a few samples (ranging from 1 to 16). If the I2SLINK is enabled with an empty Transmit FIFO, a Transmit Under-run error bit will be set in the Status register. For further details, see Section 14.6.3. This is hence an optional step, which prevents such an error. If Step 3 is not executed, then Programmed I/O must clear the Transmit Under-run status bit by setting bit 5 of the Interrupt Clear Register. For further details, see Section 14.6.5.
- 4. The following control bits can be simultaneously programmed in the I2SC's Serial Audio Controller Global Control register (SACR0):
 - a. Enable I2SLINK by setting the ENB bit (bit-0) of SACR0.
 - b. Maintain BITCLK direction as programmed in Step1. Modifying BITCLK direction will glitch the clock and affect I2SLINK activity.
 - c. Program transmit and receive threshold levels by programming the TFTH and RFTH bits of SACR0[11:8] and SACR0(15:12), respectively. See Section 14.6.1.2, regarding permitted threshold levels.

Once the I2SLINK is enabled, frames filled with 0s will be transmitted if the Transmit FIFO is still empty. This will set a Transmit Under-run status bit in SASR0. Step 3 can be executed to avoid this error condition. Valid data is sent across the I2SLINK after filling the Transmit FIFO with at least one sample. One sample consists of a 32-bit value with 16 bits each dedicated to a left and a right value.

Enabling the I2SLINK will also cause zeros to be recorded by the I2SC until the CODEC sends in valid data.

Enabling the I2SLINK also enables transmit and receive DMA Requests.

14.3.2 Disabling and Enabling Audio Replay

Audio transmission is enabled automatically when the I2SC is enabled. Transmission, or replay, can be stopped by asserting the DRPL bit of the SACR1 Register. For more details, see Section 14.6.2.

Asserting the DRPL bit in SACR1 has the following effects:

- All I2SLINK replay activity is disabled. The frame or data sample, in the midst of which the
 replay is disabled, will have invalid data (some data bits will be over-written with zeros). To
 avoid this, disable replay only after the transfer of valid data. In this case, frames with zeros
 are transmitted.
- 2. Transmit FIFO pointers are reset to zero.
- 3. Transmit FIFO fill-level is reset to zero.
- 4. Zeros are transmitted across the I2SLINK.
- 5. Transmit DMA requests are disabled.

14.3.3 Disabling and Enabling Audio Record

Audio recording is enabled automatically when the I2SC is enabled. Recording can also be stopped by asserting the DREC bit of the SACR1 Register. For more details, see Section 14.6.2.



Asserting the DREC bit in SACR1 has the following effects:

- I2SLINK recording activity is disabled. The frame or data sample, in the midst of which the
 recording is disabled, could have invalid data (some data bits will be over-written with zeros).
 To avoid this, disable record only after the transfer of valid data.
- 2. Receive FIFO pointers are reset to zero.
- 3. Receive FIFO fill-level is reset to zero.
- 4. Any read operations by the DMA/CPU are returned with zeros.
- 5. Receive DMA requests are disabled.

14.3.4 Transmit FIFO Errors

A status bit is set during Transmit Under-run conditions. If enabled, this can trigger an interrupt. For further details, see Section 14.6.3, Section 14.6.6 and Section 14.6.5. During Transmit Under-run conditions, the last valid sample is continuously sent out across the I2SLINK. Transmit Under-run can occur under the following conditions:

- 1. Valid transmit data is still available in memory, but the DMA controller starves the Transmit FIFO, as it is busy servicing other higher-priority peripherals.
- 2. The DMA controller has transferred all valid data from memory to the Transmit FIFO.

During the second condition, the last valid sample is continuously sent across the I2SLINK until the I2SC is turned off by disabling the SACR0[ENB] bit.

14.3.5 Receive FIFO Errors

A status bit is set during Receive Over-run conditions. If enabled, this can trigger an interrupt. For further details, see Section 14.6.3, Section 14.6.6 and Section 14.6.5. During Receive Over-run conditions, data sent by the CODEC is lost (will not be recorded).

14.3.6 Trailing Bytes

When the CODEC has completed transmitting valid data, zeros will be recorded by the I2SC, and this will continue until the unit is turned off by disabling the SACR0[ENB] bit.

If the total buffer size of the received data is less than a factor of the receive threshold, zeross will be recorded. A receive DMA request is made when the programmed threshold is reached.

14.4 Serial Audio Clocks and Sampling Frequencies

The BITCLK is the rate at which audio data bits enter or leave the I2SLINK. If BITCLK is an output, SYSCLK is used by the CODEC to run delta sigma ADC operations.

BITCLK can be supplied either by the CODEC or by an internal PLL. If supplied internally, BITCLK and SYSCLK are configured as output pins, and both are supplied to the CODEC. If BITCLK is supplied by the CODEC, then it is configured as an input pin. In this case, the SYSCLK's GPIO pin can be used for an alternate function.



The BITCLK, as shown in Table 14-2, is different for different sampling frequencies. If the BITCLK is chosen as an output, the Audio Clock Divider Register divides the 147.46MHz PLL clock to generate the SYSCLK. The SYSCLK is further divided by four to generate the BITCLK. The sampling frequency is the frequency of the SYNC signal, which is generated by dividing the BITCLK by 64. See Section 14.6.4, for further details about the register.

A sampling rate of 48kHz supports MPEG2 and MPEG4. A rate of 44.1kHz supports MP3.

Table 14-2. Supported Sampling Frequencies

Audio Clock Divider	SYSCLK =	BITCLK =	SYNC or Sampling frequency =
Register (31:0)	147.6 MHz / (SADIV)	SYSCLK / 4	BITCLK / 64
0x0000_000C	12.288 MHz	3.072 MHz	48.000 kHz (closest std = 48 kHz)
0x0000_000D	11.343 MHz	2.836 MHz	44.308 kHz (closest std = 44.1 kHz)
0x0000_001A	5.671 MHz	1.418 MHz	22.154 kHz (closest std = 22.05 kHz)
0x0000_0024	4.096 MHz	1.024 MHz	16.000 kHz (closest std = 16.00 kHz)
0x0000_0034	2.836 MHz	708.92 kHz	11.077 kHz (closest std = 11.025 kHz)
0x0000_0048	2.048 MHz	512.00 kHz	8.000 kHz (closest std = 8.00 kHz)

14.5 Data Formats

14.5.1 FIFO and Memory Format

FIFO buffers are 16 levels deep and 32 bits wide. This stores 32 samples per channel in each direction.

Audio data is stored with two samples (Left + Right) per 32-bit word, even if samples are smaller than 16 bits. The Left channel data occupies bits [15:0], while the Right channel data uses bits [31:16] of the 32-bit word. Within each 16-bit field, the audio sample is left-justified, with unused bits packed as zeroes on the right-hand (LSB) side.

In memory, the mapping of stereo samples is the same as in the FIFO buffers. However, single-channel audio occupies a full 32-bit word per sample, using either the upper or lower half of the word, depending on whether it's considered a Left or Right sample.

14.5.2 I²S and MSB-Justified Serial Audio Formats

I²S and MSB-Justified are similar protocols for digitized stereo audio transmitted over a serial path.

The BITCLK supplies the serial audio bit rate, the basis for the external CODEC bit-sampling logic. Its frequency is 64 times the audio sampling frequency. Divided by 64, the resulting 8 kHz to 48 kHz signal signifies timing for Left and Right serial data samples passing on the serial data paths. This Left/Right signal is sent to the CODEC on the SYNC pin. Each phase of the Left/Right signal is accompanied by one serial audio data sample on the data pins SDATA_IN and SDATA_OUT.



Figure 14-1 and Figure 14-2 provide timing diagrams that show formats for I²S and MSB-justified modes of operations.

Data is transmitted and received in frames of 64 BITCLK cycles. Each frame consists of a Left sample and a Right sample. Each frame holds 16-bits of valid sample data (shown in the figures) and 16-bits of padded zeros (not shown in the figures). The transmit and receive FIFOs only hold valid sample data (not padded zero data).

In the Normal I²S mode, the SYNC is low for the Left sample and high for the Right sample. Also, the MSB of each data sample lags behind the SYNC edges by one BITCLK cycle.

In the MSB-Justified mode, the SYNC is high for the Left sample and low for the Right sample. Also, the MSB of each data sample is aligned with the SYNC edges.

Figure 14-1. I²S Data Formats (16 bits)

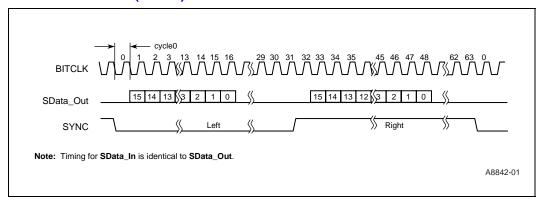
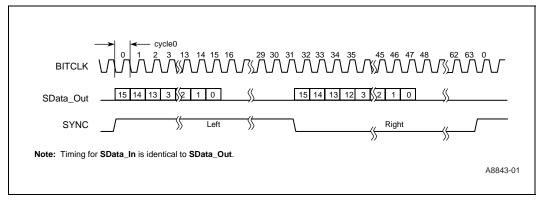


Figure 14-2. MSB-Justified Data Formats (16 bits





14.6 Registers

The I²S Controller registers are all 32-bit addressable, ranging from 0x4040_0000 through 0x404F_FFFF.

The I²S Controller has the following types of registers:

- Control registers are used to program common control, alternate mode specific control.
- The Data Register is used for Transmit and Receive FIFO accesses.
- The Status Register signals the state of the FIFO buffers and the status of the interface that is selected by the common control register.
- The Interrupt Registers include the Interrupt Mask Register, the Interrupt Clear Register, and the Interrupt Test Register.

14.6.1 Serial Audio Controller Global Control Register (SACR0)

SACRO, shown in Table 14-3, controls common I²S functions.

The ENB bit controls the I2SLINK, as follows:

- Clearing ENB to zero does the following:
 - disables any I2SLINK activity
 - resets all Receive FIFO pointers and also the counter that controls the I2SLINK
 - resets the Receive FIFO
 - does not affect the Transmit FIFO
 - the output pin SYNC will not toggle
 - de-asserts all DMA requests
 - any read accesses to the Data Register (SADR), by the processor, or by the DMA controller is returned with zeros
 - disables all interrupts.
- Setting ENB to one does the following:
 - enables I2SLINK activity
 - enables DMA requests.

Note: If ENB is toggled in the middle of a normal operation, the RST bit must also be set and cleared to reset all I2SC registers.

Note: The SACR0[ENB] control signal crosses clock domains. It is registered in an internal clock domain that is much faster than the BITCLK domain. It takes four BITCLK cycles and four internal clock cycles before SACR0[ENB] is conveyed to the slower BITCLK domain. If the control setting is modified at a rate faster than (4 BITCLK + 4 internal clock) cycles, the last updated value in this time frame is stored in a temporary register and is transferred to the BITCLK domain.



Table 14-3. SACR0 Bit Definitions

10	DIC	' '	T J.	U,	•	,,,,	Dit	DC		J 113	•																				
			P			al Ad 40_0		SS			Se				Contr ol Reg			Glob	al					I ² S	Со	ntro	ller				
Bit	31	30	29	28	2	7 26	25	24	23 22	21	2	20 19	18	1	7 16	15	1	14 13	12	11	10	9	8	7	6	5	4	3	2	1	0
							r	esei	ved								F	RFTH			Т	FTH		poviosor	DOA IDCD	STRF	EFWR	RST	BCKD	reserved	ENB
Reset	0	0	0	0	0	0	0	0	0 0	0	(0 0	0	C	0 0	0		1 1	1	0	1	1	1	0	0	0	0	0	0	0	0
		В	its			Na	me											D	esci	ipti	on										
		31	:16			-	_		reserv	ed																					
		15	5:12			RF	тн		Recei thresh Recei	old	va	lue m	inus	1.																	he
	Receive DMA Request asserted whenever the Receive FIFO has >= (RFTH+1) entries. Transmit FIFO interrupt or DMA threshold. Set to value 0 – 15. This value must be set to the threshold value minus 1. Transmit DMA Request asserted whenever the Transmit FIFO has < (TFTH+1) entries. 7:6 — reserved)																
		7	':6			-	_		reserv	ed																					
									Selec	Tra	เทร	mit o	r Re	cei	ive FI	FO	fo	r EFV	VR b	ase	d s	speci	al pu	urpo	se f	func	tion				
			5			ST	RF		0 = T 1 = F See T	ece	ive	FIF) is	sel	lected																
									This b	it er	nah	oles a	sne	cia	al nuri	ากระ	ا د	FIFO	Write	e/Re	ad	l fund	rtion								
			4			EF	WR		0 = S 1 = S See T	pec pec	ial ial	purp	ose I ose I	FIF	-0 W -0 W	rite/I	26	ead F	unct	ion i	s c	disab	led	•							
					<u> </u>				Reset							rioto	r	0.000	ont t	hio r		iotor	/C A	CD	٦١.						
			3			R	ST		0 = N 1 = F	ot r	es	et							ергі	1115 1	ey	listei	(SA	CK	J).						
									This b	it sp	ec	cifies	nput	/o	utput	dire	ct	tion of	BIT	CLK	(:										
			2			ВС	KD		0 = Ii 1 = C												ver	n out	to t	he C	OD	EC.					
			1			-			reserv	ed															_						
									Enabl	e I ² S	S f	unctio	n:																		
			0			EI	NB		0 = 1	SLI	NŁ	Κis d	sabl	ed	t																

14.6.1.1 Special purpose FIFO Read/Write function

1 = I2SLINK is enabled

As shown in Table 14-4, EFWR and STRF can be programmed for special purpose FIFO accesses. Under normal operating conditions, the processor or the DMA controller can only write to the Transmit FIFO and only read the Receive FIFO. Programming these bits allows the processor or the DMA controller to read and write both FIFOs.



Table 14-4. FIFO Write/Read table

EFWR	STRF	Description
0	х	Normal CPU/DMA Write/Read condition: A write access to the Data Register writes a Transmit FIFO entry. A read access to the Data Register reads out a Receive FIFO entry. I2SLINK reads from the Transmit FIFO and writes to the Receive FIFO.
1	0	 CPU or DMA only writes and reads Transmit FIFO: A write access to the Data Register writes a Transmit FIFO entry. A read access to the Data Register reads out a Transmit FIFO entry. I2SLINK cannot read the Transmit FIFO but can write to the Receive FIFO.
1	1	 CPU or DMA only writes and reads Receive FIFO: A write access to the Data Register writes a Receive FIFO entry. A read access to the Data Register reads out a Receive FIFO entry. I2SLINK can read the Transmit FIFO but cannot write to the Receive FIFO.

14.6.1.2 Suggested TFTH and RFTH for DMA servicing

The DMA controller can only be programmed to transfer 8, 16, or 32 bytes of data. This corresponds to 2, 4, or 8 FIFO samples. Table 14-5 shows the recommended TFTH and RFTH values to prevent Transmit FIFO over-run errors and Receive FIFO under-run errors.

Table 14-5. TFTH and RFTH Values for DMA Servicing

DMA Transfer Size	# of FIFO entries	TFTH	Value	RFTH	Value
		Min	Max	Min	Max
8 Bytes	2	0	14	1	15
16 Bytes	4	0	12	3	15
32 Bytes	8	0	8	7	15

14.6.2 Serial Audio Controller I²S/MSB-Justified Control Register (SACR1)

SACR1, shown in Table 14-6, specifically controls the I2S and MSB-Justified modes.

SACR1 bits DRPL, DREC, and AMSL cross clock domains. They are registered in an internal clock domain that is much faster than the BITCLK domain. It takes 4 BITCLK cycles and 4 internal clock cycles before these controls are conveyed to the slower BITCLK domain. If the above control settings are modified at a rate faster than (4 BITCLK + 4 internal clock) cycles, the last updated value in this time frame is stored in a temporary register and is transferred to the BITCLK domain.



Table 14-6. SACR1 Bit Definitions

			F		ical Ad 040_0											² S/MS gister	SB	-					I ² S	S Co	ntro	ller			
Bit	31	30	29	28	27 26	3 2	5 2	24 23	22	21 20	19	18	17	16	15	14 1	3	12	11	10	9	8	7	6	5	4	3	2 1	0
											rese	erved	ł												ENLBF	DRPL	DREC	reserved	AMSL
Reset	0	0	0	0	0 0	()	0 0	0	0 0	0	0	0	0	0	0)	0	0	0	0	0	0	0	0	0	0	0 0	0
		Ві	its		N	am	е										De	scr	ipti	on									
		31	:6			_		res	erve	ed																			
		į	5		ΕN	NLB	F	0 =	: I ²	I ² S/M S/MSB S/MSB	Inte	rface	e Lo	ор Е	3ack	Fund	ctic	on is	dis										
		4	4		D	RPI	L	0 =	: Re	e Repla eplayin eplayin	g Fu	nctic	n is	s ena	able	d	SB	-Ju:	stifie	ed Ir	nterf	ace	-						
		;	3		DI	RE	2	0 =	R	e Reco ecordin ecordin	g Fu	nctio	on i	s en	able	d	SE	3-Ju	stifi	ed Ir	nteri	face	:						
		2	:1			_		res	erve	ed																			
		()		Al	MS	L	0 =	: Se	Alternelect I ² elect M	S Op	erat	ion	Mod	de) Op	erat	tion:	:							

14.6.3 Serial Audio Controller I²S/MSB-Justified Status Register (SASR0)

SASR0, shown in Table 14-7, is used for recording the status of the FIFOs and I2SLINK.

Only 4 bits are assigned for TFL and RFL. Actual fill levels are interpreted as follows:

```
Actual_TFL(4:0) = {~TNF, TFL(3:0)}

Actual_RFL(4:0) calculation:
   if (RFL(3:0) == 4'b0)
        Actual_RFL(4:0) = {RNE, RFL(3:0)}
   else
        Actual_RFL(4:0) = {1'b0, RFL(3:0)}
```

This is a read-only register. Ignore reads from reserved bits.



Та	ble	14-7	7. S	AS	R0 F	Bit	De	finitic	ns																					
					I Add		6		5	Serial Ju					ller I Reg			B-					I ² S	Со	ntro	ller				
Bit	31	30 2	9 28	3 27	26	25	24	23 22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
						r	ese	rved								R	FL			Т	FL		reserved	ROR	TUR	RFS	TFS	BSY	RNE	TNF
Reset	0	0 (0 0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
		Bits	,		Nai	me											D	escr	ripti	on										
		31:16	6		_	_		reserv	ed																					
		15:12	2		RF	FL		Receiv	e F	IFO L	.eve	el: n	umb	er	of er	ntrie	es ir	n Red	ceiv	e FI	FO									
		11:8	3		TF	FL		Transr	nit F	IFO I	Lev	el: r	num	ber	of e	ntri	es i	n Tra	ansn	nit F	IFC)								
		7			_	_		reserv	ed																					
		6			RC	DR		Receive $0 = R$ $1 = I^2$ Can in Cleare	ecei S at terru	ve FI temp upt pr	FO ted	has dat	s no ta wi or if l	rite bit6	to fu	ıll R Seri	ece al A	eive f Judio	Inte) (Ir erruj	ot M	lask	Reg	jiste	er is	set.				
		5			TU	JR		Transr $0 = T_1$ $1 = I^2$ Can in Cleare	ans S at	mit F temp upt pr	IFC ted	ha dat	s no a re or if l	ad bit \$	from 5 of 5	an Ser	em ial <i>F</i>	npty ⁻ Audic	Tran	smi erru	t FII pt N	/lask		gisto	er is	set	t.			
		4			RF	-s		Receiv 0 = R 1 = R Can in Cleare	ecei ecei terru	ve FI ve FI upt pr	FO FO	lev lev	el be el is or if l	elov at bit 4	v RF or al	oov Ser	e R ial <i>F</i>	FL th Audic	res Int	hold erru	l. pt N	/lask	Re	_		s set	t.			
		3			TF	-s		Transr 0 = Tr 1 = Tr Can in Cleare	ans ans terru	mit F mit F upt pr	IFC IFC) lev) lev essc	/el e /el is or if l	xce at bit 3	eds or b of s	elo Ser	w T ial <i>F</i>	FL th Audic	nres Int	hold erru	ł pt N	/lask	Re	giste		s set	i.			
		2						$I^{2}S Bu$ $0 = I^{2}$ $1 = I^{2}$	S is	idle d	or d	lisal rans	oled smitt	ting	or r	ece	ivin	ıg a f	ram	е										
		1			RN	ΝE		Receiv 0 = R 1 = R	ecei	ve FI	FΟ	is e	empi		ty															
								Transn	nit F	IFO i	not	full:																		

14.6.4 Serial Audio Clock Divider Register (SADIV)

0 = Transmit FIFO is full 1 = Transmit FIFO is not full

SADIV, shown in Table 14-8, is used for generating six different BITCLK frequencies and hence six different sampling frequencies.

0

TNF



The reset value, 0x0000001A, defaults to a sampling frequency of 22.05 kHz.

Note: Setting this register to values other than those shown in Section 14.2 is not allowed and will cause unpredictable activity.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 14-8. SADIV Bit Definitions

			F			Add 0_00		S			Se	rial	Auc	dio C	Cloc	k Di	vide	r Re	egist	ter					I ² S	Со	ntro	ller				
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												res	serv	ed														S	ADI	V		
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	0	1	0
		Bi	ts			Na	me												De	escr	iptic	n										
		31	:7			-	-		res	erve	d																					
		6:	:0			SA	DIV		000 000 001 010 011	110 110 10 010 010	ilock 00 – 01 – 10 – 00 – 00 –	BIT BIT BIT BIT	CLI CLI CLI CLI	K of K of K of K of K of	3.0° 2.8° 1.4 1.0° 708	72M 36 M 18M 24M 3.92	1Hz MHz MHz MHz MHz KH:	z	are:													

14.6.5 Serial Audio Interrupt Clear Register (SAICR)

SAICR, shown in Table 14-9, is the Interrupt Control Register. This is only an addressable location and no data is actually stored. These addressable locations are used only for clearing status register (SASR0) bits. Each bit position corresponds to an interrupt source bit position in the Status register.

This is a write-only register. Write zeros to reserved bits.

Table 14-9. SAICR Bit Definitions

			F	Phys 0x4		Add 0_00		S			Se	rial .	Aud	io Ir	iterr	upt	Clea	ar R	egis	ter					I ² S	Co	ntro	ller				
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												res	serv	ed												ROR	TUR		res	serv	ed	
Reset	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r	r
		Bi	ts			Na	me												De	escr	ipti	on										
		31	:7			-	-		res	erve	ed																					
		(3			RC	DR		Cle	ar F	Rece	eive	FIF	O 0	verr	un lı	nter	rupt	and	IRC)R s	tatu	s bit	in S	SAS	SR0.						
		Ę	5			ΤL	JR		Cle	ar T	ran	smit	FIF	O u	nde	r-ru	n İn	terru	ıpt a	and	TUF	R sta	tus	bit i	n S	ASR	₹0.					
		4	:0			_	-		res	erve	d																					



14.6.6 Serial Audio Interrupt Mask Register (SAIMR)

Writing a one to the corresponding bit position in the SAIMR, shown in Table 14-10, enables the corresponding interrupt signal.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 14-10. SAIMR Bit Descriptions

			P			Add		S			Se	rial <i>i</i>	Aud	io Ir	terr	upt	Mas	k R	egis	ter					I ² S	Со	ntro	ller				
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												res	serv	ed												ROR	TUR	RFS	TFS	res	serv	ed
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
·		Bi	ts			Na	me												De	scr	iptic	on										
Í		31	:7			_	_		res	erve	ed																					
Í		6	3			RC	OR		Ena	able	Re	ceiv	e FI	FO	Ove	rrur	COI	nditi	on b	ase	d in	terrı	upt.									
•		Ę	5			Τl	JR		Ena	able	FIF	ΟU	Inde	er-ru	n co	ndi	ion	bas	ed ii	nter	rupt											
•		4	1			RF	FS		Ena	able	Re	ceiv	e FI	FO	Ser	vice	Red	ques	st ba	sec	linte	errup	ot.									
•		3	3			TF	-S		Ena	able	Tra	nsm	it F	IFO	Ser	vice	Re	que	st ba	ase	d int	erru	pt.									
•		2:	:0			-			res	erve	ed																					

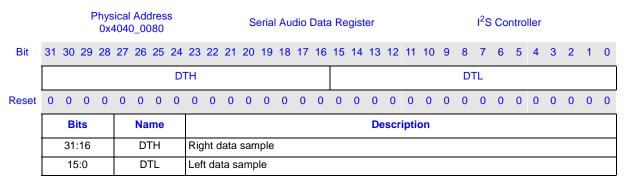
14.6.7 Serial Audio Data Register (SADR)

Writing a 32-bit sample to SADR, shown in Table 14-11, updates the data into the Transmit FIFO. Reading this register flushes a 32-bit sample from the Receive FIFO.

Figure 14-3 illustrates data flow through the FIFOs and SADR.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 14-11. SADR Bit Descriptions





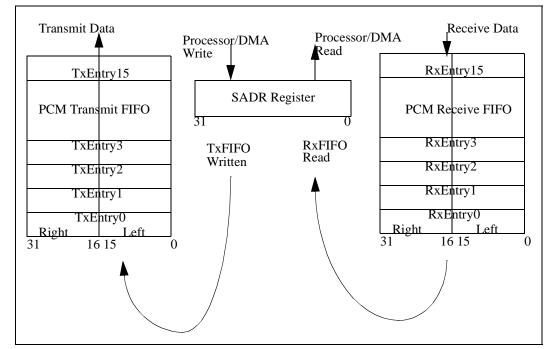


Figure 14-3. Transmit and Receive FIFO Accesses Through the SADR

14.7 Interrupts

The following SASR0 status bits, if enabled, interrupt the processor:

- Receive FIFO Service DMA Request (RFS)
- Transmit FIFO Service DMA Request (TFS)
- Transmit Under-run (TUR)
- Receive Over-run (ROR).

Note: For further details, see Section 14.6.3.

14.8 I²S Controller Register Summary

All registers are word addressable (32 bits wide) and hence increment in units of 0x00004. All I2SC registers are mapped in the address range of 0x4040_0000 through 0x4040_0080, as shown in Table 14-12.



Table 14-12. Register Memory Map

Address (paddr(9:0)	Register name	Description
0x4040_0000	SACR0	Global Control Register
0x4040_0004	SACR1	Serial Audio I ² S/MSB-Justified Control Register
0x4040_0008	_	reserved
0x4040_000C	SASR0	Serial Audio I ² S/MSB-Justified Interface and FIFO Status Register
0x4040_0014	SAIMR	Serial Audio Interrupt Mask Register
0x4040_0018	SAICR	Serial Audio Interrupt Clear Register
0x4040_001C through 0x4040_005C	_	reserved
0x4040_0060	SADIV	Audio clock divider register. See Section 14.4.
0x4040_0064 through 0x4040_007C	_	reserved
0x4040_0080	SADR	Serial Audio Data Register (TX and RX FIFO access register).



MultiMediaCard Controller

15.1 Overview

The PXA255 processor MultiMediaCard (MMC) controller acts as a link between the software used to access the processor and the MMC stack (a set of memory cards). The MMC controller is designed to support the MMC system, a low-cost data storage and communications system. A detailed description of the MMC system is available through the MMC Association's web site at www.mmca.org. The processor's MMC controller is based on the standards outlined in *The MultiMediaCard System Specification Version 2.1* with the exception that one- and three-byte data transfers are not supported and the maximum block length is 1023.

The MMC controller supports the translation protocol from a standard MMC or Serial Peripheral Interface (SPI) bus to the MMC stack. The software used to access the processor must indicate whether to use MMC or SPI mode as the protocol to communicate with the MMC controller.

The MMC controller features:

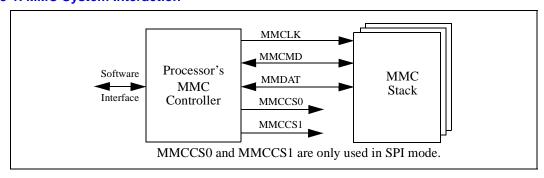
- Data transfer rates up to 20 Mbps
- A response FIFO
- Dual receive data FIFOs
- Dual transmit FIFOs
- Support for two MMCs in either MMC or SPI mode

The MMC controller contains all card-specific functions, serves as the bus master for the MMC system, and implements the standard interface to the card stack. The controller handles card initialization; CRC generation and validation; and command, response, and data transactions.

The MMC controller is a slave to the software and consists of command and control registers, a response FIFO, and data FIFOs. The software has access to these registers and FIFOs and generates commands, interprets responses, and controls subsequent actions.

Figure 15-1 shows a block diagram of the interaction of a typical MMC stack, the MMC controller, and a software.

Figure 15-1. MMC System Interaction





The MMC bus connects the card stack to the controller. The software and controller can turn the MMC clock on and off. The card stack and the controller communicate serially through the command and data lines and implement a message-based protocol. The messages consist of the following tokens:

Command: a 6-byte command token starts an operation. The command set includes card
initialization, card register reads and writes, data transfers, etc. The MMC controller sends the
command token serially on the MMCMD signal. The format for a command token is shown in
Table 15-1.

Table 15-1. Command Token Format

Bit Position	Width (bits)	Value	Description
47	1	0	start bit
46	1	1	transmission bit
[45:40]	6	х	command index
[39:8]	32	х	argument
[7:1]	7	х	CRC7
0	1	1	end bit

- Response: a response token is an answer to a command token. Each command has either a specific response type or no response type. The format for a response token varies according to the response expected and the card's mode. Response token formats are detailed *The MultiMediaCard System Specification Version 2.1*.
- Data: data is transferred serially between the controller and the card in 8-bit blocks at rates up to 20 Mbps. The format for the data token depends on the card's mode. Table 15-2 shows the data token format for MMC mode and Table 15-3 shows the data token format for SPI mode.

Table 15-2, MMC Data Token Format

Stream Data	Block Data	Description
1	0	start bit
x	х	data
no CRC	х	CRC7
1	1	end bit

Table 15-3, SPI Data Token Format

Value	Description
11111110	start byte
x	data
x	CRC16

In MMC mode, all operations contain command tokens and most commands have an associated response token. Read and write commands also have a data token. Command and response tokens are sent and received on the bidirectional MMCMD signal and data tokens are sent and received on



the bidirectional MMDAT signal. A typical MMC mode command timing diagram with and without a response is shown in Figure 15-2 while Figure 15-3 shows a typical MMC mode timing diagram for a sequential read or write

Figure 15-2. MMC Mode Operation Without Data Token

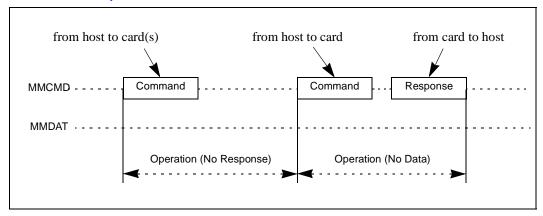
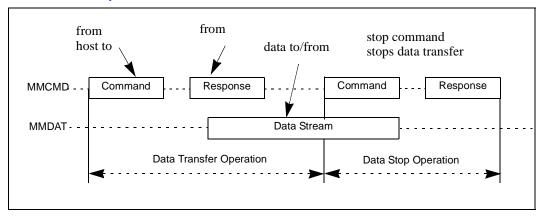


Figure 15-3. MMC Mode Operation With Data Token



In SPI mode, not all commands are available. The available commands have both a command and response token. The MMCMD and MMDAT signals are no longer bidirectional in SPI mode. The MMCMD is an output and the MMDAT is an input with respect to the processor. The command and data tokens to be written are sent on the MMCMD signal and the response and read data tokens are received on the MMDAT signal. Figure 15-4 shows a typical SPI mode timing diagram without a data token. Figure 15-5 and Figure 15-6 show SPI mode read and write timing diagrams, respectively.



Figure 15-4. SPI Mode Operation Without Data Token

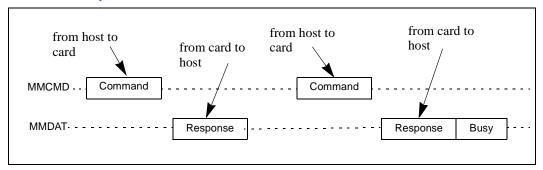


Figure 15-5. SPI Mode Read Operation

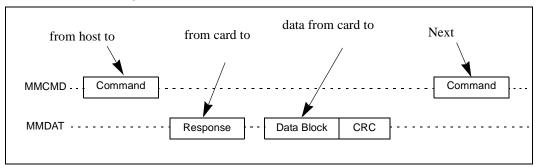
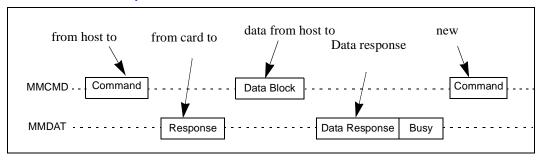


Figure 15-6. SPI Mode Write Operation



Note: One- and three-byte data transfers are not supported with this controller. Data transfers of 10 or more bytes are supported for stream writes only.

Refer to *The MultiMediaCard System Specification* for detailed information on MMC and SPI modes of operation.

15.2 MMC Controller Functional Description

The software must read and write the MMC controller registers and FIFOs to initiate communication to a card.



The MMC controller is the interface between the software and the MMC bus. It is responsible for the timing and protocol between the software and the MMC bus. It consists of control and status registers, a 16-bit response FIFO that is eight entries deep, two 8-bit receive data FIFOs that are 32 entries deep, and two 8-bit transmit FIFOs that are 32 entries deep. The registers and FIFOs are accessible by the software.

The MMC controller also enables minimal data latency by buffering data and generating and checking CRCs.

Refer to Section 15.4 for examples.



15.2.1 Signal Description

The MMC controller signals are MMCLK, MMCMD, MMDAT, MMCCS0, and MMCCS1. Table 15-4 describes each signal's function.

Table 15-4. MMC Signal Description

Signal Name	Input/Output	Description
MMCLK	Output	Clock signal to MMC
MMCMD	BiDirectional	Command line
MMDAT	BiDirectional	Data line
MMCCS0	Output	Chip Select 0 (used only in SPI mode)
MMCCS1	Output	Chip Select 1 (used only in SPI mode)

The MMCLK, MMCCS0, and MMCCS1 signals are routed through alternate functions within the GPIO. Refer to Section 4.1, "General-Purpose I/O" on page 4-1 for a description of the process used to assign these signals to a specific GPIO. Even though there are several GPIO assigned to each signal, each signal must be programmed to one of the possible GPIOs. Refer to Section 4.1.2, "GPIO Alternate Functions" on page 4-2 for a complete description of the GPIO alternate functions.

15.2.2 MMC Controller Reset

The MMC controller can only be reset by a hard or soft reset of the processor. The ways to reset the processor and the MMC controller can be found in Section 2.6, "Reset" on page 2-6. All registers and FIFO controls are set to their default values after any reset.

15.2.3 Card Initialization Sequence

After reset, the MMC card must be initialized by sending 80 clocks to it on the MMCLK signal. To initialize the MMC card, set the MMC_CMDAT[INIT] bit to a 1. This sends 80 clocks before the current command in the MMC_CMD register. This function is useful for acquiring new cards that have been inserted on the bus. Chip selects are not asserted during the initialization sequence.

After the 80-clock initialization sequence, the software must continuously send CMD1 (see Table 15-18 for command definitions) by loading the appropriate command index into the MMC_CMD register until the card indicates that the power-up sequence is complete. The software can then assign an address to the card or put it into SPI mode.

15.2.4 MMC and SPI Modes

After reset, the MMC card is in the MMC mode. The card may remain in MMC mode or be configured to SPI mode by setting the MMC_SPI register bits. The following sections briefly describe each mode as it pertains to the MMC controller.



15.2.4.1 MMC Mode

In MMC mode, the MMCMD and MMDAT signals are bidirectional and require external pullups. The command and response tokens are sent and received via the MMCMD signal and data is read and written via the MMDAT signal.

After an MMC card is powered on, it is assigned a default relative card address (RCA) of 0x0001. The software assigns different addresses to each card during the initialization sequence described in Section 15.2.3. A card is then addressed by its new relative address in the argument portion of the command token that is protected with a 7-bit CRC (see Table 15-1). For a description of the identification process when multiple cards are connected to a system, refer to the Card Identification Process as described in *The MultiMediaCard System Specification*.

There are five formats for the response token, including a no response token. The response token length is 48 or 136 bits and may be protected with a 7-bit CRC. Details of the response token can be found in *The MultiMediaCard System Specification*.

In write data transfers, the data is suffixed with a 5-bit CRC status token from the card. After the CRC status token, the card may indicate that it is busy by pulling the MMDAT line low.

The start address for a read operation can be any random byte address in the valid address space of the card memory. For a write operation, the start address must be on a sector boundary and the data length must be an integer multiple of the sector length. A sector is the number of blocks that will be erased during the write operation and is fixed for each MMC card. A block is the number of bytes to be transferred.

The MMC mode supports these data transfer modes:

- Single block read/write: in single block mode, a single block of data is transferred. The starting address is specified in the command token of the read or write command used. The software must set the block size in the controller by entering the number of bytes to be transferred in the MMC_BLKLEN register. The data block is protected with a 16-bit CRC that is generated by the sending unit and checked by the receiving unit. The CRC is appended to the data after the last data bit is transferred.
- Multiple block read/write: in multiple block mode, multiple blocks of data are transferred.
 Each block is the same length as specified by the software in the MMC_BLKLEN register.
 The blocks of data are stored or retrieved from contiguous memory addresses starting at the address specified in the command token. The software specifies the number of blocks to transfer in the MMC_NOB register. Each data block is protected by appending a 16-bit CRC. Multiple block data transfers are terminated with a stop transmission command.
- Stream read/write: in stream mode, a continuous stream of data is transferred. The starting address is specified in the command token of the read or write command used. The data stream is terminated with a stop transmission command. For write transfers, the stop transmission command must be transmitted with the last six bytes of data. This ensures that the correct amount of data is written to the card. For read transfers, the stop transmission command may occur after the data transmission has occurred. There is no CRC protection for data in this mode.



15.2.4.2 SPI Mode

SPI mode is an optional secondary communication protocol. In SPI mode, the MMCMD and MMDAT lines are unidirectional and only single block data transfers are allowed. The MMCMD signal is an output from the controller and sends the command token and write data to the MMC card. The MMDAT signal is an input to the controller and receives the response token and read data from the MMC card.

Note: When the card is in SPI mode, the only way to return to MMC mode is by toggling the power to the card.

Card addressing is implemented with hardware chip selects, MMCCS1 and MMCCS0. All command, response, and data tokens are 8-bits long and are transmitted immediately following the assertion of the respective chip select.

The command token is protected with a 7-bit CRC. The card always sends a response to a command token. The response token has four formats, including an 8-bit error response. The length of the response tokens is one, two, or five bytes.

SPI mode offers a non protected mode. In this mode, CRC bits of the command, response, and data tokens are still required in the tokens but these bits are ignored by the card and the controller.

In write data transfers, the data is suffixed with an 8-bit CRC status token from the card. As in MMC mode, the card may indicate that it is busy by pulling the MMDAT line low after the status token. In read data transfers, the card may respond with the data or a data error token one byte long.

15.2.5 Error Detection

The MMC controller detects these errors on the MMC bus and reports them in the status register (MMC STAT):

- Response CRC error: a CRC error was calculated on the command response.
- Response time out: the response did not begin before the specified number of clocks.
- Write data CRC error: the card returned a CRC status error on the data.
- Read data CRC error: a CRC error was calculated on the data.
- Read time out: the read data operation did not begin before the specified number of clocks.
- SPI data error: a read data error token was detected In SPI mode.

15.2.6 Interrupts

The MMC controller generates interrupts to signal the status of a command sequence. The software is responsible for masking the interrupts appropriately, verifying the interrupts, and performing the appropriate action as necessary.

Interrupts and masking are described in Section 15.5.11 and Section 15.5.12. The CMDAT[DMA_EN] bit will also mask the MMC_I_MASK[RXFIFO_RD_REQ] and MMC_I_MASK[TXFIFO_WR_REQ] interrupt bits.



15.2.7 Clock Control

Both the MMC controller and the software can control the MMC bus clock (MMCLK) by turning it on and off. This helps to control the data flow to prevent under runs and overflows and also conserves power. The software can also change the frequency at any time to achieve the maximum data transfer rate specified for a card's identification frequency.

The MMC controller has an internal frequency generator that may start, stop, and divide the MMC bus clock. The software may start and stop the clock by setting the appropriate bits in the MMC_STRPCL register. The MMCLK frequency is controlled by the value written in the MMC_CLKRT register.

To write any MMC controller register for the next command sequence, software must:

- 1. Stop the clock.
- 2. Write the registers.
- 3. Restart the clock.

Software must not stop the clock when it attempts to read the receive FIFOs or write the transmit FIFOs. When the clock stops, it resets the pointers in the FIFOs and any data left in the FIFOs can not be transmitted or accessed. When the receive FIFOs are empty and the MMC STAT[DATA TRAN DONE] is set, software may stop the clock.

The software can specify the clock divisor of the 20 Mhz clock by setting the MMC_CLKRT register. The clock rate may be set as:

- 20 Mhz
- 1/2 of 20 Mhz, 10 Mhz
- 1/4 of 20 Mhz, 5 Mhz
- 1/8 of 20 Mhz, 2.5 Mhz
- 1/16 of 20 Mhz, 1.25 Mhz
- 1/32 of 20 Mhz, 625 khz
- 1/64 if 20 Mhz, 312.5 khz

The controller can also turn the clock off automatically. If both receive FIFOs become full during data reads, or one receive FIFO is being read by the software and the other receive FIFO becomes full, or both transmit FIFOs become empty during data writes, or one transmit FIFO is being written by the software and the other transmit FIFO is empty, the controller will automatically turn the clock off to prevent data overflows and under runs. For read data transfers, the controller turns the clock back on after a receive FIFO has been emptied. For write data transfers, the controller turns the clock back on after the transmit FIFO is no longer empty.

Warning: Stopping the clock while data is in the transmit or receive FIFOs will cause unpredictable results.

If the software stops the clock at any time, it must wait for the MMC_STAT[CLK_EN] status bit to be cleared before proceeding.



15.2.8 Data FIFOs

The controller FIFOs for the response tokens, received data, and transmitted data are MMC_RES, MMC_RXFIFO, and MMC_TXFIFO, respectively. These FIFOs are accessible by the software and are described in the following paragraphs.

15.2.8.1 Response Data FIFO (MMC_RES)

The response FIFO, MMC_RES, contains the response received from an MMC card after a command is sent from the controller. MMC RES is a read only, 16-bit, and 8-entry deep FIFO.

The FIFO will hold all possible response lengths. Responses that are only one byte long are located on the MSB of the 16-bit entry in the FIFO. The first half-word read from the response FIFO is the most significant half-word of the received response.

The FIFO does not contain the response CRC. The status of the CRC check is in the status register, MMC_STAT.

15.2.8.2 Receive Data FIFO, MMC RXFIFO

The two receive data FIFOs are read only by the software and are readable on a single byte basis. They are dual FIFOs, where each FIFO is 32 entries of 1-byte data. Access to the FIFOs is controlled by the controller and depends on the status of the FIFOs.

Both FIFOs and their controls are cleared to a starting state after a system reset and at the beginning of all command sequences.

The FIFOs swap between the software and MMC bus. At any time, while the software has read access to one of the FIFOs, the MMC bus has write access to the other FIFO.

For purposes of an example, the FIFOs are called RXFIFO1 and RXFIFO2. After a reset or at the beginning of a command sequence, both FIFOs are empty and the software has read access to RXFIFO1 and the MMC has write access to RXFIFO2. When RXFIFO2 becomes full and RXFIFO1 is empty, the FIFOs swap and the software has read access to RXFIFO2 and the MMC has write access to RXFIFO1. When RXFIFO1 becomes full and RXFIFO2 is empty, the FIFOs swap and the software has read access to RXFIFO1 and the MMC has write access to RXFIFO2.

This swapping process continues through out the data transfer and is transparent to both the software and the MMC controller.

If at any time both FIFOs become full and the data transmission is not complete, the controller turns the MMCLK off to prevent any overflows. When the clock is off, data transmission from the card stops until the clock is turned back on. After the software has emptied the FIFO that it is connected to, the controller turns the clock on to continue data transmission.

The full status of the FIFO that the software is connected to is registered in the MMC_STAT[RECV_FIFO_FULL] bit.

The receive FIFO is readable on byte boundaries and the FIFO read request is only asserted once per FIFO access (once per 32 bytes available). Therefore, 32 bytes must be read for each request, except for the last read which may be less than 32 bytes.



If the DMA is used, it must be programmed to do 1-byte reads of 32-byte bursts. The last read can be less than a 32-byte burst. Some examples are:

• Receive 96 bytes of data:

Read 32 bytes three times.

For the DMA, use three descriptors of 32 bytes and 32-byte bursts.

• Receive 98 bytes of data:

Read 32 bytes three times, then read two more bytes.

For the DMA, use three descriptors of 32 bytes and 32-byte bursts and one descriptor of two more bytes and 8-, 16-, or 32- byte bursts.

• Receive 105 bytes:

Read 32 bytes three times, then read nine more bytes.

For the DMA, use three descriptors of 32 bytes and 32-byte bursts and one descriptor of nine or more bytes and 16- or 32-byte bursts.

15.2.8.3 Transmit Data FIFO, MMC TXFIFO

The two transmit data FIFOs are written only by the software and are writable on a single byte basis. They are dual FIFOs, where each FIFO is 32 entries of one byte data. Access to the FIFOs is controlled by the controller and depends on the status of the FIFOs.

Both FIFOs and their controls are cleared to a starting state after a system reset and at the beginning of all command sequences.

The FIFOs swap between the software and MMC bus. At any time, while the software has write access to one of the FIFOs, the MMC bus has read access to the other FIFO.

For purposes of an example, the FIFOs are called TXFIFO1 and TXFIFO2. After a reset or at the beginning of a command sequence, both FIFOs are empty and the software has write access to TXFIFO1 and the MMC has read access to TXFIFO2. When TXFIFO1 becomes full and TXFIFO2 is empty, the FIFOs swap and the software has write access to TXFIFO2 and the MMC has read access to TXFIFO1. When TXFIFO2 becomes full and TXFIFO1 is empty, the FIFOs swap and the software has write access to TXFIFO1 and the MMC has read access to TXFIFO2.

This swapping process continues through out the data transfer and is transparent to both the software and the MMC controller.

If at any time both FIFOs become empty and the data transmission is not complete, the controller turns the MMCLK off to prevent any underruns. When the clock is off, data transmission to the card stops until the clock is turned back on. When the transmit FIFO is no longer empty, the MMC controller automatically restarts the clock.

If the software does not fill the FIFO to which it is connected, the MMC_PRTBUF[BUF_PART_FULL] bit must be set to a 1. This enables the FIFOs to swap without filling the FIFO.

The empty status of the FIFO that the software is connected to is registered in the MMC_STAT[XMIT_FIFO_EMPTY] bit.

The transmit FIFO is writable on byte boundaries and the FIFO write request is only asserted once per FIFO access (once per 32 entries available). Therefore, 32 bytes must be written for each request, except for the last write which may be less than 32 bytes.



When the DMA is used, it must be programmed to do 1-byte writes of 32-byte bursts. The last write can be less than a 32-byte burst.

If the last write is less than 32 bytes, then the MMC_PRTBUF[BUF_PART_FULL] bit must be set. When the DMA is used, the last descriptor must be programmed to allow the DMA to set an interrupt after the data is written to the FIFO. After the interrupt occurs, the software must set the MMC_PRTBUF[BUF_PART_FULL] bit.

Some examples are:

• Transmit 96 bytes of data:

Write 32 bytes three times.

For the DMA, use three descriptors of 32 bytes and 32-byte bursts.

• Transmit 98 bytes of data:

Write 32 bytes three times, then write two more bytes.

For the DMA, use three descriptors of 32 bytes and 32-byte bursts and one descriptor of two more bytes and 8-, 16- or 32-byte bursts and program the descriptor to set an interrupt, for the software to write the MMC_PRTBUF[BUF_PART_FULL] bit.

• Transmit 105 bytes:

Write 32 bytes three times, then write nine more bytes.

For the DMA, use three descriptors of 32 bytes and 32-byte bursts and one descriptor of nine more bytes and 16- or 32-byte bursts and program the descriptor to set an interrupt, for the software to write MMC_PRTBUF[BUF_PART_FULL] bit.

15.2.8.4 DMA and Program I/O

The software may communicate to the MMC controller via the DMA or program I/O.

To access the FIFOs with the DMA, the software must program the DMA to read or write the MMC FIFOs with single byte transfers, and 32-byte bursts. For example, to write 64 bytes of data to the MMC_TXFIFO, the software must program the DMA to write 64 bytes with an 8-bit port size to the MMC and for 32-byte bursts. The MMC issues a request to read the MMC_RXFIFO and a request to write the MMC_TXFIFO.

With program I/O, the software waits for the MMC_I_REG[RXFIFO_RD_REQ] or MMC_I_REG[TXFIFO_WR_REQ] interrupts before reading or writing the respective FIFO.

The CMDAT[DMA_EN] bit must be set to a 1 to enable communication with the DMA and it must be set to a 0 to enable program I/O.

15.3 Card Communication Protocol

This section discusses the software's responsibilities and the communication protocols used between the MMC and the card.



15.3.1 Basic, No Data, Command and Response Sequence

The MMC controller performs the basic MMC or SPI bus transaction. It formats the command from the command registers and generates and appends the 7-bit CRC if applicable. It then serially translates this to the MMCMD bus, collects the response data, and validates the response CRC. It also checks for response time-outs and card busy if applicable. The response data is in the MMC_RES FIFO and the status of the transaction is in the status register, MMC_STAT.

The protocol of events for the software is:

- 1. Stop the clock
- 2. Write 0x6f to the MMC_I_MASK register and wait for and verify the MMC_I_REG[CLK_IS_OFF] interrupt
- 3. Write to these registers, as necessary:
 - MMC CMD
 - MMC_ARGH
 - MMC_ARGL
 - MMC_CMDAT, this register must be written, even if there is no change to the register
 - MMC_CLKRT
 - MMC_SPI
 - MMC RESTO
- 4. Start the clock
- 5. Write 0x7b to the MMC_I_MASK register and wait for and verify the MMC_I_REG[END_CMD_RES] interrupt
- 6. Read the MMC_RES FIFO and MMC_STAT registers

Some cards may become busy as the result of a command. The software may wait for the card to become not busy by writing the MMC_I_MASK register and waiting for the MMC_I_REG[PRG_DONE] interrupt or the software can start communication to another card. The software may not access the same card again until the card is no longer busy. Refer to *The MultiMediaCard System Specification* for additional information.

15.3.2 Data Transfer

A data transfer is a command and response sequence with the addition of a data transfer to a card.

Refer to the examples in Section 15.4.

The software must follow the steps as described in Section 15.3.1. In addition, before starting the clock, the software must write these registers as necessary.

- MMC RDTO
- MMC_BLKLEN
- MMC_NOB

After the software writes the registers and starts the clock, the software must read the MMC_RES as described above and read or write the MMC_RXFIFO or MMC_TXFIFO FIFOs.



After completely reading or writing the data FIFOs, the software must wait for the appropriate interrupts. The status register, MMC_STAT, must be read to ensure that the transaction is complete and to check the status of the transaction.

When using DMA request signals, the controller indicates to the DMA when a FIFO is ready for reading or writing. It is expected that all FIFO reads and writes will empty and fill the FIFO to which it is connected. If at any time the MMC_TXFIFO is not filled (32 bytes) by the software, the software must notify the controller by setting the MMC_PRTBUF[BUF_PART_FULL]. The software can write more bytes of data than is needed into the MMC_TXFIFO, but the controller will only transmit the number of bytes in the MMC_BLKLEN register.

At the end of any data transfer or busy signal on the MMC bus, the MMC controller waits eight MMC clocks before asserting the MMC_I_REG[DATA_TRAN_DONE] interrupt to notify the software that the data transfer is complete. This guarantees that the specified minimum of eight MMC clocks occurs between a data transfer and the next command.

On write data transfers, a card may become busy while programming the data. The software may wait for the card to become not busy by writing the MMC_I_MASK register and waiting for the MMC_I_REG[PRG_DONE] interrupt or the software can start communication to another card. Refer to *The MultiMediaCard System Specification* for additional information.

The MMC controller performs data transactions in all the basic modes: single block, multiple blocks, and stream modes.

15.3.2.1 Block Data Write

In a single block data write, a block of data is written to a card. In a multiple block write, the controller performs multiple single block write data transfers on the MMC bus.

After turning the clock on to start the command sequence, the software must program the DMA to fill the MMC_TXFIFO (write 32 bytes). The software must continue to fill the FIFO until all of the data has been written to the FIFOs. The software must then wait for the transmission to complete by waiting for the MMC_I_REG[DATA_TRAN_DONE] interrupt and MMC_I_REG[PRG_DONE] interrupt. The software can then read the status register, MMC_STAT, to verify the status of the transaction.

For multiple block writes, *The MultiMediaCard System Specification* specifies that the card will continue to receive blocks of data until the stop transmission command is received. After the controller has transmitted the number of bytes specified in the MMC_NOB register, the controller will stop transmitting data. After the MMC_I_REG[DATA_TRAN_DONE] interrupt is detected, the software must setup the controller to send the stop transmission command, CMD12. Consult *The MultiMediaCard System Specification* for a description of the stop transmission command.

If both transmit FIFOs become empty during data transmission, the MMC controller turns the clock off. After a FIFO has been written, the controller turns the clock back on.

In a block data write, these parameters must be specified:

- The data transfer is a write.
- The block length if the block length is different from the previous block data transfer or this is the first time that the parameter is being specified.
- The number of blocks to be transferred.



15.3.2.2 Block Data Read

In a single block data read, a block of data is read from a card. In a multiple block read, the controller performs multiple single block read data transfers on the MMC bus.

After turning the clock on to start the command sequence, the software must program the DMA to empty the MMC_RXFIFO (read 32 bytes). The software will continue the process of emptying the FIFO until all of the data has been read from the FIFO. The software must then wait for the transmission to complete by waiting for the MMC_I_REG[DATA_TRAN_DONE] interrupt. The software can then read the status register, MMC_STAT, to verify the status of the transaction.

For multiple block reads, *The MultiMediaCard System Specification* specifies that the card will continue to send blocks of data until the stop transmission command is received. After the controller has received the number of bytes specified in the MMC_NOB register, the controller will stop receiving data. After the MMC_I_REG[DATA_TRAN_DONE] interrupt is detected, the software must set up the controller to send the stop transmission command, CMD12. Consult *The MultiMediaCard System Specification* for a description of the stop transmission command.

If both receive FIFOs become full during the data transmission, the controller turns the clock off. Once the software empties the FIFO to which it is connected, the controller turns the clock back on.

In a block data read, the following parameters must be specified:

- The data transfer is a read.
- The block length, if the block length is different from the previous block data transfer or this is the first time that the parameter is being specified.
- The number of blocks to be transferred.
- The receive data time-out period.

The controller will mark the data transaction as timed out if data is not received before the time-out period. The delay for the time-out period is defined as:

Time-out Delay =
$$\frac{(MMC_RDTO[READ_TO]) \times (128)}{10^7} sec$$

The software is required to calculate this value and write the appropriate value into the MMC_RDTO register.

15.3.2.3 Stream Data Write

The stream data write looks like the single block write except a stop transmission command is sent in parallel with the last six bytes of data.

After turning the clock on to start the command sequence, the software must start the process of filling the MMC_TXFIFO and starting the clock as describe in Section 15.3.2.1. The software must then wait for the MMC_I_REG[STOP_CMD] interrupt. This interrupt indicates that the MMC controller is ready for the stop transmission command. The software must then stop the clock, write the registers for a stop transmission command, and then start the clock. At this point, the software must wait for the MMC_I_REG[DATA_TRAN_DONE] and MMC_I_REG[PRG_DONE] interrupts.



In a stream data write, the following parameters must be specified:

- The data transfer is a write.
- The data transfer is in stream mode
- The block length, if the block length is different from the previous block data transfer or this is the first time that the parameter is being specified.
- The number of blocks to be transferred as 0xffff.

15.3.2.4 Stream Data READ

The stream data read looks like the single block read except a stop transmission command must be sent after the data transfer.

After turning the clock on to start the command sequence, the software must start the process of reading the MMC_RXFIFO as described in Section 15.3.2.2.

When it uses the DMA, the software must also configure the DMA to send an interrupt after all data has been read. After the DMA interrupt or the program has read all of the data, the software must send the stop transmission command. The MCC_STAT[DATA_TRAN_DONE] bit is not set until after software sends the stop transmission command.

In a stream data read, the following parameters must be specified:

- The data transfer is a read.
- The data transfer is in stream mode.
- The block length, if the block length is different from the previous block data transfer or this is the first time that the parameter is being specified.
- The number of blocks to be transferred as 0xffff.
- The receive data time-out period.

15.3.3 Busy Sequence

The MMC controller expects a busy signal automatically from the card after every block of data for single and multiple block write operations. It will also expect a busy at the end of a command every time that the software specifies that a busy signal is expected (i.e. a busy signal is expected after the commands for stop transmission, card select, erase, program CID, etc.). Refer to the *The MultiMediaCard System Specification*.

While a busy signal is on the MMC bus, the software can send only one of two commands:

- Send status command (CMD13).
- Disconnect command (CMD7).

If the software disconnects a card while it is in a busy state, the busy signal will be turned off and the software can connect a different card. The software may not start another command sequence on the same card while the card is busy.



15.3.4 SPI Functionality

The MMC controller can address up to two cards in SPI mode using the MMCCS[1:0] chip select signals. Once the software specifies which chip select to enable in the MMC_SPI register, the selected signal is driven active low at a falling edge of the MMC clock. The chip select remains asserted until software clears the chip select enable bit or selects another card.

Note: The clock must be stopped before writing to any registers as described in Section 15.3.1.

In SPI mode, the software has the option of performing a CRC check. The default is no CRC checking.

The command and data are sent on the MMC bus aligned to every 8 clocks as described in the SPI section of *The MultiMediaCard System Specification*.

In a read sequence, the card may return data or a data error token. If a data error token is received, the controller will stop the transmission and update the status register.

15.4 MultiMediaCard Controller Operation

The software directs all communication between the card and the controller. The operations shown in the preceding sections are valid for MMC mode only.

15.4.1 Start and Stop Clock

The set of registers is accessed by stopping the clock, writing the registers, and starting the clock.

The software stops the clock, as:

- 1. Write 0x01 in MMC STRPCL to stop the MMC clock.
- 2. Write 0x0f in MMC_I_MASK to mask all interrupts except the MMC_I_REG[CLK_IS_OFF] interrupt.
- 3. Wait for the MMC_I_REG[CLK_IS_OFF] interrupt.

To start the clock the software writes 0x02 in MMC_STRPCL.

15.4.2 Initialize

Card initialization sequences must be prefixed with 80 clock cycles. To generate 80 clock cycles before any command, the software must set the MMC_CMDAT[INIT] bit.

15.4.3 Enabling SPI Mode

To communicate with a card in SPI mode, the software must set the MMC_SPI register as follows:

- 1. MMC SPI[SPI EN] must be set to 1.
- 2. MMC_SPI[SPI_CS_EN] must be set to 1.



3. MMC_SPI[SPI_CS_ADDRESS] must be set to specify the card that the software wants to address. A 1 enables CS0 and a 0 enables CS1.

Note: When the card is in SPI mode, the only way to return to MMC mode is by toggling power to the card.

15.4.4 No Data Command and Response Sequence

For the basic no data transfer, command and response transaction, the software must:

- 1. Turn the clock off, as described in Section 15.4.1.
- 2. Write the command index in the MMC_CMD[CMD_INDEX] bits.
- 3. Write the command argument in the MMC ARGH and MMC ARGL registers.
- 4. Write the MMC_CMDAT register set as:
 - a. Write 0b00 to MMC_CMDAT[RESPONSE_FORMAT].
 - b. Clear the MMC_CMDAT[DATA_EN] bit.
 - c. Clear the MMC_CMDAT[BUSY] bit, unless the card may respond busy.
 - d. Clear the MMC_CMDAT[INIT] bit.
- 5. Write MMC_RESTO register with the appropriate value.
- 6. Write 0x1b in MMC_I_MASK to unmask the MMC_I_REG[END_CMD_RES] interrupt.
- 7. Start the clock, as described in Section 15.4.1.

The software must not make changes in the set of registers until the end of the command and response sequence, after the clock is turned on.

After the clock is turned on, the software must wait for the MMC_I_REG[END_CMD_RES] interrupt, which indicates that the command and response sequence is finished and the response is in the MMC_RES FIFO.

The software may then read the MMC_STAT register to verify the status of the transaction and then read MMC_RES FIFO. If a response time-out occurred, the MMC_RES FIFO will not contain any valid data.

15.4.5 Erase

An erase command is performed as described in Section 15.4.4 with the following additions: the BUSY_BIT in the MMC_CMDAT register must be set to a 1 after it reads the MMC_RES FIFO.

15.4.6 Single Data Block Write

In a single block write command, the software must stop the clock and set the registers as described in Section 15.4.4. These registers must be set before the clock is started:

- Set MMC_NOB register to 0x0001.
- Set MMC_BLKLEN to the number of bytes per block.



- Update the MMC_CMDAT register as:
 - Write 0x01 to MMC_CMDAT[RESPONSE_FORMAT]
 - Set the MMC_CMDAT[DATA_EN] bit.
 - Set the MMC_CMDAT[WRITE/READ] bit.
 - Clear the MMC_CMDAT[STREAM_BLOCK] bit.
 - Clear the MMC_CMDAT[BUSY] bit.
 - Clear the MMC_CMDAT[INIT] bit.
- Turn the clock on.

After it starts the clock, the software must perform these steps:

- 1. Wait for the response as described in Section 15.4.4.
- 2. Write data to the MMC_TXFIFO FIFO and continue until all of the data has been written to the FIFO.

Note: If a piece of data smaller than 32 bytes is written to the FIFO, the MMC_PRTBUF register must be set.

- 3. Set MMC_I_MASK register to 0x1e and wait for MMC_I_REG[DATA_TRAN_DONE] interrupt.
- 4. Set MMC I MASK to 0x1d.
- 5. Wait for MMC_I_REG[PRG_DONE] interrupt. This interrupt indicates that the card has finished programming. Software may wait for MMC_I_REG[PRG_DONE] or start another command sequence on a different card.
- 6. Read the MMC_STAT register to verify the status of the transaction (i.e. CRC error status).

To address a different card, the software sends a select command to that card by sending a basic no data command and response transaction. To address the same card, the software must wait for MMC_I_REG[PRG_DONE] interrupt. This ensures that the card is not in the busy state.

15.4.7 Single Block Read

In a single block read command, the software must stop the clock and set the registers as described in Section 15.4.4.



These registers must be set before the clock is started:

- Update these MMC_CMDAT register bits:
 - Set the MMC_CMDAT[RESPONSE_FORMAT] bit.
 - Set the MMC_CMDAT[DATA_EN] bit.
 - Clear the MMC CMDAT[WRITE/READ] bit.
 - Clear the MMC_CMDAT[STREAM_BLOCK] bit.
 - Clear the MMC_CMDAT[BUSY] bit.
 - Clear the MMC_CMDAT[INIT] bit.
- Set MMC NOB register to 0x0001.
- Set MMC_BLKLEN register to the number of bytes per block.
- Turn the clock on.

After it turns the clock on, the software must perform these steps:

- 1. Wait for the response as described in Section 15.4.4.
- 2. Read data from the MMC_RXFIFO FIFO, as data becomes available in the FIFO, and continue reading until all data is read from the FIFO.
- 3. Set MMC_I_MASK to 0x1e.
- 4. Wait for the MMC_I_REG[DATA_TRAN_DONE] interrupt.
- 5. Read the MMC_STAT register to verify the status of the transaction (i.e. CRC error status).

15.4.8 Multiple Block Write

The multiple block write mode is similar to the single block write mode, except that multiple blocks of data are transferred. Each block is the same length. All the registers are set as they are for the single block write, except that the MMC_NOB register is set to the number of blocks to be written.

The multiple block write mode also requires a stop transmission command, CMD12, after the data is transferred to the card. After the MMC_I_REG[DATA_TRAN_DONE] interrupt occurs, the software must program the controller registers to send a stop data transmission command.

15.4.9 Multiple Block Read

The multiple block read mode is similar to the single block read mode, except that multiple blocks of data are transferred. Each block is the same length. All the registers are set as they are for the single block read, except that the MMC_NOB register is set to the number of blocks to be read.

The multiple block read mode requires a stop transmission command, CMD12, after the data from the card is received. After the MMC_I_REG[DATA_TRAN_DONE] interrupt has occurred, the software must program the controller registers to send a stop data transmission command.



15.4.10 Stream Write

In a stream write command, the software must stop the clock and set the registers as described in Section 15.4.4. These registers must be set before the clock is started:

- Set MMC_NOB register to ffffh.
- Set MMC BLKLEN register to the number of bytes per block.
- Update MMC_CMDAT register as:
 - Write 0b01 to the MMC_CMDAT[RESPONSE_FORMAT].
 - Set the MMC_CMDAT[DATA_EN] bit.
 - Set the MMC_CMDAT[WRITE/READ] bit.
 - Set the MMC_CMDAT[STREAM_BLOCK] bit.
 - Clear the MMC_CMDAT[BUSY] bit.
 - Clear the MMC_CMDAT[INIT] bit.
- Turn the clock on.

After it turns the clock on, the software must perform these steps:

- 1. Wait for the response as described in Section 15.4.4.
- 2. Write data to the MMC_TXFIFO FIFO and continue until all of the data is written to the FIFO.

Note: When data less than 32 bytes is written to the FIFO, the MMC_PRTBUF[BUF_PART_FULL] bit must be set.

- 3. Set MMC_I_MASK to 0x77 and wait for MMC_I_REG[STOP_CMD] interrupt.
- 4. Set the command registers for a stop transaction command (CMD12).
- 5. Wait for a response to the stop transaction command as described in Section 15.4.4.
- 6. Set MMC I MASK to 0x1e.
- 7. Wait for MMC I REG[DATA TRAN DONE] interrupt.
- 8. Set MMC_I_MASK to 0x1d.
- 9. Wait for MMC_I_REG[PRG_DONE] interrupt. This interrupt indicates that the card has finished programming. Software may wait for MMC_I_REG[PRG_DONE] interrupt or start another command sequence on a different card.
- 10. Read the MMC_STAT register to verify the status of the transaction (i.e. CRC error status).

To address a different card, the software must send a select command to that card by sending a basic no data command and response transaction. To address the same card, the software must wait for MMC_I_REG[PRG_DONE] interrupt. This ensures that the card is not in the busy state.

15.4.11 Stream Read

In a stream read command, the software must stop the clock and set the registers as described in Section 15.4.4. These registers must be set before the clock is turned on:

• Set MMC_NOB register to ffffh.



- Set MMC_BLKLEN register to the number of bytes per block.
- Update the MMC_CMDAT register as:
 - Write 0x01 to the MMC_CMDAT[RESPONSE_FORMAT].
 - Set the MMC_CMDAT[DATA_EN] bit.
 - Clear the MMC_CMDAT[WRITE/READ] bit.
 - Set the MMC_CMDAT[STREAM_BLOCK] bit.
 - Clear the MMC_CMDAT[BUSY] bit.
 - Clear the MMC_CMDAT[INIT] bit.
- Turn the clock on.

After it turns the clock on, the software must perform these steps:

- 1. Wait for the response as described in Section 15.4.4.
- 2. Read data from the MMC_RXFIFO FIFO and continue until all of the data has been read from the FIFO.
- 3. Set the command registers for a stop transaction command (CMD12). If the DMA is being used, the last descriptor must set the DMA to send an interrupt to signal that all the data has been read.
- 4. Wait for a response to the stop transaction command as described in Section 15.4.4.
- 5. Set MMC_I_MASK to 0x1e.
- 6. Wait for MMC_I_REG[DATA_TRAN_DONE] interrupt.
- 7. Read the MMC_STAT register to verify the status of the transaction (i.e. CRC error status).

15.5 MMC Controller Registers

The MMC controller is controlled by a set of registers that software configures before every command sequence on the MMC bus.

Table 15-5 through Table 15-23 describe the registers and FIFOs.

15.5.1 MMC_STRPCL Register

The MMC_STRPCL, shown in Table 15-5, allows the software to start and stop the MMC bus clock. Reads from this register are unpredictable.

This is a write-only register. Write zeros to reserved bits.



Table 15-5. MMC_STRPCL Bit Definitions

			P	hysi 0x4		Add 0_00		SS				N	IMC	:_S	TRF	CL	Reg	iste	er				Μu	ıltiN	ledi	aCa	ard (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														ı	rese	rve	d														CTDDCI	0 N C L
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	escr	ipti	on										
		31	:2			_	_		res	erve	d																					
	31:2 — 1:0 strpcl									– Do – St – St	op to no op to art to serv	thin ne N he N	g /IMC	C clo	ock ock ock																	

15.5.2 MMC_Status Register (MMC_STAT)

MMC_STAT, shown in Table 15-6, is the status register for the MMC controller. The register is cleared at the beginning of every command sequence.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-6. MMC_STAT Bit Definitions (Sheet 1 of 2)

			PI		ical 4110			S					MN	IC_	STA	TR	egi	ster					Μu	ıltiN	ledi	aCa	ard (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
								r	ese	rved	ı								END_CMD_RES	PRG_DONE	DATA_TRAN_DONE	reserved		CLK_EN	RECV_FIFO_FULL	XMIT_FIFO_EMPTY	RES_CRC_ERR	SPI_READ_ERROR_TOKEN	CRC_READ_ERROR	CRC_WRITE_ERROR	TIME_OUT_RESPONSE	READ_TIME_OUT
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0	0	0
		Bi	ts			Nai	me												De	escr	ipti	on										
		31:	14			_	_		rese	erve	d																					
		1	3		EN	D_C E)_R	0 1	- (ma	nd a	and	resp							ot co			d							
		1.	2		PR	RG_I	DOI			- (ha	s n									bus;		sy								
		1	1		DA	TA_ DO		N_	Dat 0 1	_ [ansn Data Data	tra	ınsn	niss	ion t							ed										
		10	:9			_			rese	erve	d _																					



Table 15-6. MMC_STAT Bit Definitions (Sheet 2 of 2)

		ı			- Addres _0004	S			ı	` //М	C_:	STA	TR	egis	ter				ı	/lulti	Med	iaC	ard	Con	trol	ler		
Bit	31	30 29	28	27	26 25	24	23 22	2 21	20 1	9	18	17	16	15	14	13	12	11	10 9	8 (7	6	5	4	3	2	1	0
						r	eserv	ed								END_CMD_RES	PRG_DONE	DATA_TRAN_DONE	reserved	CLK EN	RECV_FIFO_FULL	XMIT_FIFO_EMPTY	RES_CRC_ERR	SPI_READ_ERROR_TOKEN	CRC_READ_ERROR	CRC_WRITE_ERROR	TIME_OUT_RESPONSE	READ_TIME_OUT
Reset	0	0 0	0	0	0 0	0	0 0	0	0	0	0	0	0	0	0			0		0	0	1	0	0	0	0	0	0
		Bits			Name											De	esci	ipti	on									
		8		C	CLK_EN			Enab MM MM	C clo																			
		7		REC	CV_FIF FULL	0_	-	ve FII Rec Rec	eive	FIF	_			I														
		6			T_FIFO)_E		mit FI Trar Trar	nsmit	FIF	-O i				′													
		5			S_CRC RR	_	1 –	No e	error C erro	on or c	the occu						se											
		4			_READ OR_TO N			ead E SPI SPI	data	err	or t								d									
		3		CR(C_REA ERROR	D_	-	Read No e CR(error	on					eive	ed d	ata											
		2			C_WRITERROR			Nrite No e Car	error	on							ssio	n of	data									
		1			E_OUT SPONS			Out R Card Car	d resp	oor	ise				ed c	ut			_									
•		0		REA	AD_TIM OUT	1E_		Time Car	d read						ed o	out												

15.5.3 MMC_CLKRT Register (MMC_CLKRT)

Card read data timed out

MMC_CLKRT, shown in Table 15-7, specifies the frequency division of the MMC bus clock. The software is responsible for setting this register.

The software can only write this register after the clock is turned off and the software has received an interrupt that indicates the clock is turned off.



This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-7. MMC_CLK Bit Definitions

			P	•	ical 4110			SS				ı	ММ	C_	CLI	K R	TR	Regi	ste	r				Mu	ultiN	/led	iaCa	ard	Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	B 1	7 '	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
														re	ese	rve	ed															K_F [2:0	RAT 0]
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	()	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Na	me													De	esc	ripti	on										
		31	:3			-	_		rese	erve	d																						
		2	:0		CLF	(_R 0		≣[2:	001 010 011 100 101 110	- 1 - 2 - 1 - 0 - 0	0 M 0 M 5 MH .5 M .25N .625	Hz z c Hz //Hz 5 M	clock clock z clo Hz c	k, , 1, ck, ck ck	/4 c 1/8 , 1/ ck,	of 2 of 16 1/3	0 M 20 of 2 2 o	//Hz MH 20 N	cloo z clo MHz MH	ck ock clod	lock												

15.5.4 MMC_SPI Register (MMC_SPI)

MMC_SPI, shown in Table 15-8, is for SPI mode only and is set by the software.

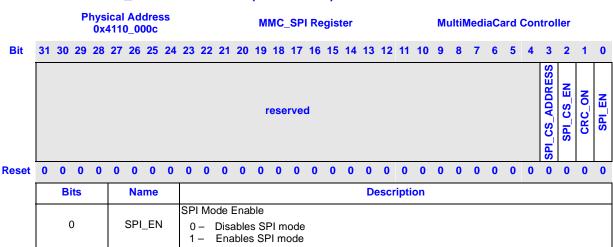
This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-8. MMC_SPI Bit Definitions (Sheet 1 of 2)

Physical Address MMC_SPI Register MultiMediaCard Controller 0x4110_000c 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 PI_CS_ADDRESS SPI_CS_EN 8 M reserved CRC Reset 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 **Bits Name Description** 31:4 reserved Specifies the relative address of the card to activate the SPI CS SPI_CS_ADD 3 0 - Enables CS1 RESS 1 - Enables CS0 SPI Chip Select Enable 2 SPI_CS_EN 0 - Disables the SPI chip select 1 - Enables the SPI chip select CRC Generation Enable 1 CRC_ON 0 - Disables CRC generation and verification Enables CRC generation and verification



Table 15-8. MMC_SPI Bit Definitions (Sheet 2 of 2)



15.5.5 MMC_CMDAT Register (MMC_CMDAT)

MMC_CMDAT, shown in Table 15-9, controls the command sequence. Writing to this register starts the command sequence on the MMC bus when the MMC bus clock is turned on.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-9. MMC CMDAT Bit Definitions (Sheet 1 of 2)

			P		ical 4110		dres 010	S				ı	мм	C_C	ME	DAT	Reg	iste	r				Mı	ultiN	/ledi	aCa	ırd (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
·											r	ese	rve	d											MMC_DMA_EN	LINI	BUSY	STREAM_BLOCK	WRITE/READ	DATA_EN	RESPONSE FORMATI1:01	
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0
		Bi	ts			Na	me												De	esci	ripti	on										
		31	8:			-	_		rese	erve	d																					
		7			MM	I_DI	AMC	Ē	0 1 Wh		Prog DM DM	grar A m	n I/C) m		, this	s bit	is a	ma	sk o	n R	XFIF	=O_	RD_	_RE	Q aı	nd T	XFI	FO_	_WR	RI	ΞQ
·		6				IN	IT		80 I 0 1		Do i	not	pred	ede	со	mma d sed							lock	(S								
		5				BU	SY								•	sign mar									com	ıma	nd.					



Table 15-9. MMC_CMDAT Bit Definitions (Sheet 2 of 2)

			P		ical 4110			SS				ı	имо	C_C	MI	DAT	Re	gist	er					Mu	ıltiN	ledi	iaCa	ard (Con	trol	ler	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	4 1	3	12	11	10	9	8	7	6	5	4	3	2	1 0
											r	ese	rve	d												MMC_DMA_EN	LINI	BUSY	STREAM_BLOCK	WRITE/READ	DATA_EN	RESPONSE_FORMAT[1:0]
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	(0	0	0	0	0	0	1	0	0	0	0	0	0 0
		В	ts			Na	me													De	scr	ipti	on									
		4	1		ST	REA O	AM_ CK	BL		_		a tra				e cu														е		
		(3		WR	ITE	/RE	AD	Rea 0 1	_	Spe		s th	at th	ne	data data																
		2	2		D	ATA	_E			_	No (tra	nsfe		rith o							les a	a da	ta tr	ans	fer					
		1	:0			SP(SE_ 1:0]	00 - 01 - 10 -	- No - Fo - Fo	res orma	spor at R	nse 1, R 2 in	in M 1b, MM	IMO R4 C i	pon C mo , and mode mode	de. d R: e. F	No in orm	t su MN nat	JPI /IC R2	oort mo	ed i ode. SPI	n SI For mo	PI m mat de	ode	:		b in	SP	mo	ode	

15.5.6 MMC_RESTO Register (MMC_RESTO)

The MMC_RESTO, shown in Table 15-10, controls the number of MMC clocks that the controller must wait after the command before it can turn on the time-out error if a response has not occurred. The default value of this register is 64.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-10. MMC_RESTO Bit Definitions

			P	hysi 0x4		Ado 0_00		SS				ı	имо	C_R	ES	го г	Regi	iste	r				Μι	ıltiN	ledi	aCa	ırd (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		reserved RES_TO															то															
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0	0	0
		Bi	ts			Na	me												De	escr	ipti	on										
		31	:7			-	_		rese	erve	d																					
		6	:0		F	RES	_TC)	Nur	nbe	r of	MM	C cl	ock	s be	fore	a r	esp	onse	e tim	ne-o	ut										



15.5.7 MMC_RDTO Register (MMC_RDTO)

MMC_RDTO, shown in Table 15-11, determines the length of time that the controller waits after a command before it turns on the time-out error if data has not been received. The length of time before a time-out occurs is measured in increments of 256 20 MHz periods and can be calculated as follows:

Time-out Delay =
$$\frac{(\text{MMC_RDTO[READ_TO]}) \times (256)}{20MHz} = \frac{(\text{MMC_RDTO[READ_TO]}) \times (128)}{10^7} \sec \frac{(\text{MMC_RDTO[READ_TO]}) \times (128)}{10^7}$$

It is the software's responsibility to calculate the value for READ_TO. The default value is 0xffff, which corresponds to 838.848 ms.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-11. MMC_RDTO Register

						_				J -																						
			P				dres 018	SS					ММ	C_F	RDT	O R	legi	ster					Mu	ultiN	/led	iaCa	ard	Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
·		30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 reserved																				R	EA	D_T	0							
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
•		В	its			Na	me												De	esci	ipti	on										
,		31	:16			_			res	erve	d																					
,		15	5:0		R	EAI	D_T	0	Spe	cifie	es th	ne le	engtl	h of	time	e be	fore	a d	ata	read	d tin	ne-o	ut									



15.5.8 MMC_BLKLEN Register (MMC_BLKLEN)

MMC_BLKLEN, shown in Table 15-12, specifies the number of bytes in a block of data.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-12. MMC_BLKLEN Bit Definitions

			P	hysi 0x4		Add 0_00		SS				N	имс	_BI	_KL	EN	Reg	jiste	er				Μu	ıltiN	ledi	aCa	rd (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 reserved																		В	LK_	LE	N								
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	its			Na	me												De	escr	ipti	on										
		31:	:10			_	_		res	erve	d																					
		9:	:0		В	BLK_	_LE	N	Nur	nbe	r of	byte	es in	the	blo	ck.																

15.5.9 MMC_NOB Register (MMC_NOB)

In block mode, MMC_NOB, shown in Table 15-13, specifies the number of blocks.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-13. MMC_NOB Bit Definitions

			P				dres 020	SS					MN	IC_	NOI	3 Re	egis	ter					Μι	ıltiN	ledi	aCa	ard (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		reserved																				M	MC.	NO	В							
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	its			Na	me												De	scr	ipti	on										
		31:	:16			_	_		rese	erve	d																					
		15	5:0		M	MC	_NC)B	Nur	nbe	r of	bloo	ks f	or a	mu	ltiple	e blo	ock 1	tran	sfer												



15.5.10 MMC_PRTBUF Register (MMC_PRTBUF)

MMC_PRTBUF, shown in Table 15-14, is used when MMC_TXFIFO is partially written. The FIFOs swap when either FIFO is full (32 bytes) or the MMC_PRTBUF register is set to a 1.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-14. MMC_PRTBUF Bit Definitions

			P		ical 411(dres 024	SS				N	IMC	_PI	RTB	UF	Reg	jiste	er				Mu	ıltiN	ledi	iaCa	ard	Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	reserved																		BUF_PART_FULL													
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
•		Bi	ts			Na	me												De	esci	ripti	on										
•		31	:1			_	_		rese	erve	d																					
									Buff			•																				
		C)		BU	_	AR7 LL	I_F	0 · 1 ·		Buff Buff							nust	be:	swa	ppe	d to	the	othe	er tr	ans	mit k	ouffe	er			
									Soft																							

15.5.11 MMC_I_MASK Register (MMC_I_MASK)

MMC_I_MASK, shown in Table 15-15, masks off the various interrupts when set to a 1.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-15. MMC_I_MASK Bit Definitions (Sheet 1 of 2)

			P		ical 1110		dres 028	ss				N	имс	:_I_	MA	SK	Reg	iste	r				Μι	ıltiM	ledi	aCa	ard (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												res	serv	ed												TXFIFO_WR_REQ	RXFIFO_RD_REQ	CLK_IS_OFF	STOP_CMD	END_CMD_RES	PRG_DONE	DATA_TRAN_DONE
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1
		Bi	ts			Na	me												De	esci	ipti	on										
		31	:7			-	_		rese	erve	d																					
·		6	6		TXF	FIFC	D_W EQ	/R_	0	_	it FI Not Mas	mas	sked		equ	est																



Table 15-15. MMC_I_MASK Bit Definitions (Sheet 2 of 2)

			PI		ical 4110			SS					имо	C_I_	MA	SK	Reg	iste	r				Μι	ıltiN	ledi	aCa	ırd (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												res	serv	/ed												TXFIFO_WR_REQ	RXFIFO_RD_REQ	CLK_IS_OFF	STOP_CMD	END_CMD_RES	PRG_DONE	DATA_TRAN_DONE
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1
		Bi	ts			Na	me												De	esci	ripti	on										
		5	5		RXI		O_F EQ	RD_		_	e FII Not Mas	ma	ske	id Re	equ	est																
		4	1		C		_IS. FF	_		_	s Off Not Mas	ma		d																		
		3	3		ST	OP.	_CN	ИD		_	for S Not Mas	ma	ske	ansa d	ctio	n Co	omn	nand	t													
		2	2		ENI		CME S)_R		_	mm Not Mas	ma	ske	spoi	nse																	
		1	1		PR	:G_	DOI	NE		_	nmi Not Mas	ma	ske																			
		()		DAT	ΓA_ DO		N_		_	ans Not Mas	ma	ske																			

15.5.12 MMC_I_REG Register (MMC_I_REG)

MMC_I_REG, shown in Table 15-16, shows the currently requested interrupt. The FIFO request interrupts, TXFIFO_WR_REQ, and RXFIFO_RD_REQ are masked off with the MMC_DMA_EN bit in the MMC_CMDAT register. The software is responsible for monitoring these bits in program I/O mode.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 15-16. MMC_I_REG Bit Definitions

Physical Address	MMC I DEC Demistes	MultiMadia Cand Cantuallan
0x4110_002c	MMC_I_REG Register	MultiMediaCard Controller

				UX	• • • •		120																									
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													serv													TXFIFO_WR_REQ	RXFIFO_RD_REQ	CLK_IS_OFF	STOP_CMD	END_CMD_RES	PRG_DONE	DATA_TRAN_DONE
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

Bits	Name	Description
31:5	_	reserved
6	TXFIFO_WR_ REQ	Transmit FIFO Write Request 0 - No Request for data write to MMC_TXFIFO FIFO 1 - Request for data write to MMC_TXFIFO FIFO Cleared after each write but immediately set again unless there are no entries left in the FIFIO.
5	RXFIFO_RD_ REQ	Receive FIFO Read Request 0 - No Request for data read from MMC_RXFIFO FIFO 1 - Request for data read from MMC_RXFIFO FIFO Cleared after each read but immediately set again unless the FIFO is empty.
4	CLK_IS_ OFF	Clock Is Off 0 – MMC clock has not been turned off 1 – MMC clock has been turned off, due to stop bit in STRP_CLK register Cleared by the MMC_STAT[CLK_EN] bit when the clock is started.
3	STOP_CMD	For stream mode writes. 0 - MMC is not ready for the stop transmission command 1 - MMC is ready for the stop transmission command Cleared when CMD12 is loaded in the MMC_CMD register and the clock is started.
2	END_CMD_R ES	End Command Response 0 – MMC has not received the response 1 – MMC has received the response or a response time-out has occurred Cleared by the MMC_STAT[END_CMD_RES] bit.
1	PRG_DONE	Programming Done 0 – Card has not finished programming and is busy 1 – Card has finished programming and is no longer busy Cleared by the MMC_STAT[PRG_DONE] bit.
0	DATA_TRAN_ DONE	Data Transfer Done 0 - Data transfer is not complete 1 - Data transfer has completed or a read data time-out has occurred Cleared by the MMC_STAT[DATA_TRAN_DONE] bit.



15.5.13 MMC_CMD Register (MMC_CMD)

MMC CMD, shown in Table 15-17, specifies the command number.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-17. MMC_CMD Register

Physical Address MultiMediaCard Controller MMC_CMD Register 0x4110 0030 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 reserved CMD_INDEX **Bits** Name **Description** 31:8 reserved 7 reserved Start bit for command sequence and cannot be changed. 6 reserved Transmission bit in command sequence and cannot be changed. 5:0 CMD_INDEX | Command index (see Table 15-18)

Table 15-18. Command Index Values (Sheet 1 of 3)

CMD	COMM	MODE	ABBREVIATION
000000	CMD0	MMC/SPI	GO_IDLE_STATE
000001	CMD1	MMC/SPI	SEND_OP_COND
000010	CMD2	MMC	ALL_SEND_CID
000011	CMD3	MMC	SET_RELATIVE_ADDR
000100	CMD4	MMC	SET_DSR
000101	CMD5	reserved	
000110	CMD6	reserved	
000111	CMD7	MMC	SELECT/DESELECT_CARD
001000	CMD8	reserved	
001001	CMD9	MMC/SPI	SEND_CSD
001010	CMD10	MMC/SPI	SEND_CID
001011	CMD11	MMC	READ_DAT_UNTIL_STOP
001100	CMD12	MMC	STOP_TRANSMISSION
001101	CMD13	MMC/SPI	SEND_STATUS
001110	CMD14	reserved	
001111	CMD15	MMC	GO_INACTIVE_STATE
010000	CMD16	MMC/SPI	SET_BLOCKLEN
010001	CMD17	MMC/SPI	READ_SINGLE_BLOCK
010010	CMD18	MMC	READ_MULTIPLE_BLOCK
010011	CMD19	reserved	



Table 15-18. Command Index Values (Sheet 2 of 3)

CMD INDEX	COMM	MODE	ABBREVIATION
010100	CMD20	MMC	WRITE_DAT_UNTIL_STOP
010101	CMD21	reserved	
010110	CMD22	reserved	
010111	CMD23	reserved	
011000	CMD24	MMC/SPI	WRITE_BLOCK
011001	CMD25	MMC	WRITE_MULTIPLE_BLOCK
011010	CMD26	MMC	PROGRAM_CID
011011	CMD27	MMC/SPI	PROGRAM_CSD
011100	CMD28	MMC/SPI	SET_WRITE_PROT
011101	CMD29	MMC/SPI	CLR_WRITE_PROT
011110	CMD30	MMC/SPI	SEND_WRITE_PROT
011111	CMD31	reserved	
100000	CMD32	MMC/SPI	TAG_SECTOR_START
100001	CMD33	MMC/SPI	TAG_SECTOR_END
100010	CMD34	MMC/SPI	UNTAG_SECTOR
100011	CMD35	MMC/SPI	TAG_ERASE_GROUP_START
100100	CMD36	MMC/SPI	TAG_ERASE_GROUP_END
100101	CMD37	MMC/SPI	UNTAG_ERASE_GROUP
100110	CMD38	MMC/SPI	ERASE
100111	CMD39	MMC	FAST_IO
101000	CMD40	MMC	GO_IRQ_STATE
101001	CMD41	reserved	
101010	CMD42	MMC/SPI	LOCK_UNLOCK
101011	CMD43	reserved	
101100	CMD44	reserved	
101101	CMD45	reserved	
101110	CMD46	reserved	
101111	CMD47	reserved	
110000	CMD48	reserved	
110001	CMD49	reserved	
110010	CMD50	reserved	
110011	CMD51	reserved	
110100	CMD52	reserved	
110101	CMD53	reserved	
110110	CMD54	reserved	
110111	CMD55	MMC/SPI	APP_CMD
111000	CMD56	MMC/SPI	GEN_CMD
·			· · · · · · · · · · · · · · · · · · ·



Table 15-18. Command Index Values (Sheet 3 of 3)

CMD	COMM	MODE	ABBREVIATION
111001	CMD57	reserved	
111010	CMD58	SPI	READ_OCR
111011	CMD59	SPI	CRC_ON_OFF
111100	CMD60	MMC	reserved for manufacturer
111101	CMD61	MMC	reserved for manufacturer
111110	CMD62	MMC	reserved for manufacturer
111111	CMD63	MMC	reserved for manufacturer

15.5.14 MMC_ARGH Register (MMC_ARGH)

MMC_ARGH, shown in Table 15-19, specifies the upper 16 bits of the argument for the current command.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-19. MMC_ARGH Bit Definitions

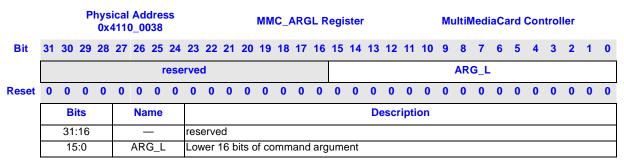
			Pl	hysi 0x4		Ad()_0(SS					ММ	C_A	RG	H R	egi	ster					Μι	ıltiN	ledi	aCa	ırd (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		reserved																						AR	3_H							
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bits Name Description																														
	31:16 — reserved																															
15:0 ARG_H Upper 16 bits of command argument																																

15.5.15 MMC_ARGL Register (MMC_ARGL)

MMC_ARGL, shown in Table 15-20, specifies the lower 16 bits of the argument in the current command.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-20. MMC_ARGL Bit Definitions





15.5.16 **MMC_RES FIFO**

MMC_RES FIFO, shown in Table 15-21, contains the response after a command. It is 16 bits wide by eight entries. The RES FIFO does not contain the 7-bit CRC for the response. The status for CRC checking and response time-out status is in the status register, MMC_STAT.

The first half-word read from the response FIFO is the most significant half-word of the received response.

This is a read-only register. Ignore reads from reserved bits.

Table 15-21. MMC_RES, FIFO Entry

			P	hysi 0x4			dres 03c	SS					MM	C_R	ES	FIF	0 E	ntry	,				М	ultiN	led	iaCa	ard (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		reserved																				RI	ESP	ON	SE_	DA	TA					
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X	X
		В	its			Na	me												De	escr	ipti	on										
	31:16 — reserve																															
	15:0 RESPONSE_DATA Two by											of re	espc	nse	dat	а																

15.5.17 MMC_RXFIFO FIFO

MMC_RXFIFO, shown in Table 15-22, consists of two dual FIFOs, where each FIFO is eight bits wide by 32 entries deep. This FIFO holds the data read from a card. It is a read only FIFO to the software, and is read on 8-bit boundaries. The eight bits of data are read on a 32-bit boundary and occupying the least significant byte lane (7:0).

This is a read-only register. Ignore reads from reserved bits.

Table 15-22. MMC RXFIFO, FIFO Entry

			PI			Ad 0_00		SS					MN	IC_	RXF	IFC	En	try					Mu	ultiN	ledi	iaCa	ard (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•	reserved READ_DATA t 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																															
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	X	X	X	X	X	X	X	X
•		В	ts			Na	me												De	scr	ipti	on										
•	31:8 — reserved																															
7:0 READ_DATA One byte of read data																																



15.5.18 MMC_TXFIFO FIFO

MMC_TXFIFO, shown in Table 15-23, consists of two dual FIFOs, where each FIFO is eight bits wide by 32 entries deep. This FIFO holds the data to be written to a card. It is a write only FIFO to the software, and is written on boundaries eight bits wide. The eight bits of data are written on a 32-bit APB and occupy the least significant byte lane (7:0).

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 15-23. MMC_TXFIFO, FIFO Entry

			P	hysi 0x4		Ad()_ 0 (SS					MN	/IC_	TXF	IFO	En	try					Μι	ıltiN	ledi	aCa	ard (Con	trol	ler		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			reserved WRITE_DATA 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0																													
Reset	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	X	X	X	X	X	X	X	X
		Bits Name Description																														
		31:16 — reserved																														
		7:0 WRITE_DATA One byte of write data																														

15.6 MultiMediaCard Controller Register Summary

Table 15-24 lists the address, name, and description of the MMC Controller Registers.

Table 15-24. MMC Controller Registers (Sheet 1 of 2)

Address	Name	Description
0x4110_0000	MMC_STRPCL	Control to start and stop MMC clock
0x4110_0004	MMC_STAT	MMC status register (read only)
0x4110_0008	MMC_CLKRT	MMC clock rate
0x4110_000c	MMC_SPI	SPI mode control bits
0x4110_0010	MMC_CMDAT	Command/response/data sequence control
0x4110_0014	MMC_RESTO	Expected response time out
0x4110_0018	MMC_RDTO	Expected data read time out
0x4110_001c	MMC_BLKLEN	Block length of data transaction
0x4110_0020	MMC_NOB	Number of blocks, for block mode
0x4110_0024	MMC_PRTBUF	Partial MMC_TXFIFO FIFO written
0x4110_0028	MMC_I_MASK	Interrupt Mask
0x4110_002c	MMC_I_REG	Interrupt Register (read only)
0x4110_0030	MMC_CMD	Index of current command
0x4110_0034	MMC_ARGH	MSW part of the current command argument
0x4110_0038	MMC_ARGL	LSW part of the current command argument



Table 15-24. MMC Controller Registers (Sheet 2 of 2)

Address	Name	Description
0x4110_003c	MMC_RES	Response FIFO (read only)
0x4110_0040	MMC_RXFIFO	Receive FIFO (read only)
0x4110_0044	MMC_TXFIFO	Transmit FIFO (write only)



Network SSP Serial Port

This chapter describes the signal definitions and operation of the Intel® PXA255 Processor Network Synchronous Serial Protocol (NSSP) serial port.

The NSSP is configured differently than the SSPC.

16.1 Overview

The NSSP is a synchronous serial interface that connects to a variety of external analog-to-digital (A/D) converters, telecommunication CODECs, and many other devices that use serial protocols for data transfer. The NSSP provides support for the following protocols:

- Texas Instruments (TI) Synchronous Serial Protocol*
- Motorola Serial Peripheral Interface* (SPI) protocol
- National Semiconductor Microwire*
- Programmable Serial Protocol (PSP)

The NSSP operates as full-duplex devices for the TI Synchronous Serial Protocol*, SPI*, and PSP protocols and as half-duplex devices for the Microwire* protocol.

The FIFOs can be loaded or emptied by the CPU using programmed I/O or DMA burst transfers.

16.2 Features

- Supports the TI Synchronous Serial Protocol*, the Motorola SPI* protocol, National Semiconductor Microwire*, and a Programmable Serial Protocol (PSP)
- Two independent transmit and receive FIFOs, each 16 samples deep by 32-bits wide
- Sample sizes from four to 32-bits
- Maximum bit rate of 13 Mbps in slave of clock mode, requires using DMA
- · Master-mode and slave-mode operation
- Receive-without-transmit operation



16.3 Signal Description

Table 16-1 lists the external signals between the SSP serial ports and external device. If the port is disabled, its pins are available for GPIO use. See Section 4.1, "General-Purpose I/O" for details on configuring pin direction and Section 4.2, "Interrupt Controller" for Interrupt capabilities.

Table 16-1. SSP Serial Port I/O Signals

Name	Direction	Description
NSSPSCLK	Input/Output	NSSPSCLK is the serial bit clock used to control the timing of a transfer. NSSPSCLK is generated internally (master mode) or is supplied externally (slave mode) as indicated by SSCR1[SCLKDIR] as defined in Table 16-4.
NSSPSFRM	Input/Output	NSSPSRM is the serial frame indicator that indicates the beginning and the end of a serialized data word. SSPSRM is generated internally (master mode) or is supplied externally (slave mode) as indicated by SSCR1[SFRMDIR] as defined in Table 16-4.
NSSPTXD	Output	NSSPTXD is the transmit data (serial data out) serialized data line. It is available on two GPIO pins, GPIO[83] or GPIO[84]. See Section 4.1, "General-Purpose I/O" for details.
NSSPRXD	Input	NSSPRXD is the receive data (serial data in) serialized data line. It is available on two GPIO pins, GPIO[83] or GPIO[84]. See Section 4.1, "General-Purpose I/O" for details.

The Network SSP can output either NSSPTXD and NSSPRXD on either GPIO[83] or GPIO[84]. This allows a system to dynamically change the direction of transfer for this port. The NSSP can change direction if enabled, but it must be idle.

16.4 Operation

The SSP controller transfers serial data between the PXA255 processor and an external device through FIFOs. The PXA255 processor CPU initiates the transfers using programmed I/O or DMA bursts to and from memory. Separate transmit and receive FIFOs and serial data paths permit simultaneous transfers in both directions to and from the external device, depending on the protocols chosen.

Programmed I/O transfers data directly between the CPU and the SSP Data Register (SSDR). DMA transfers data between memory and the SSP Data Register (SSDR). Data written to the SSP Data Register (by either the CPU or DMA) is automatically transmitted by the transmit FIFO. Data received by the receive FIFO is automatically sent to the SSP Data Register.

16.4.1 Processor and DMA FIFO Access

The CPU or DMA accesses data through the SSP transmit and receive FIFOs. A CPU access takes the form of programmed I/O, transferring one FIFO entry per access. The FIFO are seen as one 32-bit location by the processor. CPU accesses are normally triggered by an SSSR interrupt and are always 32-bits wide. CPU writes to the FIFOs ignore bits beyond the programmed FIFO data size (EDSS/DSS value); and CPU reads return zeroes in the MSBs down to the programmed data size.



The FIFOs can also be accessed by DMA bursts (in multiples of one, two or four bytes) depending upon the EDSS value. When SSCR0[EDSS] is set, DMA bursts must be in multiples of four bytes (the DMA must have the SSP configured as a 32-bit peripheral). When SSCR0[EDSS] is cleared, DMA bursts must be in multiples of one or two bytes (the DMA's DCMD[WIDTH] register must be at least the SSP data size programmed into the SSCR0[EDSS] and SSCR0[DSS]. If the DMA DCMD[WIDTH] field is configured for 1 byte width, the DMA burst size must be 8 or 16.

For writes, the SSP takes the data from the transmit FIFO, serializes it, and sends it over the serial wire (SSPTXD) to the external device. Receive data from the external device (on SSPRXD) is converted to parallel words and stored in the receive FIFO.

When exceeded, a programmable trigger threshold generates an interrupt or DMA service request that, if SSCR1[TIE] or SSCR1[TSRE] are enabled, signal the CPU or DMA, respectively, to refill the transmit FIFO. Similarly, a programmable trigger threshold generates an interrupt or DMA service request that, if SSCR1[RIE] or SSCR1[RSRE] are enabled, signal the CPU or DMA, respectively, to empty the receive FIFO.

The receive and transmit FIFOs are differentiated by whether the access is a read or a write transfer. Reads automatically target the receive FIFO, while writes write data to the transmit FIFO. From a memory-map perspective, both reads and writes are at the same address. The FIFOs are 16 samples deep by one word wide.

16.4.2 Trailing Bytes in the Receive FIFO

When the number of samples in the receive FIFO is less than the trigger threshold and no additional data is received, the remaining bytes are called trailing bytes. Trailing bytes must be handled by the processor. Trailing bytes are identified via a time-out mechanism and the existence of data within the receive FIFO.

16.4.2.1 Time-out

A time-out condition exists when the receive FIFO is idle for the period of time defined by the Time-Out Register (SSTO). When a time-out occurs, the receiver time-out interrupt (SSSR[TINT]) is set. If the time-out interrupt is enabled (SSCR1[TINTE] set) a time-out interrupt occurs to signal the processor that a time-out condition has occurred. The time-out timer is reset after receiving a new sample or after the processor reads the receive FIFO. Once SSSR[TINT] is set it must be cleared by writing a one to it. If the time-out interrupt is enabled, clearing SSCR1[TINTE] also causes the time-out interrupt to be de-asserted.

16.4.2.2 Removing Trailing Bytes

In this case, no receive DMA service request is generated. To read out the trailing bytes, have the software wait for the time-out interrupt and then read all remaining entries as indicated by SSSR[RFL] and SSSR[RNE].

Note: The time-out interrupt must be enabled by setting SSCR1[TINTE].

16.4.3 Data Formats

Four pins transfer data between the PXA255 processor and external CODECs or modems. Although four serial-data formats exist, each has the same basic structure and in all cases the pins are used as follows:



- SSPSCLK-Defines the bit rate at which serial data is driven onto and sampled from the port.
- SSPSFRM—Defines the boundaries of a basic data unit, comprised of multiple serial bits.
- SSPTXD-The serial data path for transmitted data, from system to peripheral.
- SSPRXD-The serial data path for received data, from peripheral to system.

A data frame can contain from four to 32-bits, depending on the selected protocol. Serial data is transmitted most significant bit first. Four protocols are supported: TI Synchronous Serial Protocol*, SPI, Microwire*, and a PSP.

The SSPSFRM function and use varies between each protocol.

- For the TI Synchronous Serial Protocol*, SSPSFRM is pulsed high for one (serial) data period at the start of each frame. Master and slave modes are supported. TI Synchronous Serial Protocol* is a full-duplex protocol.
- For the SPI* protocol, SSPSFRM functions as a chip select to enable the external device (target of the transfer) and is held active-low during the data transfer (during continuous transfers, the SSPSFRM signal can be either held low or pulsed depending upon the value of SSCR1_x[SPH]). Master and slave modes are supported. SPI* is a full-duplex protocol.
- For the Microwire* protocol, SSPSFRM functions as a chip select to enable the external device (target of the transfer) and is held active-low during the data transfer. Slave mode is not supported for Microwire*. SSPSFRM for Microwire* is also held low during continuous transfers. Microwire* is a half-duplex protocol.
- For the PSP, SSPSFRM is programmable in direction, delay, polarity, and width. Master and slave modes are supported. PSP can be programmed to be either full or half duplex.

The SSPSCLK function and use varies between each protocol.

- For TI Synchronous Serial Protocol*, data sources switch transmit data on the rising edge of SSPSCLK and sample receive data on the falling edge. Master and slave modes are supported.
- For SPI*, the SSP lets programmers select which edge of SSPSCLK to use for switching transmit data and for sampling receive data. In addition, users can move the phase of SSPSCLK, shifting its active state one-half cycle earlier or later at the start and end of a frame. Master and slave modes are supported.
- For Microwire*, both data sources switch (change to the next bit) transmit data on the falling edge of SSPSCLK and sample receive data on the rising edge. Slave mode is not supported for Microwire*.
- For PSP, the protocol allows for the configuration of which edge of the SSPSCLK is used for switching transmit data and the edge for sampling receive data. In addition, the idle state for SSPSCLK can be controlled and the number of active clocks that precede and follow the data transmission. Master and slave modes are supported.

Microwire* uses a half-duplex, master-slave messaging protocol. At the start of a frame, the controller transmits a one or two-byte control message to the peripheral; no data is sent by the peripheral. The peripheral interprets the message and if the message is a read request, the peripheral responds with the requested data, one clock after the last bit of the request message. Return data—part of the same frame—can be from four to 16-bits in length. The total frame length is 13 to 33 bits. The SSPSCLK is active during the entire frame.

Note: The serial clock (SSPSCLK), if driven by the SSP, toggles only while an active data transfer is underway, unless receive-without-transmit mode is enabled by setting SSCR1[RWOT] and the frame format is not Microwire*, in which case the SSPSCLK toggles regardless of whether



transmit data exist within the transmit FIFO. At other times, SSPSCLK holds in an inactive or idle state as defined by the protocol.

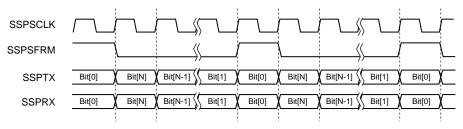
16.4.3.1 TI Synchronous Serial Protocol* Details

When outgoing data in the SSP controller is ready to transmit, SSPSFRM asserts for one clock period. On the following clock, data to be transmitted is driven on SSPTXD one bit at a time, the most significant bit first. For receive data, the peripheral similarly drives data on the SSPRXD pin. The word length can be from four to 32-bits. All output transitions occur on the rising edge of SSPSCLK while data sampling occurs on the falling edge. At the end of the transfer, the SSPTXD signal either retains the value of the last bit sent (LSB) or clears depending on the serial form and the value of the SSPSP[ETDS] (See Figure 16-1 through Figure 16-8). If the SSP is disabled or reset, SSPTXD is forced low.

Figure 16-1 shows the TI Synchronous Serial Protocol for when back-to-back frames are transmitted. Figure 16-2 shows the TI Synchronous Serial Protocol for a single transmitted frame. Once the transmit FIFO contains data, SSPSFRM is pulsed high for one SSPSCLK period and the value to be transmitted is transferred from the transmit FIFO to the transmit logic serial shift register. On the next rising edge of SSPSCLK, the most significant bit of the four to 32-bit data frame is shifted to the SSPTXD pin. Likewise, the MSB of the received data is shifted onto the SSPRXD pin by the off-chip serial slave device. Both the SSP and the off-chip serial slave device then latch each data bit into the serial shifter on the falling edge of each SSPSCLK. The received data is transferred from the serial shifter to the receive FIFO on the first rising edge of SSPSCLK after the last bit has been latched.

For back-to-back transfers, the start of one frame is the completion of the previous frame. The MSB of one transfer immediately follows the LSB of the preceding with no "dead" time between them. When the SSP is a master to the frame sync (SSPSFRM) and a slave to the clock (SSPSCLK), at least three extra clocks are needed at the beginning and end of each block of transfers to synchronize internal control signals (a block of transfers is a group of back-to-back continuous transfers).

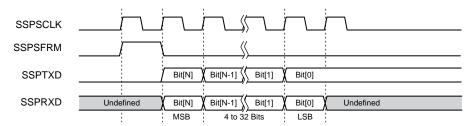
Figure 16-1. Texas Instruments Synchronous Serial Frame* Protocol (multiple transfers)



A9650-01



Figure 16-2. Texas Instruments Synchronous Serial Frame* Protocol (single transfers)



A9518-02

16.4.3.2 SPI Protocol Details

The SPI protocol has four possible sub-modes, depending on the SSPSCLK edges selected for driving data and sampling received data and on the selection of the phase mode of SSPSCLK (see Section 16.4.3.2.1 for complete descriptions of each mode).

When the SSP is disabled or in idle mode, SSPSCLK and SSPTXD are low and SSPSFRM is high. When transmit data is available to send, SSPSFRM goes low (one clock period before the first rising edge of SSPSCLK) and stays low for the remainder of the frame. The most significant bit of the serial data is driven onto SSPTXD one half-cycle later. Halfway into the first bit period, SSPSCLK asserts high and continues toggling for the remaining data bits. Data transitions on the falling edge of SSPSCLK. Four to 32 bits can be transferred per frame.

With the assertion of SSPSFRM, receive data is simultaneously driven from the peripheral on SSPRXD, MSB first. Data transitions on SSPSCLK falling edges and is sampled by the controller on rising edges. At the end of the frame, SSPSFRM is de-asserted high one clock period (one half clock cycle after the last falling edge of SSPSCLK) after the last bit latched at its destination and the completed incoming word is shifted into the incoming FIFO. The peripheral can drive SSPRXD to a high-impedance state after sending the last bit of the frame. SSPTXD retains the last value transmitted when the controller goes into idle mode, unless the SSP is disabled or reset (which forces SSPTXD low).

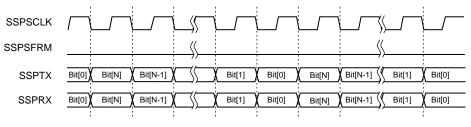
For back-to-back transfers, frames start and complete similar to single transfers, except SSPSFRM does not de-assert between words. Both transmitter and receiver are configured for the word length and internally track the start and end of frames. There are no dead bits; the least significant bit of one frame is followed immediately by the most significant bit of the next.

When using the SPI protocol, the SSP can either be a master or a slave device. However, the clock and frame direction must be the same. For example, the SSCR1[SCLKDIR] and SSCR0[SFRMDIR] must both be set or both be cleared.

Figure 16-3 shows when back-to-back frames are transmitted for the Motorola SPI* frame protocol. Figure 16-4 shows one of the four possible configurations for the Motorola SPI* frame protocol for a single transmitted frame.



Figure 16-3. Motorola SPI* Frame Protocol (multiple transfers)

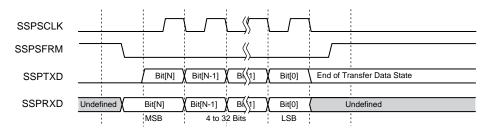


A9651-01

Note: When configured as either master or slave (to clock or frame) the SSP continues to drive SSPTXD with the last bit of data sent (the LSB). If SSCR0[SSE] is cleared, SSPTXD goes low. The state of SSPRXD is undefined before the MSB and after the LSB is transmitted. For minimum power consumption, this pin must not float.

Note: The phase and polarity of SSPSCLK can be configured for four different modes. This example shows just one of those modes (SSCR1[SPO] = 0, SSCR1[SPH] = 0).

Figure 16-4. Motorola SPI* Frame Protocol (single transfers)



A9519-02

Note: When configured as either master or slave (to clock or frame) the SSP continues to drive SSPTXD with the last bit of data sent (the LSB). If SSCR0[SSE] is cleared, SSPTXD goes low. The state of SSPRXD is undefined before the MSB and after the LSB is transmitted. For minimum power consumption, this pin must not float.

16.4.3.2.1 Serial Clock Phase (SPH)

The phase relationship between the SSPSCLK and the serial frame (SSPSFRM) pins when the Motorola SPI* protocol is selected is controlled by SSCR1[SPH].

When SPH is cleared, SSPSCLK remains in its inactive or idle state (as determined by SSCR1[SPO]) for one full cycle after SSPSFRM is asserted low at the beginning of a frame. SSPSCLK continues to transition for the rest of the frame. It is then held in its inactive state for one-half of an SSPSCLK period before SSPSFRM is de-asserted high at the end of the frame.



When SPH is set, SSPSCLK remains in its inactive or idle state (as determined by SSCR1[SPO]) for one-half cycle after SSPSFRM is asserted low at the beginning of a frame. SSPSCLK continues to transition for the remainder of the frame and is then held in its inactive state for one full SSPSCLK period before SSPSFRM is de-asserted high at the end of the frame.

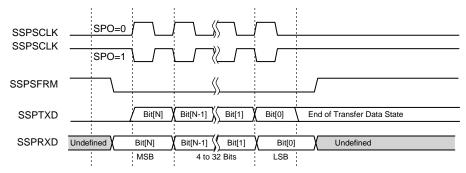
The combination of the SSCR1[SPO] and SSCR1[SPH] settings determine when SSPSCLK is active during the assertion of SSPSFRM and which SSPSCLK edge transmits and receives data on the SSPTXD and SSPRXD pins.

When programming SSCR1[SPO] and SSCR1[SPH] to the same value (both set or both cleared), transmit data is driven on the falling edge of SSPSCLK and receive data is latched on the rising edge of SSPSCLK. When programming SSCR1[SPO] and SSCR1[SPH] to opposite values (one set and the other cleared), transmit data is driven on the rising edge of SSPSCLK and receive data is latched on the falling edge of SSPSCLK.

Note: SSCR1[SPH] is ignored for all data frame formats except for the Motorola SPI* protocol.

Figure 16-6 shows the pin timing for all four programming combinations of SSCR1[SPO] and SSCR1[SPH]. The SSCR1[SPO] inverts the polarity of the SSPSCLK signal and SSCR1[SPH] determines the phase relationship between SSPSCLK and SSPSFRM, shifting the SSPSCLK signal one-half phase to the left or right during the assertion of SSPSFRM.

Figure 16-5. Motorola SPI* Frame Protocols for SPO and SPH Programming (multiple transfers)



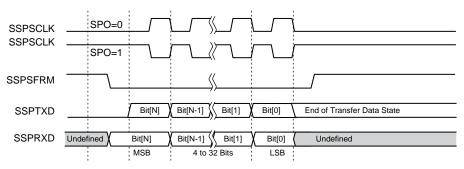
A9652-01

Note: When configured as either master or slave (to clock or frame) the SSP continues to drive SSPTXD with the last bit of data sent (the LSB). If SSCR0[SSE] is cleared, SSPTXD goes low. The state of



SSPRXD is undefined before the MSB and after the LSB is transmitted. For minimum power consumption, this pin must not float.

Figure 16-6. Motorola SPI* Frame Protocols for SPO and SPH Programming (single transfers)



A9520-02

Note: When configured as either master or slave (to clock or frame) the SSP continues to drive SSPTXD with the last bit of data sent (the LSB). If SSCR0[SSE] is cleared, SSPTXD goes low. The state of SSPRXD is undefined before the MSB and after the LSB is transmitted. For minimum power consumption, this pin must not float.

16.4.3.3 Microwire* Protocol Details

The Microwire* protocol is similar to SPI, except transmission is half-duplex instead of full-duplex and it uses master-slave message passing. While in the idle state or when the SSP is disabled, SSPSCLK and SSPTXD are low and SSPSFRM is high.

Each serial transmission begins with SSPSFRM asserting low, followed by an eight or 16-bit command word sent from the controller to the peripheral on SSPTXD. The command word data size is selected by the Microwire* transmit data size bit (SSCR1[MWDS]). SSPRXD is controlled by the peripheral and remains in a high-impedance state. SSPSCLK asserts high midway into the command's most significant bit and continues toggling at the bit rate.

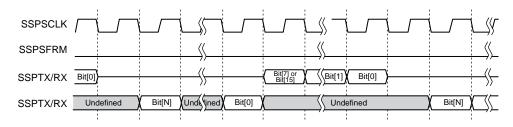
One bit-period after the last command bit, the peripheral returns the serial-data requested most significant bit first on SSPRXD. Data transitions on the falling edge of SSPSCLK and is sampled on the rising edge. The last falling edge of SSPSCLK coincides with the end of the last data bit on SSPRXD and SSPSCLK remains low after that (if it is the only word or the last word of the transfer). SSPSFRM de-asserts high one-half clock period later.

The start and end of a series of back-to-back transfers are like those of a single transfer; however, SSPSFRM remains asserted (low) throughout the transfer. The end of a data word on SSPRXD is followed immediately by the start of the next command byte on SSPTXD with no dead time.

When using the Microwire* protocol, the SSP can function only as a master (frame and clock are outputs). Therefore, both SSCR1[SCLKDIR] and SSCR0[SFRMDIR] must both be cleared. Figure 16-7 shows the National Semiconductor Microwire* frame protocol with eight-bit command words when back-to-back frames are transmitted. Figure 16-8 shows the National Semiconductor Microwire* frame protocol with eight-bit command words for a single transmitted frame.



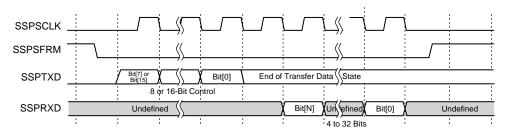
Figure 16-7. National Semiconductor Microwire Frame Protocol (multiple transfers)



A9653-01

Note: When configured master the SSP continues to drive SSPTXD with the last bit of data sent (the LSB) or it drives zero, depending on the status of SSPSP[ETDS]. If SSCR0[SSE] is cleared, SSPTXD goes low. The state of SSPRXD is undefined before the MSB and after the LSB is transmitted. For minimum power consumption, this pin must not float.

Figure 16-8. National Semiconductor Microwire* Frame Protocol (single transfers)



A9521-02

Note: When configured master the SSP continues to drive SSPTXD with the last bit of data sent (the LSB) or it drives zero, depending on the status of SSPSP[ETDS]. If SSCR0[SSE] is cleared, SSPTXD goes low. The state of SSPRXD is undefined before the MSB and after the LSB is transmitted. For minimum power consumption, this pin must not float.

16.4.3.4 **PSP Details**

The PSP provides programmability for several parameters that determine the transfer timings between data.

There are four possible serial clock sub-modes, depending on the SSPSCLK edges selected for driving data and sampling received data and the selection of idle state of the clock.

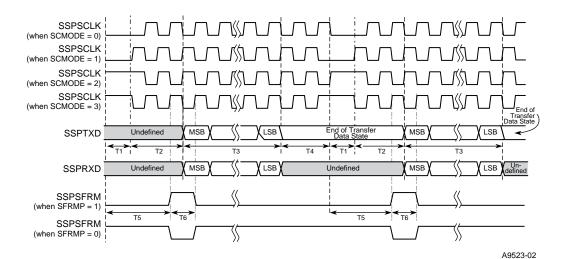
For the PSP, the idle and disable modes of the SSPTXD, SSPSCLK, and SSPSFRM are programmable via SSPSP[ETDS], SSPSP[SCMODE] and SSPSP[SFRMP]. When transmit data is ready, the SSPSCLK remains in its idle state for the number of serial clock (SSPSCLK) clock periods programmed within the start delay (SSPSP[STRTDLY]) field. SSPSCLK then starts toggling, SSPTXD remains in the idle state for the number of cycles programmed within the dummy start field (SSPSP[DMYSTRT]). The SSPSFRM signal asserts after the number of half-



clocks programmed in the field SSPSP[SFRMP]. The SSPSFRM remains asserted for the number of half-clocks programmed within SSPSP[SFRMWDTH]. Four to 32-bits can be transferred per frame. Once the LSB transfers, the SSPSCLK continues toggling based on the dummy stop field (SSPSP[DMYSTOP]). SSPTXD either retains the last value transmitted or is forced to zero, depending on the value programmed within the end of transfer data state field (SSPSP[ETDS]), when the controller goes into idle mode, unless the SSP is disabled or reset (which forces SSPTXD low). Refer to Table 16-2 for more information.

With the assertion of SSPSFRM, receive data is simultaneously driven from the peripheral on SSPRXD, MSB first. Data transitions on SSPSCLK based on the serial clock mode selected and are sampled by the controller on the opposite edge. When the SSP is a master to the frame sync (SSPSFRM) and a slave to the clock (SSPSCLK), at least three extra clocks are needed at the beginning and end of each block of transfers to synchronize internal control signals (a block of transfers is a group of back-to-back continuous transfers).

Figure 16-9. Programmable Serial Protocol (multiple transfers)



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Figure 16-10. Programmable Serial Protocol (single transfers)

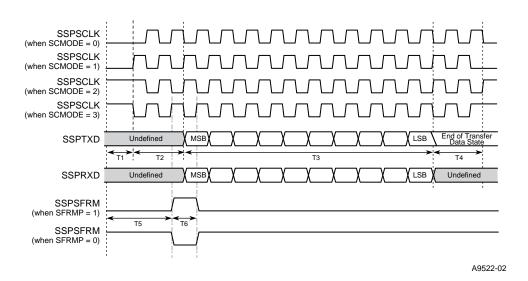


Table 16-2. Programmable Serial Protocol (PSP) Parameters

Symbol	Definition	Range	Units
_	Serial clock mode (SSPSP[SCMODE])	(Drive, Sample, SSPSCLK Idle) 0 - Fall, Rise, Low 1 - Rise, Fall, Low 2 - Rise, Fall, High 3 - Fall, Rise, High	_
_	Serial frame polarity (SSPSP[SFRMP])	High or Low	_
T1	Start delay (SSPSP[STRTDLY])	0 - 7	Clock period
T2	Dummy start (SSPSP[DMYSTRT])	0 - 3	Clock period
Т3	Data size (SSCR0[EDSS] and SSCR0[DSS])	4 - 32	Clock period
T4	Dummy stop (SSPSP[DMYSTOP])	0 - 3	Clock period
T5	SSPSFRM delay (SSPSP[SFRMDLY]	0 - 88	Half clock period
T6	SSPSFRM width (SSPSP[SFRMWDTH]	1 - 44	Clock period
	End of transfer data state (SSPSP[ETDS])	Low or [bit 0]	_

Note: The SSPSFRM delay must not extend beyond the end of T4. SSPSFRM Width must be asserted for at least 1 SSPSCLK, and must be deasserted before the end of the T4 cycle (i.e. in terms of time, not bit values, $(T5 + T6) \le (T1 + T2 + T3 + T4)$, $1 \le T6 < (T2 + T3 + T4)$, and $(T5 + T6) \ge (T1 + T1)$ to ensure that SSPSFRM is asserted for at least 2 edges of the SSPSCLK). While the PSP can be programmed to generate the assertion of SSPSFRM during the middle of the data transfer (after the MSB was sent), the SSP is not able to receive data in frame slave mode (SSCR1[SFRMDIR] is



set) if the assertion of frame is not before the MSB is sent (For example, T5 <= T2 if SSCR1[SFRMDIR] is set). Transmit Data transitions from the "End of Transfer Data State" to the next MSB value upon the assertion of frame. The start delay field should be programmed to 0 whenever SSPSCLK or SSPSFRM is configured as an input.

16.4.4 Hi-Z on SSPTXD

The PXA255 processor NSSP supports placing SSPTXD into Hi-Z during idle times instead of driving SSPTXD.

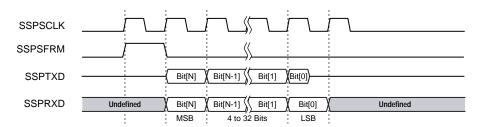
SSCR1[TTE] enables Hi-Z on SSPTXD. SSCR1[TTELP] controls when SSPTXD is placed into Hi-Z.

16.4.4.1 TI Synchronous Serial Port

When SSCR1[TTE] is 0, the SSP behaves as described in Section 16.4.3.1.

If SSCR1[TTE] is 1 and SSCR1[TTELP] is 0, SSPTXD is driven with the MSB at the first rising edge of SSPSCLK after SSPSFRM is asserted. SSPTXD is Hi-Z after the falling edge of SSPSCLK for the LSB (1 clock edge after the clock edge that starts the LSB). Figure 16-11shows the pin timing for this mode.

Figure 16-11. TI SSP with SSCR[TTE]=1 and SSCR[TTELP]=0

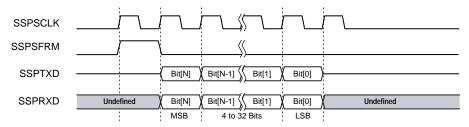


A9974-01

If SSCR1[TTE] is 1 and SSCR1[TTELP] is 1, SSPTXD is driven with the MSB at the first rising edge of SSPSCLK after SSPSFRM is asserted. SSPTXD is Hi-Z at the next rising edge of SSPSCLK after the LSB (2 clock edges after the clock edge that starts the LSB). Figure 16-12 shows the pin timing for this mode.



Figure 16-12. TI SSP with SSCR[TTE]=1 and SSCR[TTELP]=1



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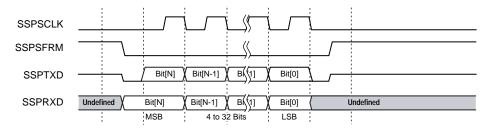
Note: If SSPSCLK is an input, the device driving SSPSCLK must provide another clock edge to cause the TXD line to go to Hi-Z.

16.4.4.2 Motorola SPI

When SSCR1[TTE] is 0, the SSP behaves as described in Section 16.4.3.2.

If SSCR1[TTE] is 1, SSPTXD is driven only when SSPSFRM is 0. When SSPSFRM is 1, SSPTXD is Hi-Z. During the time between the last falling edge and SSPSFRM rising, SSPSP[EDTS] controls the value driven on SSPTXD. Figure 16-13 shows the pin timing for this mode.

Figure 16-13. Motorola SPI with SSCR[TTE]=1



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Note: SSCR1[TTELP] must be 0 for Motorola SPI.

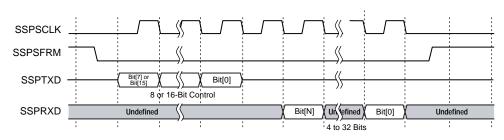
16.4.4.3 National Semiconductor Microwire

When SSCR1[TTE] is 0, the SSP behaves as described in Section 16.4.3.3.

If SSCR1[TTE] is 1, SSPTXD is driven at the same clock edge that the MSB is driven. SSPTXD is Hi-Z after the next rising edge of SSPSCLK for the LSB (1 clock edge after the clock edge that starts the LSB). Figure 16-14 shows the pin timing for this mode.



Figure 16-14. National Semiconductor Microwire with SSCR1[TTE]=1



A9977-01

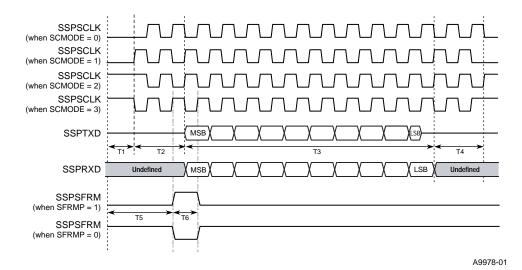
Note: SSCR1[TTELP] must be 0 for National Semiconductor Microwire.

16.4.4.4 Programmable Serial Protocol

When SSCR1[TTE] is 0, the SSP behaves as described in Section 16.4.3.4.

If SSCR1[TTE] is 1 and SSCR1[TTELP] is 0, SSPTXD is driven at the same clock edge that the MSB is driven. If the SSP is a slave to frame SSPTXD is Hi-Z on the clock edge after the edge that starts the LSB. Figure 16-15 shows the pin timing for this mode.

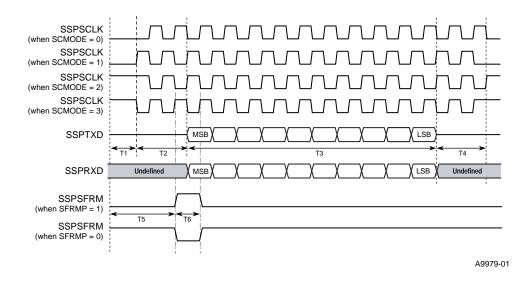
Figure 16-15. PSP mode with SSCR1[TTE]=1 and SSCR1[TTELP]=0 (slave to frame)



If the SSP is a master to frame, SSPTXD is Hi-Z two clock edges after the clock edge that drives the LSB. This occurs even if the SSP is a master of clock and this clock edge does not appear on the SSPSCLK. Figure 16-16 shows the pin timing for this mode.

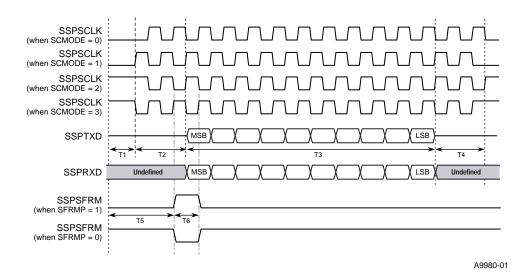


Figure 16-16. PSP mode with SSCR1[TTE]=1 and SSCR1[TTELP]=0 (master to frame)



SSCR1[TTELP] can only be set to 1 in PSP mode if the SSP is a slave to frame. If SSCR1[TTE] is 1 and SSCR1[TTELP] is 1 and the SSP is a slave to frame, SSPTXD is driven at the same clock edge that the MSB is driven. SSPTXD is Hi-Z two clock edges after the clock edge that starts the LSB. This occurs even if the SSP is a master of clock and this clock edge does not appear on the SSPSCLK. If the SSP is a slave of clock, then the device driving SSPSCLK must provide another clock edge. Figure 16-17 shows the pin timing for this mode.

Figure 16-17. PSP mode with SSCR1[TTE]=1 and SSCR1[TTELP]=1 (must be slave to frame)





16.4.5 FIFO Operation

Two separate and independent FIFOs are present for transmit (to peripheral) and receive (from peripheral) serial data. FIFOs are filled or emptied by programmed I/O or DMA bursts.

16.4.5.1 Using Programmed I/O Data Transfers

The PXA255 processor can perform FIFO filling and emptying in response to an interrupt from the FIFO logic. Each FIFO has a programmable trigger threshold at which an interrupt is triggered. When the number of entries in the receive FIFO exceeds the value in SSCR1[RFT], an interrupt is generated (if enabled). This interrupt signals the CPU to empty the receive FIFO. When the number of entries in the transmit FIFO is less than or equal to the value of (SSCR1[TFT] + 1), an interrupt is generated (if enabled). This interrupt signals the CPU to refill the transmit FIFO.

Reading the SSP Status Register (see Section 16.5.3) shows whether the FIFO is full, empty or how many samples it contains.

16.4.5.2 Using DMA Data Transfers

The DMA controller can be programmed to transfer data to and from the SSP FIFOs. To prevent overruns of the transmit FIFO or underruns of the receive FIFO when using the DMA, take care when setting the transmit and receive trigger thresholds.

The programming model for using the DMA is as:

- Program the total number of transmit and receive byte lengths, burst sizes, and peripheral width. Program DCMD[WIDTH] to 0b01 for SSP formats of 8 bits or less; to 0b10 for SSP formats of 9 to 16 bits; to 0b11 for SSP formats of more than 16 bits. When DCMD[WIDTH] is 0b01 (1 byte), then the DMA burst size must be configured for 8 or 16 bytes per burst.
- Set the preferred values in the SSP control registers.
- Set the SSE bit in the SSP Control Register 0 to enable the SSP (see Section 16.5.1).
- Set the run bits in the DMA Command Register.
- Wait for both the DMA transmit and receive interrupt requests.
- If the transmit/receive byte length is not an even multiple of the transfer burst size, a trailingbyte condition may occur as described within Section 16.4.2.
- In full-duplex formats where the SSP always receives the same number of data samples as it transmits, the DMA channel must be set up to transmit and receive the same number of bytes.

16.4.6 Baud-Rate Generation

When the SSP is configured as the master of the SSPSCLK (as determined by SSCR1[SCLKDIR]), the baud rate (or serial bit-rate clock SSPSCLK) is generated internally by dividing the 3.6864 MHz clock by a programmable divider (SSCR0[SCR]).

This generates baud rates up to a maximum of 3.68 Mbits per second. When driven by an external clock, SSPSCLK can be driven up to 13 MHz, generating baud rates up to 13 Mbits per second. At these fast baud rates, using polled/interrupt mode is insufficient to keep the FIFO filled. You must use DMA mode.



16.5 Register Descriptions

Each SSP consists of seven registers: three control, one data, one status, one time-out, and one test.

- The SSP control registers (SSCR0, SSCR1) configure the baud rate, data length, frame format, data-transfer mechanism, and port enabling. They also permit setting the FIFO trigger threshold that triggers an interrupt.
- Access all registers using aligned words.

Note: Write the SSP registers after a reset but before the SSP is enabled.

- The SSP Time-Out (SSTO) register programs the time-out value used to signal a specified period of receive FIFO inactivity.
- While in PSP mode, the SSP Programmable Serial Protocol (SSPSP) register programs the parameters used in defining the data transfer.
- The data register is mapped as one 32-bit location, which physically points to either of two 32-bit registers: one register is for writes of data transfers to the transmit FIFO and the other register is for reads that take data from the receive FIFO. A write cycle or burst write puts successive words into the SSP write register and then into the transmit FIFO. A read cycle or burst read takes data from the SSP read register and the receive FIFO reloads it with available data bits it has stored.

Do not increment the address using read and write DMA bursts.

Besides showing the state of the FIFO buffers, the status register shows whether the
programmable trigger threshold has been passed and whether a transmit or receive FIFO
service request is active. The status register also shows how full the FIFO is. Flag bits indicate
when the SSP is actively transmitting data, when the transmit FIFO is not full, and when the
receive FIFO is not empty. The SSSR[ROR] bit signals an overrun of the receive FIFO In this
case newly received data is discarded.

When programming registers, reserved bits must be written as zeroes and read as undefined.

16.5.1 SSP Control Register 0 (SSCR0)

SSCR0, shown in Table 16-3, contains bit fields that control various functions within the SSP. Before enabling the SSP (via SSE) the desired values for this register must be set.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 16-3. SSCR0 Bit Definitions (Sheet 1 of 2)

		0x4	1140	0_00	00							S	SCF	₹0							N	etw	ork	SSF	o Se	rial	Por	t		
31	30 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			re	serv	ed					EDSS							SCR						SSE	reserved	HAT H			DSS		
?	? ?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0	0	0	0	0	0	0	0	?	0	0	0	0	0	0
	Bits			Nar	ne												De	scri	otior	1										
	31:21			_	-		res	erve	d																					
							EX	ΓΕΝ	DED	DAT	A SIZ	ZE S	SELE	ECT	:															
	20			ED	20		Use	ed in	conj	uncti	on w	ith D	SS	to s	sele	ct th	ne si	ze of	the	data	trar	nsm	ittec	lan	d red	ceiv	ed b	y th	e S	SP.
	20			EDS	55		0 –	Pre	-app	ende	d to	the	DSS	S va	alue.	Se	ts th	ne DS	SS ra	ange	fro	m 4	-16-	bits	i.					
							1 –	Pre	-app	ende	d to	the	DSS	S va	alue.	Se	ts th	ne DS	SS ra	ange	fro	m 1	7-32	-bit	S.					
							THE	E SE	RIA	_ CL(OCK	RA	ГЕ:																	
							by S	SSC	R1[S	oit rat CLK 1 (a	DIR]	. Th	ne m	axi	mun	n bi	t rate	e is 3	.686	64 M	bps	. Th	e cl	ock	is di	vide				
	19:8			SC	R		SS((Ma	CR1	SCL um o	gnor KDIF f 13 I O fill	R]) ar MHz)	nd tr . At	ansı the	mis se f	sion ast l	da bau	ta ra d ra	tes a	ire d	eter	mine	ed b	y th	e è	ctern	al c	lévic		t to	
							NO	TE:		ware ses tl														led I	beca	ause	e doi	ng s	0	
							Ser	ial b	it rat	e = S	SP (Cloc	k / (SCF	₹ +	1)														
							SYI	NCH	RON	IOUS	SEI	RIAL	_ PC	DRT	EN	AΒ	LE/C	DISA	BLE:											
							stop	pec	by p	disa rogra	amm	ers	to m	ninir	nize	ро	wer	cons	ump	tion.			-							
							frar	ne b	eing	d du trans SSF	mitte	ed o	r rec	ceiv	ed.	Cle	aring	g SS												
	7			SS	E		NO	TE:	SS7 regi	r resona, a ster in the contract of the contra	nd S s res -con	SPS et b figu	SP ce efor- ring	onti e re the	rol re e-en: SS	egis abli CR(ters ng th	are ne S SCR	prop SP b 1, or	erly y se SSF	re-c tting PSP	onfi SS	igure SE. <i>F</i>	ed a	nd t , SS	hat E m	the s	SSS be c	R lea	red
							0 –	SS	Рор	eratio	on di	sabl	ed																	
							1 –	SS	Рор	eratio	on er	nable	ed																	
	6			_	-		rese	erve	d																					
							FR	٩ME	FOF	RMAT	:																			
							SEI	EC.	TS w	hich	fram	e fo	rma	t to	use															
	F. 1			ΓР	_		0b0	0 –	Seria	ıl Per	iphe	ral li	nterf	face) *															
	5:4			FR	.Г		0b0	1 –	TI Sy	nchr	onou	s S	erial	Pro	otoc	ol*														
							0b1	0 –	Micro	wire	*																			
							0b1	1 – 1	Prog	ramm	nable	Se	rial F	Prot	toco	I														



Table 16-3. SSCR0 Bit Definitions (Sheet 2 of 2)

		0x4	4140	00_	00							SSC	R0							Ne	etwo	ork :	SSF	Se	rial	Por	t		
31	30 29	28	27	26	25	24	23	22	21	20	19 18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	(
			re	serv	/ed					EDSS						SCR						SSE	reserved	FRF			DSS		
?	? ?	?	?	?	?	?	?	?	?	0	0 0	0	0	0	0	0	0	0	0	0	0	0	?	0	0	0	0	0	(
	Bits			Nar	ne											De	scri	otion											_
							DAT	TA S	IZE S	SELE	CT:																		_
							SSF	P. Th	ne co																				
		Used in conjunction with EDSS to select the size of the data transmitted and received by SSP. The concatenated 5-bit value of EDSS and DSS provides a data range from four to 3 in length. For the Microwire* protocol, DSS and EDSS are used to determine the receive data size. size of the transmitted data is either eight or 16-bits (determined by SSCR1[MWDS]) and EDSS bit is ignored. The EDSS and DSS fields are ignored for Microwire* transmit data s MWDS (alone) configures this. However, for all modes (including Microwire*) EDSS and DSS are used to determine the receive data size. When data is programmed to be less than 32 bits, the FIFO must be programmed right-just.															the size	•											
									ata is		gramme Data S		be				bits DSS		FIFO			oe p	rog	ramı	med	righ	nt-ju	stifi	e
							1		0b00		17-bit				0		0b00		rese			nde	fine	·d					_
							1		0b00		18-bit				0		0b00		rese		•								
							1		0b00		19-bit				0		0b00		rese										
	3:0			DS	SS		1		0b00	11	20-bit	data			0		0b00		4-bi										
							1		0b01	00	21-bit	data			0		0b01	100	5-bi	t da	ta								
							1		0b01	01	22-bit	data			0		0b01	101	6-bi	t da	ta								
							1		0b01	10	23-bit	data			0		0b01	110	7-bi	t da	ta								
							1		0b01	11	24-bit	data			0		0b01	111	8-bi	t da	ta								
							1		0b10	00	25-bit	data			0		0b10	000	9-bi	t da	ta								
							1		0b10	01	26-bit	data			0		0b10	001	10-l	oit d	ata								
							1		0b10	10	27-bit	data			0		0b10	010	11-k	oit d	ata								
							1		0b10	11	28-bit	data			0		0b10	011	12-l	oit d	ata								
							1		0b11	00	29-bit	data			0		0b11	100	13-l	oit d	ata								
							1		0b11	01	30-bit	data			0		0b11	101	14-l	oit d	ata								
							1		0b11	10	31-bit	data			0		0b11	110	15-l	oit d	ata								
							1		0b11	11	32-bit	data			0		0b11	111	16-l	oit d	ata								

16.5.2 SSP Control Register 1 (SSCR1)

SSCR1, shown in Table 16-4, contains bit fields that control various SSP functions. Before enabling the port (using SSCR0[SSE]), the desired values for this register must be set.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 16-4. SSCR1 Bit Definitions (Sheet 1 of 2)

Ta	ble	16	-4.	SS	CR1	Bit	De	fin	itio	ns	(Sh	eet	1 o	of 2)																
				0x0	4140_0)004	ŀ						\$	SSC	R1							Ne	two	ork S	SP	Sei	rial	Por	t		
Bit	31	30	29	28	27 26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	TTELP	TTE	EBCEI	SCFR	reserved	SCLKDIR	SFRMDIR	RWOT	reserved	TSRE	RSRE	TINTE		reserved		STRF	EFWR		RF	т			TF	т		MWDS	SPH	SPO	LBM	Ħ	RIE
Reset	0	0	0	0	? ?	0	0	0	0	0	0	0	?	?	?	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bi	ts			Nam	ie											ſ	Desc	riptio	on										
									TRA	NSN	/IT H	II-Z	_AT	ER I	PHA	SE:	:														
		3	1		ד	ΓΤΕL	_P		Clock This For LSE For edg If S: 0 —	k edge only TI State of State	ge) la v occur SP for reser form ter th CLK i	ater formation and a final section and the LS and the L	than vith t, the the SB (inp Z tin	the is molected SSF the ut, the ning	rmal TIS ean (). P is LSB he c	SSP s the state of	forr e Save ave ores ce d	nat, SPT to frent lrivir	d belo	the F s Hi- , this cloc SPSC ow fo	PSP f Z afte mea k). CLK r r TT	former the ans the muster	nat if ne ris the s	f the sing SSP	SS edg TXI an	P is ge a D is	a s fter Hi-z	lave the Z tw	to LSE o cl	fran 3 (T ock ge.	ne. The
			If SSPSCLK is an input, the device driving 0 - SSPTXD Hi-Z timing is as described 1 - SSPTXD Hi-Z timing is extended by SSP is a slave to frame. TRANSMIT HI-Z ENABLE: This bit controls whether or not SSPTXD is																		,				,						
		30	0			TTE			This For SSF clock For SSF edg For star edg this 0 -	Micro Micro PTXE kk ede SPI, TI SS PSFF e afte PSP ts the clock	contro owire o is H ge th SSP SP fo RM is er the form LSE er the	ols w * SS ii-Z a at st TXE rma ass e clo nat, i ii3. Fo e clo ge do	wheth SPT in the service of the serv	her of the the the the Hi-Z SPT; dans edge ess SP for edge and a drive	or nor nor nor nor nor nor nor nor nor n	iven tt ris 3). enev is dr Hi- t sta a si at if the ear (n at ing ver striver Starts lave LSE con S	SSP with the toff SSF SSP SSP is P is	same le of SPFRM th the father fame le same le	e clock SSPS I is do MSI MSI MIIIng . SSI mas	eass at tedge	ge to form of the following th	that r the ed. First I SSF Hi-Z me,	the Me LSE rising PSCL Z on t	MSI 3 (1 J ec K f he TX	B is clood dge for the sam	driv ck e of S he L ne c s Hi-	en, edge SPS SB	SCL (1 d	er the Karlock	fter k hat
											JNT E		-				-														
		29	9		E	EBCI	ΕI		Disa whe inte 0 –	ables on the rnal b Bit o	bit ce SSI oit co count	ount is a unte	erre a sla er ha or ev	or in ave t as re vents	nterr to cl ach s wi	upts ock ed (Il ge	s. SS or f). ener	SSR ram ate	R will : ie and an inf	d the terru	SSF pt.	de									rs
		28	8		;	SCF	R		SLA SCI sou In m Mas 0 –	TR in rce is naste ster n	slave s runr r moe node ck inp	e moning de (Sont)	REE ode con SCL 7:	E RU (SC tinuo KDI	JNN LKE ousl R cl	IING DIR s y. eare	ed)	mus this	st be bit is busly	clea igno runn	red ing	f the		out clo	ock	fror	m th	e e	xter	nal	

reserved

27:26



Table 16-4. SSCR1 Bit Definitions (Sheet 2 of 2)

la	ble	16-	4. S	S	CR1	Bit	De	tını	itioi	ns (She	et	20	12)																	
			0x	04	140_0	004							5	SCR	1							Ne	etwo	ork	SSF	Se	rial	Por	t		
Bit	31	30 2	9 28	3 2	2 7 26	25	24	23	22	21	20	19	18	17 1	6 ′	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	TTELP	TTE	SCFR		reserved	SCLKDIR	SFRMDIR	RWOT	reserved	TSRE	RSRE	TINTE		reserved		STRF	EFWR		RI	FT			ті	FT		MWDS	SPH	SPO	LBM	TIE	RIE
Reset	0	0	0 0		? ?	0	0	0	0	0	0	0	?	? '	?	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
•		Bits	•	T	ı	lam	е											-	Desc	ripti	on										
		25			sc	CLKE	DIR		Dete 0 –	ermir Mas SSF Slav uses	ter nes were neset we me	heth node .K. ode, odet	ner te, the	TE CL he po port port a ine w SSPF	rt is ge icts nen	s th ner s as	e mate	naste s SS	er or SPSC e, rec	slav CLK i ceive	inter	nally SPS	, ac CLK	cts a	s th m a	e m n ex	aste	er, a	nd d	rive e aı	
		24			SF	RMI	DIR		Dete Whe exte drivi befo to w	ermir en SF ernal ing fr ore as vait a TE:	RMI device ame sserti certa When over Altern	wheth DIR is the must have the mode the	ner t is se Vhei st wa ramou mou e GFI GPI ly, the	TION: the SS et, the the the hit at I e. (No unt of PIO al O dire he GP e port	por eas ex tim terr ctic IO	rt ac t is t th tterr e b nate on t pin ner	cts cor cor ne e nal efo efo is a rate	as the figure as the clock reas n reaso	ne sla ired a valen k cyc ssert on is exan outpu	ave as a at of 1 cles a ting f selenple, t if S	and r slave 10 Sare n rame cted the FRM inter	ecei e to SPS eed e). for t GPI IDIR	the CLI ed,	the fran KS a the port in is	SS ne, after exte , thi an	PSF the e ernal s bit inpu	RM extended ablin dev	signernal ernal ng th vice s pre SFR	nal f dev ie po just ecec MD	rom vice ort nee lenc IR=	eds ce :1.
	RECEIVE WITH OUT TRANSMIT: Puts the SSP into a mode similar to half duplex. This allows the port to receive data transmitting data (half-duplex only). When RWOT is set, the port continues to clock in receive data, regardless of data ex the transmit FIFO. Data is sent/received immediately after the port enable bit (SSCR is set. In this mode, if there is no data to send, the DMA service requests and interruthe transmit FIFO must be disabled (clear TSRE and TIE). If the transmit FIFO is en zeroes are transmitted which must be discarded by the external peripheral. The transmit FIFO underrun condition does not occur when RWOT is set. When RW enabled, SSSR[BUSY] remains active (set to 1) until software clears the RWOT bit. 0 — Transmit/Receive mode. 1 — Receive With Out Transmit mode.														exis CR0 errup em	sting [SS ots f oty,	g in SE]) for all														

16.5.3 SSP Programmable Serial Protocol Register (SSPSP)

reserved

SSPSPx, shown in Table 16-5, contains bit fields used to program the various programmable serial-protocol parameters. The contents of these registers are ignored if the PSP is not selected.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 16-5. SSPSP Bit Definitions (Sheet 1 of 2)

			C)x4	140_	002	C			Ì			S	SPS	SP							Ne	etwo	ork	SSF	Se	rial	Poi	rt		
Bit	31	30 29	9 2	28	27 2	26 2	25 2	24 23	22	21	20	19	18	17	16	15	14 1	3	12	11	10	9	8	7	6	5	4	3	2	1	0
		re	ese	erve	ed			DMYSTOP	reserved		SFI	RMW	/DTI	н				SF	RMI	DLY			DMYSTRT			STRTDLY		ETDS	SFRMP	ACOMO.	
Reset	?	? ?		?	?	? '	?	0 0	?	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bits			ı	lam	ie											De	scrip	otior	1										
		31:25	;			_		res	erve	d																					
		24:23	3		DM	YST	TOF	De las	term t bit	(bit 0	the n	ransı	mitte	ed (S											CLK	is a	ctiv	e fo	llow	ing t	the
																	sig	nal													
		22:16 SFRMWDTH (Start Delay) _{max} + (Dummy start) _{max} + (Data size) _{max} + (Dummy Stop) _{max} . In slave mode (SSCR1[SFRMDIR] set), this field is ignored, however the incoming frame smust be asserted for at least 1 SSPSCLK duration. In PSP mode, the incoming frame signal must be deasserted for at least 1 SSPSCLK afte assertion (before the next sample is transferred). SERIAL FRAME DELAY:															Ū														
		15:9			SF	RMI	DLY	De the	term tran	_ FRA ines i sfer. arts b	the n	umb prog	er o	ed v	/alu	e se	ts th	e n	umb	er of	half	SSI	PSC	LK	cycl						
		8:7			DM	IYST	TRT	De	term	Y ST/ ines ((D) o	the n	umb						es	after	STF	RTDI	_Y tł	nat p	orec	ede	the	traı	nsm	itted	I	
		6:4			ST	RTE	DLY	De dat	term a tra		the n	iumb ne st	er o art c	f SS lelay	SPS(y fie	CLk ld n	nust l	oe p	orogi	ramn	ned t	0 0	whe	enev	er S	SSP	SCL	_K c	r	eer	1
	data transfers. The start delay field must be programmed to 0 whenever SSPSCLK or SSPSFRM is configured as an input (SSCR1[SCLKDIR] = 1 or SSCR1[SFRMDIR] = 1). END OF TRANSFER DATA STATE: Determines the state of SSPTXD at the end of a transfer. When cleared, the state of SSF forced to 0 after the last bit (bit 0) of the frame is sent and remains 0 through the next idle When set, the state of SSPTXD retains the value of the last bit sent (bit 0) through the neriod. 0 - Low 1 - Last Value <bit 0=""></bit>														dle	peri	od.														
		2			S	FRM	ЛΡ	De In (SS 0 -	term dle r SCR SS	FRA ines mode 1[SFI SPSF	the a or w RMD RM i	ctive hen IR] s s act	sta the et), ive I	te of SSF this low.	f the is bit i	disa	bled	, S	SPSI	FRM	is in	its	inac	tive					mod	de	



Table 16-5. SSPSP Bit Definitions (Sheet 2 of 2)

				0>	414	0_0	02C							•	SSP	SP							N	etwo	ork	SSF	P Se	erial	Po	rt		
Bit	31	30	29	28	3 27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
·			re	ser	ved			DMYSTOP		reserved		SFI	RMV	/DT	Н				SF	RMI	DLY			DMYSTRT	5		STRTDLY		ETDS	SFRMP	SCMODE	
Reset																0	0	0	0													
•	? ? ? ? ? ? 0 0 ? 0 0 0 0 0 0 0 0 0 0 0																															
		,	1:0		;	SCM	1OD	E	Selection Select	ects oper SSP\$ 00 - [01 - [0 - [BIT- one ation SCLK Data Data Data Data	of for is si and Drive Drive Drive	ur se mila I on en (F en (F	erial r to whice allir Risin	cloc how th ed ng), I g), I	sk m SS dge: Data Data	ode CR ² s da a Sa a Sa	1[SF amp amp ampl	PO] a s driv led (led (led (and S ven a (Risir Fallir Fallir	SSCI and s ng), I ng), I	R1[Sample dle state dle st	SPH oled Stat Stat Stat	tog e (L e (L e (H	ow) ow) ow) ligh)	er d	•	•		,	e sta	ate

16.5.4 SSP Time Out Register (SSTO)

The SSTO register, shown in Table 16-6, specifies the time-out value used to signal a period of inactivity within the receive FIFO.

This is a read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 16-6. SSTO Bit Definitions

				0X	414	0_0	028								SS	го							N	etw	ork	SSI	P Se	erial	Ро	rt		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
•			r	ese	rve	d													Т	IMEC	DUT											
Reset	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
•		Bi	ts			Na	me												De	scri	otior	1										
•		31:	24			_	_		res	erve	d																					
•									TIM	IEO	JT:																					
		23	:0		Т	IME	OU	Т				o set INT]				t int	erva	ıl. V	/hen	the	TIMI	EOL	JT v	alue	is c	clear	ed,	no t	ime	-out	occ	urs
										time		t inte	rval	is gi	ven	by t	he e	equa	ation	: Tim	ne-ou	ut In	terv	al =	(TIN	ИΕС)UT) / P	erip	hera	l Cl	ock

16.5.5 SSP Interrupt Test Register (SSITR)

SSITR, shown in Table 16-7 on page 16-25, contains bit fields used for testing purposes only.

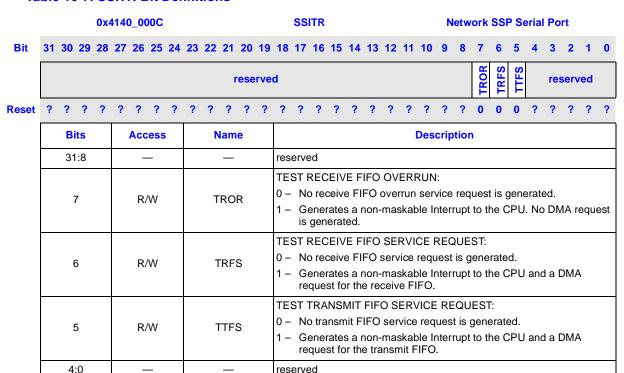
Setting bits in this register causes the SSP controller to generate interrupts and DMA requests if enabled. This is useful in testing the port's functionality.



Setting any of these bits also causes the corresponding status bit(s) to be set in the SSP Status Register (SSSR). The interrupt or service request caused by the setting of one of these bits remains active until the bit is cleared.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 16-7. SSITR Bit Definitions



16.5.6 SSP Status Register (SSSR)

SSSR, shown in Table 16-8 contains bit fields that signal overrun errors and the transmit and receive FIFO service requests. Each of these hardware-detected events signals an interrupt request to the interrupt controller. The status register also contains flags that indicate:

- When the SSP is actively transmitting data
- When the transmit FIFO is not full
- When the receive FIFO is not empty

One interrupt signal is sent to the interrupt controller for each SSP. These events can cause an interrupt:

- Receiver time-out,
- Receive FIFO overrun,
- · Receive FIFO request
- Transmit FIFO request.



Bits that cause an interrupt signal the request as long as the bit is set. The interrupt clears when the bits clear. Read and write bits are called status bits (status bits are referred to as sticky and once set by hardware, they must be cleared by software); Read-only bits are called flags. Writing a 1 to a status bit clears it; writing a 0 has no effect. Read-only flags are set and cleared by hardware; writes have no effect. The reset state of read-write bits is zero and all bits return to their reset state when SSCR0[SSE] is cleared. Additionally, some bits that cause interrupts have corresponding mask bits in the control registers and are indicated in the section headings that follow.

Set the desired values for this register before enabling the SSP (via SSCR0[SSE]).

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 16-8. SSSR Bit Definitions (Sheet 1 of 3)

															- 1																	
				0x4	4140	0_0	800								SSS	SR							N	etw	ork	SSF	Se	rial	Por	t		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
			r	ese	rve	d			BCE	CSS	TUR	reserved	TINIT		reserved			į	KFL			Æ			ROR	RFS	TFS	BSY	RNE	INF	reserved	
Reset	?	?	?	?	?	?	?	?	0	0	0	?	0	?	?	?	1	1	1	1	0	0	0	0	0	0	0	0	0	1	?	?
		В	its			Na	me												De	scrip	otior	1										
•		31	:24			_	_		rese	erve	d																					
									BIT	CO	UNT	ERR	OR:																			

Bits	Name	Description
31:24	_	reserved
		BIT COUNT ERROR:
23	BCE	Indicates that the SSP has detected the SSPSFRM signal asserted at an incorrect time. This bit will cause an interrupt if SSCR1[BCE] is set. The SSP will ignore the current sample and the next sample in order to re-synchronize with the master.
		Write one to clear this bit.
		0 – SSPSFRM has not been asserted out of synchronization.
		1 – SSPSFRM has been asserted out of synchronization.
		CLOCK SYNCHRONIZATION STATUS:
		A read-only bit that indicates the SSP is busy synchronizing the control signals. This bit is only valid when the SSP is a slave to frame.
22	CSS	Software must wait until this bit is a 0 before allowing an external device to assert the SSPSFRM signal.
		0 – The SSP is ready for slave operations.
		1 – The SSP is busy synchronizing slave mode signals.
		TRANSMIT FIFO UNDER RUN:
21	TUR	Indicates that the transmitter tried to send data from the transmit FIFO when the transmit FIFO was empty. When set, an interrupt is generated to the CPU that cannot be locally masked by any SSP register bit. Setting TUR does not generate any DMA service request. To clear TUR, software sets it. TUR remains set until cleared by software writing a one to it which also reset its Interrupt request. Writing a zero to this bit does not affect TUR.
2.		TUR can be set only when the port is a slave to the FRAME signal (SSCR1[SFRMDIR] set) and is not set if the port is in receive-without-transmit mode (SSCR1[RWOT] set).
		Write one to clear this bit.
		0 - Transmit FIFO has not experienced an under run
		1 – Attempted read from the transmit FIFO when the FIFO was empty, request interrupt.
20		reserved



Ta	ble	16-	-8.	SS	SR	Bi	t D)efi	niti	ion	s (S	hee	et 2	of	3)																	
				0x4	4140	_00	800								SS	SSR								N	etw	ork	SSF	° Se	erial	Po	rt	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	7 16	15	14	1 13	1	12	11	10	9	8	7	6	5	4	3	2	1 0
			r	ese	rved	1			BCE	CSS	TUR	eserved	F		reserved				RFL				Ē			ROR	RFS	TFS	BSY	RNE	TNF	reserved
Reset	?	?	?	?	?	?	?	?	0	0	0	?	0	?	?		1	1	1		1	0	0	0	0	0	0	0	0	0	1	? ?
		Bit	S			Nar	me												D	esc	crip	tion										
	-							-	RE	CEI	VER	TIME	Ξ-OL	II TL	NTI	ERRI	JPT	:														
		RECEIVER TIME-OUT INTERRUPT: Indicates that the receive FIFO has been idle (no samples received) for the period of time defined by the value programmed within SSTO. This interrupt can be masked by SSCR1[TINTE]. Write one to clear this bit. 0 - No Receiver Time-out pending 1 - Receiver Time-out pending 18:16 - reserved															е															
	-	18:	Write one to clear this bit. 0 - No Receiver Time-out pending 1 - Receiver Time-out pending																													
	0 - No Receiver Time-out pending 1 - Receiver Time-out pending															fer to																
		11:	.8			TF	FL		Nur	mbei en tl		alid (entrie	es (r	min	nus 1) ad, th	,		,							d pro	ogra	ımm	ners	mus	st re	fer to
		7				RC)R		Indi com disc exis Who regi rese Writ	icate nplet carde sts. en s ister ets it te or	tely fi ed. T set, and bit. T ts into te to	at the illed. This property of the second contract the second cont	e Rec Whe proce errup settin ot rec ar this	ceive en n ess is ot is ng of ques s bit	gelof Rost.	ogic a	a is r ted f ted t does	to to no	eiver all no the Cot ge	d, F ew CPU ene	ROF dat U the erate	R is a real at case any	ass ceiv anno DN	erted red u ot be	d an until e loc servi	nd th at le	ne ne east ma	ewly one	rec e em d by	eive pty any	ed da FIFO	
		6				RF	-S		Indi num thre set, requ	icate nber esho , an uest ries oma	es that of value of the left o	at the alid e is cle rupt is SCR1 the v	e rece entrie eared is ge 1[RS value ared	eive es in d wh nera RE] e of S	e FI hen ate] is SS		equi eive as fe en S Afte [RF] SRE	ires ewe SSC r th T],	IFO er en CR1 ne CI RFS and S	is e trie [RII PU S (a: SSC	equates the sequent of the sequent tended to the sequent tended to the sequent tended to the sequent tended tended to the sequent tended tended to the sequent tended ten	al to nan t s set DMA the s	or (he i . Se rea serv	grea trigg etting ads ice i	ter to ger the g RF the requ	than hres FS s FIF(lest both	the shold signa O su or ir n be	e rec d va als a uch t nteri set	eive lue. a DM that rupt)	FIF Wh IA s it ha	O ti en F ervi	rigger RFS is ce

0 - Receive FIFO level exceeds RFT trigger threshold or the SSP is disabled 1 - Receive FIFO level is at or above RFT trigger threshold, request Interrupt



Ta	Table 16-8. SSSR Bit Definitions (Sheet 3 of 3)																														
	0x4140_0008								SSSR											Ne	Network SSP Serial Port										
Bit	31	30 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
	reserved							BCE	CSS	TUR	reserved	TNI		reserved			į	RFL			Ħ			ROR	RFS	TFS	BSY	RNE	TNF	reserved	1000
Reset	?	? ?	?	?	?	?	?	0	0	0	?	0	?	?	?	1	1	1	1	0	0	0	0	0	0	0	0	0	1	?	?
		Bits			Name				Description																						
		5	TRANSMIT FIFO SERVICE REQUEST: Indicates that the transmit FIFO requires service to prevent an underrun. TFS is set when the number of valid entries in the transmit FIFO is equal to or lesser than the transmit FIFO trigge threshold. It is cleared when it has fewer entries than the trigger threshold value. When TFS is set, an Interrupt is generated when SSCR1[TIE] is set. Setting TFS signals a DMA service request if SSCR1[TSRE] is set. After the CPU or DMA fills the FIFO such that it has at least a many entries as the value of SSCR1[TFT], TFS (and the service request or interrupt) is automatically cleared. SSCR1[TSRE] and SSCR1[TIE] must not both be set. 0 — Transmit FIFO level exceeds TFT trigger threshold or the SSP is disabled 1 — Transmit FIFO level is at or below TFT trigger threshold, request Interrupt															er is													
	4				BS	SY	i 6 6	SSP BUSY: Indicates that the port is actively transmitting or receiving data and is cleared when the port is idle or disabled. This bit does not generate an Interrupt. Software must wait for the Tx Fifo to empty first and then wait for the BSY bit to be cleared at the end of a data transfer. 0 - SSP is idle or disabled 1 - SSP currently transmitting or receiving a frame																							
	3				RI	NE	 	RECEIVE FIFO NOT EMPTY: Indicates that the receive FIFO contains one or more entries of valid data. It is cleared when it no longer contains any valid data. This bit does not generate an Interrupt. When using programmed I/O, this bit can be polled to remove remaining bytes of data from the receive FIFO since CPU Interrupt requests are made only when the receive FIFO trigger threshold has been met or exceeded. 0 — Receive FIFO is empty. 1 — Receive FIFO is not empty.																							
	2				11	NF	I i	TRANSMIT FIFO NOT FULL: Indicates that the transmit FIFO contains one or more entries that do not contain valid data. TNF is cleared when the FIFO is completely full. This bit does not generate an Interrupt. When using programmed I/O, this bit can be polled to fill the transmit FIFO over its trigger threshold.																							

SSP Data Register (SSDR) 16.5.7

1:0

0 - Transmit FIFO is full 1 - Transmit FIFO is not full

reserved

SSDR, shown in Table 16-9, is a single address location that read and write data transfers access. SSDR represents two physical registers: the first is temporary storage for data on its way out through the transmit FIFO. The other register is temporary storage for data coming in through the receive FIFO.



As the system accesses the register, FIFO control logic transfers data automatically between the registers and FIFOs as fast as the system moves it. Unless attempting a write to a full transmit FIFO, data in the FIFO shifts up or down to accommodate the new word(s). Status bits show users whether the FIFO is full, above the programmable trigger threshold, below the programmable trigger threshold or empty.

For transmit data transfers, the register can be written by the system processor anytime it falls below its trigger threshold when using programmed I/O.

When a data size of less than 32-bits is selected, do not left-justify data written to the transmit FIFO. Transmit logic left-justifies the data and ignores any unused bits. Received data of less than 32-bits is automatically right-justified in the receive FIFO.

When the SSP is programmed for the Microwire* protocol and the size of the Transmit data is eight bits (SSCR1[MWDS] cleared), the most significant bits are ignored. Similarly, if the size for the Transmit data is 16 bits (SSCR1[MWDS] set), the most significant 16 bits are ignored. SSCR0[DSS] controls the Receive data size.

Both FIFOs are cleared when the port is reset, or by clearing SSCR0[SSE].

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 16-9. SSDR Bit Definitions

	0x4140_0010									SSDR											Network SSP Serial Port											
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		DATA																														
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?
	Bits Name							Description																								
	31:0					DA	TΑ	TRANSMIT/RECEIVE DATA:																								
									Dat	a wo	ord to	be v	vritte	n to	o/rea	ad fr	om	Tra	nsm	it/rec	eive	FIF	C									

16.6 Network SSP Serial Port Register Summary

Table 16-10 shows the registers associated with the NSSP and their physical addresses.

Table 16-10. NSSP Register Address Map

Physical Address	Name	Description
0x4140_0000	NSSCR0	NSSP Control register 0
0x4140_0004	NSSCR1	NSSP Control register 1
0x4140_0008	NSSSR	NSSP Status register
0x4140_000C	NSSITR	NSSP Interrupt Test register
0x4140_0010	NSSDR	NSSP Data Write Register / Data Read register
0x4140_0028	NSSTO	NSSP Time Out register
0x4140_002C	NSSPSP	NSSP Programmable Serial Protocol





Hardware UART

This chapter describes the signal definitions and operations of the PXA255 processor hardware UART (HWUART) port.

The HWUART interface pins are available via either the PCMCIA general purpose I/O (GPIO) pins or the BTUART pins. Refer to Section 4.1.2, "GPIO Alternate Functions" for more information. When using the HWUART through the PCMCIA pins, they are driven at the same voltage level as the memory interface. Because the PCMCIA signal nPWE is used for variable-latency input / output (VLIO), VLIO cannot be used while the HWUART interface is using the PCMCIA pins.

The HWUART is configured differently than the other UARTs. The HWUART supports full hardware flow control.

17.1 Overview

The HWUART contains a UART and a slow infrared transmit encoder and receive decoder that conforms to the IrDA Serial Infrared (SIR) Physical Layer Link Specification.

The UART performs serial-to-parallel conversion on data characters received from a peripheral device or a modem and parallel-to-serial conversion on data characters received from the processor. The processor can read the UART's complete status during functional operation. Status information includes the type and condition of transfer operations and error conditions (parity, overrun, framing, or break interrupt) associated with the UART.

The HWUART operates in FIFO or non-FIFO mode. In FIFO mode, a 64-byte transmit FIFO holds data from the processor until it is transmitted on the serial link and a 64-byte receive FIFO buffers data from the serial link until it is read by the processor. In non-FIFO mode, the transmit and receive FIFOs are bypassed.

The HWUART can also use direct memory access (DMA) to transfer data to and from the HWUART.

17.2 Features

Software can program interrupts to meet its requirements. This minimizes the number of computations required to handle the communications link. The UART operates in an environment that is controlled by software and can be polled or is interrupt driven. The HWUART has these features:

- Functionally compatible with the 16550A and 16750 UART specifications. The UART supports not only the 16550A and 16750 industry standards but these additional functions as well:
 - DMA requests for transmit and receive data services
 - Slow infrared asynchronous interface



- Non-Return-to-Zero (NRZ) encoding/decoding function
- 64 byte transmit/receive FIFO buffers
- Programmable receive FIFO trigger threshold
- Auto baud-rate detection
- Auto flow
- Maximum baud rate of 921,600 bps.
- Ability to add or delete standard asynchronous communications bits (start, stop, and parity) in the serial data
- Independently controlled transmit, receive, line status, and data set interrupts
- Programmable baud rate generator that allows the internal clock to be divided by 1 to 2¹⁶-1 to generate an internal 16X clock. This clock can be used to drive the internal transmitter and receiver logic.
- Modem control functions clear to send (nCTS) and request to send (nRTS)
- Autoflow capability controls data I/O without generating interrupts:
 - nRTS (output) controlled by UART receiver FIFO
 - nCTS (input) from modem controls UART transmitter
- Fully programmable serial-interface:
 - 5-, 6-, 7-, or 8-bit characters
 - Even, odd, and no parity detection
 - 1, 1.5, or 2 stop bit generation
 - Baud rate generation up to 921 Kbps
 - False start bit detection.
- 64-byte transmit FIFO
- 64-byte receive FIFO
- · Complete status reporting capability
- Ability to generate and detect line breaks
- Internal diagnostic capabilities that include:
 - Loopback controls for communications link fault isolation
 - Break, parity, and framing error simulation
- Fully prioritized interrupt system controls
- Separate DMA requests for transmit and receive data services
- Slow infrared asynchronous interface that conforms to the Infrared Data Association (IrDA) standard



17.3 Signal Descriptions

Table 17-1 lists and describes each external signal that is connected to the UART module. The pins are connected to the PXA255 processor through GPIOs.

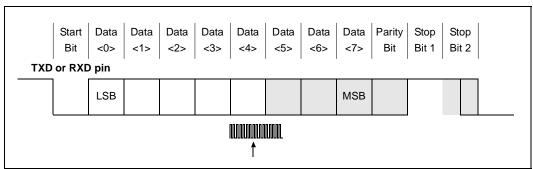
Table 17-1. UART Signal Descriptions

Name	Туре	Description
RXD	Input	SERIAL INPUT – Serial data input to the receive shift register. In infrared mode, it is connected to the infrared receiver input.
TXD	Output	SERIAL OUTPUT – Serial data output to the communications link-peripheral, modem, or data set. The TXD signal is set to the logic 1 state upon a Reset operation. It is connected to the output of the infrared transmitter in infrared mode.
nCTS	Input	CLEAR TO SEND – When low, indicates that the modem or data set is ready to exchange data. Non-Autoflow Mode: When not in autoflow mode, bit 4 (CTS) of the Modem Status register (MSR) indicates the state of nCTS. Bit 4 is the complement of the nCTS signal. Bit 0 (DCTS) of the Modem Status register indicates whether the nCTS input has changed state since the previous reading of the Modem Status register. When the CTS bit of the MSR changes state and the modem status interrupt is enabled, an interrupt is generated. nCTS has no effect on the transmitter. The user can program the UART to interrupt the processor when DCTS changes state. The programmer can then stall the outgoing data stream by starving the transmit FIFO or disabling the UART with the Interrupt Enable register (IER). NOTE: If UART transmission is stalled by disabling the UART, the user will not receive an MSR interrupt when nCTS reasserts. This is because disabling the UART also disables interrupts. To get around this, either use Auto CTS in Autoflow Mode, or program the nCTS GPIO pin to interrupt. Autoflow Mode: In autoflow mode, the UART transmit circuity checks the state of nCTS before transmitting each byte. If nCTS is high, no data is transmitted.
nRTS	Output	REQUEST TO SEND – When low, signals the modem or the data set that the UART is ready to exchange data. Non-Autoflow Mode: The nRTS output signal can be asserted by setting bit 1 (RTS) of the Modem Control register to 1. The RTS bit is the complement of the nRTS signal. Autoflow Mode: nRTS is automatically asserted by the autoflow circuitry when the receive buffer exceeds its programmed trigger threshold. It is deasserted when enough bytes are removed from the buffer to lower the data level back to the trigger threshold.

17.4 Operation

The format of a UART data frame is shown in Figure 17-1.

Figure 17-1. Example UART Data Frame





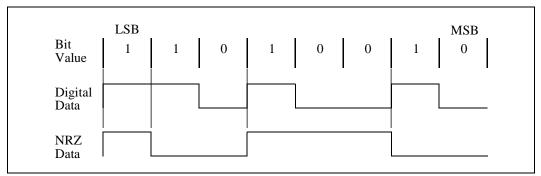
Receive data sample counter frequency is 16 times the value of the bit frequency. The 16X clock is created by the baud rate generator. Each bit is sampled three times in the middle. Shaded bits in Figure 17-1 are optional and can be programmed by software.

Each data frame is between seven and 12 bits long, depending on the size of the data programmed, whether parity is enabled, and the number of stop bits. A data frame begins by transmitting a start bit that is represented by a high to low transition. The start bit is followed by data that is five to eight bits wide and begins with the least significant bit (LSB). The data bits are followed by an optional parity bit. The parity bit is set if even parity is enabled and the data byte has an odd number of ones or if odd parity is enabled and the data byte has an even number of ones. The data frame ends with one, one and a half, or two stop bits, as programmed by software. The stop bits are represented by one, one and a half, or two successive bit periods of a logic one.

The UART has two FIFOs: one transmit and one receive. The transmit FIFO is 64 bytes deep and eight bits wide. The receive FIFO is 64 bytes deep and 11 bits wide. Three bits are used for tracking errors.

The UART can use NRZ coding to represent individual bit values. NRZ coding is enabled when Interrupt Enable register (IER) bit 5 (IER[5]) is set to high. A one is represented by a line transition and a zero is represented by no line transition. Figure 17-2 shows the data byte 0b 0100 1011 in NRZ coding. The LSB is transmitted first.

Figure 17-2. Example NRZ Bit Encoding – (0b0100 1011



17.4.1 Reset

The UART is disabled on reset. To enable the UART, use software to program the GPIO registers then set IER[UUE]. When the UART is enabled, the receiver waits for a frame start bit and the transmitter sends data if it is available in the Transmit Holding register. Transmit data can be written to the Transmit Holding register before the UART is enabled. In FIFO mode, data is transmitted from the FIFO to the Transmit Holding register before it goes to the pin.

When the UART unit is disabled, the transmitter or receiver finishes the current byte and stops transmitting or receiving data. When the UART is enabled, data in the FIFO is not cleared and transmission resumes.

17.4.2 FIFO Operation

The UART has a transmit FIFO and a receive FIFO each holding 64 characters of data. There are three separate methods for moving data into and out of the FIFOs: interrupts, polling, and DMA.



17.4.2.1 FIFO Interrupt Mode Operation

17.4.2.1.1 Receive Interrupt

For a receive interrupt to occur, the receive FIFO and receive interrupts must be enabled. The Interrupt Identification register (IIR) bits 1 and 2 (IIR[IID]) change to show that receive data is available when the FIFO reaches its trigger threshold. IIR[IID] changes to show the next waiting interrupt when the FIFO drops below the trigger threshold. A change in IIR[IID] triggers an interrupt to the core. Software reads IIR[IID] to determine the cause of the interrupt.

The receiver line status interrupt (IIR = 0xC6) has the highest priority and the received data available interrupt (IIR = 0xC4) is lower. The line status interrupt occurs only when the character at the front of the FIFO has errors.

The data ready bit (DR in the Line Status register) is set when a character is transferred from the shift register to the receive FIFO. The DR bit is cleared when the FIFO is empty.

17.4.2.1.2 Character Timeout Interrupt

A character timeout interrupt occurs when the receive FIFO and receive timeout interrupt are enabled and these conditions exist:

- At least one character is in the FIFO.
- The most recently received character was received more than four continuous character times ago. If two stop bits are programmed, the second is included in this interval.
- The most recent FIFO read was performed more than four continuous character times ago.

After the processor reads one character from the receive FIFO or a new start bit is received, the timeout interrupt is cleared and the timeout is reset. If a timeout interrupt has not occurred, the timeout is reset when a new character is received or the processor reads the receive FIFO.

17.4.2.1.3 Transmit Interrupt

Transmit interrupts can only occur when the transmit FIFO and transmit interrupt are enabled. The transmit data request interrupt occurs when the transmit FIFO is at least half empty. The interrupt is cleared when the Transmit Holding register (THR) is written or the IIR is read.

17.4.2.2 FIFO Polled Mode Operation

When the FIFOs are enabled, clearing both IER[DMAE] and IER[4:0] places the serial port in FIFO polled operating mode. The receiver and the transmitter are controlled separately. Either one or both can be in polled mode. In polled mode, software checks receiver and transmitter status via the Line Status register (LSR). The processor polls the following bits for receive and transmit data service:

- **Receive Data Service** The processor checks the data ready (LSR[DR]) bit which is set when one or more bytes remain in the receive FIFO or Receive Buffer register (RBR).
- Transmit Data Service The processor checks the transmit data request (LSR[TDRQ]) bit which is set when transmitter needs data.

The processor can also check the transmitter empty (LSR[TEMT]) bit, which is set when the transmit FIFO and Transmit Holding register are empty.



17.4.2.3 **FIFO DMA Mode Operation**

The UART has two DMA requests: one for transmit data service, and one for receive data service. DMA requests are generated in FIFO mode only. The requests are activated by setting IER[DMAE].

- Data Transmit Data Service When IER[DMAE] is set, if the transmit FIFO is less than half full, the transmit-DMA request is generated. The DMA controller then writes data to the FIFO. For each DMA request, the DMA controller can send 8, 16, or 32 bytes of data to the FIFO. The actual number of bytes to be transmitted is programmed in the DMA controller.
- Data Receive Data Service When IER[DMAE] is set, the receive-DMA request is generated when the receive FIFO reaches its trigger threshold with no errors in its entries. The DMA controller then reads data from the FIFO. For each DMA request, the DMA controller can read 8, 16, or 32 bytes of data from the FIFO. The actual number of bytes to be read is programmed in the DMA controller along with the bus width.

17.4.2.4 **DMA Receive Programming Errors**

If the DMA channel stops prematurely, because it encounters the end of a descriptor chain or other error, the processor must be notified, since the DMA controller can no longer service the UARTs FIFOs. If this occurs, the processor must correct the situation by programming another descriptor or by servicing the FIFOs via interrupt or polling modes as previously described. The DMA must set DCSR[StopIrqEn] to generate an interrupt if a stopped channel occurs.

17.4.2.5 **DMA Error Handling**

An error interrupt is used when DMA requests are enabled. The interrupt is generated when LSR bit 7 is set to 1. This happens when a receive DMA request is not generated and the receive FIFO has an error. The error interrupt tells the processor to handle the data in the receive FIFO through programmed I/O. The error interrupt is enabled when DMA requests are enabled and cannot be masked. Receiver line status interrupts occur when the error is at the front of the FIFO.

When DMA requests are enabled and an interrupt occurs, software must first read the LSR to

verify an error interrupt exists, then check the IIR for the source of the interrupt. If an interrupt occurs and LSR[FIFOE] is clear, software must read the ISR to determine the error condition. When the last error byte is read from the FIFO, DMA requests are automatically enabled. Software is not required to check for the error interrupt if DMA requests are disabled. Error interrupts only occur when DMA requests are enabled.

If an error occurs while in DMA mode:

- receive-DMA requests are disabled
- error interrupt IIR[IID] is generated

The processor must now read the error bytes through programmed I/O. When all errors have been removed from the FIFO, the receive DMA requests are once again enabled automatically by the UART.

If an error occurs when the receive FIFO trigger threshold has been reached, such that a receive DMA request is set, users need to wait for the DMA to finish the transfer before reading out the error bytes through programmed I/O. If not, FIFO underflow could occur.



Note: Ensure that the DMA controller has completed the previous receive DMA requests before the error interrupt handler begins to clear the errors from the FIFO. If not, FIFO underflow could occur.

17.4.2.6 Removing Trailing Bytes In DMA Mode

When the number of entries in the receive FIFO is less than its trigger threshold, and no additional data is received, the remaining bytes are called trailing bytes. The remaining bytes must then be removed via the processor as described in Section 17.4.2.1.

17.4.3 Autoflow Control

Autoflow control uses the clear to send (nCTS) and request to send (nRTS) signals to automatically control the flow of data between the UART and external modem. When autoflow is enabled, the remote device is not allowed to send data unless the UART asserts nRTS low. If the UART deasserts nRTS while the remote device is sending data, the remote device is allowed to send one additional byte after nRTS is deasserted. An overflow could occur if the remote device violates this rule. Likewise, the UART is not allowed to transmit data unless the remote device asserts nCTS low. This feature increases system efficiency and eliminates the possibility of a receive FIFO overflow error due to long interrupt latency.

Autoflow mode can be used in two ways: full autoflow, automating both nCTS and nRTS; and half autoflow, automating only nCTS. Full autoflow is enabled by setting MCR[AFE] and MCR[RTS] to 1. Auto-nCTS-only mode is enabled by setting MCR[AFE] and clearing MCR[RTS].

When in full autoflow mode, nRTS is asserted when the UART FIFO is ready to receive data from the remote transmitter. This occurs when the amount of data in the receive FIFO is below the programmable trigger threshold value. When the amount of data in the receive FIFO reaches the programmable trigger threshold, nRTS is deasserted. It is asserted once again when enough bytes are removed from the FIFO to lower the data level below the trigger threshold.

When in full or half-autoflow mode, nCTS is asserted by the remote receiver when the receiver is ready to receive data from the UART. The UART checks nCTS before sending the next byte of data and will not transmit the byte until nCTS is low. If nCTS goes high while the transfer of a byte is in progress, the transmitter sends the byte.

Note: Autoflow mode can be used only in conjunction with FIFO mode.

17.4.4 Auto-Baud-Rate Detection

The HWUART supports auto-baud-rate detection. When enabled, the UART counts the number of 14.7456 MHz clock cycles within the start-bit pulse. This number is then written into the Auto-Baud-Count register (ACR, see Table 17-13) and is used to calculate the baud rate. When the ACR is written, a auto-baud-lock interrupt is generated (if enabled), and the UART automatically programs the Divisor Latch registers (Section 17.5.3) with the appropriate baud rate. If preferred, the processor can read the ACR and use this information to program the Divisor-Latch registers with a baud rate calculated by the processor. After the baud rate has been programmed, it is the responsibility of the processor to verify that the predetermined characters (usually AT or at) are being received correctly. For the autobaud rate detection circuit to work correctly, the first data bit transmitted after the start bit must be a logic '1'. If a logic '0' is transmitted, the auto-baud circuit counts the zero as part of the start bit, resulting in an incorrect baud rate being programmed into the Divisor Latch Register Low (DLL) and Divisor Latch Register High (DLH) registers.



If the UART is to program the Divisor Latch registers, you can choose one of two methods for auto-baud calculation: table-based and formula-based. Set Auto-Baud Control register (ABR), bit 3 (ABR[ABT]) to select which method you want to use (refer to Section 17.5.8). When the formula method is used, any baud rate defined by the parameters in Section 17.5.3 can be programmed by the UART. The formula method works well for higher baud rates, but could possibly fail below 28.8 kbps if the remote transmitter's actual baud rate differs by more than one percent of its target. The table method is more immune to such errors as the table rejects uncommon baud rates and rounds to the common ones. The table method allows any baud rate defined by the formula in Section 17.5.3 above 28.8 kbps. Below 28.8 kbps the only baud rates which can be programmed by the UART are 19200, 14400, 9600, 4800, 1200, and 300 baud.

When the baud rate is detected, the auto-baud circuitry disables itself by clearing ABR, bit 0 (ABR[ABE]). If users want to re-enable auto-baud detection, ABR[ABE] must be set.

Note: Auto-baud rate detection is not supported with IrDA (slow infrared) mode.

See Section 17.5.8 for more information on auto-baud.

17.4.5 Slow Infrared Asynchronous Interface

The Slow Infrared (SIR) interface supports two-way wireless communication that uses infrared transmission. The SIR provides a transmit encoder and receive decoder to support a physical link that conforms to the Infrared Data Association, *Serial Infrared Physical Layer Link Specification*, October 17, 1995, Version 1.1.

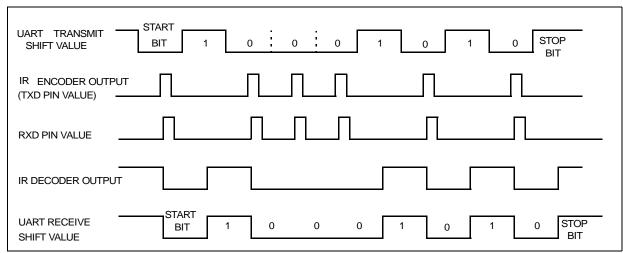
The SIR interface does not contain the actual IR LED driver or the receiver amplifier. The I/O pins attached to the SIR only have digital CMOS level signals. The SIR supports two-way communication, but full duplex communication is not possible because reflections from the transmit LED enter the receiver. The SIR interface supports frequencies up to 115.2 Kbps. Because the input clock is 14.7456 MHz, the baud divisor must be eight or more.

17.4.5.1 Operation

The SIR modulation technique works with 5-, 6-, 7-, or 8-bit characters with an optional parity bit. The data is preceded by a zero value start bit and ends with one or more stop bits. The encoding scheme is to set a pulse 3/16 of a bit wide in the middle of every zero bit and send no pulses for bits that are ones. The pulse for each zero bit must occur, even for consecutive bits with no edge between them.



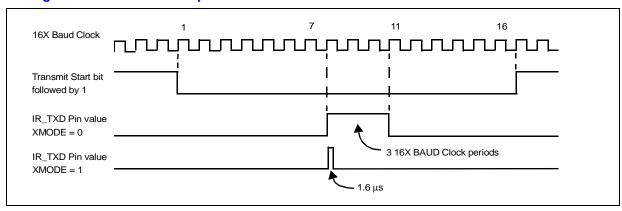




The top line in Figure 17-3 shows an asynchronous transmission as it is sent from the UART. The second line shows the pulses generated by the IR encoder at the TXD pin. A pulse is generated in the middle of the START bit and any data bit that is a zero. The third line shows the values received at the RXD input pin. The fourth line shows the receive decoder's output. The receive decoder drives the receiver data line low when it detects a pulse. The bottom line shows how the UART's receiver interprets the decoder's action. This last line is the same as the first, but it is shifted half a bit period.

When XMODE is cleared, each zero bit has a pulse width of 3/16 of a bit time. When XMODE is set, a pulse of $1.6 \,\mu s$ is generated in the middle of each zero bit. The shorter infrared pulse generated when XMODE is set reduces the LED's power consumption. At 2400 bps, the LED is normally on for $78 \,\mu s$ for each zero bit that is transmitted. When XMODE is set, the LED is on only $1.6 \,\mu s$ (as show in Figure 17-4).

Figure 17-4. XMODE Example.



To prevent transmitter LED reflection feed back to the receiver, disable the IR receiver decoder when the IR transmit encoder transmits data and disables the IR transmit encoder when the IR receiver decoder receives data. The RCVEIR and XMITIR bits in the Infrared Selection Register (ISR) must not be set at the same time (refer to Section 17.5.15).



17.5 Register Descriptions

17.5.1 Receive Buffer Register (RBR)

In non-FIFO mode, the RBR, shown in Table 17-2, holds the character(s) received by the UART's Receive Shift Register. If the RBR is configured to use fewer than eight bits, the bits are right-justified and the most significant bits (MSB) are zeroed. Reading the register empties the register and clears LSR[DR] (refer to Section 17.5.11, "Line Status Register (LSR)" on page 17-19).

In FIFO mode, the RBR latches the value of the data byte at the front of the FIFO (see Table 17-2).

This is a write-only register. Write zeros to reserve bits.

Table 17-2. RBR Bit Definitions

7:0

		0					dres OLA	ss B=0)			Re	ad E	3uff	er F	Regi	ster	(RI	BR)			PX	125	5 Pr	осе	sso	r H	ardv	vare	e UA	ART	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
										r	ese	rve	d														By	e 0				
		Bi	ts			Na	me												De	escr	ipti	on										
		31	:8						rese	erve	d																					

17.5.2 Transmit Holding Register (THR)

Byte 0

Byte 0

In non-FIFO mode, the THR, shown in Table 17-3, holds the next data byte(s) to be transmitted. When the transmit shift register (TSR) is emptied, the contents of the THR are loaded in the TSR and the LSR[TDRQ] is set to 1 (refer to Section 17.5.15).

In FIFO mode, a write to the THR puts data into the end of the FIFO. The data at the front of the FIFO is loaded to the TSR when that register is empty.

This is a write-only register. Write zeros to reserve bits.

Table 17-3. THR Bit Definitions

		(P 0x41	hys 60_)		Ti	rans	mit	Hol	ldin	g Ro	egis	ter	(ТН	R)		PXA	\25	5 Pr	осе	esso	r H	ardv	ware	e UA	RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												rese	rve	d														By	te 0			
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0	0
•		В	its			Na	me												De	escr	ipti	on										
		3′	1:8			_			rese	erve	d																					
		7	:0			Byt	te 0		Byt	e 0																						

17.5.3 Divisor Latch Registers (DLL and DLH)

The HWUART contains a programmable baud rate generator that can take the 14.7456 MHz fixed input clock and divide it by a number that is between 1 and 2^{16} –1. The baud rate generator output frequency is 16 times the baud rate. Two 8-bit latches store the divisor in a 16-bit binary format.



Load these divisor latches during initialization to ensure that the baud rate generator operates properly. If each divisor latch is loaded with a 0, the 16X clock stops. The divisor latches are accessed with a word write.

The baud rate of the data shifted in to or out of a UART is given by the formula:

$$BaudRate = \frac{14.7456 \text{ MHz}}{(16xDivisor)}$$

For example, if the divisor is 24, the baud rate is 38400 bps.

The divisor's reset value is 0x0002.

Table 17-4 and Table 17-5 describe the DLL and DLH registers.

These are read/write registers. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 17-4. DLL Bit Definitions

		0		hysi 60_)		Di	viso	r La	atch	Re	gist	er l	_ow	(DL	L)		PXA	25	5 Pr	осе	SSC	r H	ard	ware	e UA	RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	t														DI	LL			
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	1	0
		Bi	its			Nai	me												De	scr	ipti	on										
		31	:8			_	_		rese	erve	ed																					
		7:	:0			DI	_L		Low	v by	te c	omp	are	valı	ue to	ge	nera	ate l	oaud	d rat	e											

Table 17-5. Divisor Latch Register High (DLH) Bit Definitions

		C		hysi 60_)		Di	viso	r La	itch	Re	gist	er F	ligh	(DL	.H)		PXA	\25	5 Pr	осе	sso	r Ha	ardv	vare	∌ U <i>F</i>	RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											ı	ese	rve	d														DL	.н			
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0	0
		В	ts			Na	me												De	scr	ipti	on										
		31	:8			_	_		res	erve	ed																					
		7	:0			DI	_H		Hig	h by	rte d	comp	oare	val	ue t	o ge	ner	ate l	bau	d ra	te											

17.5.4 Interrupt Enable Register (IER)

The IER, shown in Table 17-6, enables the five types of interrupts that set a value in the Interrupt Identification register (IIR) (refer to Section 17.5.5). To disable an interrupt, software must clear the appropriate bit in the IER. Software can enable some interrupts by setting the appropriate bit.

The character timeout indication interrupt is separated from the received data available interrupt to ensure that the processor and the DMA controller do not service the receive FIFO at the same time. When a character timeout indication interrupt occurs, the processor must handle the data in the receive FIFO through programmed I/O.



Enabling DMA requests also enables a separate error interrupt. For additional information see Section 17.4.2.5.

Set bit 7 of the IER to enable DMA requests. The IER also contains the unit enable and NRZ coding enable control bits. Bits 7 through 4 are used differently from the standard 16550A register definition.

Note: MCR[OUT2] is a global interrupt enable, and must be set to enable UART interrupts.

Note: To ensure that the DMA controller and programmed I/O do not access the same FIFO, software must not set DMAE (bit 7) while TIE (bit 1) or RAVIE (bit 0) are set to 1.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 17-6. IER Bit Definitions

			P	•			dres				1	nter	rup	t En	abl	e Re	egis	ter	(IER	!)		PX/	25	5 Pr	осе	SSO	r Ha	ardv	vare	UA	RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											ı	ese	rve	d											DMAE	UUE	NRZE	RTOIE	MIE	RLSE	TE.	RAVIE
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0	0

Bits	Name	Description
31:8	_	reserved
		DMA REQUESTS ENABLE:
7	DMAE	0 = DMA requests are disabled 1 = DMA requests are enabled
		UART Unit Enable:
6	UUE	0 = the unit is disabled 1 = the unit is enabled
		NRZ CODING ENABLE.
5	NRZE	NRZ encoding/decoding is only used in UART mode, not in infrared mode. If the slow infrared receiver or transmitter is enabled, NRZ coding is disabled.
		0 = NRZ coding disabled 1 = NRZ coding enabled
		RECEIVER TIME OUT INTERRUPT ENABLE (Source IIR[TOD]):
4	RTOIE	0 = Receiver data Time out Interrupt disabled 1 = Receiver data Time out Interrupt enabled
		MODEM INTERRUPT ENABLE (Source IIR[IID]):
3	MIE	0 = Modem Status Interrupt disabled 1 = Modem Status Interrupt enabled
		RECEIVER LINE STATUS INTERRUPT ENABLE (Source IIR[IID]):
2	RLSE	0 = Receiver Line Status Interrupt disabled 1 = Receiver Line Status Interrupt enabled
		TRANSMIT DATA REQUEST INTERRUPT ENABLE (Source IIR[IID]):
1	TIE	0 = Transmit FIFO Data Request Interrupt disabled 1 = Transmit FIFO Data Request Interrupt enabled
		RECEIVER DATA AVAILABLE INTERRUPT ENABLE (Source IIR[IID]):
0	RAVIE	 0 = Receiver Data Available (trigger threshold reached) Interrupt disabled 1 = Receiver Data Available (trigger threshold reached) Interrupt enabled



17.5.5 Interrupt Identification Register (IIR)

The UART prioritizes interrupts in four levels (see Table 17-7) and records them in the IIR. The IIR stores information that indicates that a prioritized interrupt is pending and identifies the source of the interrupt. The Interrupt Identification Register (IIR) bit definitions are shown in Table 17-8.

If additional data is received before a receiver time out interrupt is serviced, the interrupt is deasserted.

Read IIR to determine the type and source of UART interrupts. To be 16550 compatible, the lower 4 bits of the IIR are priority encoded, shown in Table 17-9. If two or more interrupts represented by these bits occur, only the interrupt with the highest priority is displayed. The autobaud lock interrupt is not priority encoded. It asserts/deasserts independently of the lower 4 bits.

IIR[nIP] indicates the existence of an interrupt in the lower four bits of the IIR. A low signal on this bit indicates an encoded interrupt is pending. If this bit is high, no encoded interrupt is pending, regardless of the state of the other 3 bits. nIP has no effect or association with IIR[ABL], which asserts/deasserts independently of nIP.

Table 17-7. Interrupt Conditions

Priority Level	Interrupt origin
1 (highest)	Receiver line status – One or more error bits were set.
2	Received data is available. In FIFO mode, trigger threshold was reached. In non-FIFO mode, RBR has data.
2	Receiver timeout occurred. Occurs only in FIFO mode, when data is in the receive FIFO but no data has been sent for a set time period.
3	Transmitter requests data. In FIFO mode, the transmit FIFO is at least half empty. In non-FIFO mode, the THR has been transmitted.
4 (lowest)	Modem status – One or more modem input signal has changed state.

This is a read-only register. Ignore reads from reserved bits.

Table 17-8. IIR Bit Definitions (Sheet 1 of 2)

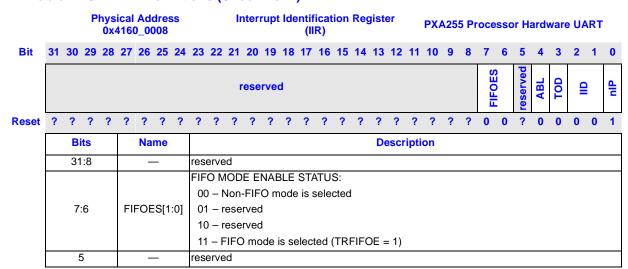




Table 17-8. IIR Bit Definitions (Sheet 2 of 2)

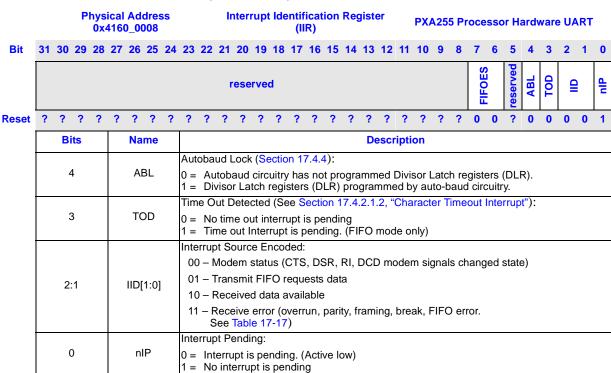


Table 17-9 shows the priority, type, and source of the Interrupt Identification register interrupts. It also gives the reset condition used to deassert the interrupts. Bits (0-3) of the IIR register represent priority encoded interrupts. Bits (4-7) do not.

Table 17-9. Interrupt Identification Register Decode (Sheet 1 of 2)

Interr	upt	: ID	bit	S			Interrupt SET/RESET Fund	etion
	3	2	1	0	Priority	Туре	Source	RESET Control
nIP	0	0	0	1	-	None	No interrupt is pending.	_
IID[11]	0	1	1	0	Highest	Receiver Line Status	Overrun error, parity error, framing error, break interrupt.	Reading the Line Status register.
							Non-FIFO mode – Receive buffer is full.	Non-FIFO mode – Reading the Receiver Buffer register.
IID[10]	0	1	0	0	Second Highest	Received Data Available.	FIFO mode – Trigger threshold was reached.	FIFO mode – Reading bytes until receiver FIFO drops below trigger threshold or setting RESETRF bit in FIFO Control register (FCR).
TOD	1	1	0	0	Second Highest	Character Timeout indication.	FIFO mode only: At least 1 character is left in the receive buffer indicating trailing bytes.	Reading the receiver FIFO or setting RESETRF bit in FCR.
IIDI041	0	0	1	0	Third	Transmit FIFO	Non-FIFO mode: Transmit Holding register empty	Reading the IIR (if the source of the interrupt) or writing into the Transmit Holding register.
IID[01]	U	0	1	0	Highest	Data Request	FIFO mode: transmit FIFO has half or less than half data.	Reading the IIR (if the source of the interrupt) or writing to the transmitter FIFO.



Table 17-9. Interrupt Identification Register Decode (Sheet 2 of 2)

Interr	rupt	ID	bit	S			Interrupt SET/RESET Func	tion
	3	2	1	0	Priority	Туре	Source	RESET Control
IID[00]	0	0	Fourth			Modem Status	Clear to send, data set ready, ring indicator, received line signal detect.	Reading the Modem Status register.
							Non Prioritized Interrupts:	
ABL		4	1		None	Autobaud Lock indication.	Autobaud circuitry has locked onto the baud rate.	Reading the IIR register

17.5.6 FIFO Control Register (FCR)

The FCR, shown in Table 17-10, is a write-only register that is located at the same address as the IIR, which is a read-only register. The FCR enables/disables the transmitter/receiver FIFOs, clears the transmitter/receiver FIFOs, and sets the receiver FIFO trigger threshold.

This is a write-only register. Write zeros to reserved bits.

Table 17-10. FCR Bit Definitions (Sheet 1 of 2)

Physical Address 0x4160_0008 FIFO Control Register (FCR) PXA255 Processor Hardware UART

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

reserved

reserved

Bits	Name	Description
31:8	_	reserved
		Interrupt Trigger Level (threshold) – When the number of bytes in the receiver FIFO equals the interrupt trigger threshold programmed into this field and the received data available interrupt is enabled via the IER, an interrupt is generated and appropriate bits are set in the IIR. The receive DMA request is also generated when the trigger threshold is reached.
7:6	ITL	0b00 – 1 byte or more in FIFO causes interrupt (not valid in DMA mode)
		0b01 – 8 bytes or more in FIFO causes interrupt and DMA request
		0b10 – 16 bytes or more in FIFO causes interrupt and DMA request
		0b11 – 32 bytes or more in FIFO causes interrupt and DMA request
5:4	_	reserved
3	TIL	Transmitter Interrupt Level – Determines when interrupts or DMA requests are sent from the transmit FIFO.
3	IIL	0 = Interrupt/DMA request when FIFO is half empty. 1 = Interrupt/DMA request when FIFO is empty
2	RESETTF	Reset Transmitter FIFO – When RESETTF is set to 1, all the bytes in the transmitter FIFO are cleared. The TDRQ bit in the LSR is set and the IIR shows a transmitter requests data interrupt, if the TIE bit in the IER is set. The Transmitter Shift register is not cleared and it completes the current transmission.
		0 = Writing 0 has no effect 1 = The transmitter FIFO is cleared



Table 17-10. FCR Bit Definitions (Sheet 2 of 2)



Bits	Name	Description
1	RESETRF	Reset Receiver FIFO – When RESETRF is set to 1, all the bytes in the receiver FIFO are cleared. The DR bit in the LSR is reset to 0. All the error bits in the FIFO and the FIFOE bit in the LSR are cleared. Any error bits, OE, PE, FE or BI, that had been set in LSR are still set. The receiver shift register is not cleared. If the IIR had been set to received data available, it is cleared.
		0 = Writing 0 has no effect 1 = The receiver FIFO is cleared
0	TRFIFOE	Transmit and Receive FIFO Enable – TRFIFOE enables/disables the transmitter and receiver FIFOs. When TRFIFOE = 1, both FIFOs are enabled (FIFO Mode). When TRFIFOE = 0, the FIFOs are both disabled (non-FIFO Mode). Writing a 0 to this bit clears all bytes in both FIFOs. When changing from FIFO mode to non-FIFO mode and vice versa, data is automatically cleared from the FIFOs. This bit must be 1 when other bits in this register are written or the other bits are not programmed. 0 = FIFOs are disabled 1 = FIFOs are enabled

17.5.7 Receive FIFO Occupancy Register (FOR)

The FOR, shown in Table 17-11, shows the number of bytes currently remaining in the receive FIFO. It can be used by the processor to determine the number of trailing bytes to remove. The FOR is incremented once for each byte of data written to the receive FIFO and decremented once for each byte read.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 17-11. FOR Bit Definitions

			P	hys 0x4	ical 1160			SS			F	IFO	Occ	cupa	anc	y Re	egis	ter	(FO	R)		PX	\25	5 Pr	OCE	esso	or Ha	ardı	vare	e U <i>A</i>	\RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
		30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 reserved																	Byte	e Co	ount	i										
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0
	Bits Name																		De	esci	ipti	on										
		31	:7			-	_		res	erve	ed																					
		6	:0		B۱	vte	Cou	nt	Nur	mbe	r of	byte	es ((0-64) rei	mair	ning	in t	he r	ece	ver	FIF)									



17.5.8 Auto-Baud Control Register (ABR)

The ABR, shown in Table 17-12, controls the functionality and options for auto-baud-rate detection within the UART. Through this register, users can enable or disable the auto-baud lock interrupt, direct either the processor or UART to program the final baud rate in the Divisor Latch registers, and choose between the two methods used to calculate the final baud rate.

The auto-baud circuitry counts the number of clocks in the start bit and writes this count into the Auto-Baud Count register (ACR – refer to Section 17.5.9). It then interrupts the processor if ABR[ABLIE] is set. It also automatically programs the Divisor Latch registers (DLL and DLH – refer to Section 17.5.3) if ABR[ABUP] bit is set.

See Section 17.4.4 for more information on auto-baud rate.

Note: Auto-baud rate detection is not supported with slow infrared Mode.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 17-12. ABR Bit Definitions

			P				dres	SS			Αι	itob	auc	l Co	ntro	ol R	egis	ster	(AE	BR)		PX	A25	5 Pr	осе	sso	r H	ard	ware	e U	ART	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													ı	rese	rve	d													ABT	ABUP	ABLIE	ABE
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0

Bits	Name	Description
31:4	_	reserved
		AutoBaud Table. Uses a table to select low-speed baud rates instead of a formula. Only valid if ABR[ABUP] is set.
3	ABT	 Table used to calculate baud rates which limits UART to choosing common baud rates as shown in Section 17.4.4, "Auto-Baud-Rate Detection" on page 17-7 Formula used to calculate baud rates allowing all possible baud rates to be chosen by UART.
2	ABUP	AutoBaud UART Program. Controls if the UART automatically programs DLL and DLH or if the processor must program DLL and DLH based on the results in the ACR.
2	ABOF	0 = Processor Programs Divisor Latch registers 1 = UART Programs Divisor Latch registers
		Autobaud Lock Interrupt Enable.
1	ABLIE	0 = Autobaud Lock Interrupt disabled (Source IIR[ABL]) 1 = Autobaud Lock Interrupt enabled (Source IIR[ABL])
		Autobaud Enable.
0	ABE	0 = Autobaud Disabled 1 = Autobaud Enabled

17.5.9 Auto-Baud Count Register (ACR)

The ACR, shown in Table 17-13, stores the number of 14.7456 MHz clock cycles within a start bit pulse. This value is then used by the processor or the UART to calculate the baud rate. If auto-baud mode (ABR[ABE] – Section 17.5.8) and auto-baud interrupts (ABR[ABLIE]) are enabled, the UART interrupts the processor with the auto-baud lock interrupt (IIR[ABL] – Section 17.5.5) after it has written the count value into the ACR. The value is written regardless of the state of the auto-baud UART program bit (ABR[ABUP]).



This is a read-only register. Ignore reads from reserved bits.

Table 17-13. ACR Bit Definitions

			P	hysi 0x4	ical 1160						A	utol	oau	d Co	oun	t Re	gis	ter (ACI	R)		PXA	\25	5 Pr	осе	sso	r Ha	ardv	vare	e UA	RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
·							r	ese	rve	d													Co	unt	Val	ue						
Reset	? ? ? ? ? ? ? ? ?										?	?	?	?	?	?	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		Bits Name																	De	scr	ipti	on										
		31:	:16			-	_		rese	erve	d																					
•		15	5:0			AC	CR		Nur	nbe	r of	14.7	7456	6 MF	lz c	lock	сус	les	with	in a	sta	rt bit	pul	se.								

17.5.10 Line Control Register (LCR)

The LCR., shown in Table 17-14 specifies the format for the asynchronous data communications exchange. The serial data format consists of a start bit, five to eight data bits, an optional parity bit, and one, one and a half, or two stop bits. The LCR has bits that allow access to the divisor latch and bits that can cause a break condition.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.

Table 17-14. LCR Bit Definitions (Sheet 1 of 2)

						Ad 0_00						Lin	e C	onti	rol F	Regi	iste	r (L	CR)			PX/	A25	5 Pı	oce	SSC	or Ha	ardv	ware	U.A	RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											r	ese	rve	d											DLAB	SB	STKYP	EPS	PEN	STB	MLS	
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0	0

Bits	Name	Description
31:8	_	reserved
		DIVISOR LATCH ACCESS BIT
7	DLAB	Must be set to access the divisor latches of the baud rate generator during a READ or WRITE operation. Must be cleared to access the receiver buffer, the Transmit Holding register, or the IER.
		 0 = Access Transmit Holding register (THR), Receive Buffer register (RBR) and IER. 1 = Access Divisor Latch registers (DLL and DLH)
		SET BREAK
6	SB	Causes a break condition to be transmitted to the receiving UART. Acts only on the TXD pin and has no effect on the transmitter logic. In FIFO mode, wait until the transmitter is idle, LSR[TEMT]=1, to set and clear SB.
		0 = No effect on TXD output 1 = Forces TXD output to 0 (space)
		STICKY PARITY
5	STKYP	Forces the bit value at the parity bit location to be the opposite of the EPS bit, rather than the parity value. This stops parity generation. If PEN = 0, STKYP is ignored.
		0 = No effect on parity bit 1 = Forces parity bit to be opposite of EPS bit value



Table 17-14. LCR Bit Definitions (Sheet 2 of 2)

			P		ical 4160							Lin	e C	ont	rol F	Reg	iste	r (L	CR)			PX	125	5 Pr	осе	SSC	or Ha	ardv	vare	e UA	ıRT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
											ı	rese	rve	d											DLAB	SB	STKYP	EPS	PEN	STB	WLS	
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0	0

к	е	s	е	τ

Bits	Name	Description
		EVEN PARITY SELECT
4	EPS	Even parity select bit. If PEN = 0, EPS is ignored.
	21 0	0 = Sends or checks for odd parity 1 = Sends or checks for even parity
		PARITY ENABLE
3	PEN	Enables a parity bit to be generated on transmission or checked on reception.
	1 211	0 = No parity 1 = Parity
		STOP BITS
2	STB	Specifies the number of stop bits transmitted and received in each character. When receiving, the receiver only checks the first stop bit.
		0 = 1 Stop bit 1 = 2 Stop bits, except for 5-bit character then 1-1/2 bits
		WORD LENGTH SELECT
		Specifies the number of data bits in each transmitted or received character.
1:0	WLS[1:0]	00 – 5-bit character
1.0	WLS[1.0]	01 – 6-bit character
		10 – 7-bit character
		11 – 8-bit character

17.5.11 **Line Status Register (LSR)**

The LSR, shown in Table 17-15, provides data transfer status information to the processor.

In non-FIFO mode, LSR[4:2] show the parity error, framing error, break interrupt, and show the error status of the character that has just been received.

In FIFO mode, LSR[4:2] show the status bits of the character that is currently at the front of the FIFO.

LSR[4:1] produce a receiver line status interrupt when the corresponding conditions are detected and the interrupt is enabled. In FIFO mode, the receiver line status interrupt only occurs when the erroneous character reaches the front of the FIFO. If the erroneous character is not at the front of the FIFO, a line status interrupt is generated after the other characters are read and the erroneous character becomes the character at the front of the FIFO.

The LSR must be read before the erroneous character is read. LSR[4:1] bits are set until software reads the LSR.

See Section 17.4.2.3 for details on using the DMA to receive data.

This is a read-only register. Ignore reads from reserved bits.



Table 17-15. LSR Bit Definitions (Sheet 1 of 2)

Physical Address 0x4160_0014 Line Status Register (LSR) PXA255 Processor Hardware UART

Bits	Name	Description
31:8	_	reserved
		FIFO ERROR STATUS
7	FIFOE	In non-FIFO mode, this bit is 0. In FIFO mode, FIFOE is set to 1 when there is at least one parity error, framing error, or break indication for any of the characters in the FIFO. A processor read to the LSR does not reset this bit. FIFOE is reset when all erroneous characters have been read from the FIFO. If DMA requests are enabled (IER bit 7 is set to 1) and FIFOE is set to 1, the error interrupt is generated and no receive DMA request is generated even when the receive FIFO reaches the trigger threshold. Once the errors have been cleared, by reading the FIFO, DMA requests are automatically re-enabled. If DMA requests are not enabled (IER bit 7 is set to 0), then FIFOE set to 1 does not generate an error interrupt. 0 = No FIFO or no errors in receiver FIFO 1 = At least one character in receiver FIFO has errors
		TRANSMITTER EMPTY
6	TEMT	Set when the Transmit Holding register and the Transmitter Shift register are both empty. It is cleared when either the Transmit Holding register or the Transmitter Shift register contains a data character. In FIFO mode, TEMT is set when the transmitter FIFO and the Transmit Shift register are both empty.
		0 = There is data in the Transmit Shift register, the Transmit Holding register, or the FIFO 1 = All the data in the transmitter has been shifted out
		TRANSMIT DATA REQUEST
		Indicates that the UART is ready to accept a new character for transmission. In addition, this bit causes the UART to issue an interrupt to the processor when the transmit data request interrupt enable is set high and generates the DMA request to the DMA controller if DMA requests and FIFO mode are enabled.
5	TDRO	In non-FIFO mode, TDRQ is set when a character is transferred from the Transmit Holding register into the Transmit Shift register. The bit is cleared with the loading of the Transmit Holding register.
5	IDKQ	In FIFO mode, TDRQ is set to 1 when half of the characters in the FIFO have been loaded into the Shift register if FCR[TIL]=0, or the FIFO is empty and FCR[TIL]=1, or the RESETTF bit in FCR has been set. It is cleared when the FIFO has more data than required by FCR[TIL].
		If more than 64 characters are loaded into the FIFO, the excess characters are lost.
		0 = There is data in Transmit Holding register or FIFO waiting to be shifted out 1 = Transmit FIFO has half or less than half data (FCR[TIL]=0), or the transmit FIFO is empty (FCR[TIL]=1), or the UART is waiting for data (non-FIFO mode)
		BREAK INTERRUPT
4	ВІ	BI is set when the received data input is held low for longer than a full word transmission time (that is, the total time of start bit + data bits + parity bit + stop bits). The break indicator is reset when the processor reads the LSR. In FIFO mode, only one character equal to 0x00, is loaded into the FIFO regardless of the length of the break condition. BI shows the break condition for the character at the front of the FIFO, not the most recently received character.
		0 = No break signal has been received 1 = Break signal received



Table 17-15. LSR Bit Definitions (Sheet 2 of 2)

Physical Address 0x4160_0014

Line Status Register (LSR)

PXA255 Processor Hardware UART

Bit 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

reserved

reserved

																									Н	ш	ш					
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	1	1	0	0	0	0	0

Bits	Name	Description
		FRAMING ERROR
3	FE	FE indicates that the received character did not have a valid stop bit. FE is set when the bit following the last data bit or parity bit is detected as 0. If the LCR had been set for two stop bit mode, the receiver does not check for a valid second stop bit. The FE indicator is reset when the processor reads the LSR. The UART re-synchronizes after a framing error. To do this it assumes that the framing error was due to the next start bit, so it samples this "start" bit twice and then reads in the "data". In FIFO mode, FE shows a framing error for the character at the front of the FIFO, not for the most recently received character.
		0 = No framing error 1 = Invalid stop bit has been detected
		PARITY ERROR
2	PE	Indicates that the received data character does not have the correct even or odd parity, as selected by the even parity select bit. PE is set when a parity error is detected and is cleared when the processor reads the LSR. In FIFO mode, PE shows a parity error for the character at the front of the FIFO, not the most recently received character.
		0 = No parity error 1 = Parity error has occurred
		OVERRUN ERROR
1	OE	In non-FIFO mode, OE indicates that data in the Receive Buffer register was not read by the processor before the next character was received. The new character is lost. In FIFO mode, OE indicates that all 64 bytes of the FIFO are full and the most recently received byte has been discarded. The OE indicator is set when an overrun condition is detected and cleared when the processor reads the LSR.
		0 = No data has been lost 1 = Received data has been lost
		DATA READY
0	DR	DR is set when a complete incoming character has been received and transferred into the Receive Buffer register or the FIFO. In non-FIFO mode, DR is cleared when the receive buffer is read. In FIFO mode, DR is cleared if the FIFO is empty (last character has been read from Receive Buffer register) or the FIFO is reset with FCR[RESETRF].
		0 = No data has been received 1 = Data is available in RBR or the FIFO

17.5.12 Modem Control Register (MCR)

The MCR, shown in Table 17-16, uses the modem control pin nRTS to control the interface with a modem or data set. The MCR also controls the loopback mode. Loopback mode must be enabled before the UART is enabled.

This is a read/write register. Ignore reads from reserved bits. Write zeros to reserved bits.



Table 17-16. MCR Bit Definitions (Sheet 1 of 2)

Physical Address 0x4160_0010 Modem Control Register (MCR) PXA255 Processor Hardware UART

eset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	?	0	?
		Bi	ts			Na	me												De	scr	ipti	on										
1		- 0.4	_								.1																					_

Bits	Name	Description
31:8	_	reserved
		AUTOFLOW CONTROL ENABLE
5	AFE	0 = Auto-RTS and auto-CTS are disabled.1 = Auto-CTS is enabled. If MCR[RTS] is also set, both auto-CTS and auto-RTS is enabled.
		LOOPBACK MODE
		This bit provides a local loopback feature for diagnostic testing of the UART. When LOOP is set to a logic 1, the following occurs:
		The transmitter serial output is set to a logic 1 state.
		The receiver serial input is disconnected from the pin.
		 The output of the Transmitter Shift register is "looped back" into the Receiver Shift register input.
		 The four modem control inputs (nCTS, nDSR, nDCD, and nRI) are disconnected from the pins and the modem control output pins (nRTS and nDTR) are forced to their inactive state.
4	LOOP	Coming out of the loopback mode may result in unpredictable activation of the delta bits (bits 3:0) in the Modem Status register (MSR). It is recommended that MSR is read once to clear the delta bits in the MSR.
		Loopback mode must be configured before the UART is enabled.
		MCR[RTS] is connected to the Modem Status register CTS bit: This allows software to test CTS functionality by setting or clearing MCR[RTS]
		RTS = 1 forces CTS to 1
		RTS = 0 forces CTS to a 0
		In loopback mode, data that is transmitted is immediately received. This feature lets the processor verify the transmit and receive data paths of the UART. The transmit, receive and modem control interrupts are operational, except the modem control interrupts are activated by MCR bits, not the modem control pins. A break signal can also be transferred from the transmitter section to the receiver section in loopback mode.
		0 = Normal UART operation 1 = Loopback mode UART operation
		OUT2 SIGNAL CONTROL
		OUT2 connects the UART's interrupt output to the interrupt controller unit. When LOOP=0:
3	OUT2	0 = UART interrupt is disabled 1 = UART interrupt is enabled.
		When LOOP=1, interrupts always go to the processor.
2	_	reserved



Table 17-16. MCR Bit Definitions (Sheet 2 of 2)

											•					•																
			P	•	ical 1160			S	Modem Control Register (MCR)												PX	\25	5 Pr	осе	ssc	or Ha	ardv	ware	e UA	RT		
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
												r	ese	rved	d												AFE	LOOP	OUT2	reserved	RTS	reserved
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	?	0	?
	Bits																		De	esci	ipti	on										
			1			RT	ΓS		Cor bet Aut 0 = 1 = Aut 0 = 1 =	oflow oflow nR nR oflow Au	n ful w mo TS p TS p w mo to-R to-R	e sta I au ode oin i ode TS TS	atus toflo disa s 1 s 0 ena disa ena	of tow a able able able able able able able ab	ind d: d: d. A	uto uto	auto flow flow	wo wo	rks rks v	only with	wit bot	∃ bit h au h au gnal	to-C	CTS					set,	swit	che	Ø
	0 — reserved																															

17.5.13 Modem Status Register (MSR)

The MSR, shown in Table 17-17, provides the current state of the control lines from the modem or data set (or a peripheral device emulating a modem) to the processor. In addition to this current state information, four bits of the MSR provide change information. MSR[3:0] are set when a control input from the modem changes state. They are cleared when the processor reads the MSR.

The status of the modem control lines do not affect the FIFOs. To use these lines for flow control, IER[MIE] must be set. When an interrupt on one of the flow control pins occurs, the interrupt service routine must disable the UART. The UART continues transmission and reception of the current character and then stops. The contents of the FIFOs is preserved. If the UART is reenabled, transmission continues where it stopped.

Note: When bit 0, 1, 2, or 3 is set, a modem status interrupt is generated if IER[MIE] is set.

This is a read-only register. Ignore reads from reserved bits.

Table 17-17. MSR Bit Definitions (Sheet 1 of 2)

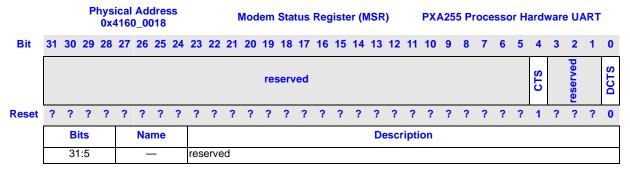




Table 17-17. MSR Bit Definitions (Sheet 2 of 2)

											•					•																
			P	hys 0x4		Ad 0_00		SS		Modem Status Register (MSR) 23 22 21 20 19 18 17 16 15 14 13 12 11												PX	A25	5 Pr	осе	SSO	r Ha	ardv	vare	e UA	RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	serv	/ed													CTS		reserved		DCTS
Reset	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	?	1	?	?	?	0
•	Bits Name					Description																										
									CLEAR TO SEND																							
			4			C ⁻	ΓS		Complement of the clear to send (nCTS) input. Equivalent to MCR[RTS] if MCR[LOOP] is set.													is										
									0 = 1 =	nC nC		pin pin																				
•	3:1 — re						reserved																									
İ	0					DC	TS		0 =	LTA No nC	cha	ange	e in	nC ⁻	ГS р			last	t rea	ad o	f MS	SR										

17.5.14 Scratchpad Register (SCR)

The SCR, shown in Table 17-18, has no effect on the UART. It is intended as a scratchpad register for use by the programmer. It is included for 16550A compatibility.

This is a read-only register. Ignore reads from reserved bits.

Table 17-18, SCR Bit Definitions

			P				dres	SS				Sc	rato	hpa	ıd R	egi	ster	(SC	CR)			PX	125	5 Pr	осе	ssc	or H	ardv	vare	e UA	RT	
Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
							reserved																		SC	CR						
Reset	? ? ? ? ? ? ? ? ? ?							?	?	?	?	?	?	?	?	?	?	?	?	?	?	0	0	0	0	0	0	0	0			
	Bits			Name														De	escr	ipti	on											
	31:8					_	_		rese	erve	d																					
	7:0			SC	CR		No effect on UART fu				fun	ctio	n																			

17.5.15 Infrared Selection Register (ISR)

Each UART can manage an IrDA module. The ISR, shown in Table 17-19, controls IrDA functions (see Section 17.4.5).

This is a read/write register. Ignore reads to reserved bits. Write zeros to reserved bits.



Table 17-19. ISR Bit Definitions

Physical Address 0x4160_0020 Infrared Selection Register (ISR) PXA255 P

PXA255 Processor Hardware UART

Bit	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
													res	serv	ed													RXPL	TXPL	XMODE	RCVEIR	XMITIR

Bits	Name	Description
31:5	_	reserved
		RECEIVE DATA POLARITY
4	RXPL	0 = SIR decoder takes positive pulses as zeros 1 = SIR decoder takes negative pulses as zeros
		TRANSMIT DATA POLARITY
3	TXPL	 0 = SIR encoder generates a positive pulse for a data bit of zero 1 = SIR encoder generates a negative pulse for a data bit of zero
		TRANSMIT PULSE WIDTH SELECT
2	XMODE	When XMODE is cleared, the UART 16X clock clocks the IrDA transmit and receive logic. When XMODE is set, receive decoder operation does not change and the transmit encoder generates 1.6 μs pulses (that are 3/16 of a bit time at 115.2 Kbps) instead of pulses 3/16 of a bit time wide.
		0 = Transmit pulse width is 3/16 of a bit time wide 1 = Transmit pulse width is 1.6 μs
		RECEIVER SIR ENABLE
1	RCVEIR	When RCVEIR is set, the signal from the RXD pin is processed by the IrDA decoder before it is fed to the UART. If RCVEIR is cleared, then all clocking to the IrDA decoder is blocked and the RXD pin is fed directly to the UART.
		0 = Receiver is in UART mode 1 = Receiver is in infrared mode
		TRANSMITTER SIR ENABLE
		When XMITIR is set to 1, the normal TXD output from the UART is processed by the IrDA encoder before it is fed to the device pin. If XMITIR is cleared, all clocking to the IrDA encoder is blocked and the UART's TXD signal is connected directly to the device pin.
0	XMITIR	When Transmitter SIR Enable is set, the TXD output pin, which is in a normally high default state, switches to a normally low default state. This can cause a false start bit unless the infrared LED is disabled before XMITIR is set.
		0 = Transmitter is in UART mode 1 = Transmitter is in infrared mode

17.6 Hardware UART Register Summary

Table 17-20 contains the register addresses for the HWUART.

Table 17-20. HWUART Register Locations (Sheet 1 of 2)

Register Addresses	DLAB Bit Value	Name	Description
0x4160_0000	0	HWRBR	"Receive Buffer Register (RBR)" (read only)
0x4160_0000	0	HWTHR	"Transmit Holding Register (THR)" (write only)
0x4160_0004	0	HWIER	"Interrupt Enable Register (IER)" (read/write)



Table 17-20. HWUART Register Locations (Sheet 2 of 2)

Register Addresses	DLAB Bit Value	Name	Description
0x4160_0008	Х	HWIIR	"Interrupt Identification Register (IIR)" (read only)
0x4160_0008	Х	HWFCR	"FIFO Control Register (FCR)" (write only)
0x4160_000C	Х	HWLCR	"Line Control Register (LCR)" (read/write)
0x4160_0010	Х	HWMCR	"Modem Control Register (MCR)" (read/write)
0x4160_0014	Х	HWLSR	"Line Status Register (LSR)" (read only)
0x4160_0018	Х	HWMSR	"Modem Status Register (MSR)" (read only)
0x4160_001C	Х	HWSPR	"Scratchpad Register (SCR)" (read/write)
0x4160_0020	Х	HWISR	"Infrared Selection Register (ISR)" (read/write)
0x4160_0024	Х	HWFOR	"Receive FIFO Occupancy Register (FOR)" (read only)
0x4160_0028	Х	HWABR	"Auto-Baud Control Register (ABR)" (read/write)
0x4160_002C	Х	HWACR	"Auto-Baud Count Register (ACR)"
0x4160_0000	1	HWDLL	"Divisor Latch Registers (DLL and DLH)" low byte (read/write)
0x4160_0004	1	HWDLH	"Divisor Latch Registers (DLL and DLH)" high byte (read/write)



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